# Coinduction - An Introductory Example

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based on: Jacobs and Rutten, A Tutorial on (Co)Algebras and (Co)Induction. EATCS Bulletin 62, 1997, p.222-259

#### coinduction

How to reason about infinite data/processes? An example:

Given 
$$zip: A^{\mathbb{N}} \times A^{\mathbb{N}} \to A^{\mathbb{N}}$$
,  $even: A^{\mathbb{N}} \to A^{\mathbb{N}}$ ,  $odd: A^{\mathbb{N}} \to A^{\mathbb{N}}$ 

such that

$$head(zip(l_1, l_2)) = head(l_1)$$

$$tail(zip(l_1, l_2)) = zip(l_2, tail(l_1))$$

$$head(even(l)) = head(l)$$

$$tail(even(l)) = even(tail(tail(l)))$$

$$odd(l) = even(tail(l))$$

$$zip(even(x), odd(x)) = x$$

Given

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$$tail(zip(l_1, l_2)) = zip(l_2, tail(l_1))$$

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$$zip(even(x), odd(x))) = x$$

$$head(zip(even(x), odd(x))) = head(even(x)) = head(x)$$

Given

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head(zip(l_1, l_2)) = head(l_1)
tail(zip(l_1, l_2)) = zip(l_2, tail(l_1))
head(even(l)) = head(l)
tail(even(l)) = even(tail(tail(l)))
odd(l) = even(tail(l))
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zip(even(x), odd(x))) = x
```

$$tail(zip(even(x), odd(x))) = zip(odd(x), tail(even(x)))$$
  
=  $zip(odd(x), even(tail(tail(x))))$   
=  $zip(even(tail(x)), odd(tail(x))) = tail(x)$ 

```
head(zip(l_1, l_2)) = head(l_1)
tail(zip(l_1, l_2)) = zip(l_2, tail(l_1))
head(even(l)) = head(l)
tail(even(l)) = even(tail(tail(l)))
odd(l) = even(tail(l))
show
zip(even(x), odd(x))) = x
```

```
tail(zip(even(x), odd(x))) = zip(odd(x), tail(even(x)))
= zip(odd(x), even(tail(tail(x))))
= zip(even(tail(x)), odd(tail(x))) = tail(x)
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head(zip(l_1, l_2)) = head(l_1)
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tail(zip(even(x), odd(x))) = zip(odd(x), tail(even(x)))
= zip(odd(x), even(tail(tail(x))))
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tail(zip(even(x), odd(x))) = zip(odd(x), tail(even(x)))
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```
tail(zip(even(x), odd(x))) = zip(odd(x), tail(even(x)))
= zip(odd(x), even(tail(tail(x))))
= zip(even(tail(x)), odd(tail(x))) = tail(x)
```

# explanation (bisimulation)

**Coinduction Proof Principle:** Two streams  $x, x' \in A^{\mathbb{N}}$  are equal iff there is a relation R with xRx' and for all y, y'

$$yRy' \Rightarrow head(y) = head(y')$$
  
 $yRy' \Rightarrow tail(y) R tail(y')$ 

**Example:** Put zip(even(x), odd(x))) R x for all  $x \in A^{\mathbb{N}}$ 

```
tail(zip(even(x), odd(x))) = zip(odd(x), tail(even(x)))
= zip(odd(x), even(tail(tail(x))))
= zip(even(tail(x)), odd(tail(x))) R tail(x)
```

# explanation (coinduction via finality)

**Def:** An object Z is **final** if for all objects X there is a unique arrow  $X \to Z$ .

**Observation:**  $A^{\mathbb{N}} \to A \times A^{\mathbb{N}}$  is the final coalgebra (for the functor  $TX = A \times X$ ).

**Fact:** To say that  $X \to A \times X$  satisfies the coinduction proof principle is equivalent to saying that  $X \to A \times X$  is the final coalgebra.

# why category theory matters

- ... simple elegant definitions (via universal properties, eg finality)
- ... category theoretic definitions are more general
- ...the right level of abstraction for many proofs (eg Lambek's lemma, Birkhoff's variety theorem)
- ... solution of domain equations
- ...duality
- ... heuristics for finding meaningful mathematical constructions
- ...CT often codes up a lot of annoying combinatorics (eg: GSOS rule format is equivalent to the naturality of a transformation between two functors)