

High Throughput Secure MPC Over Small Population in Hybrid Networks

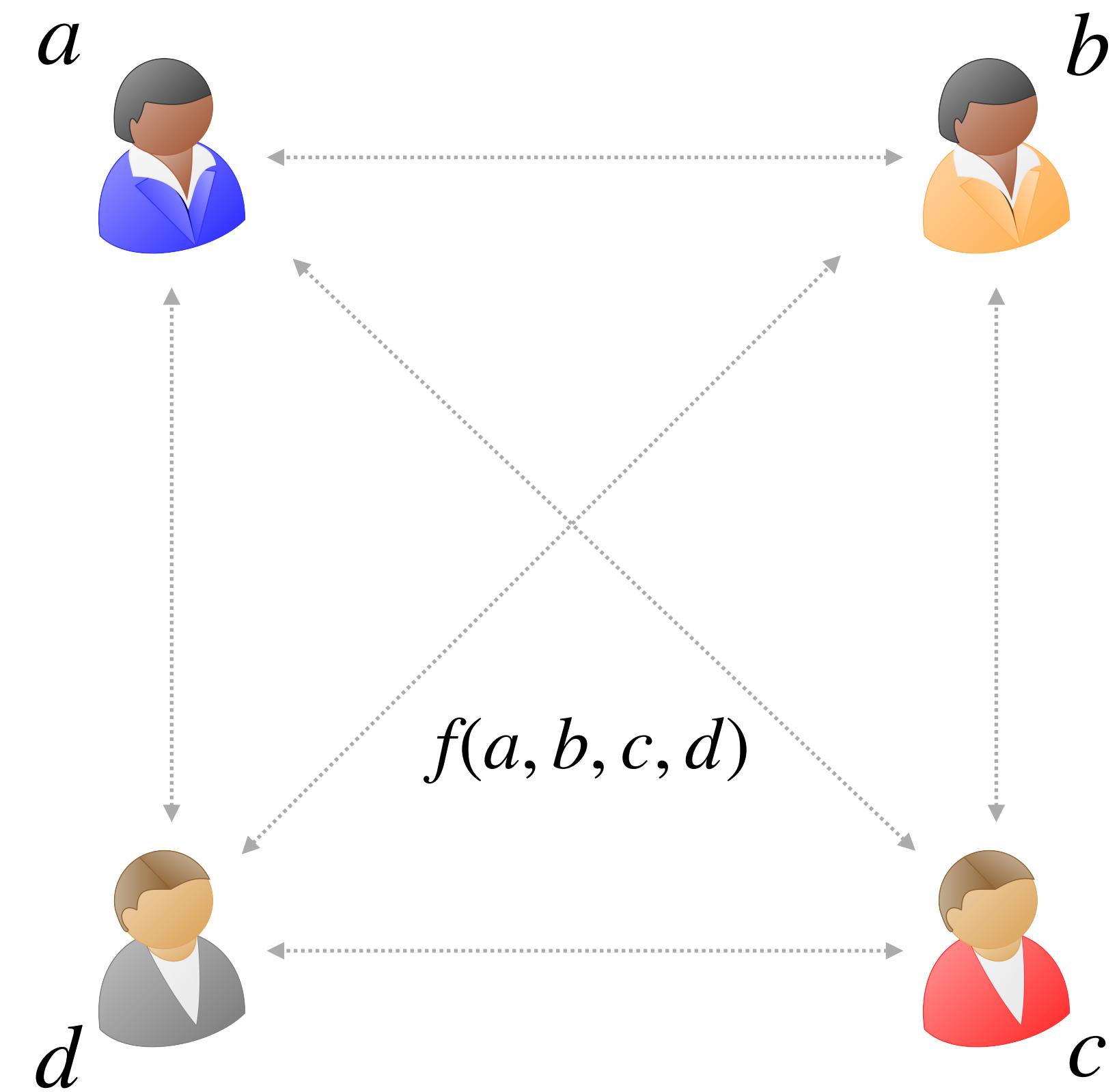
Ashish Choudhury, Aditya Hegde

INDOCRYPT 2020



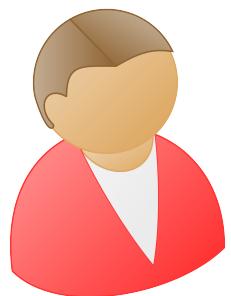
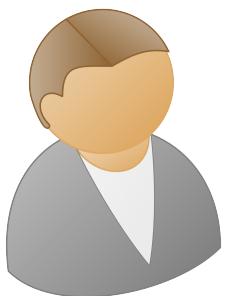
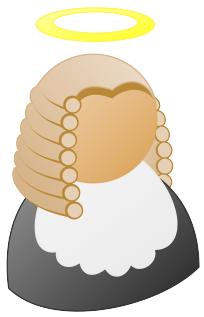
Secure Multi-Party Computation (MPC)

- Distrusting parties compute a function on private inputs



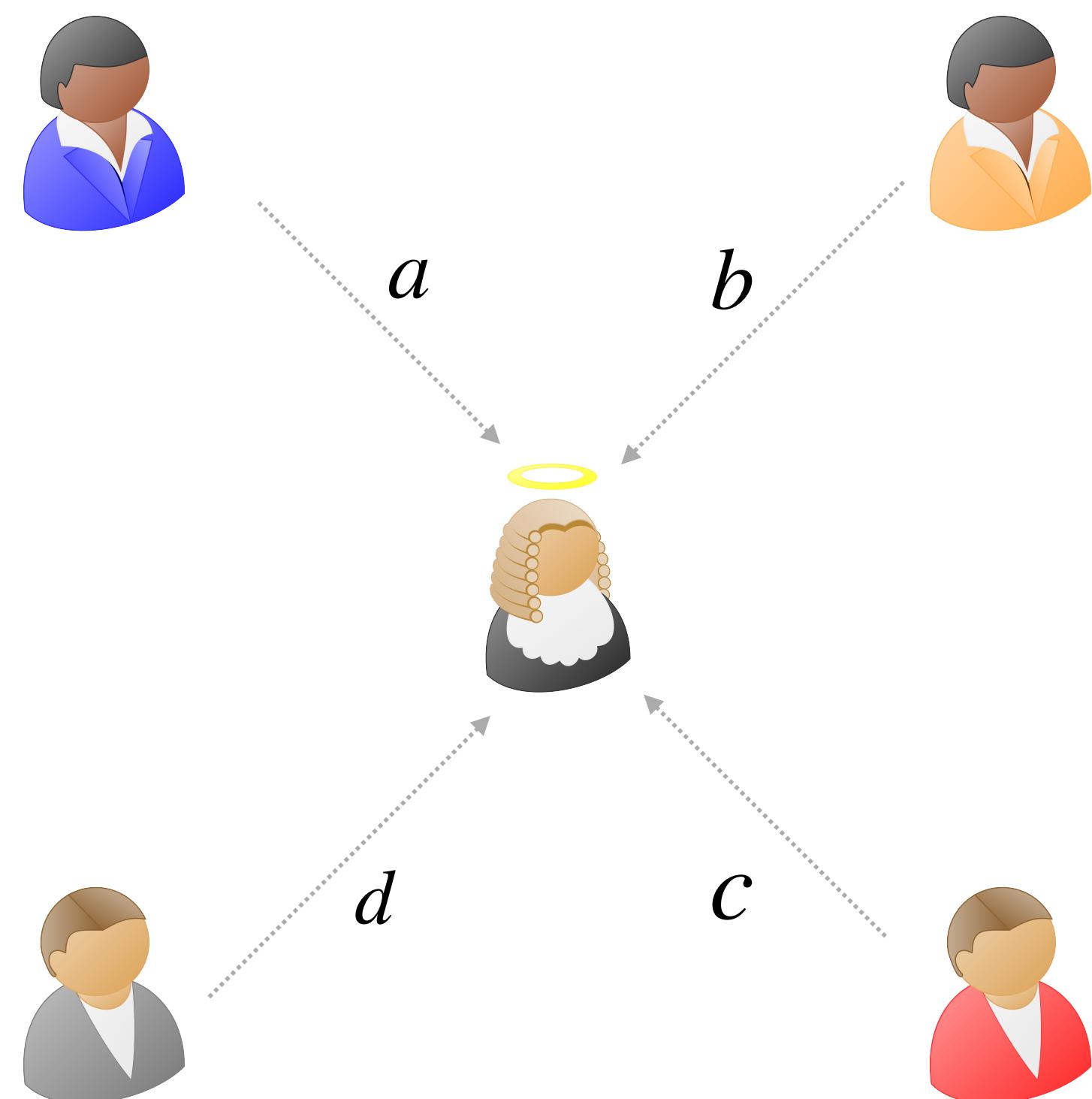
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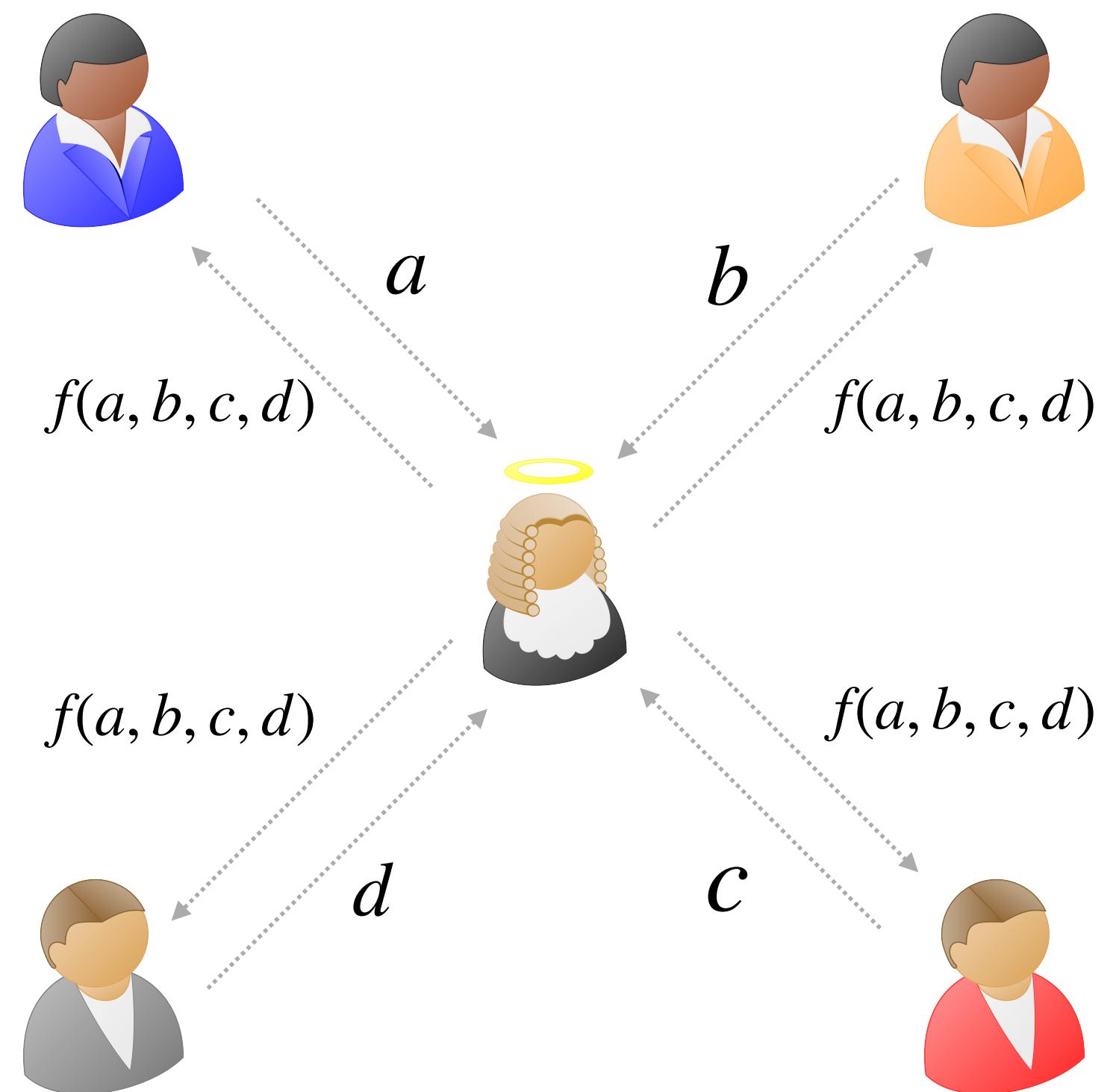
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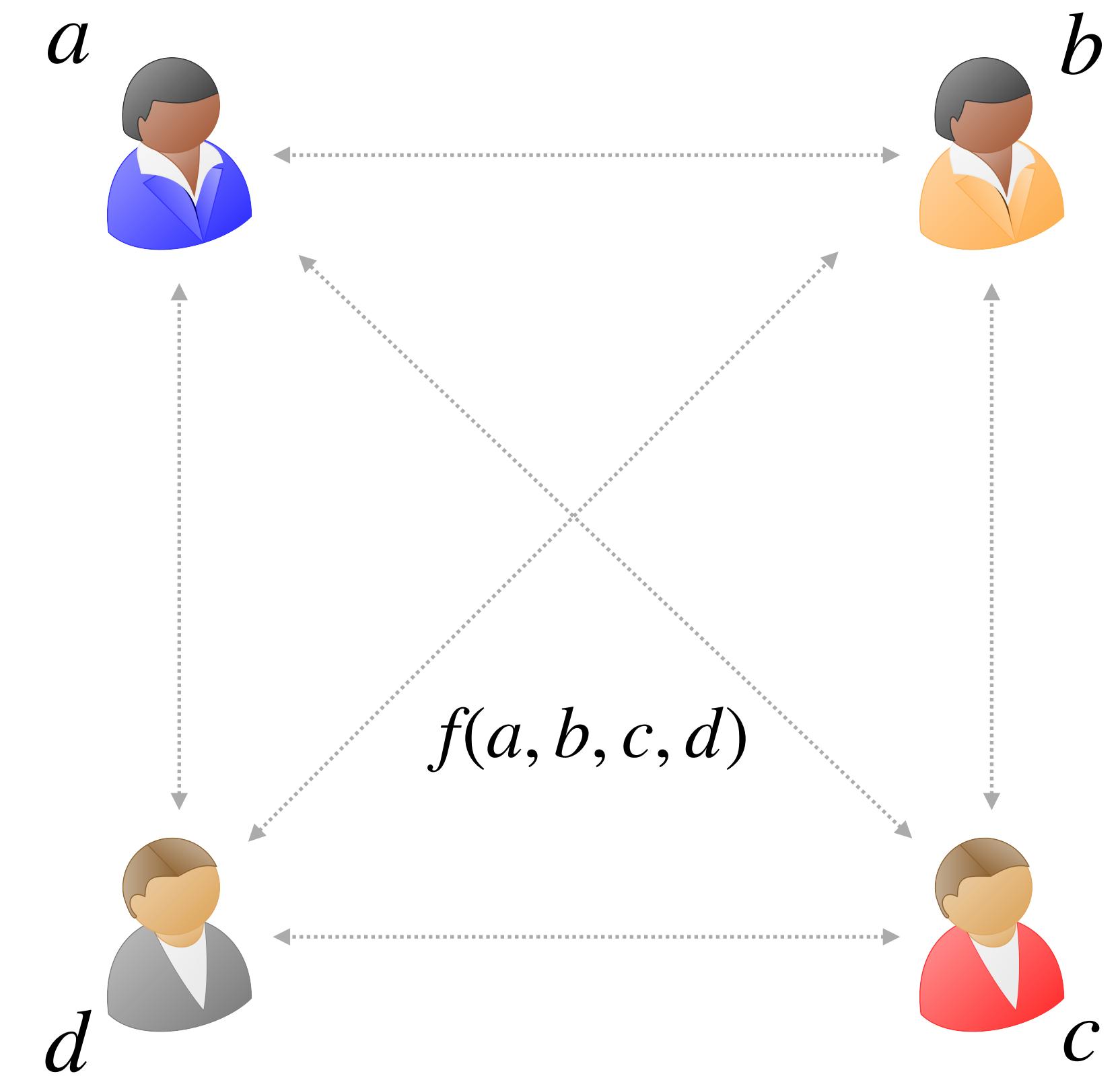
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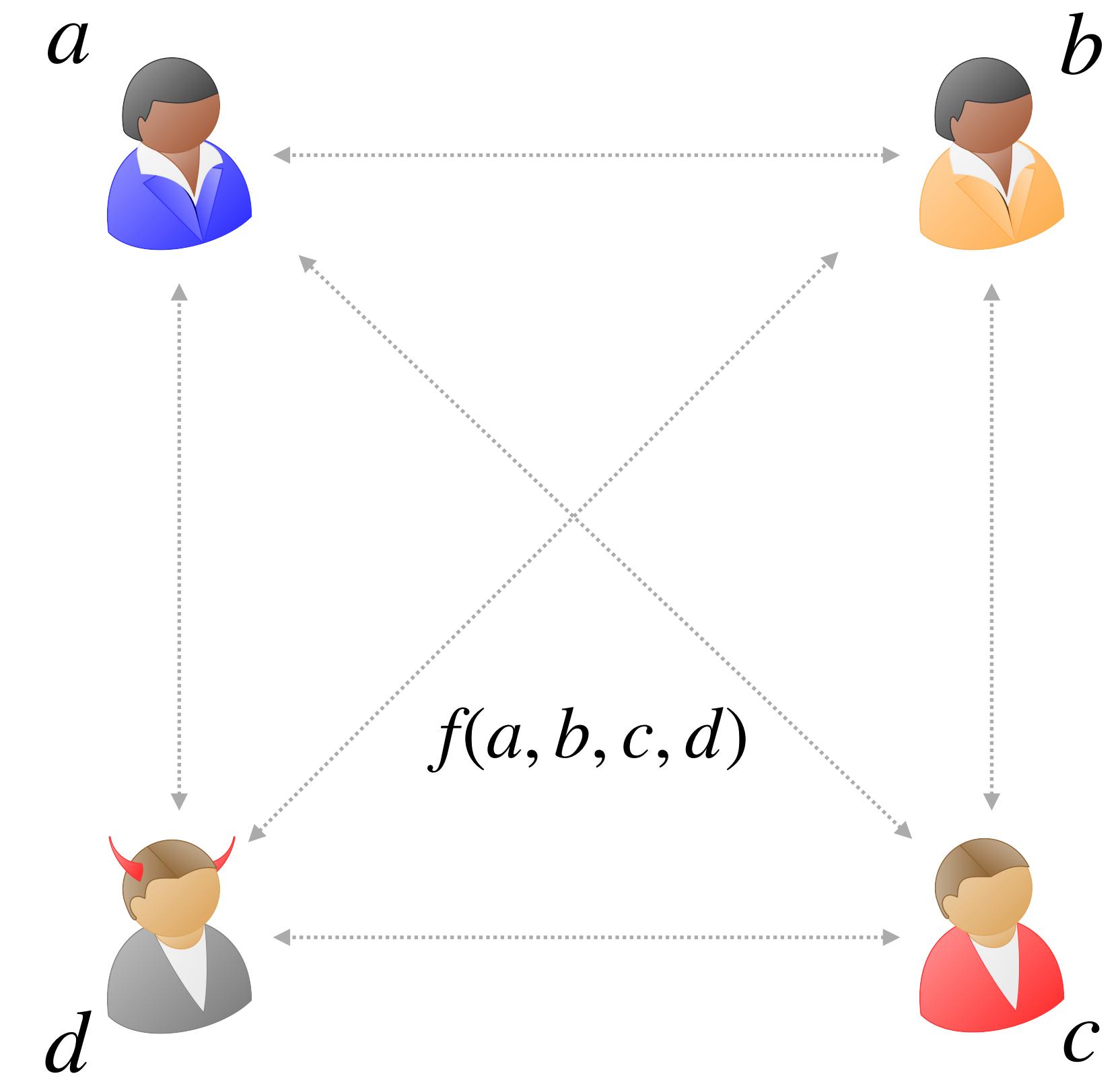
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- Setting
 - $n = 4, t = 1$
 - Malicious adversary

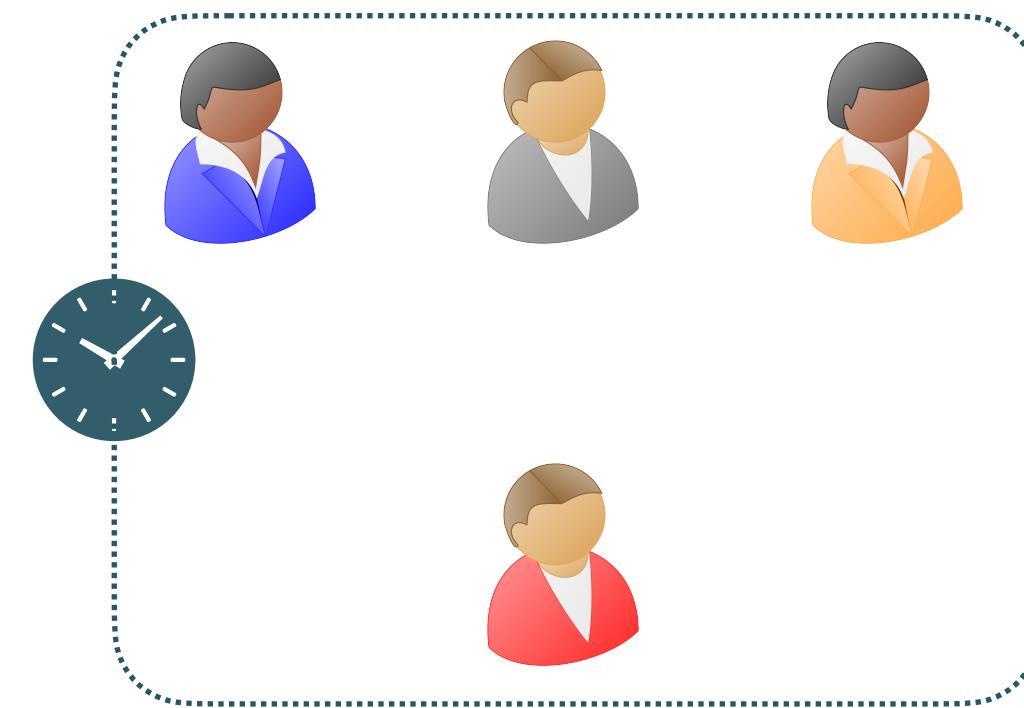


Communication Model - Synchronous and Asynchronous Networks

- Pairwise **private** and **authentic** channels

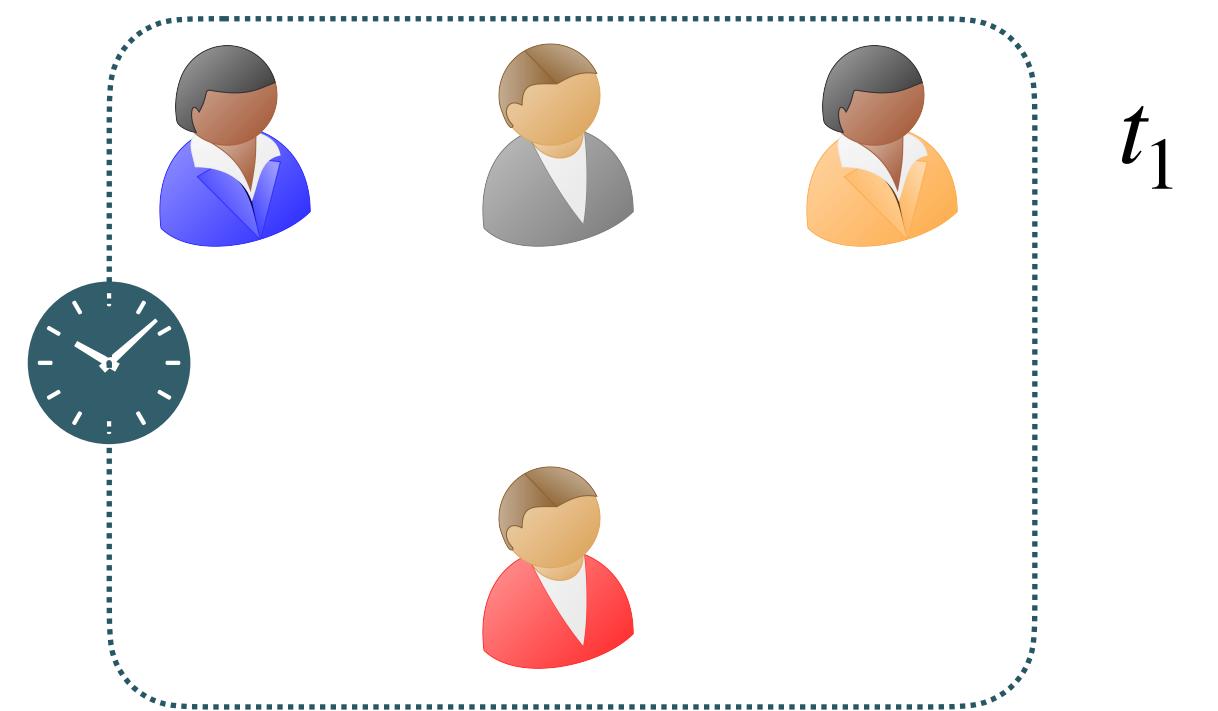
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 - Publicly known upper bound on message delay



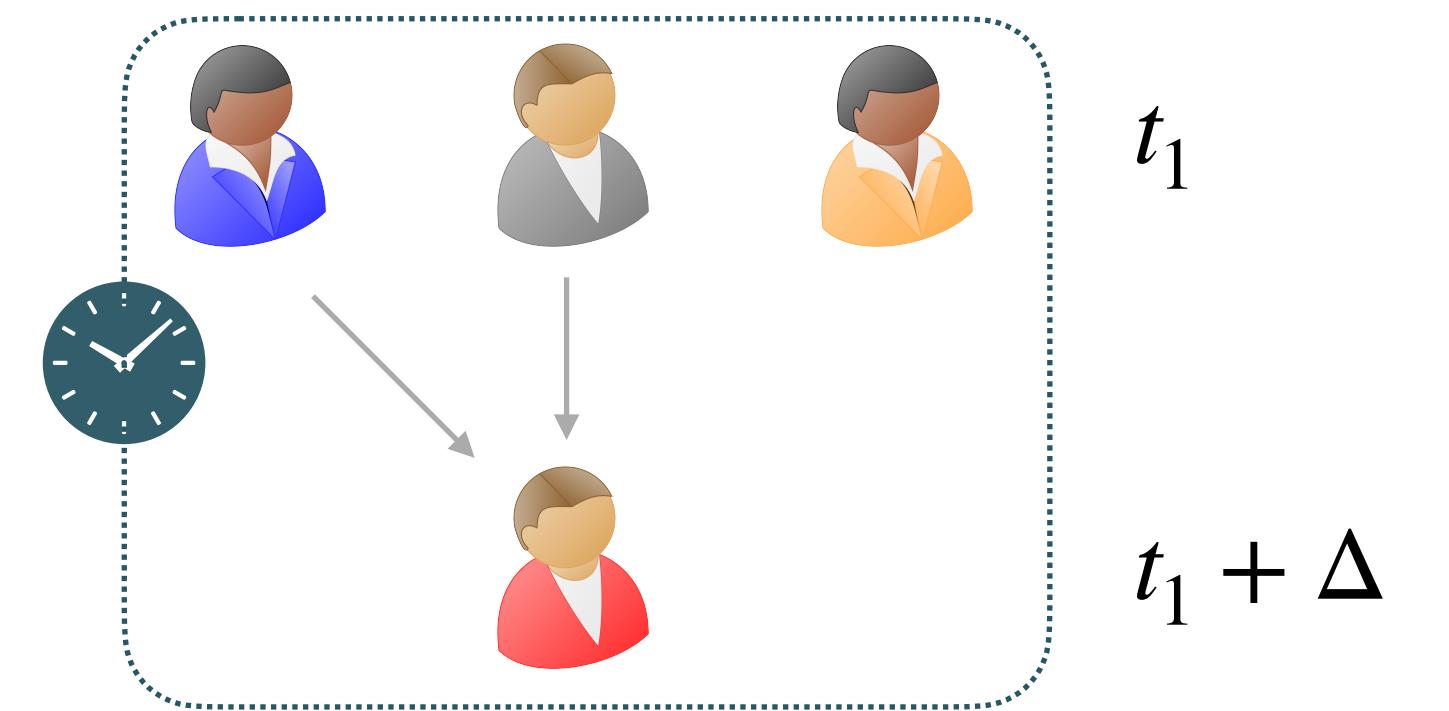
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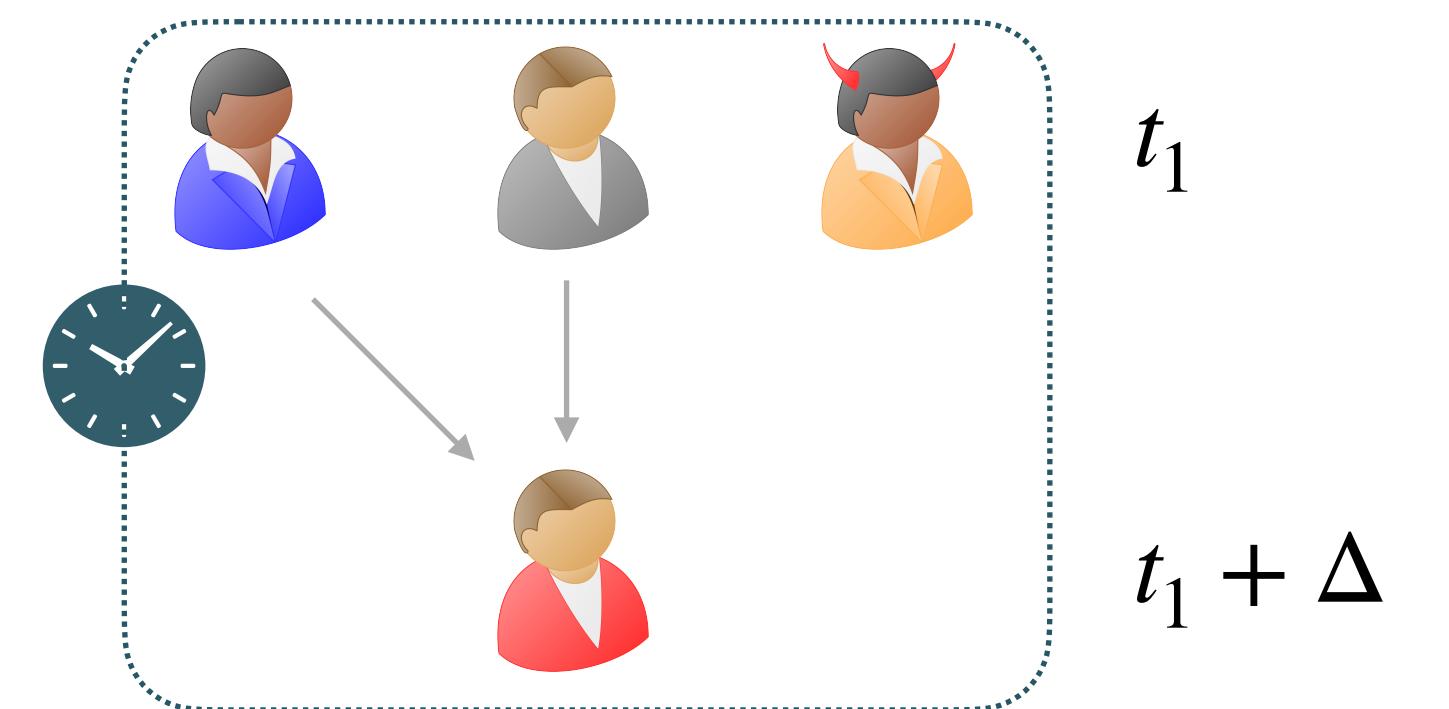
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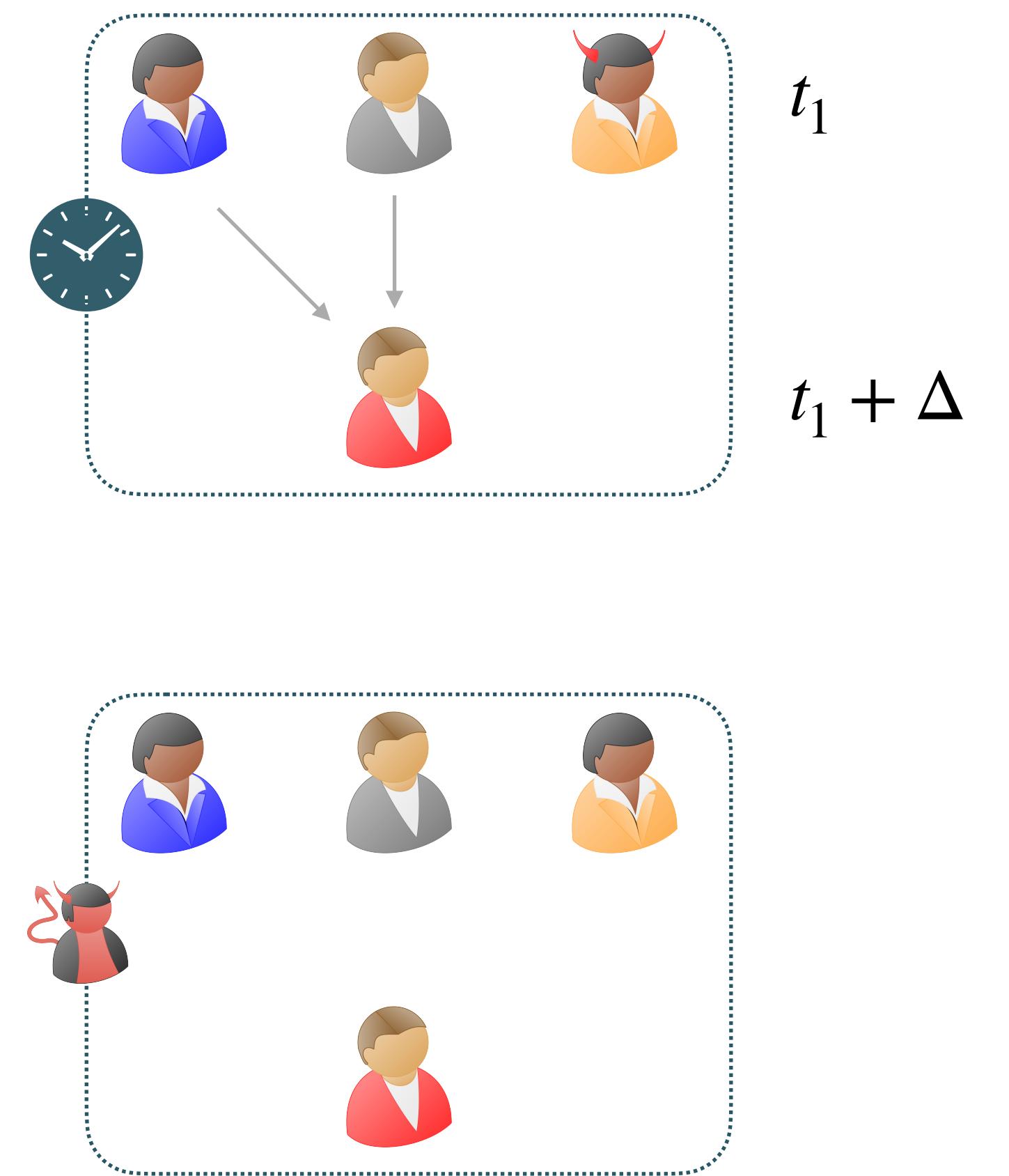
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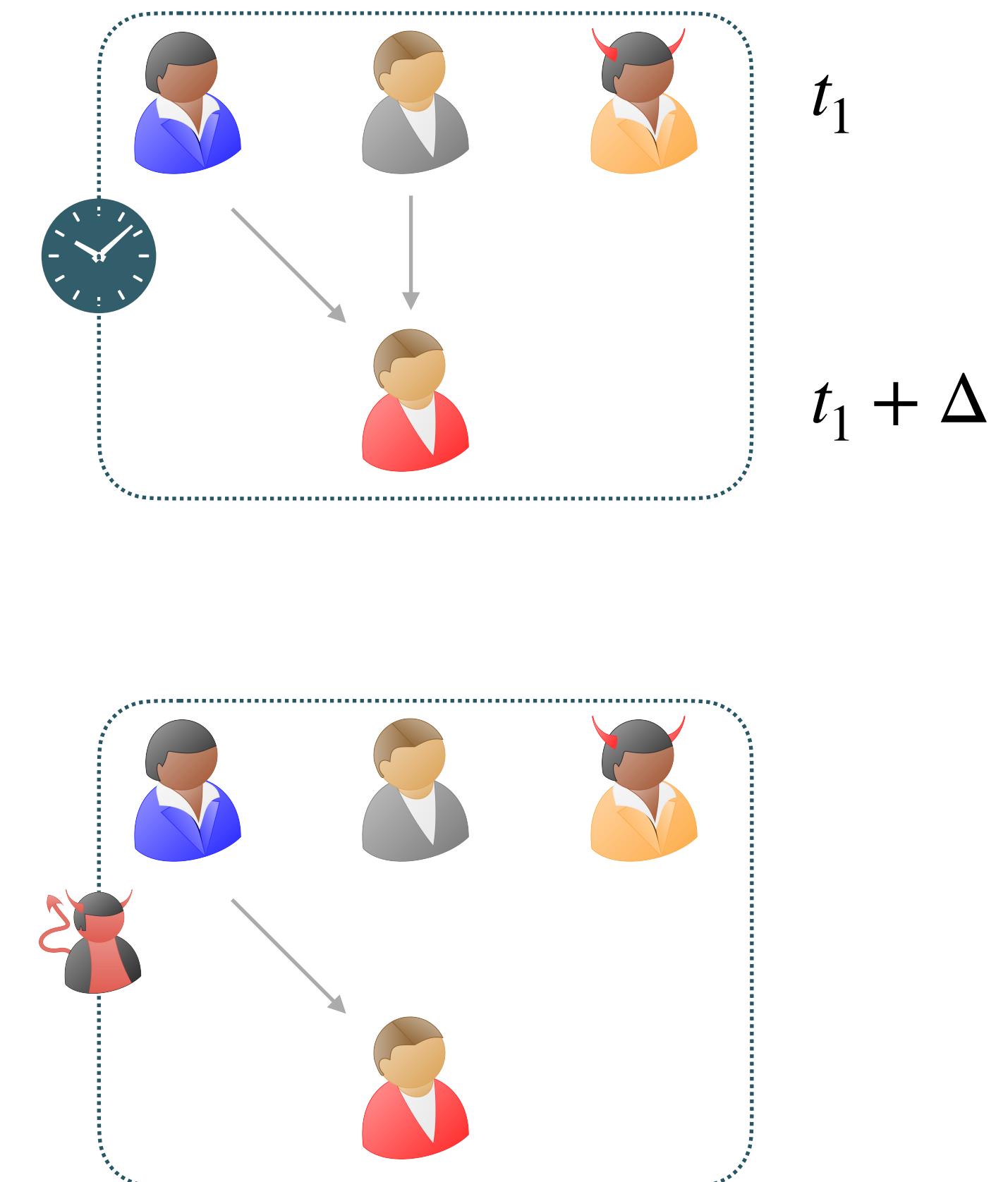
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 - No synchronisation
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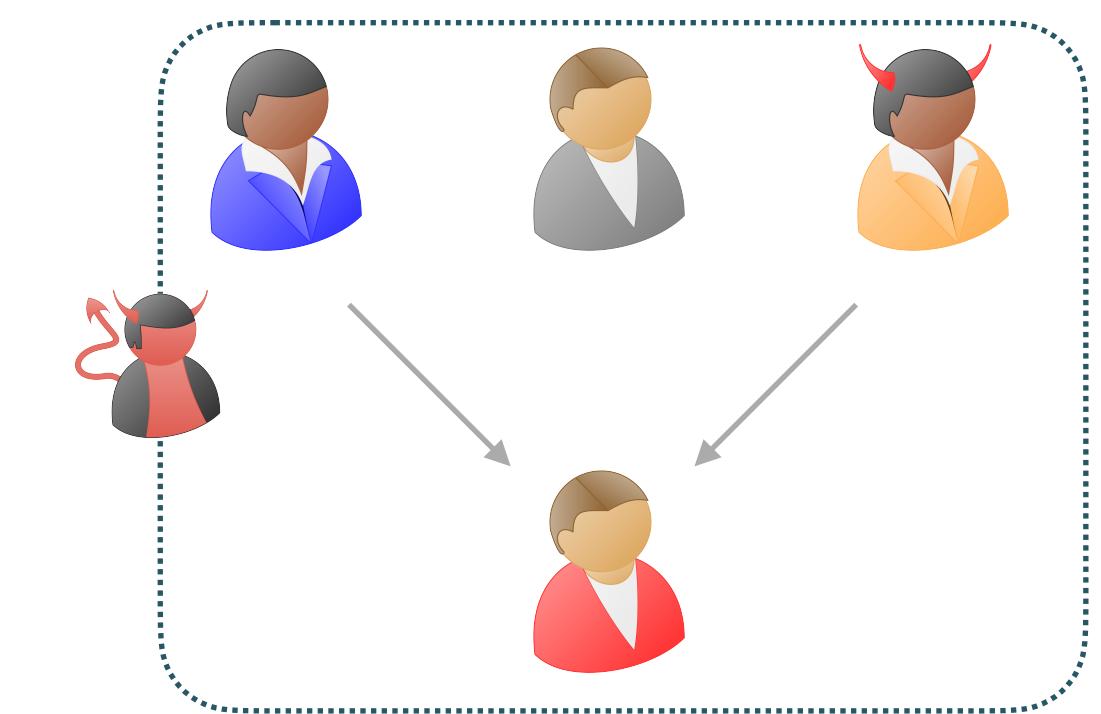
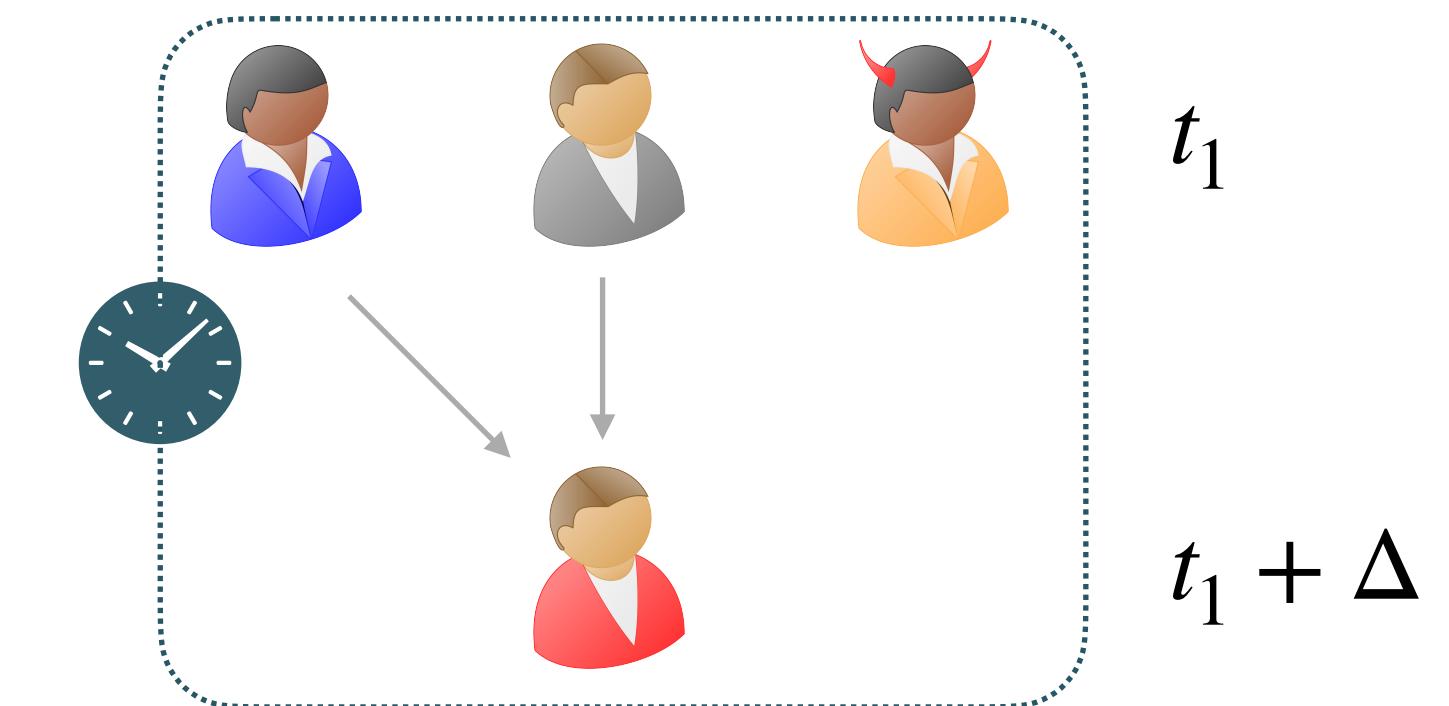
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Cannot distinguish
between delayed and
unsent message

Can only wait for
 $n - t$ messages at
each step

Communication Model - Hybrid Networks

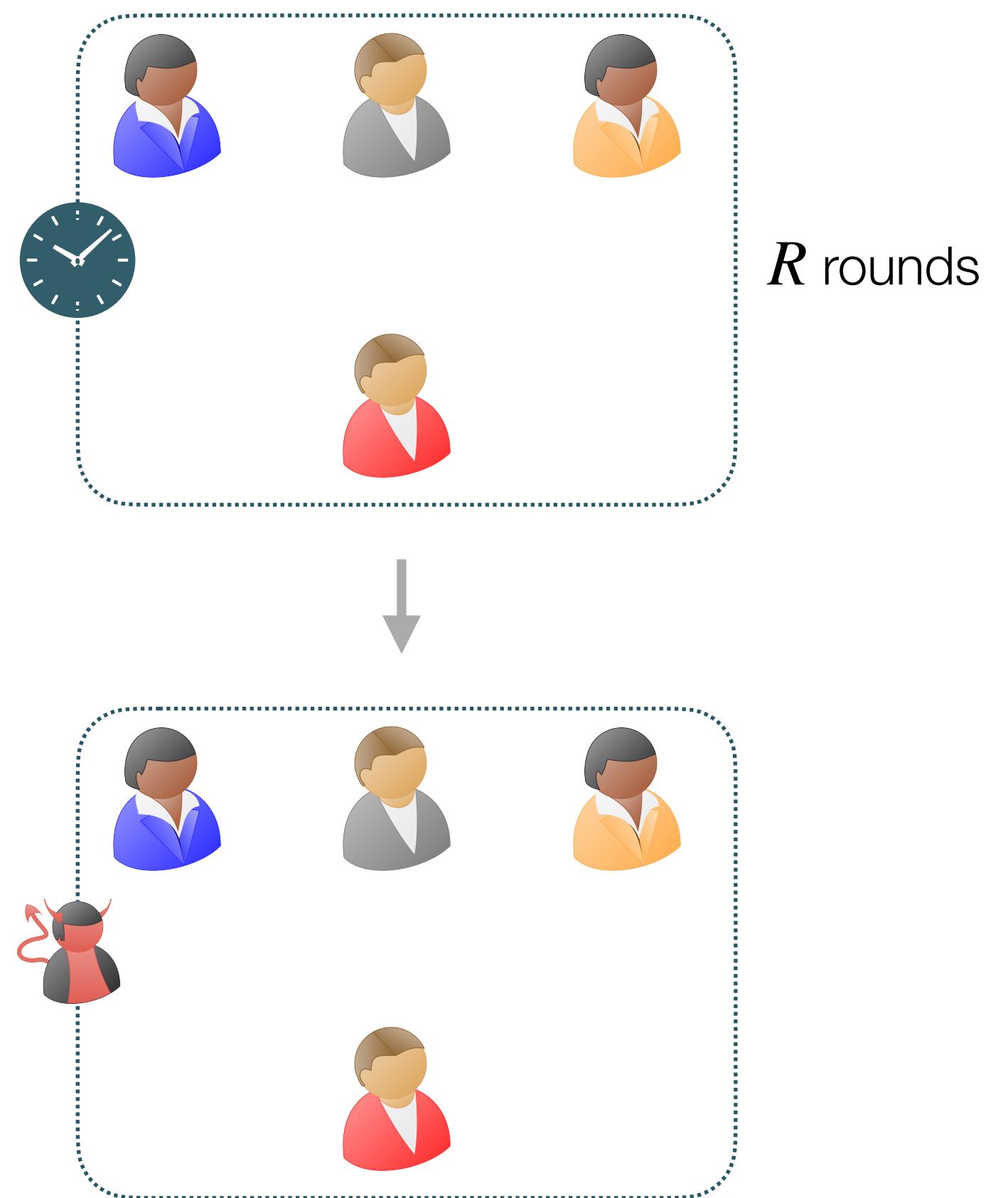
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- **Hybrid networks:** R initial synchronous rounds followed by asynchronous computation [BHN10, CHP13, PR18]
 - Assume synchronous broadcast channel in first R rounds



Our Contributions

- Perfectly secure MPC protocol over hybrid network with $R = 2$
 - First protocol in this setting
- Cryptographically secure MPC protocol over hybrid network with $R = 1$
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 - Implementation and benchmarks

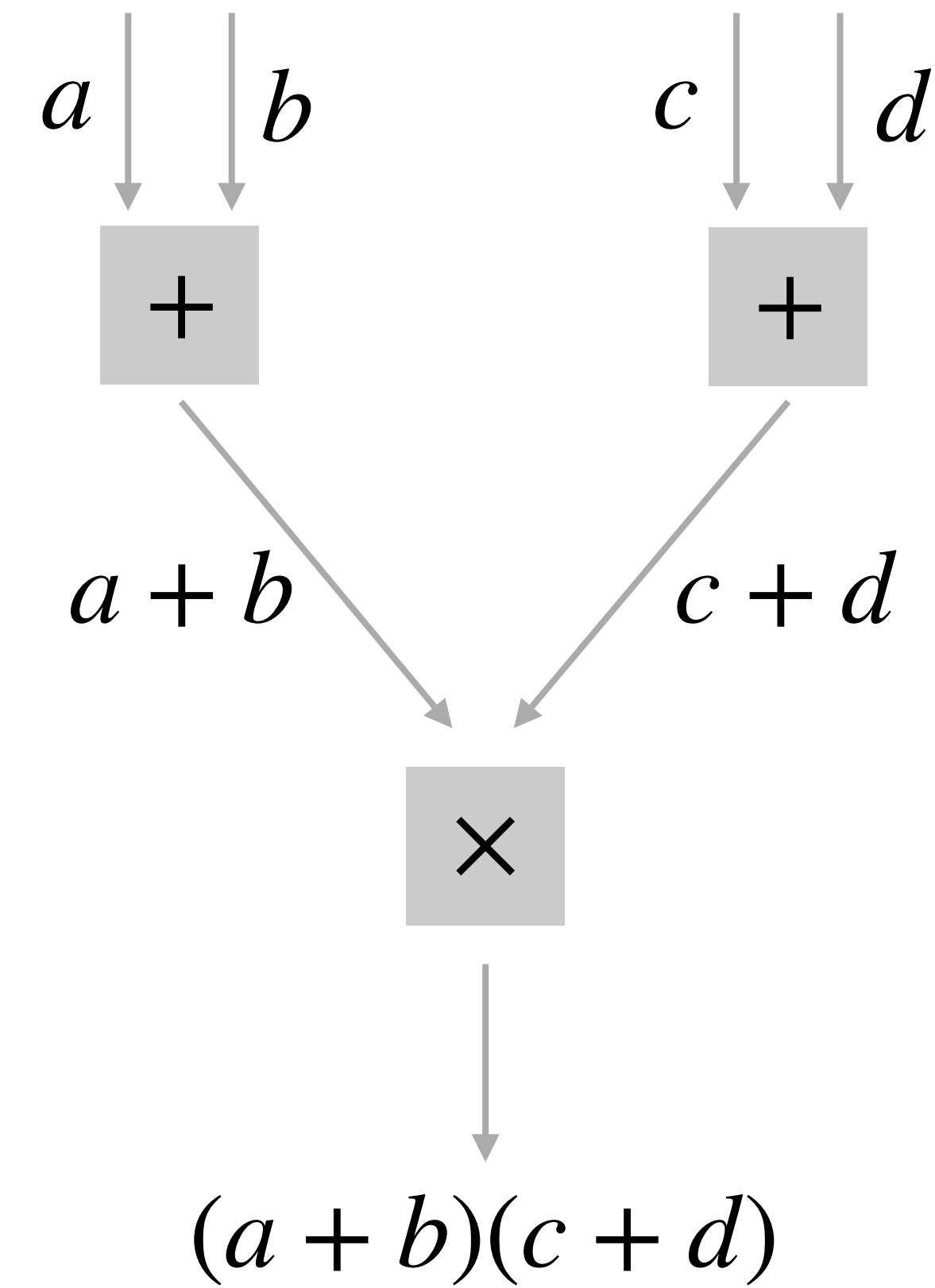
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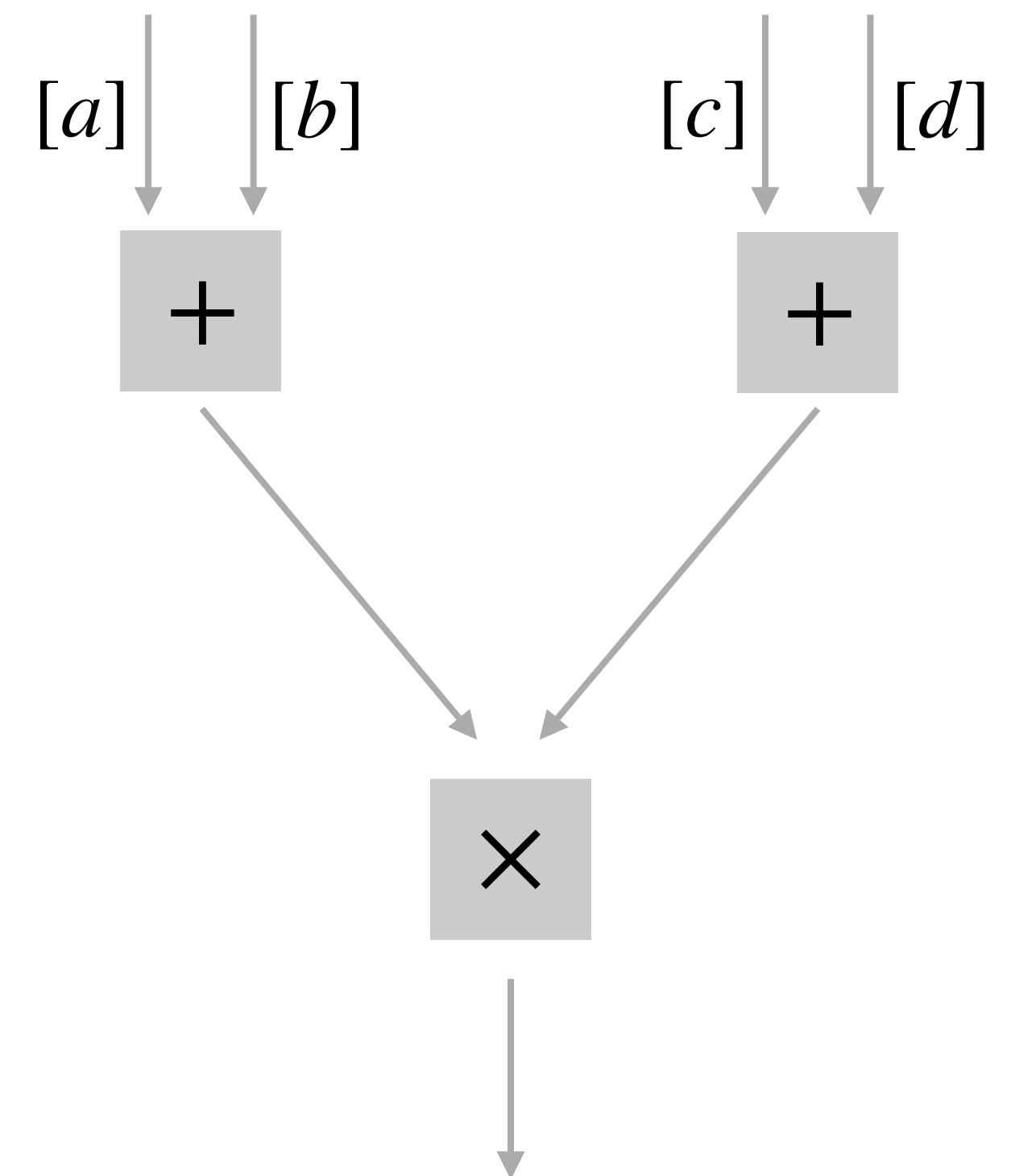
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- f represented as **arithmetic circuit** over finite field



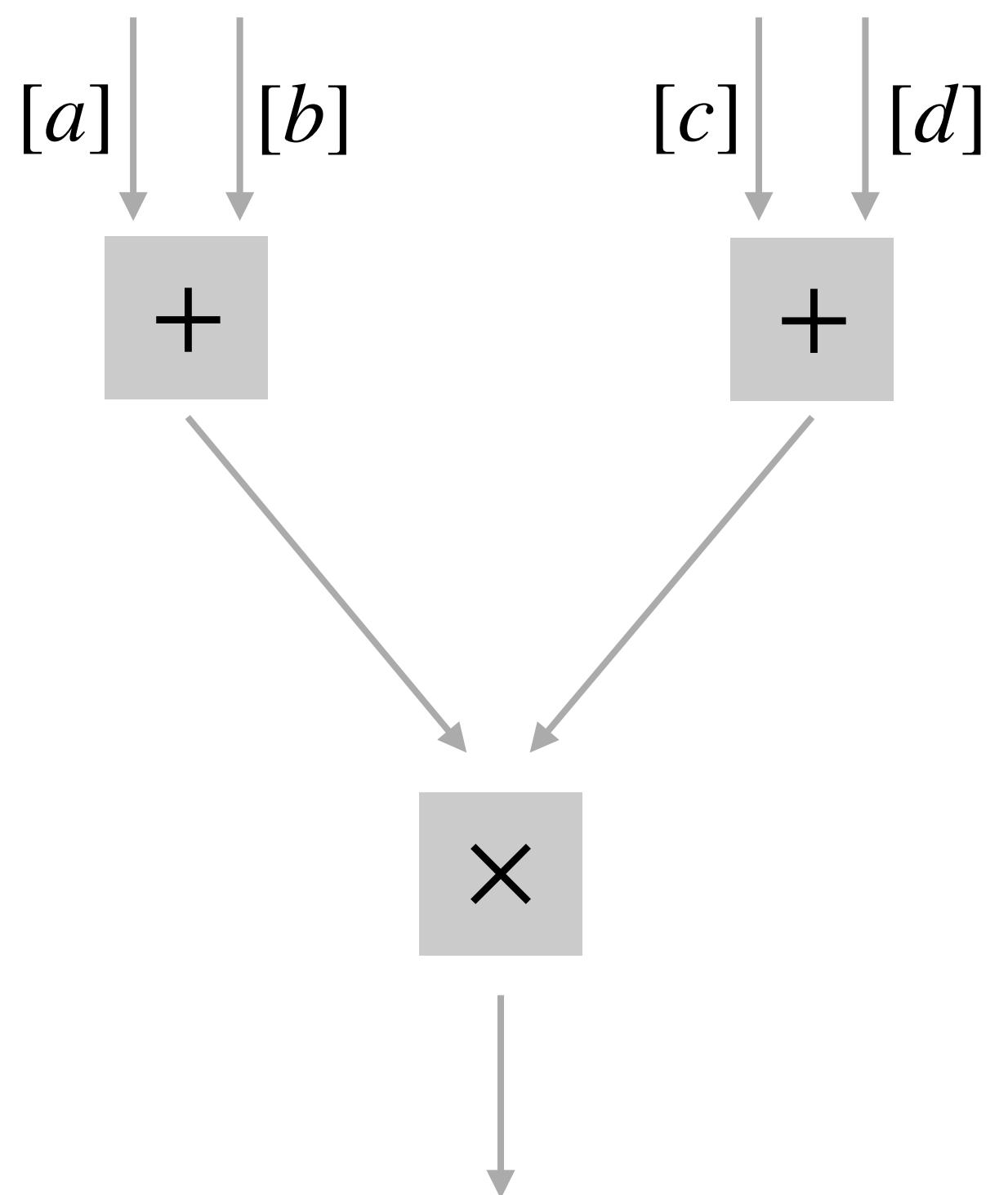
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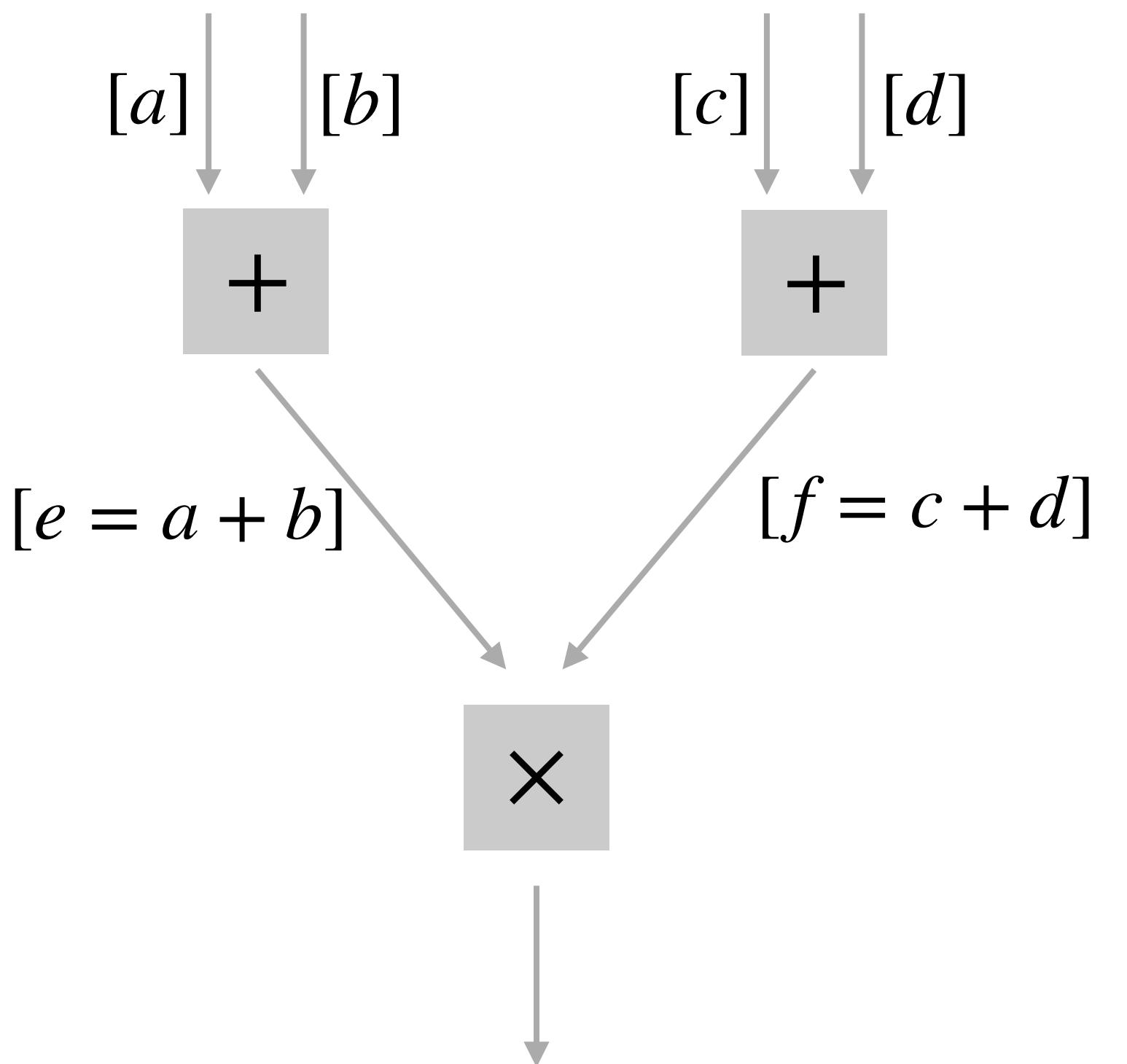
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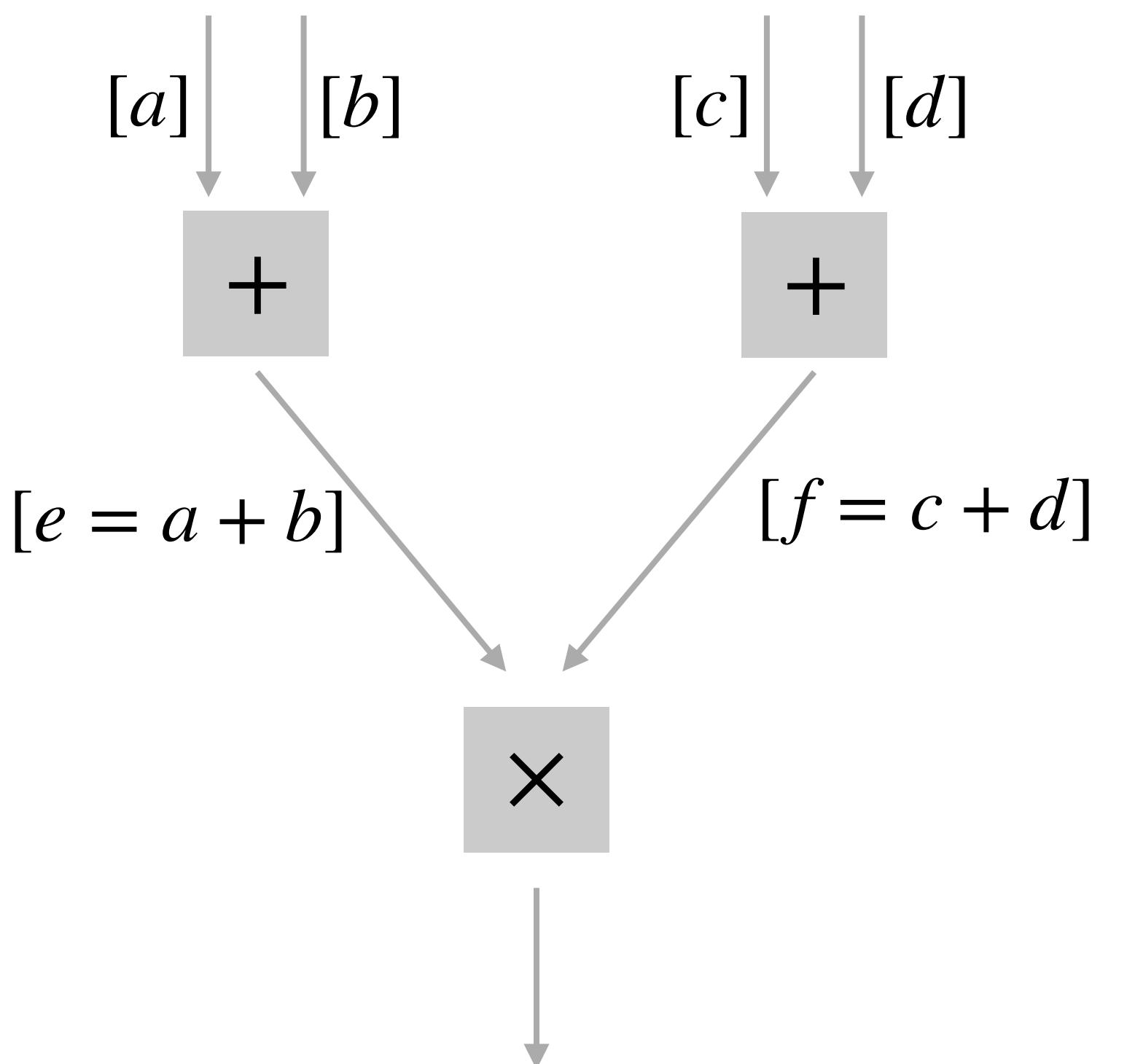
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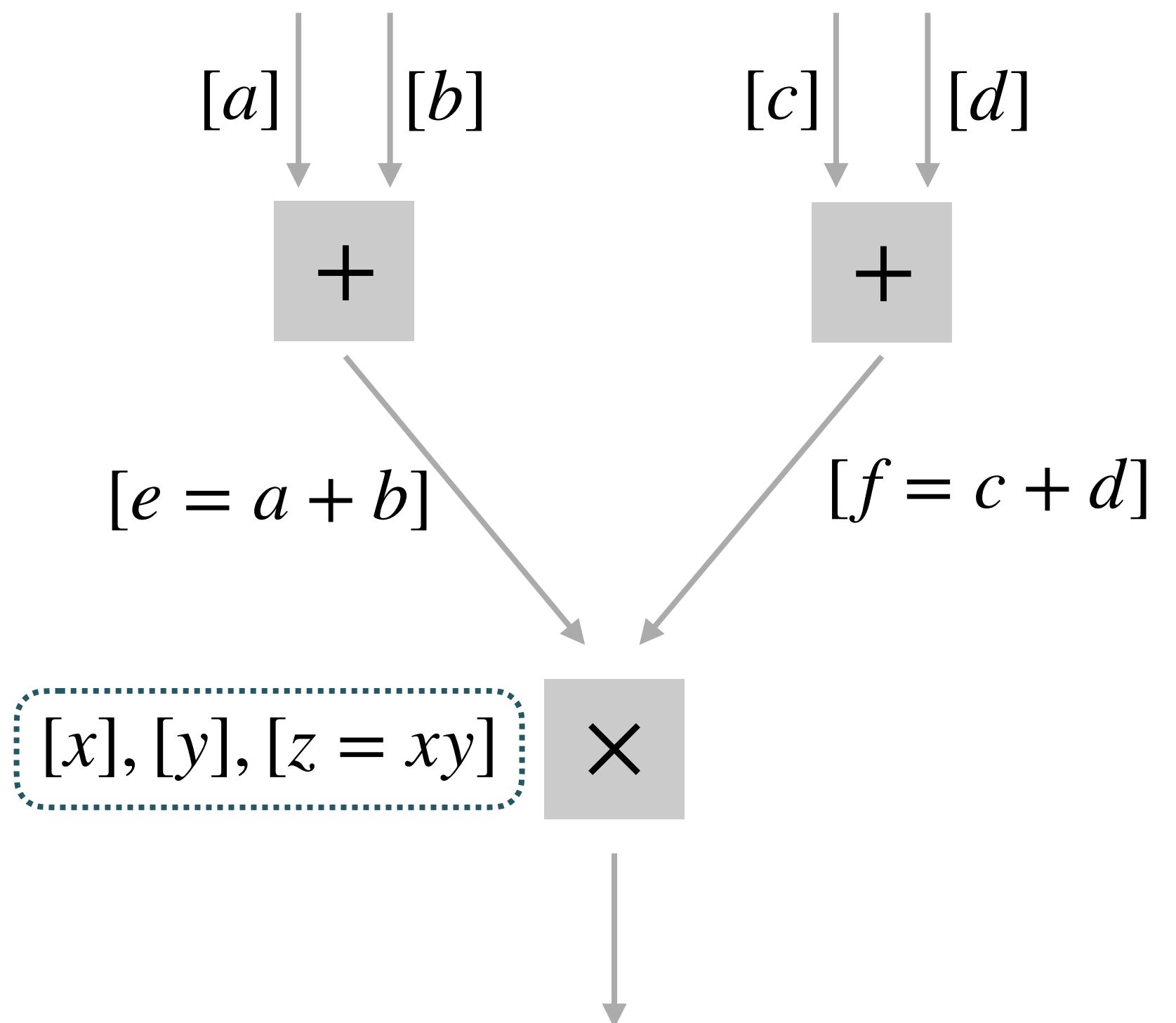
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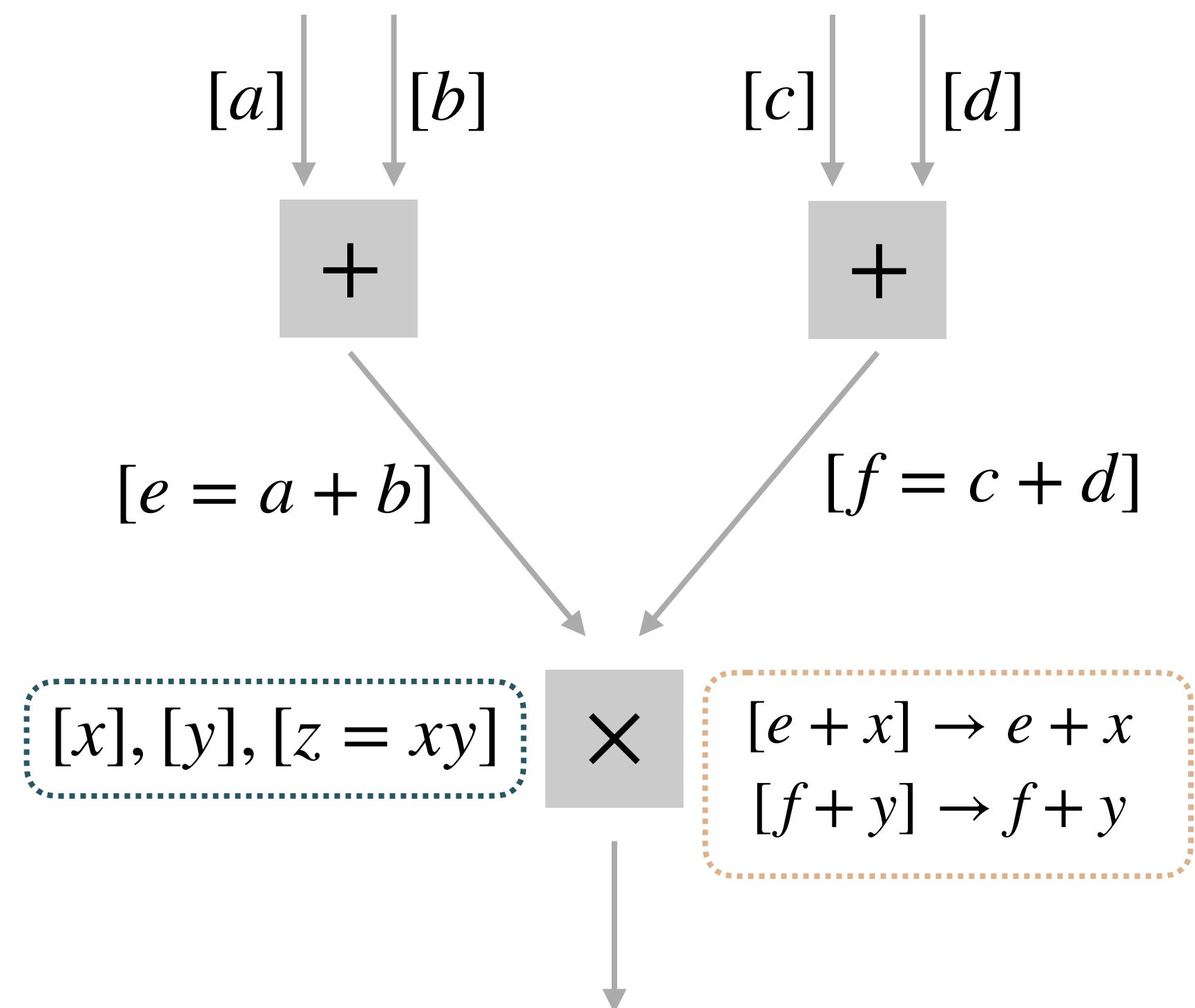
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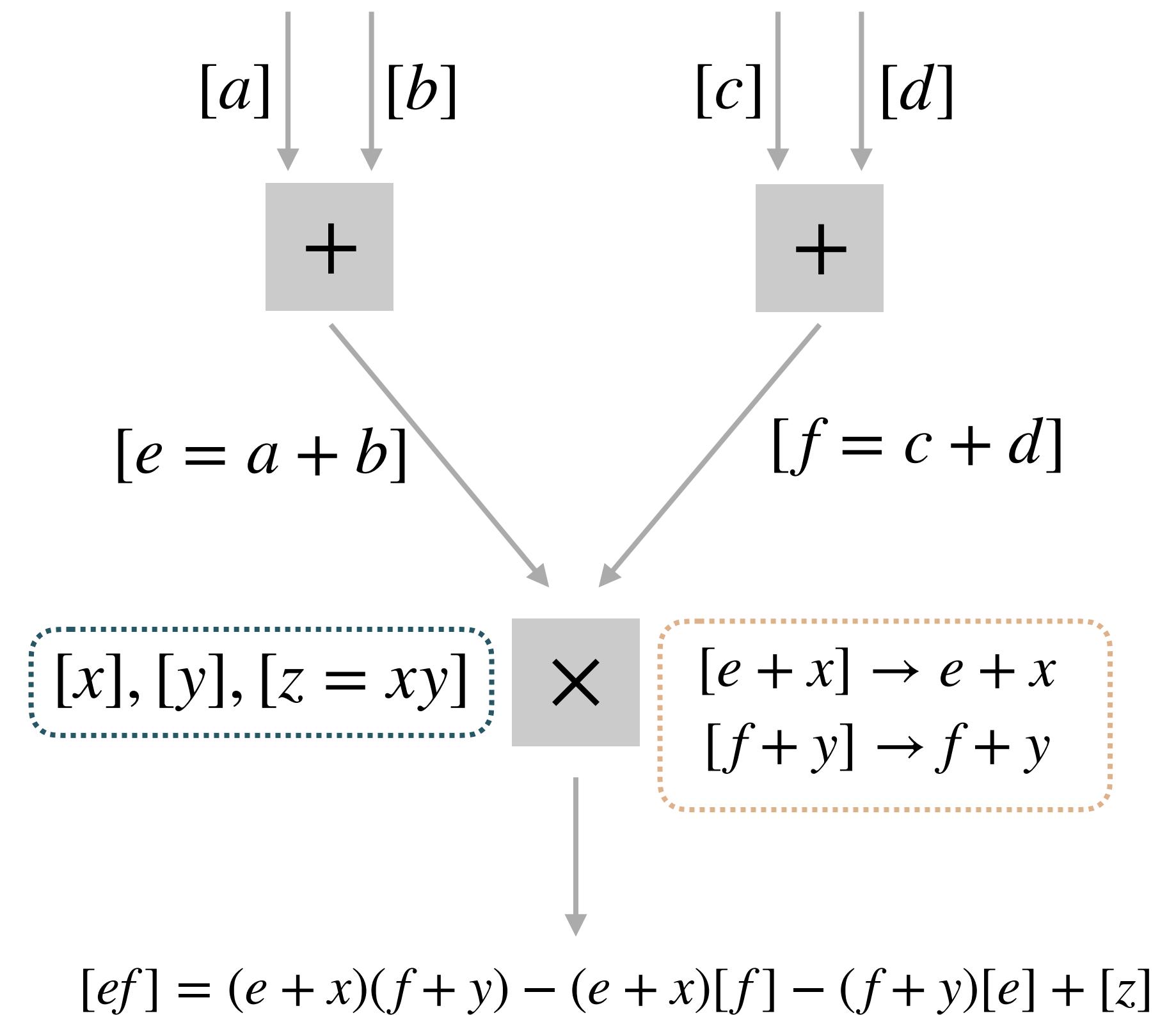
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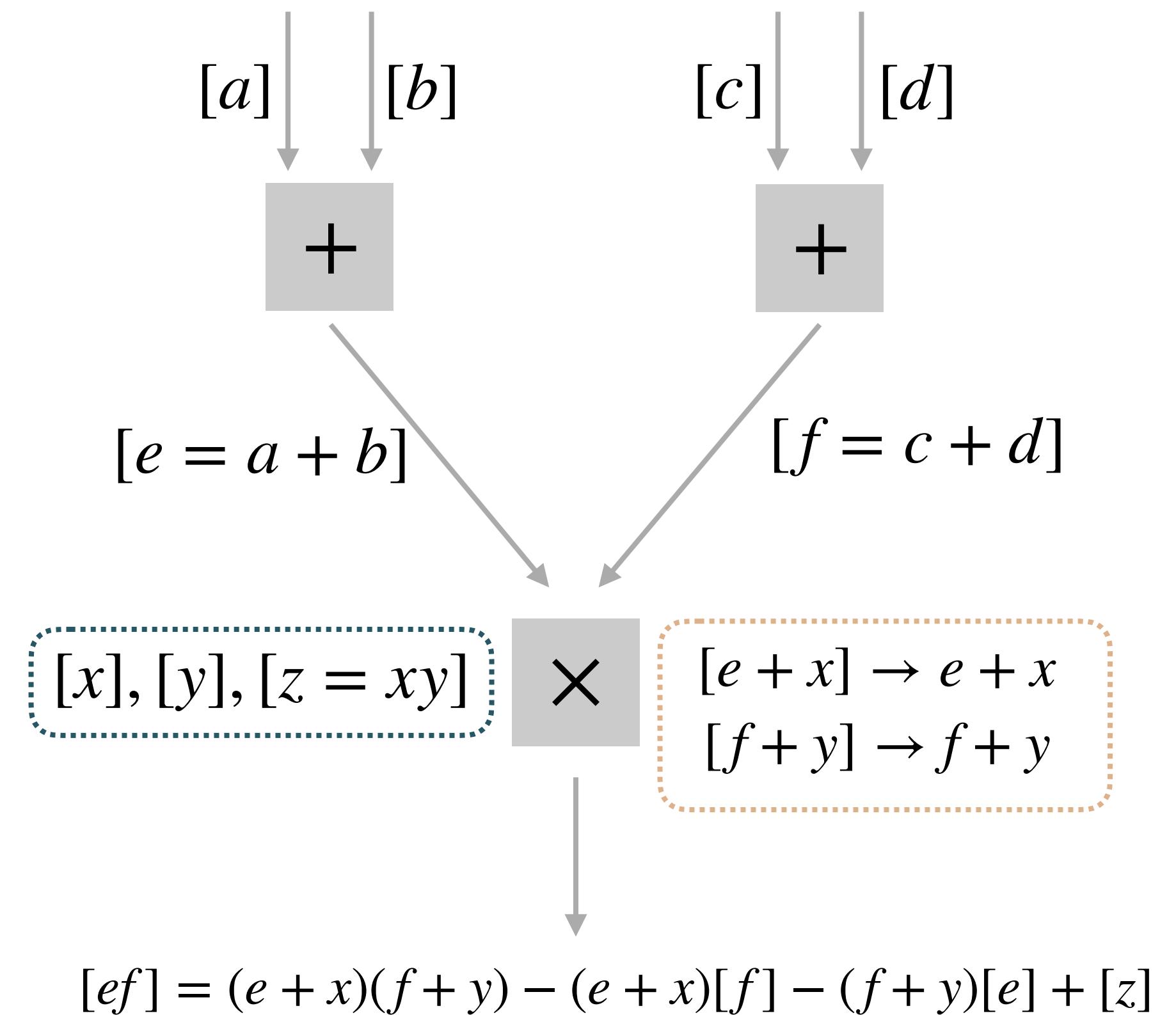
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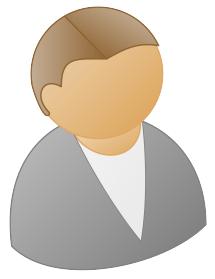
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Beaver

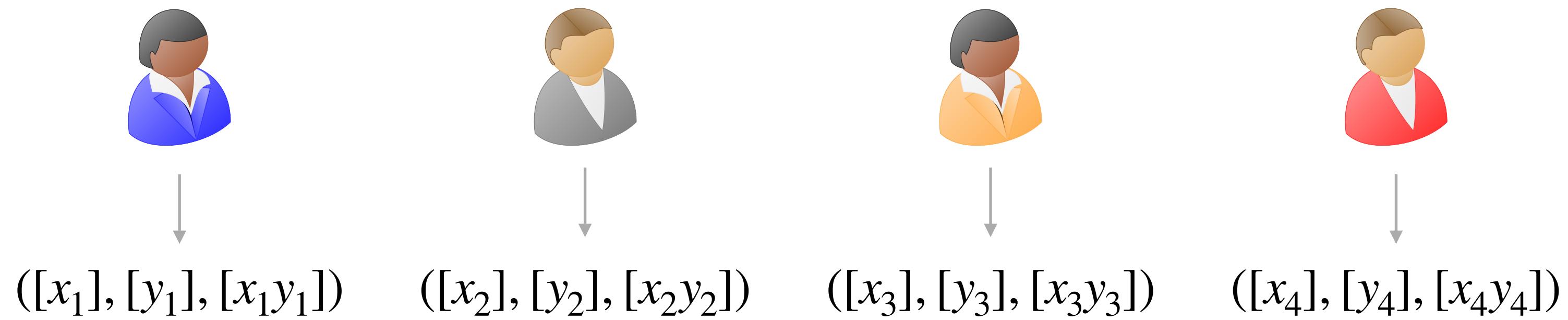
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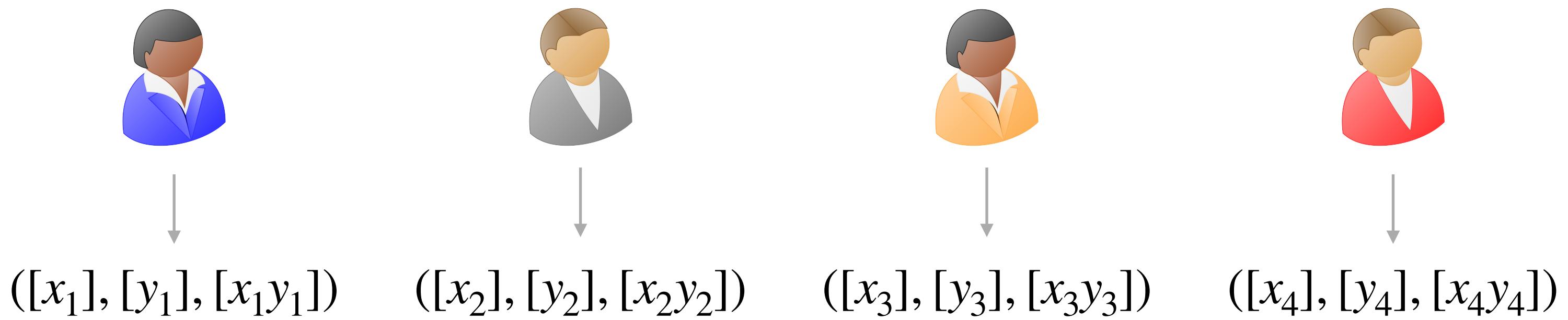
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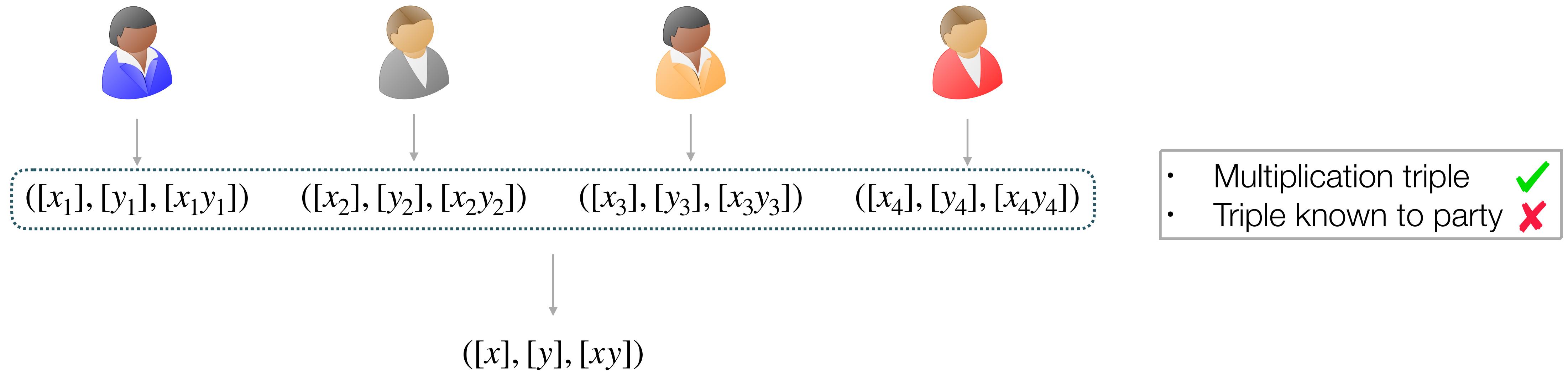
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- Multiplication triple ✓
- Triple known to party ✗

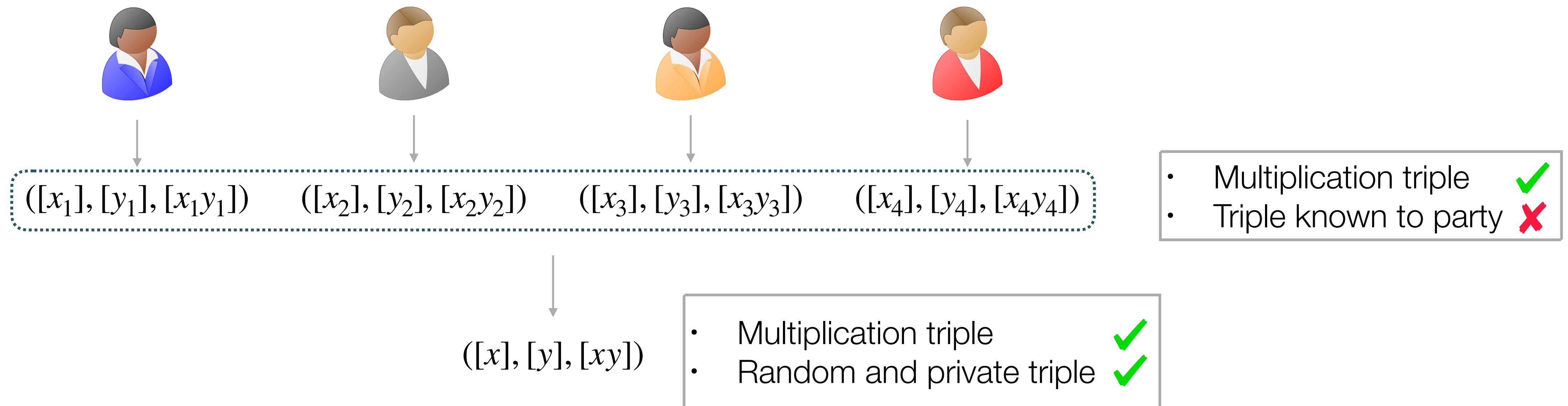
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Perfect HMPC

- Open Problem [PR18]: Perfectly secure MPC protocol over hybrid network
 - Two synchronous rounds
 - Tolerating $t < n/3$ corruptions
 - With synchronous broadcast channel
 - Guaranteed output delivery
- Input provision impossible in this setting [PR18]

Perfect HMPC - Linear Secret Sharing Scheme

- Replicated Secret Sharing [ISN89]

- $[s] = (s_1, s_2, s_3, s_4)$

- $s = s_1 + s_2 + s_3 + s_4$

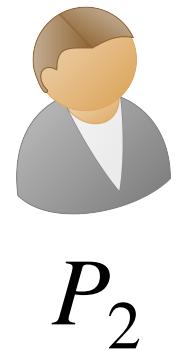
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P_1



P_2



P_3



P_4

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(s_2, s_3, s_4)

(s_3, s_4, s_1)

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- P_i does not have s_i

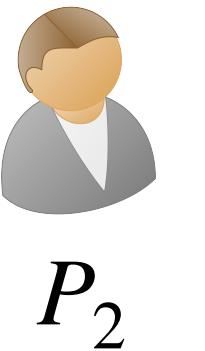
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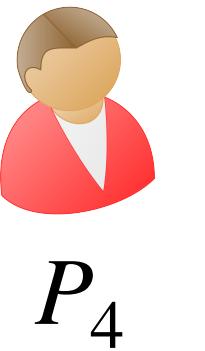
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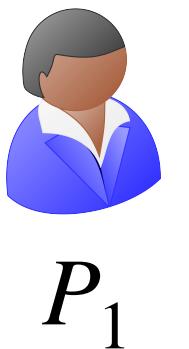
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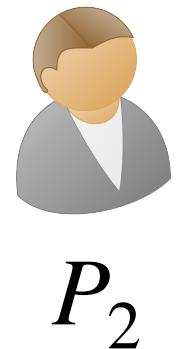
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$$(s_2, s_3, s_4) \quad (s_3, s_4, s_1) \quad (s_4, s_1, s_2) \quad (s_1, s_2, s_3)$$



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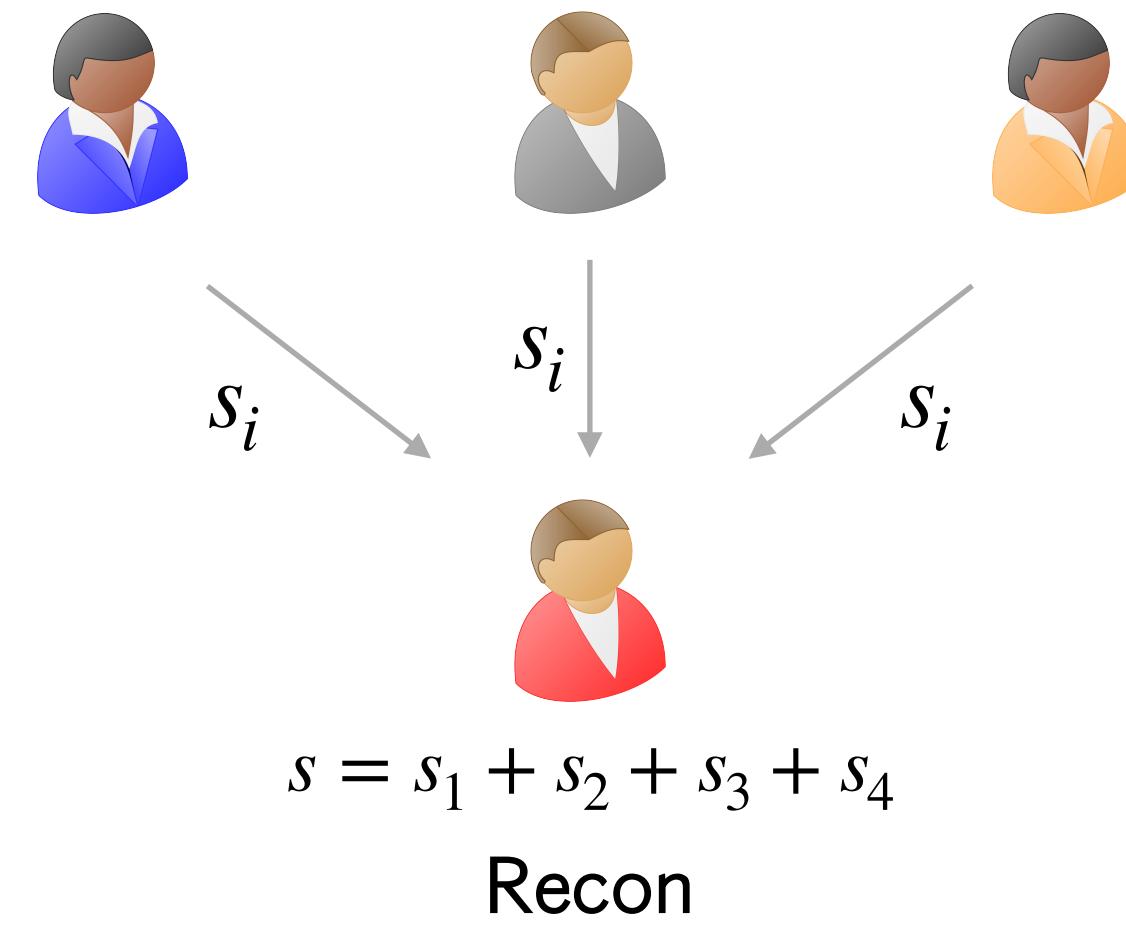
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$$[cs] = (cs_1, cs_2, cs_3, cs_4)$$
$$[s + s'] = (s_1 + s'_1, s_2 + s'_2, s_3 + s'_3, s_4 + s'_4)$$

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- Reconstruction

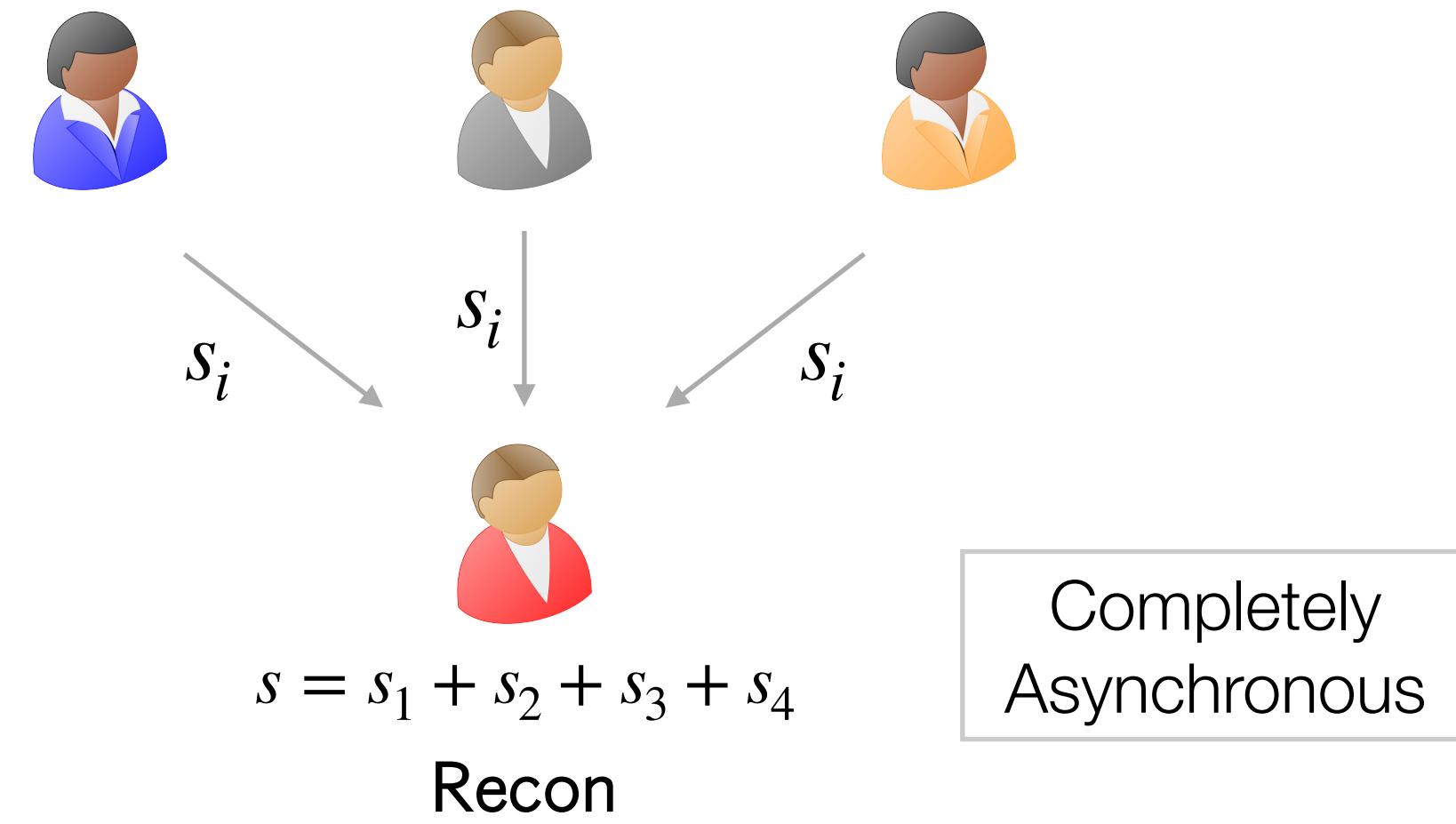
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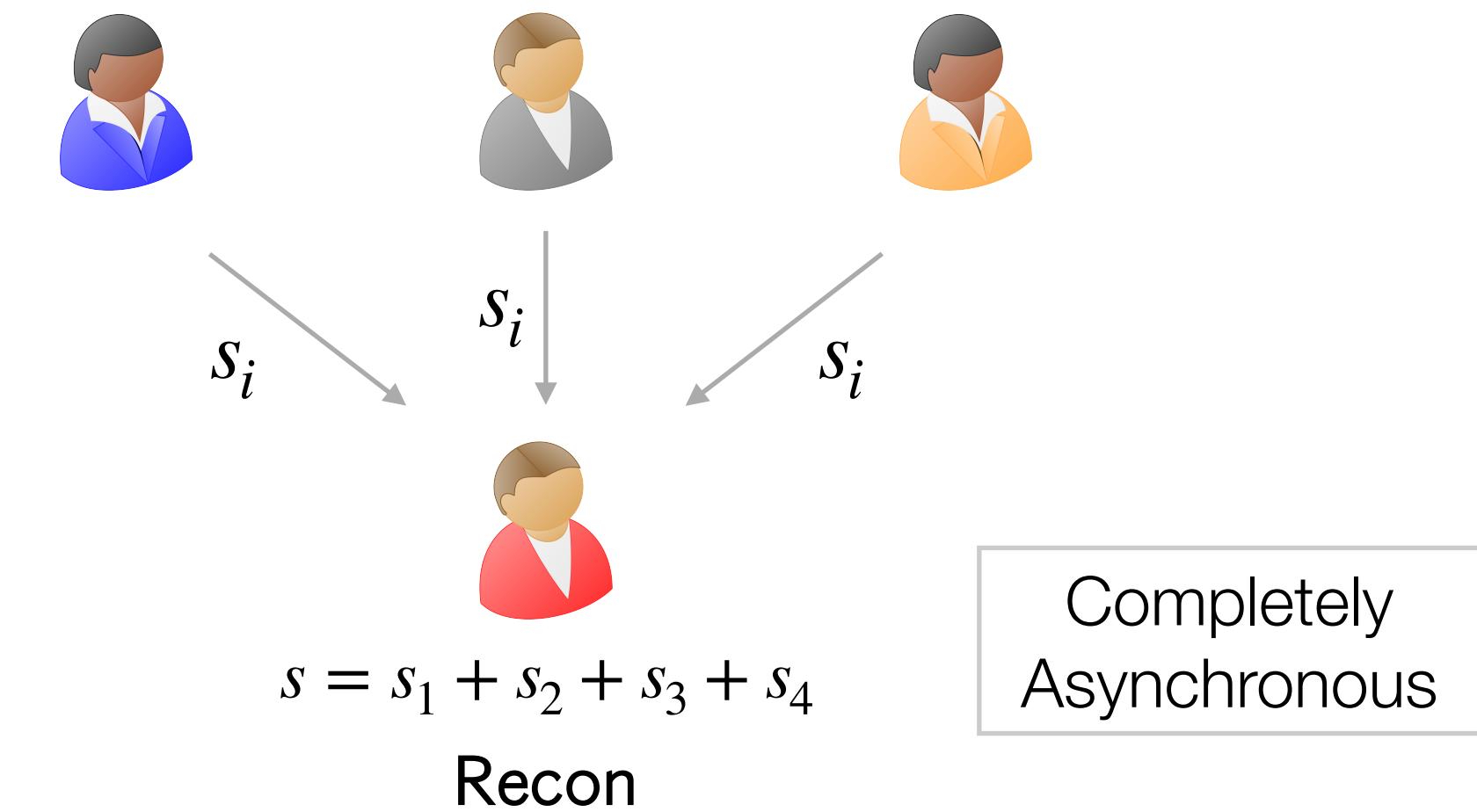
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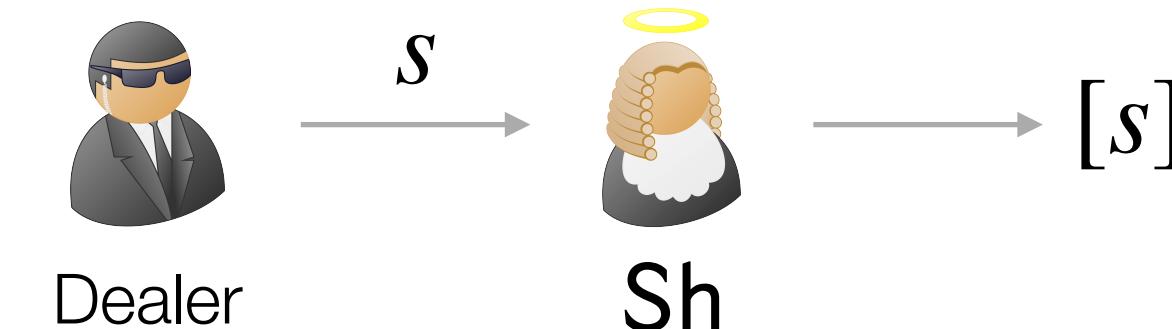
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- VSS with Party Elimination

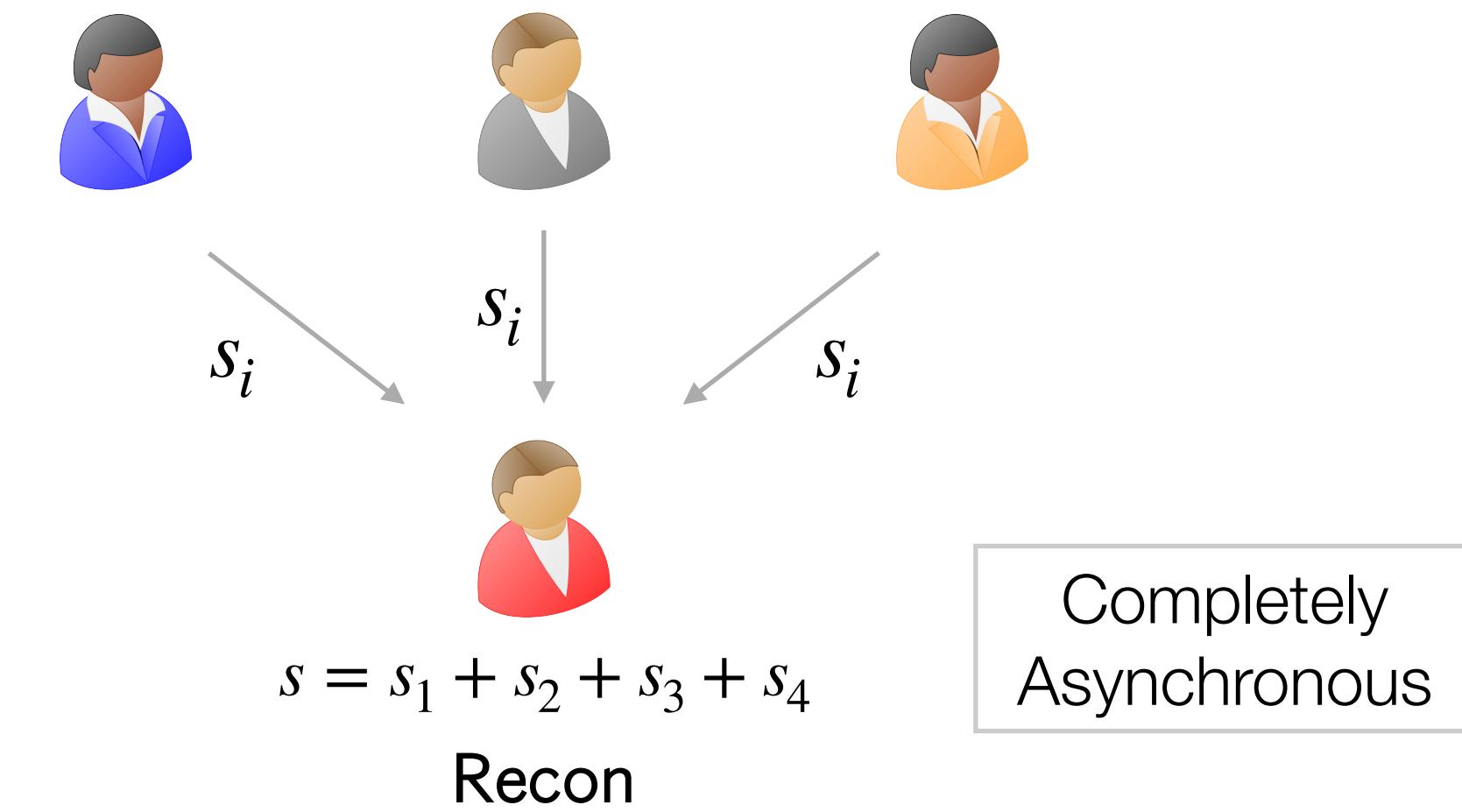
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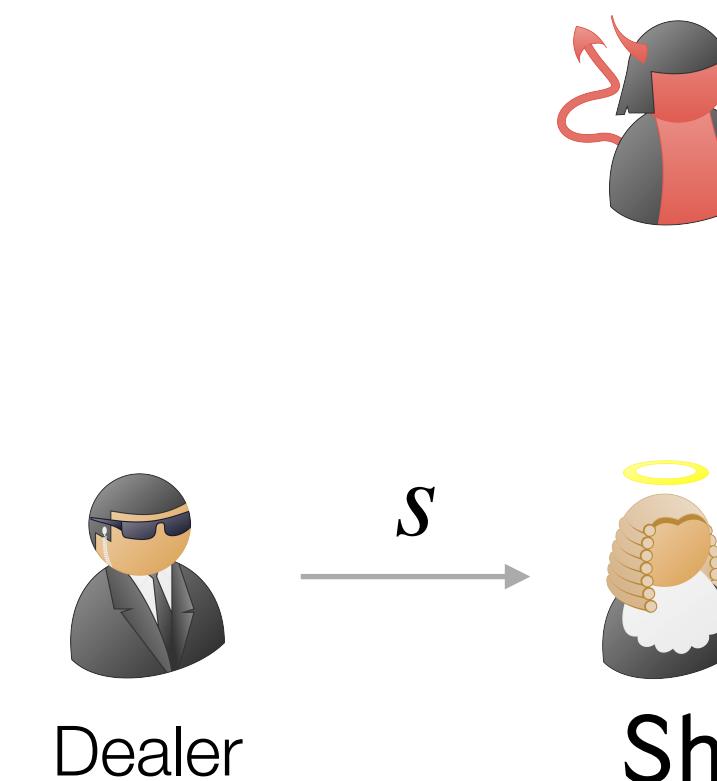
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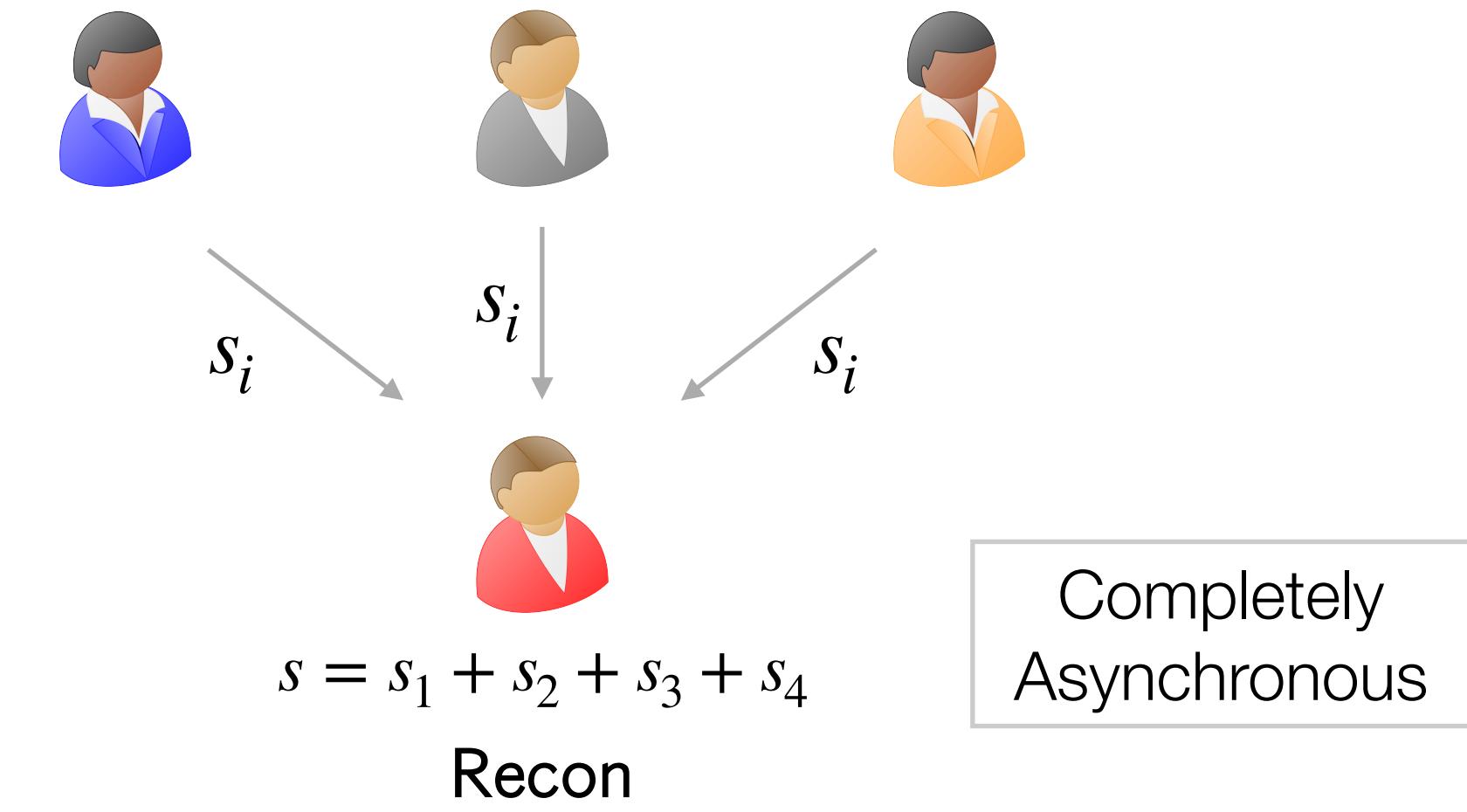
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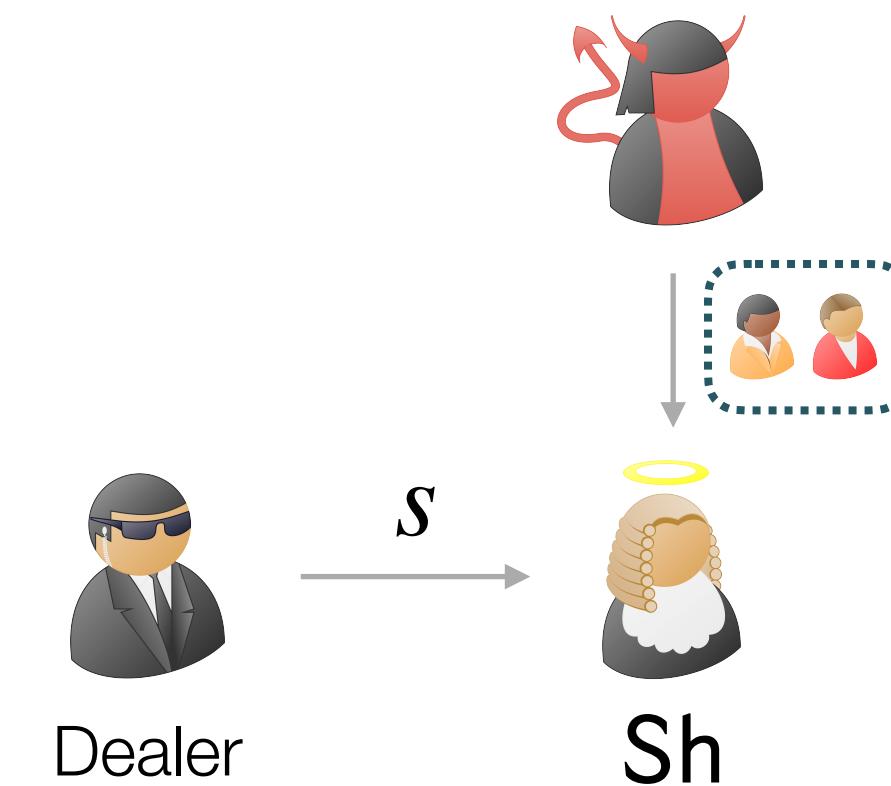
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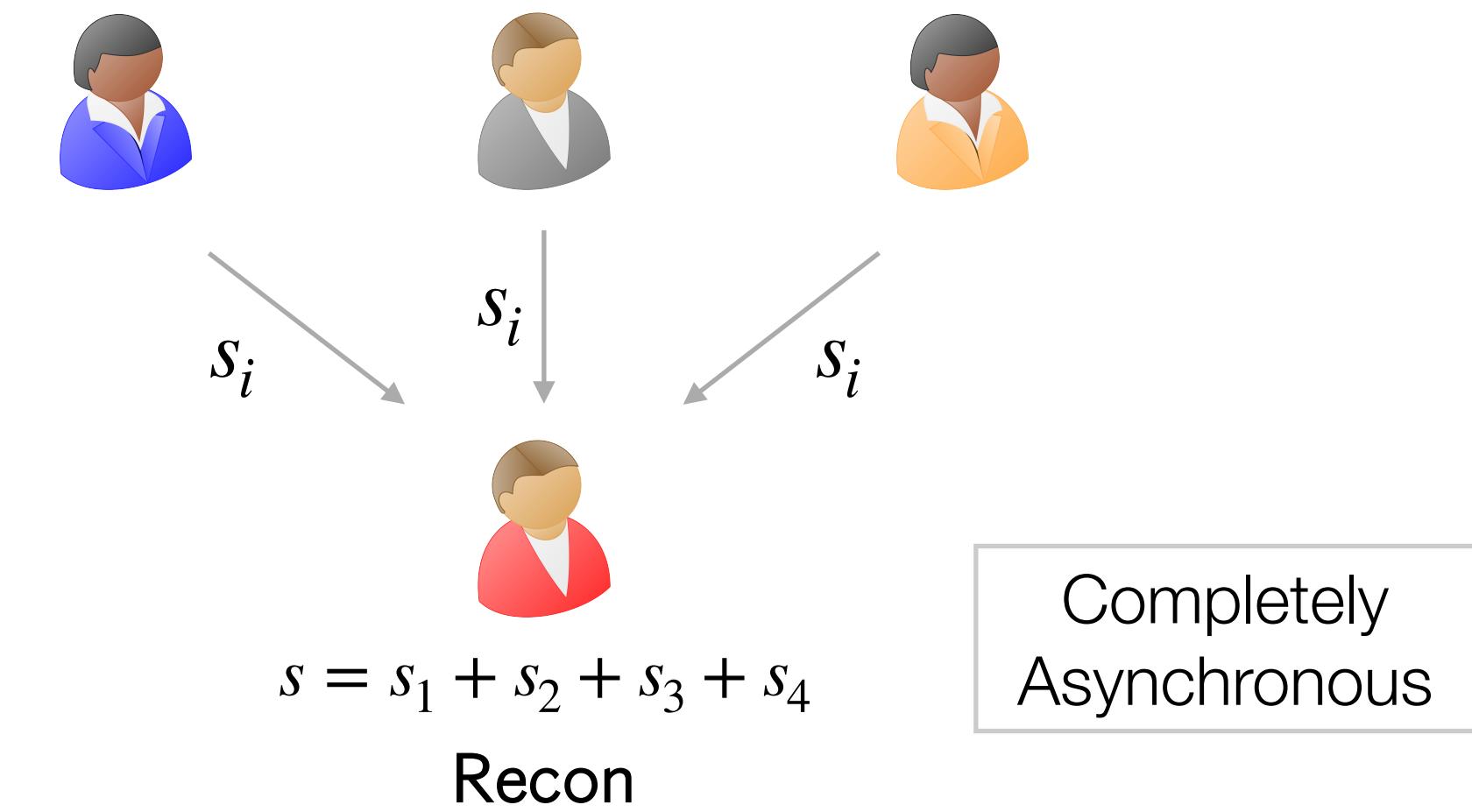
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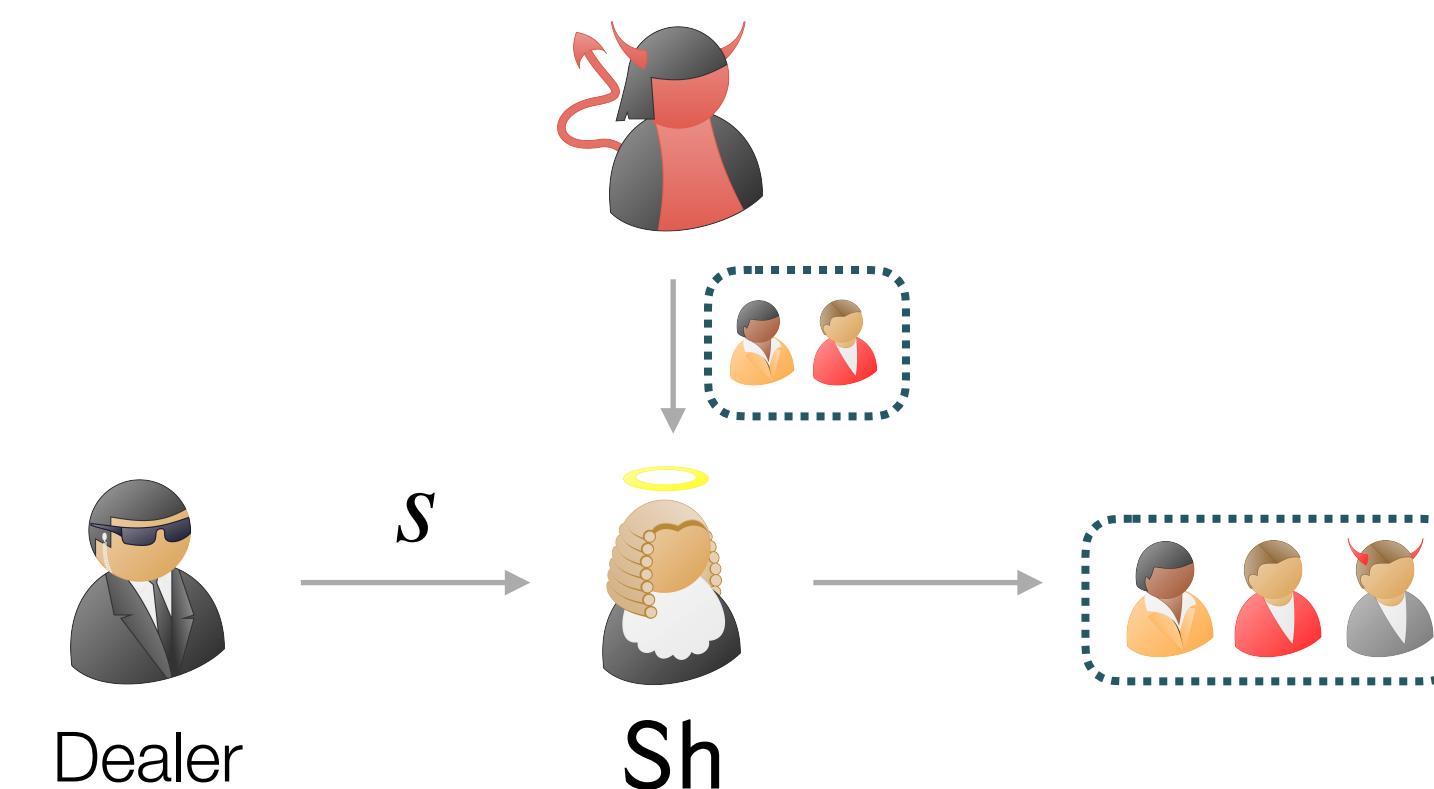
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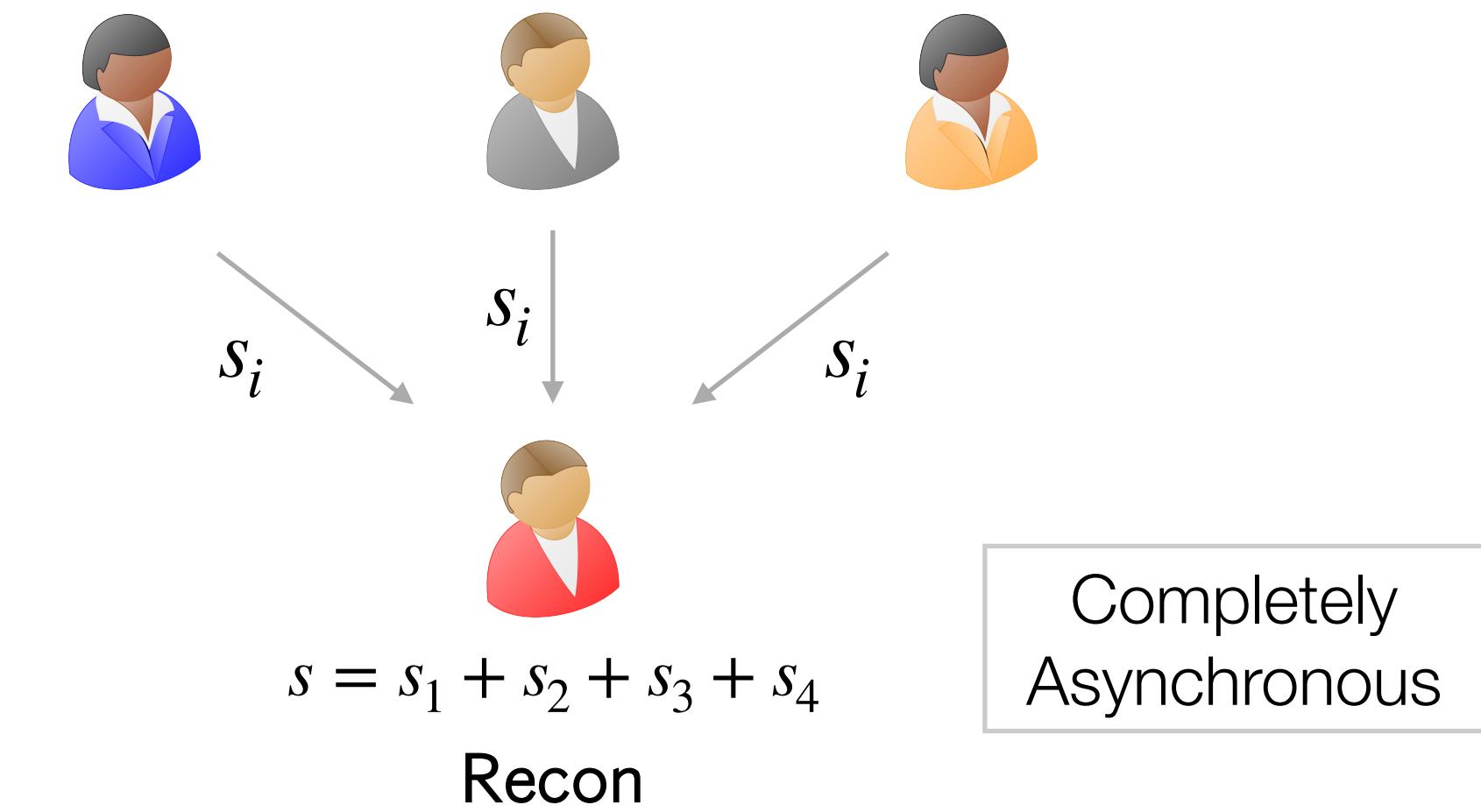
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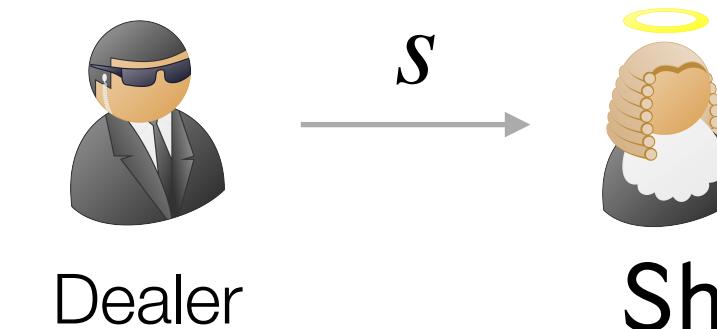
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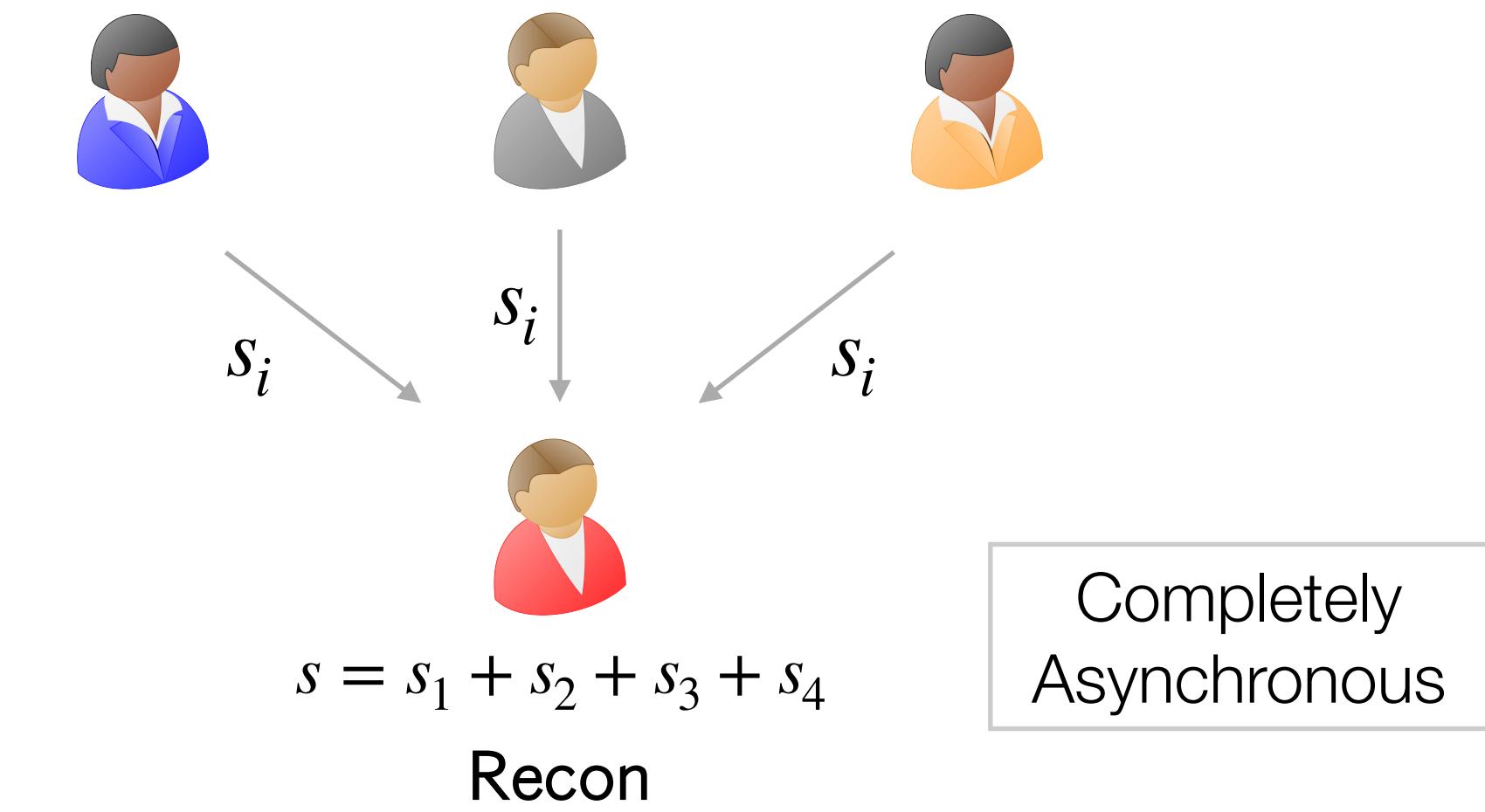
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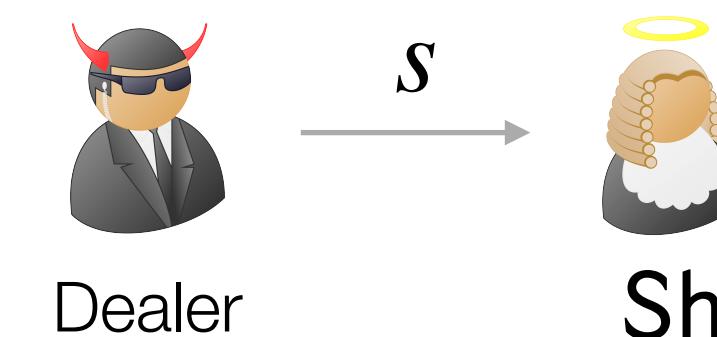
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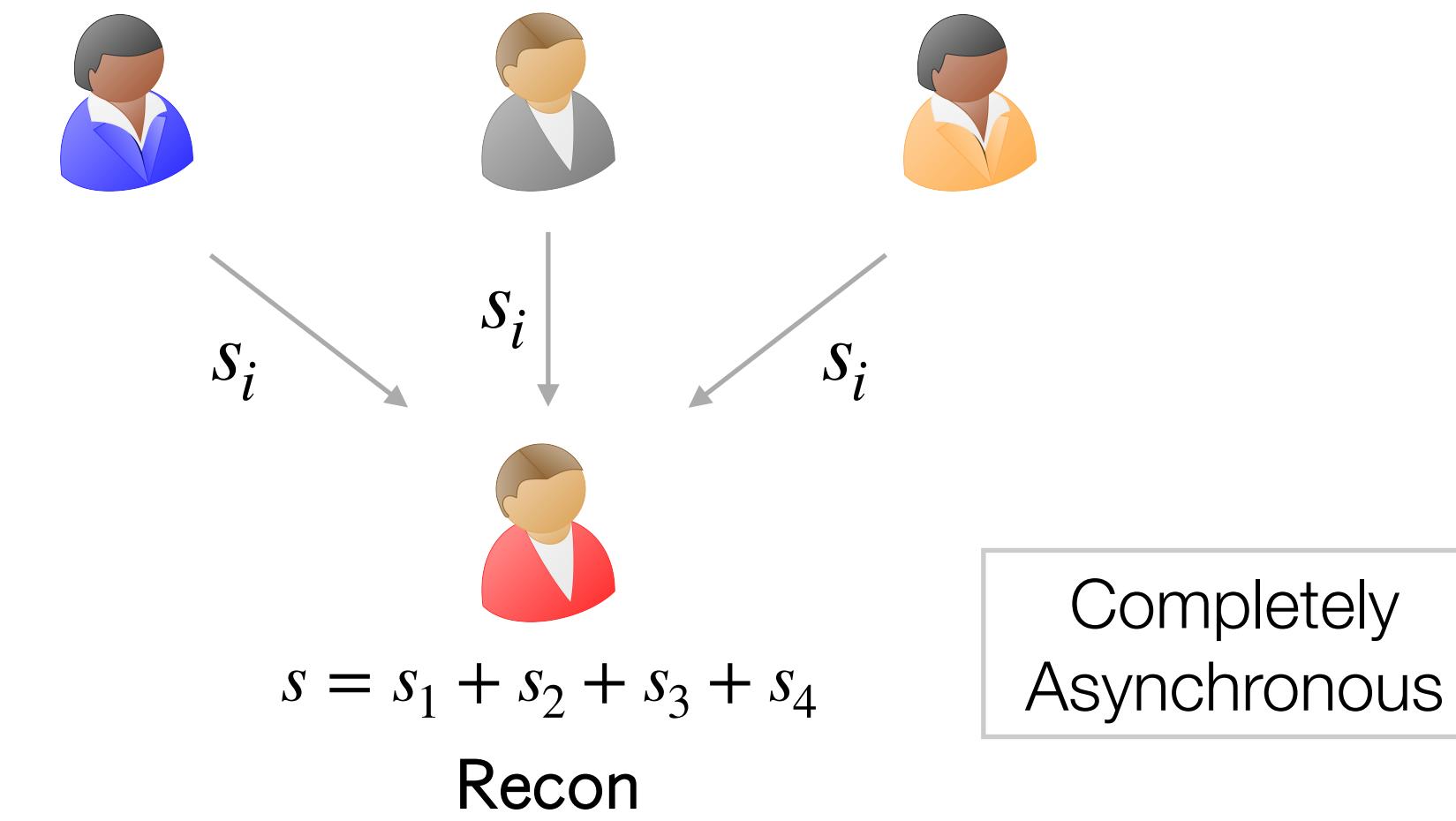
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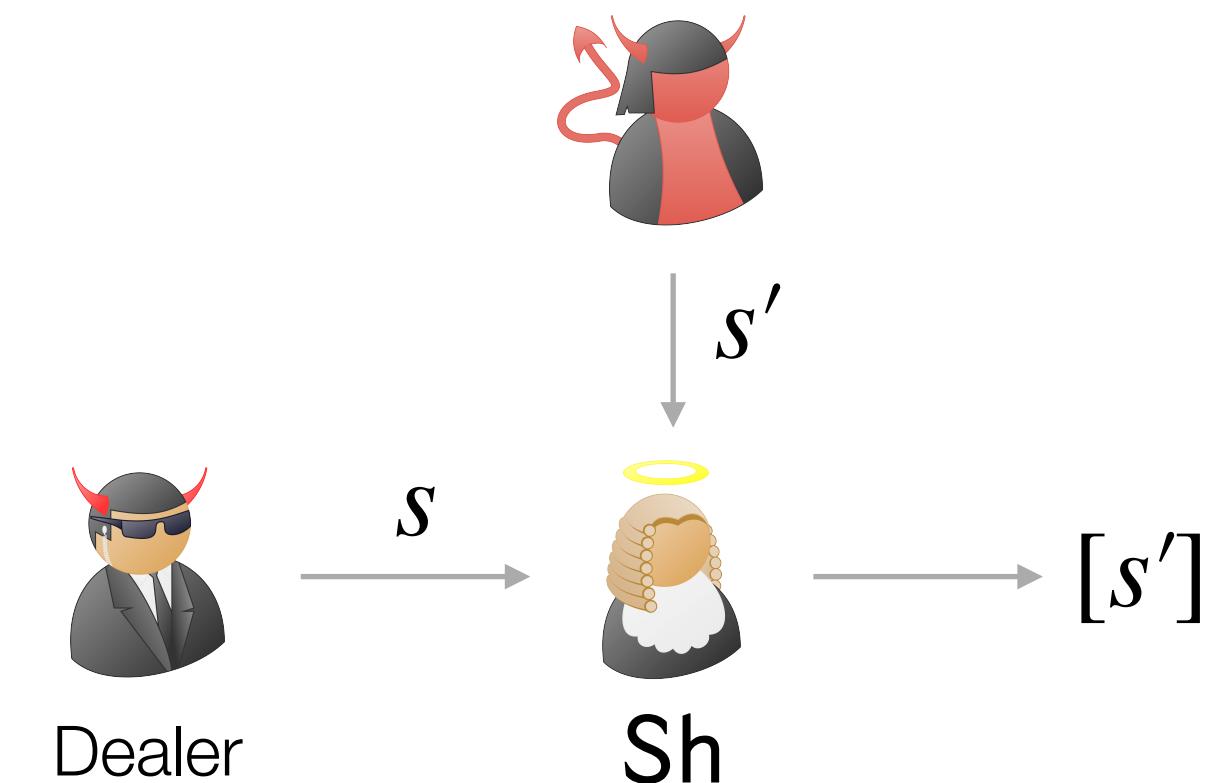
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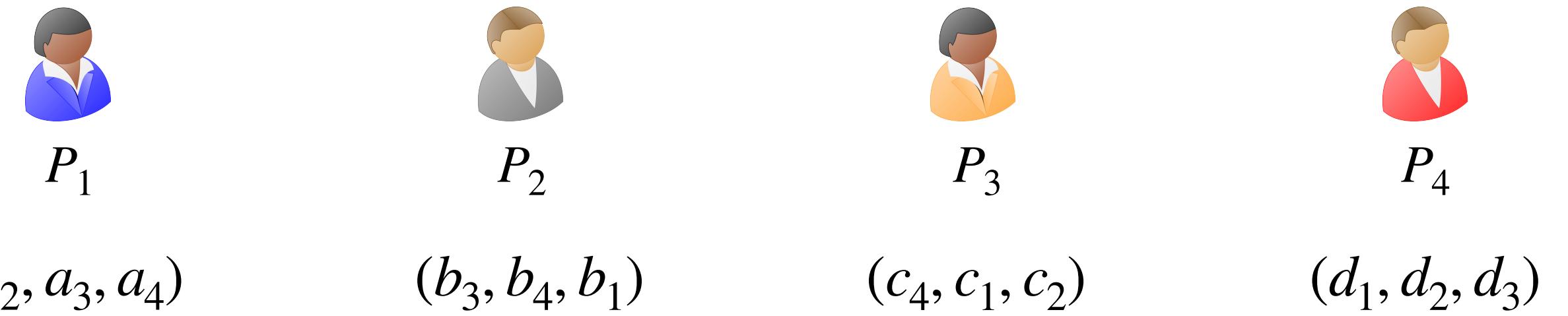
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- Round 1
 - D sends share to each party
 - Parties exchange **random pad** for each **common element** in share



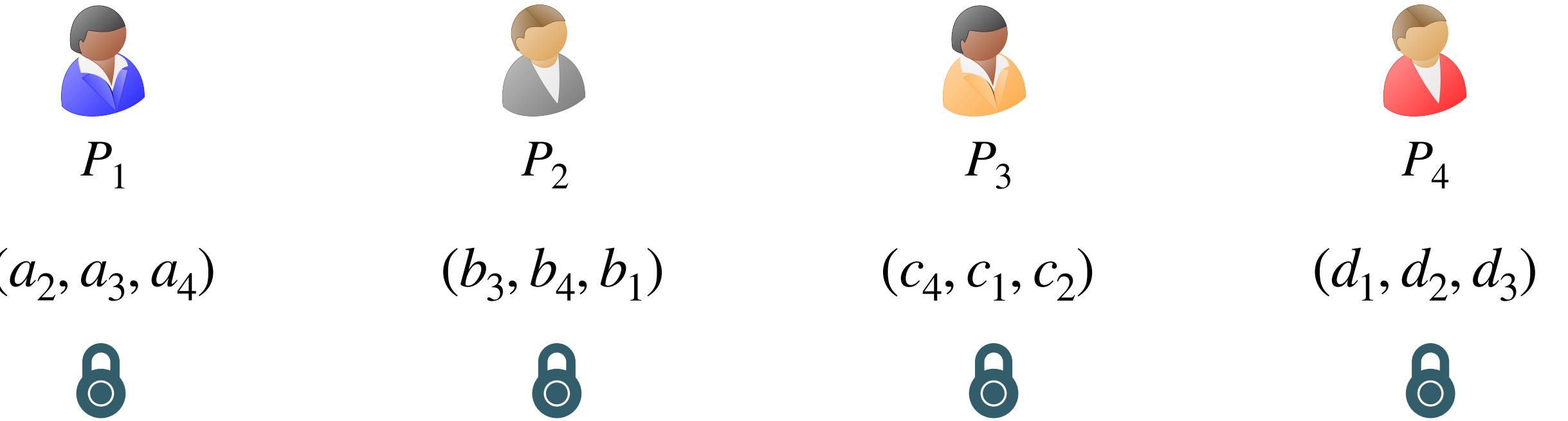
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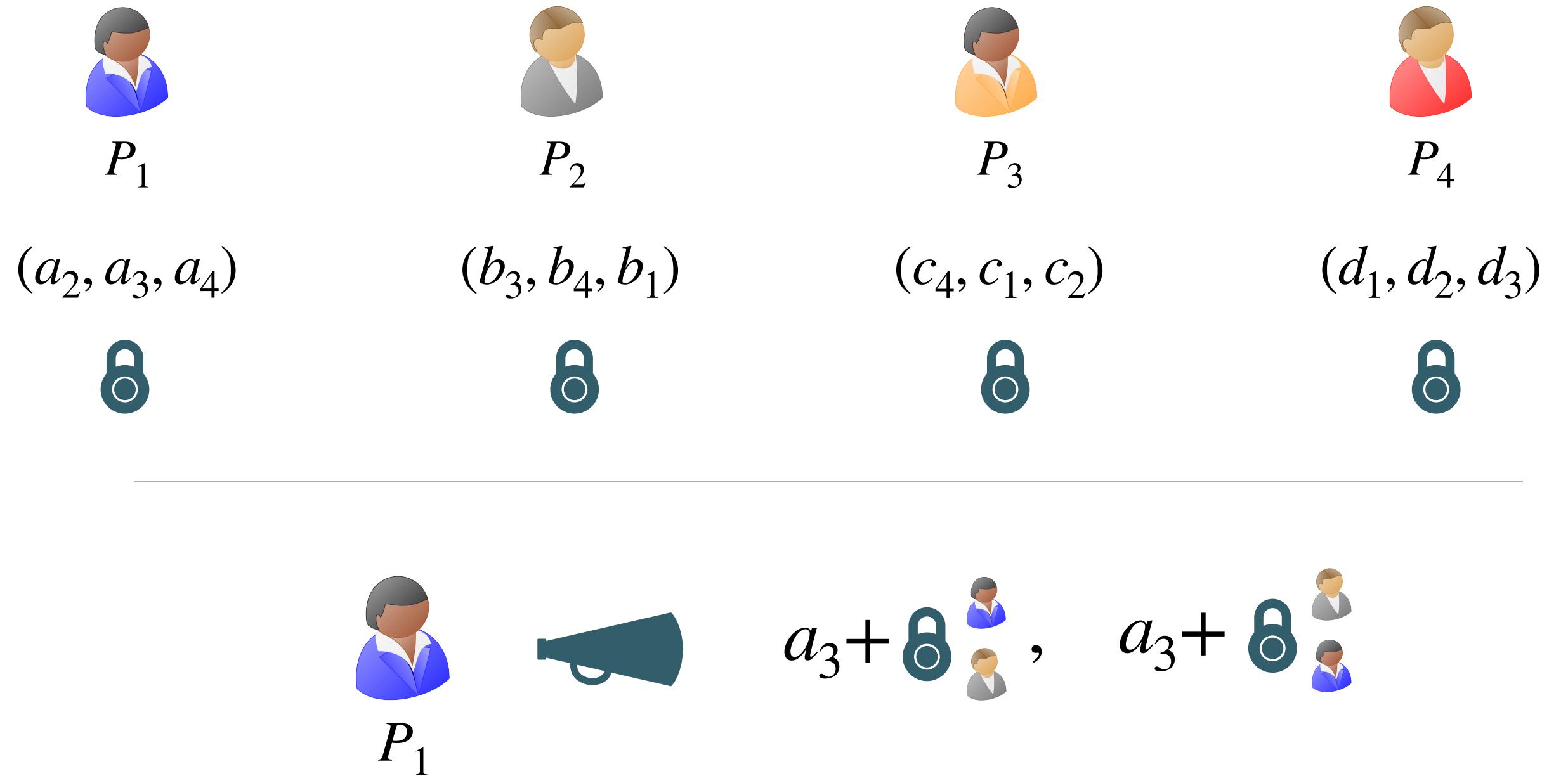
Perfect HMPC - VSS with Party Elimination

- Round 1
 - D sends share to each party
 - Parties exchange **random pad** for each **common element** in share



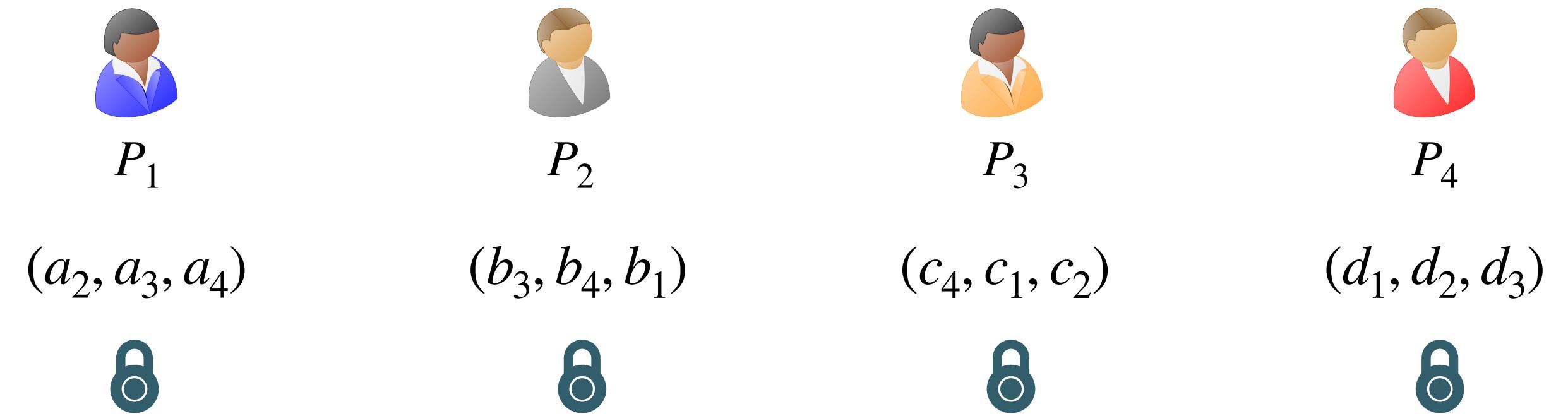
Perfect HMPC - VSS with Party Elimination

- Round 1
 - D sends share to each party
 - Parties exchange **random pad** for each **common element** in share
- Round 2
 - Parties broadcast **masked shares**



Perfect HMPC - VSS with Party Elimination

- Round 1
 - D sends share to each party
 - Parties exchange **random pad** for each **common element** in share
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- Local computation
 - If shares inconsistent, output dispute set
 - Else output with secret shares

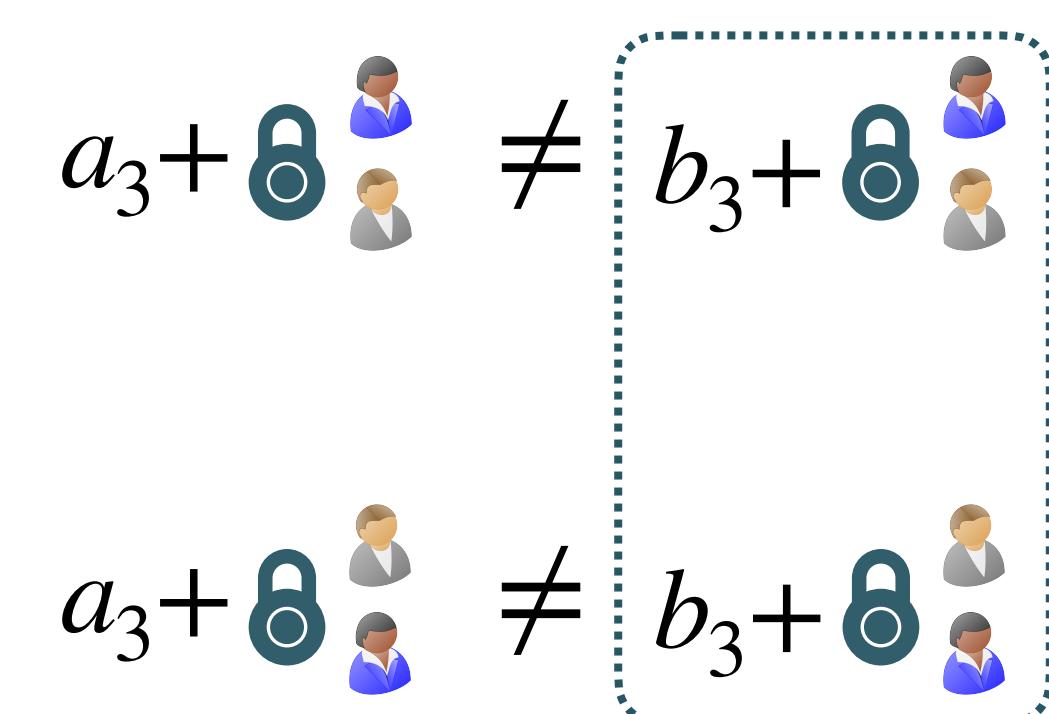
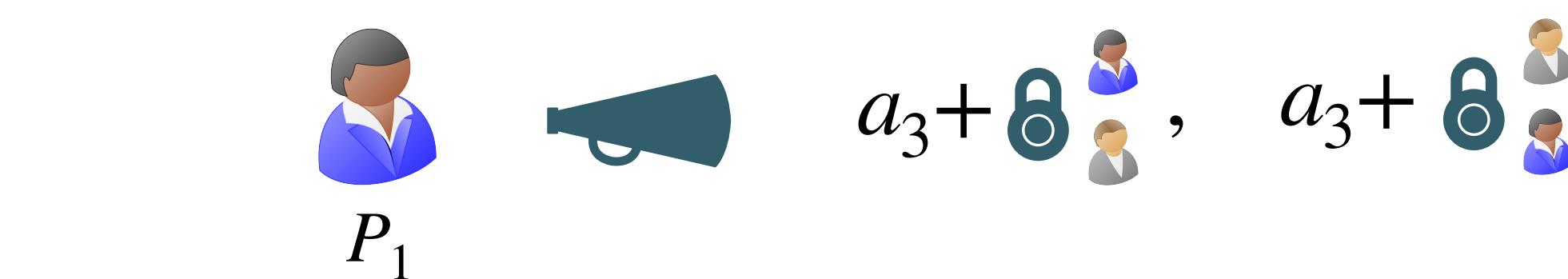
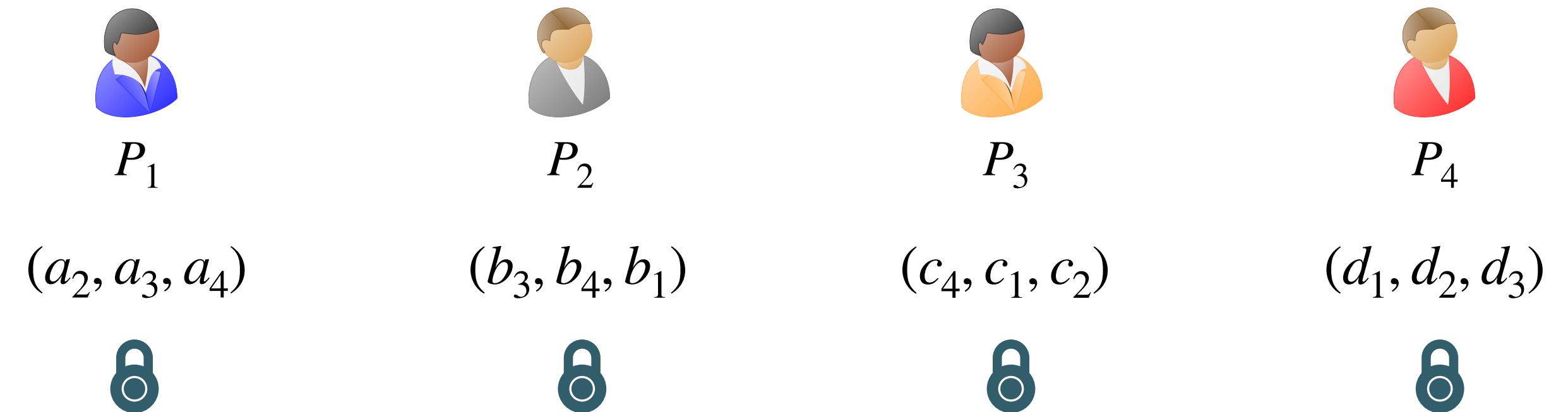


$$a_3 + \text{lock icon} \neq b_3 + \text{lock icon}$$

$$a_3 + \text{lock icon} \neq b_3 + \text{lock icon}$$

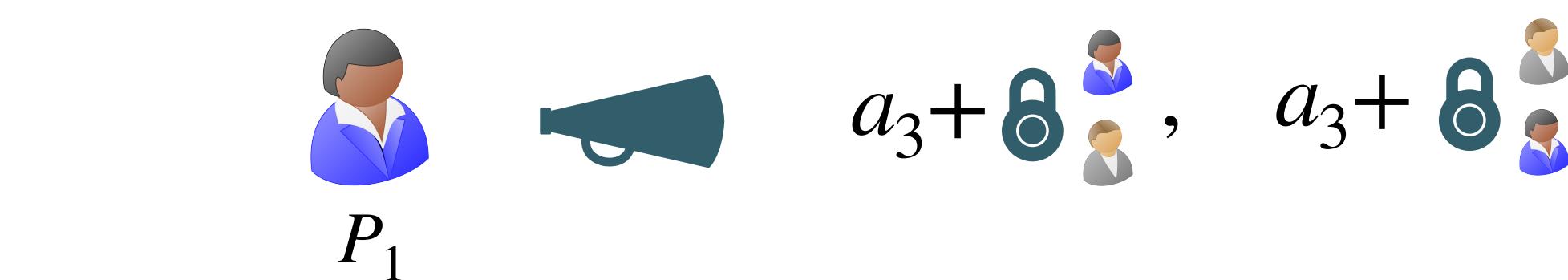
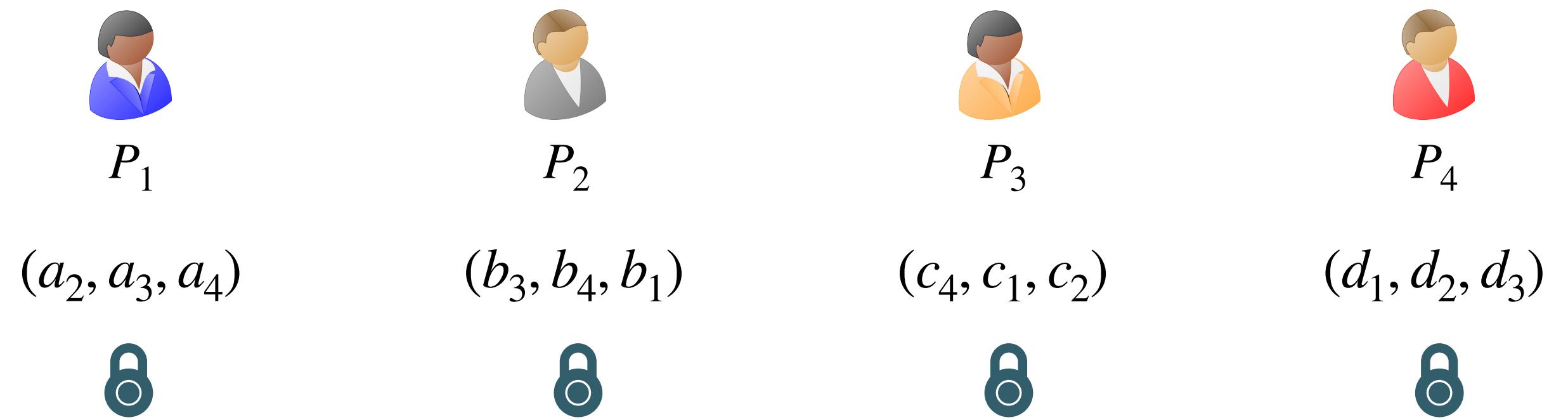
Perfect HMPC - VSS with Party Elimination

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Perfect HMPC - VSS with Party Elimination

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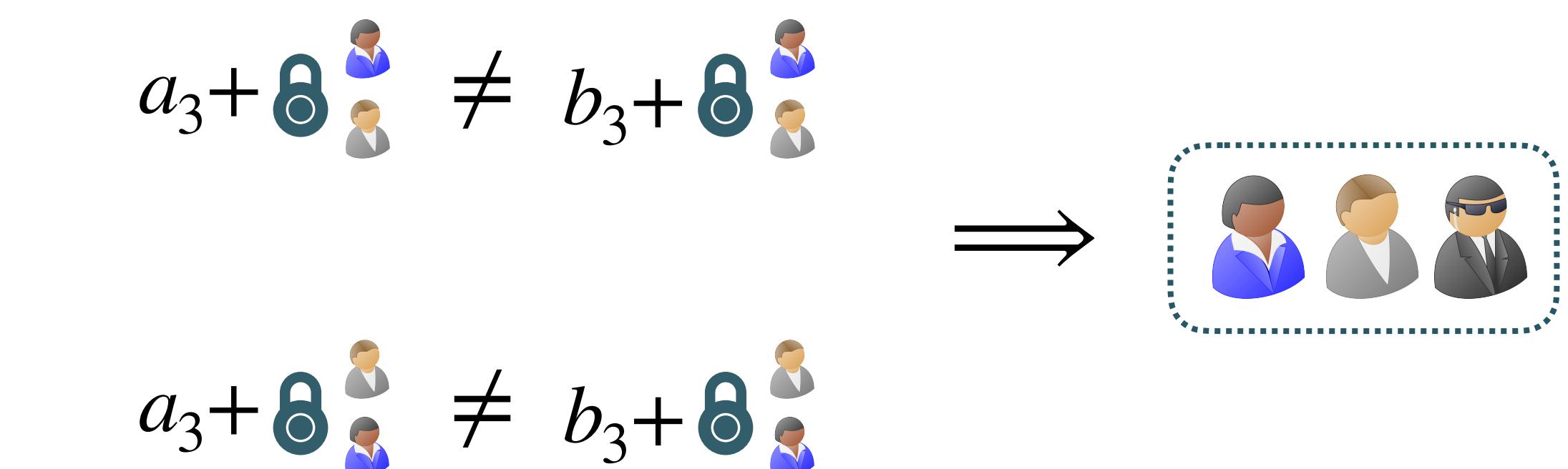
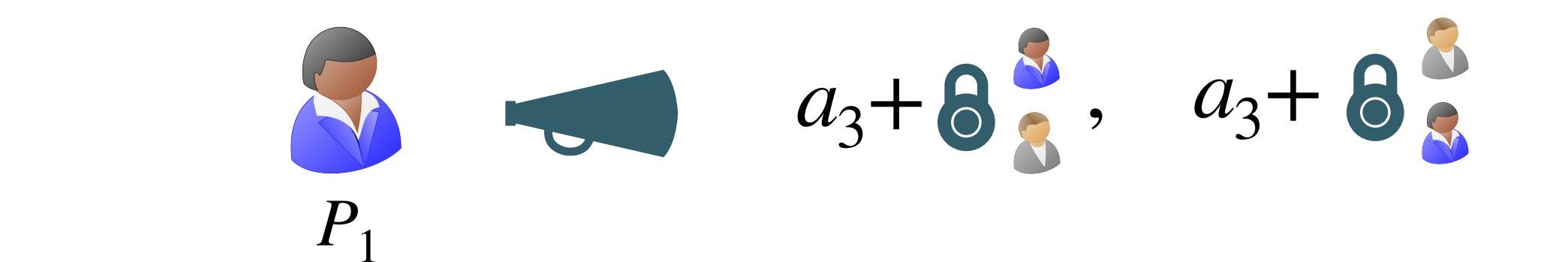
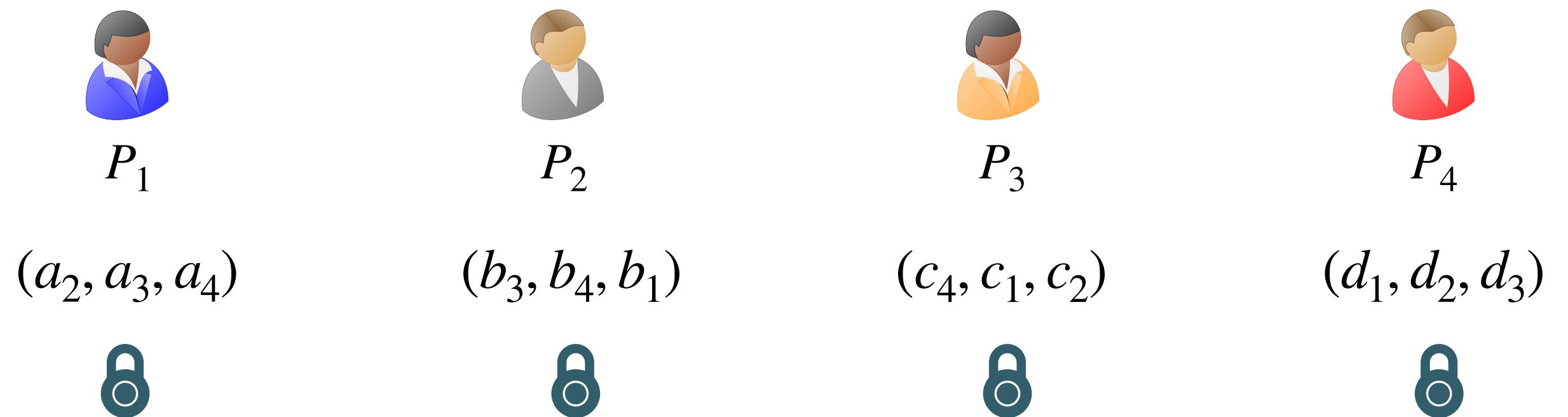


$$a_3 + \text{lock icon} \neq b_3 + \text{lock icon}$$

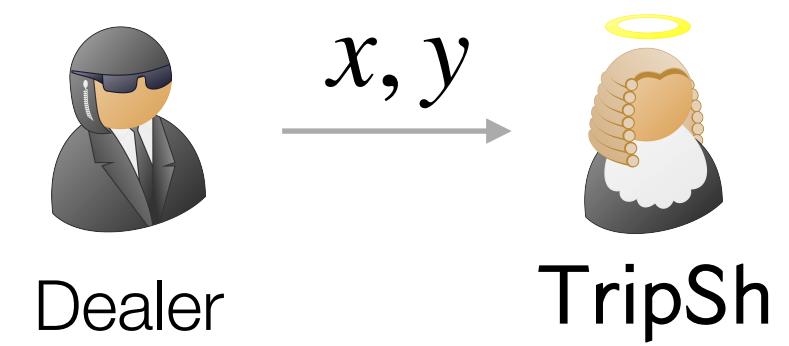
This equation compares two masked shares. On the left is $a_3 + \text{lock icon}$, and on the right is $b_3 + \text{lock icon}$. Both lock icons contain two party icons (one blue, one grey). The inequality sign (\neq) indicates that these shares are inconsistent, which triggers a dispute set.

Perfect HMPC - VSS with Party Elimination

- Round 1
 - D sends share to each party
 - Parties exchange **random pad** for each **common element** in share
- Round 2
 - Parties broadcast **masked shares**
- Local computation
 - If shares inconsistent, output dispute set
 - Else output with secret shares



Perfect HMPC - Triple Sharing with Party Elimination Functionality



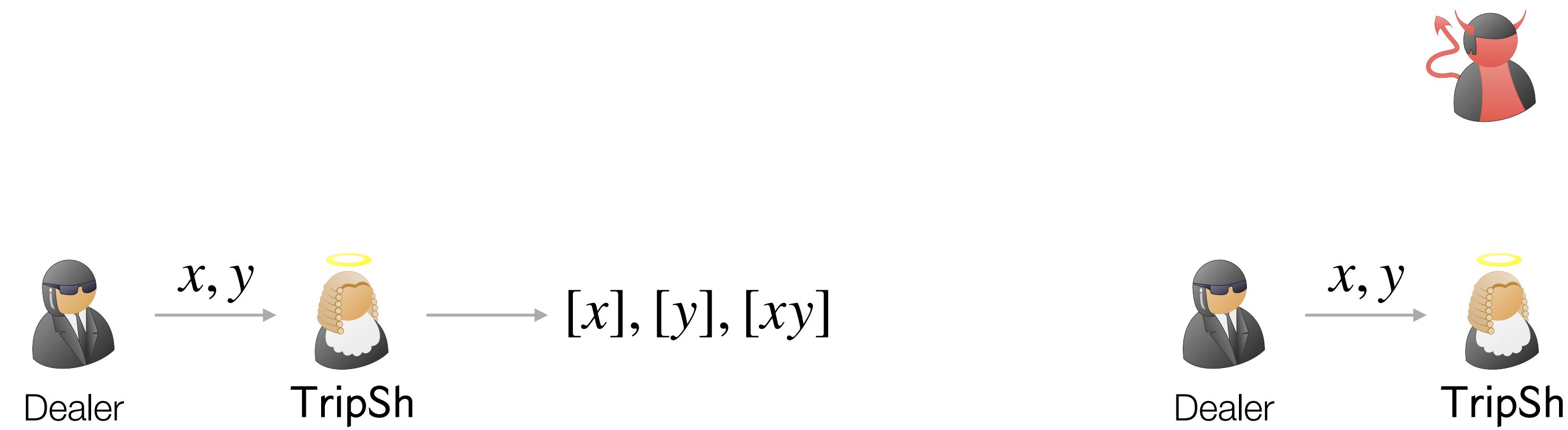
- Triple sharing with Party Elimination
 - Verified multiplication triple or dispute set

Perfect HMPC - Triple Sharing with Party Elimination Functionality



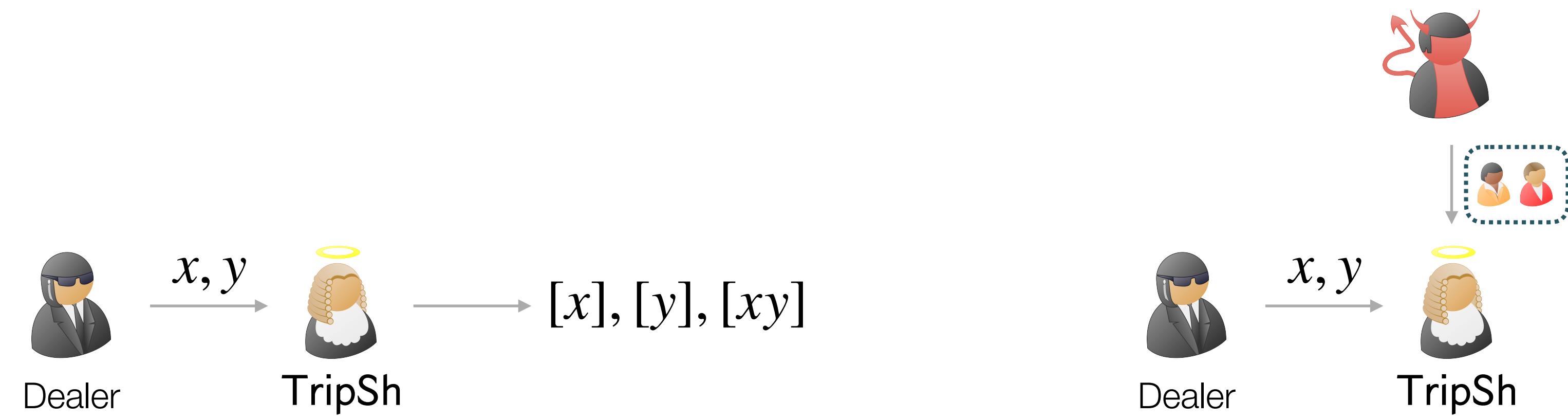
- Triple sharing with Party Elimination
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Perfect HMPC - Triple Sharing with Party Elimination Functionality



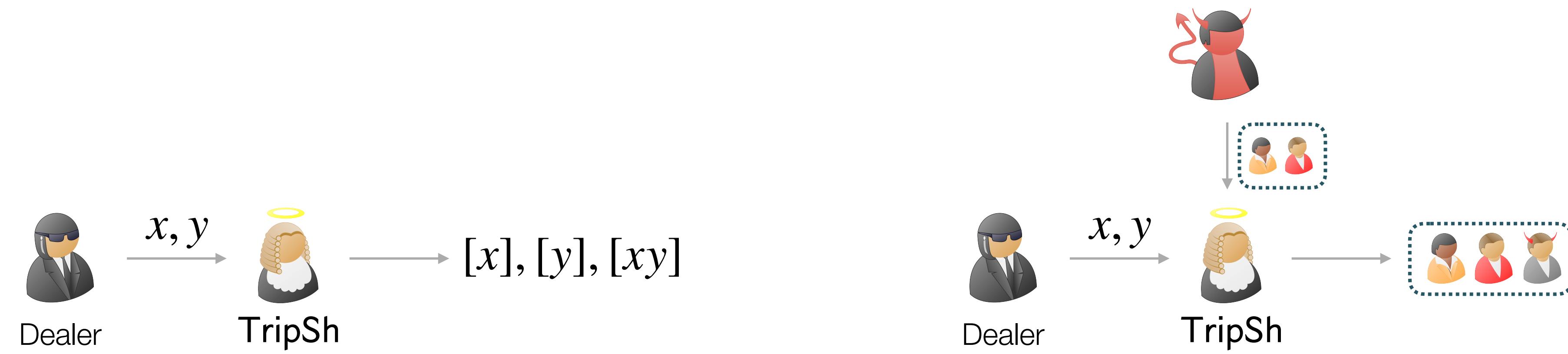
- Triple sharing with Party Elimination
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Perfect HMPC - Triple Sharing with Party Elimination Functionality



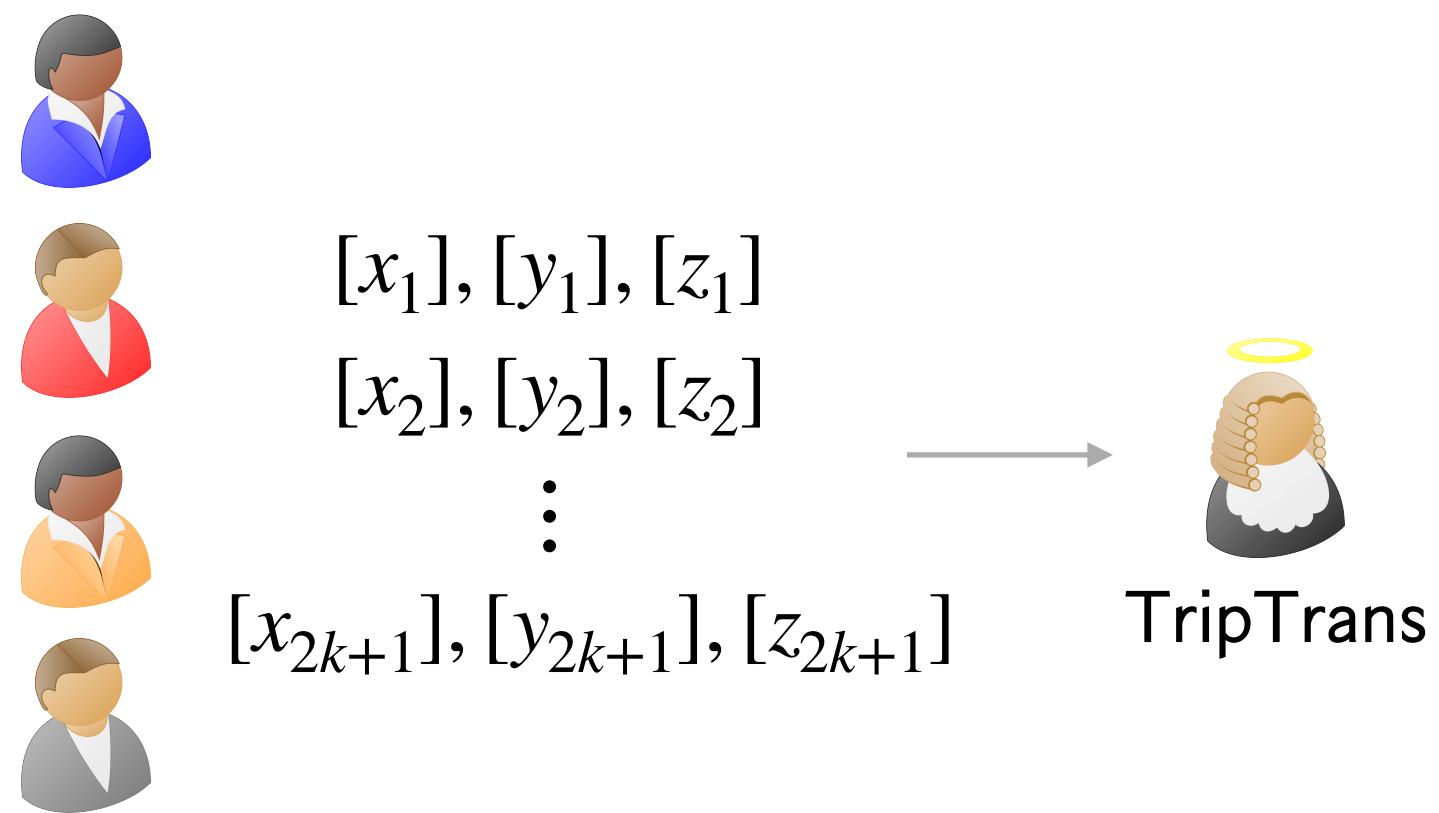
- Triple sharing with Party Elimination
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Perfect HMPC - Triple Sharing with Party Elimination Functionality



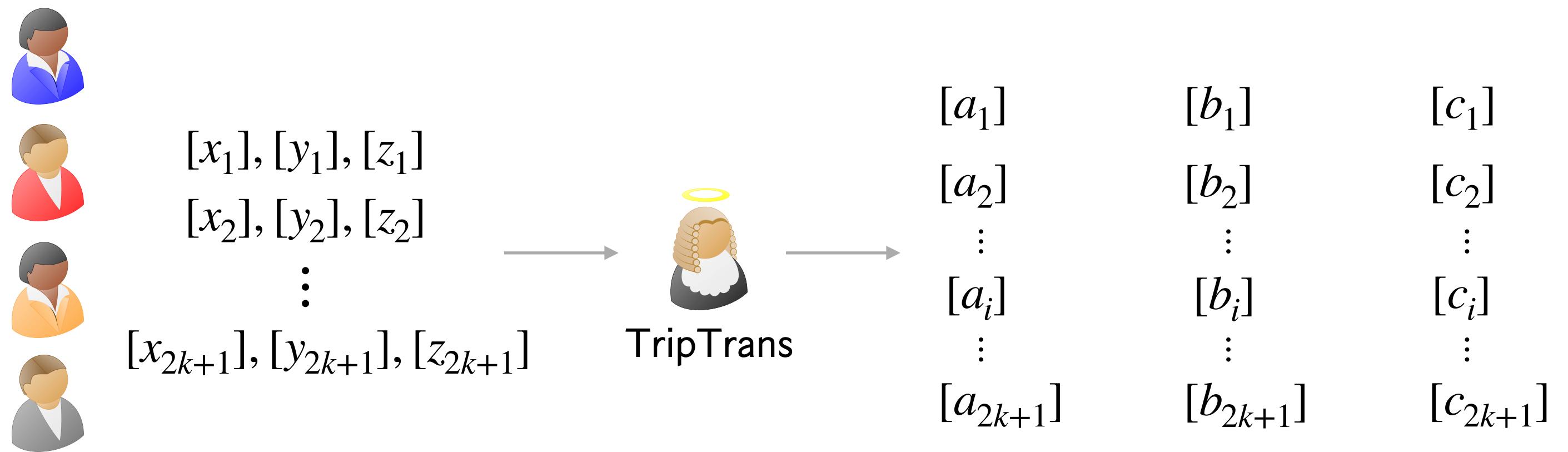
- Triple sharing with Party Elimination
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Perfect HMPC - Triple Transform Functionality



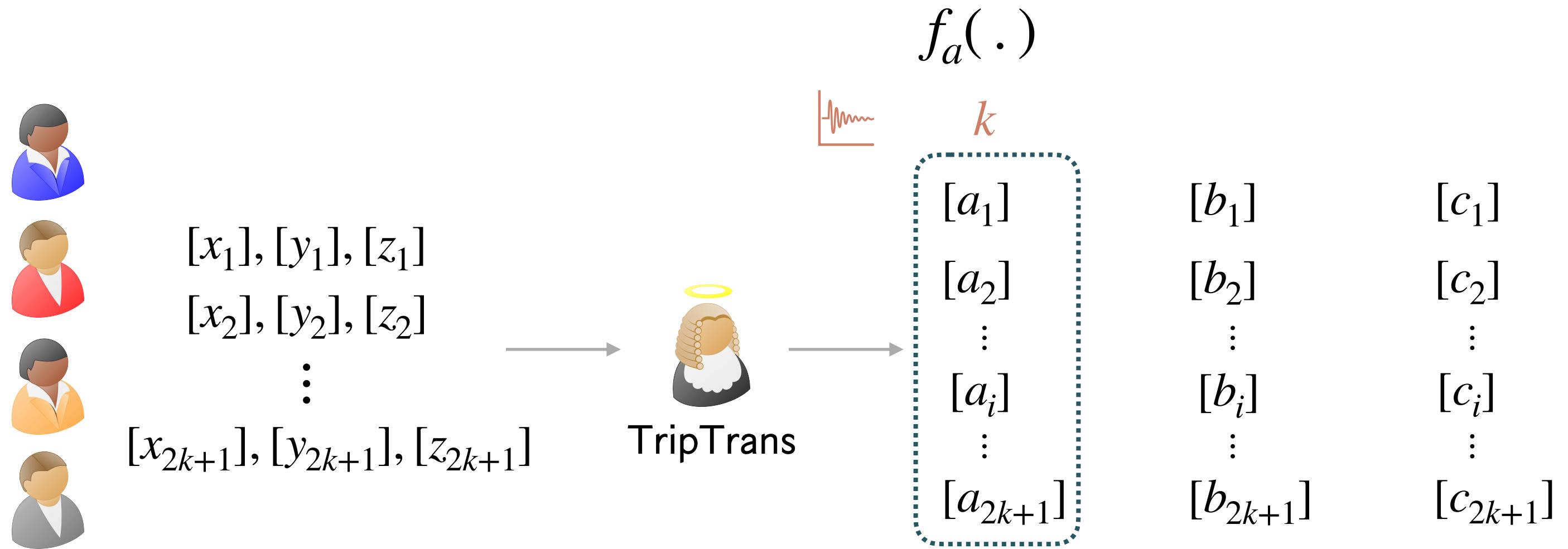
- Random triples → correlated random triples

Perfect HMPC - Triple Transform Functionality



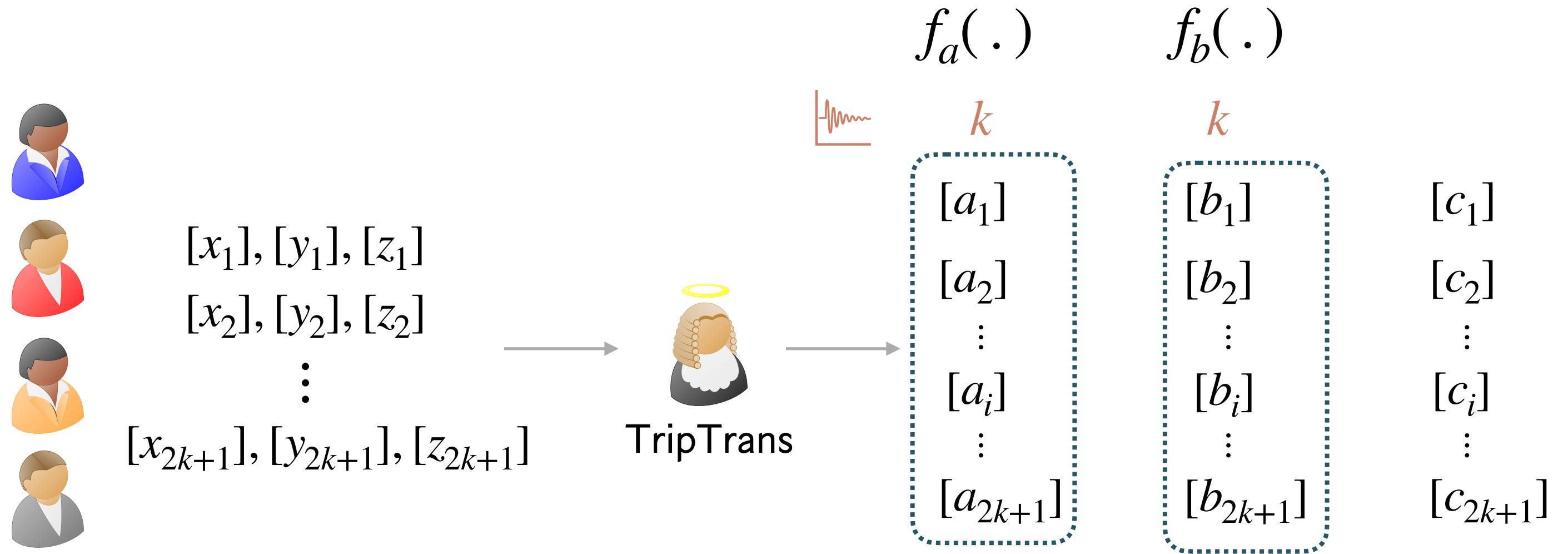
- Random triples → correlated random triples

Perfect HMPC - Triple Transform Functionality



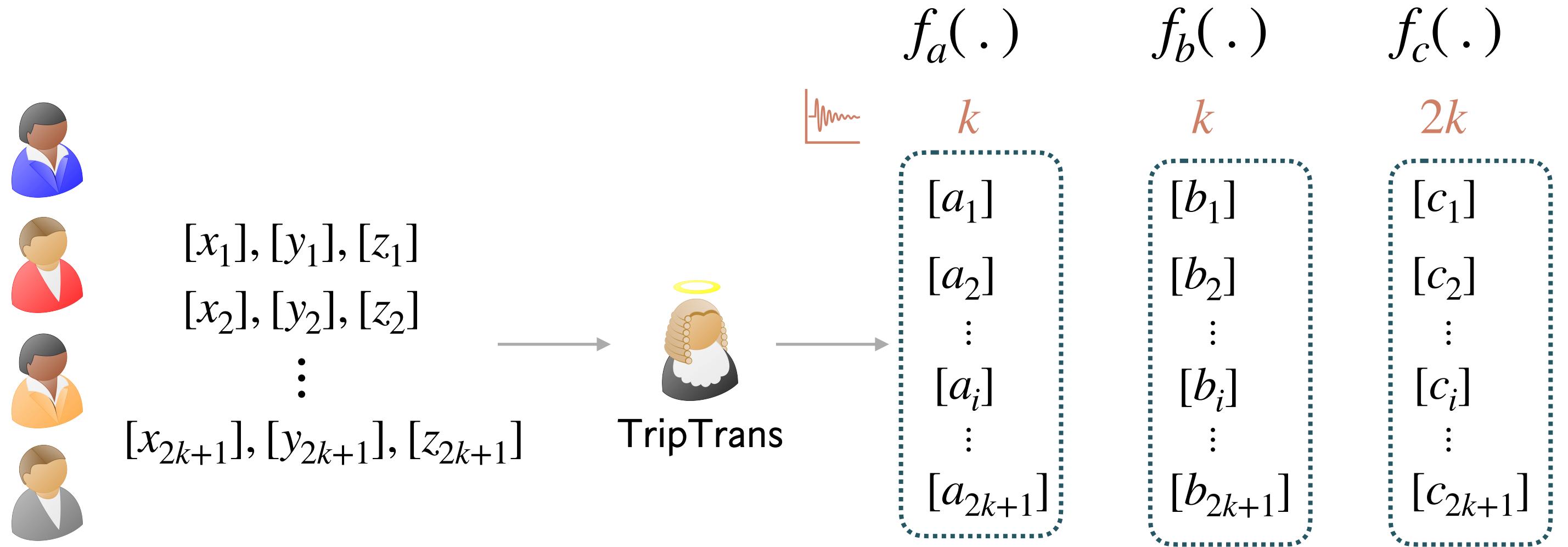
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



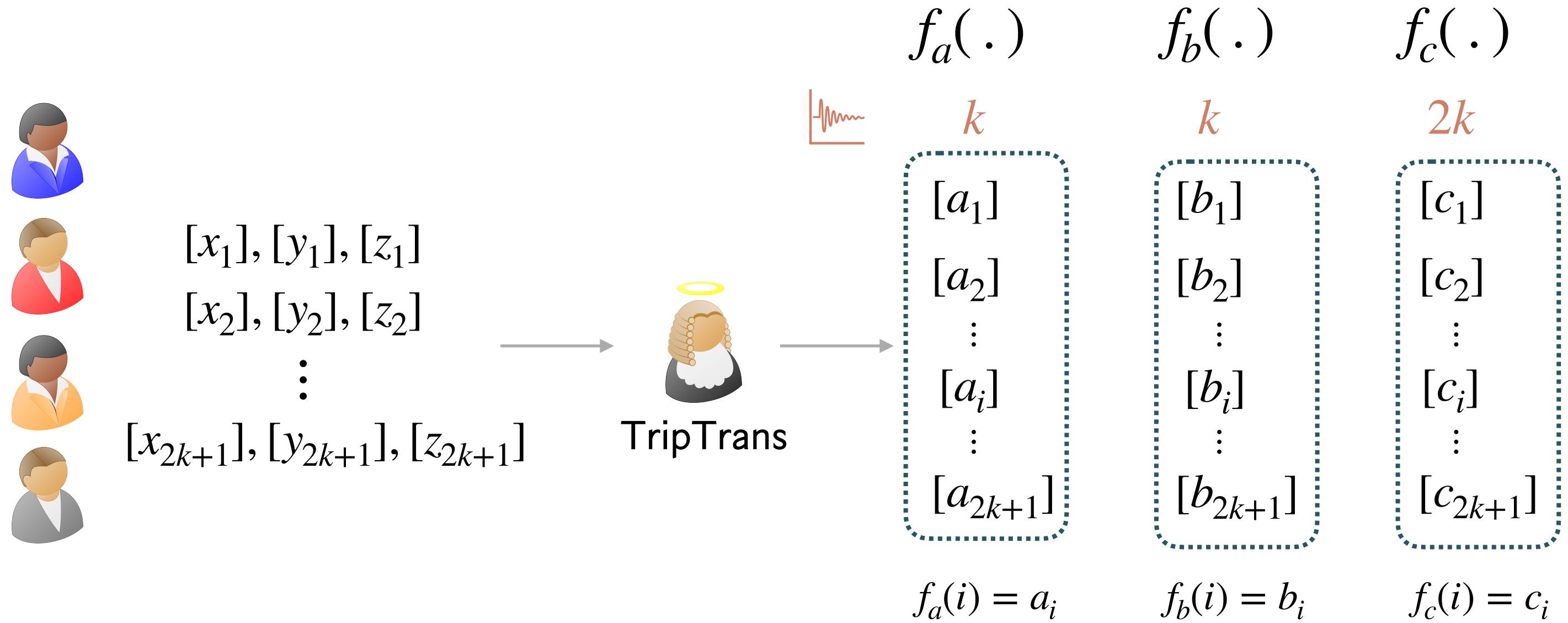
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



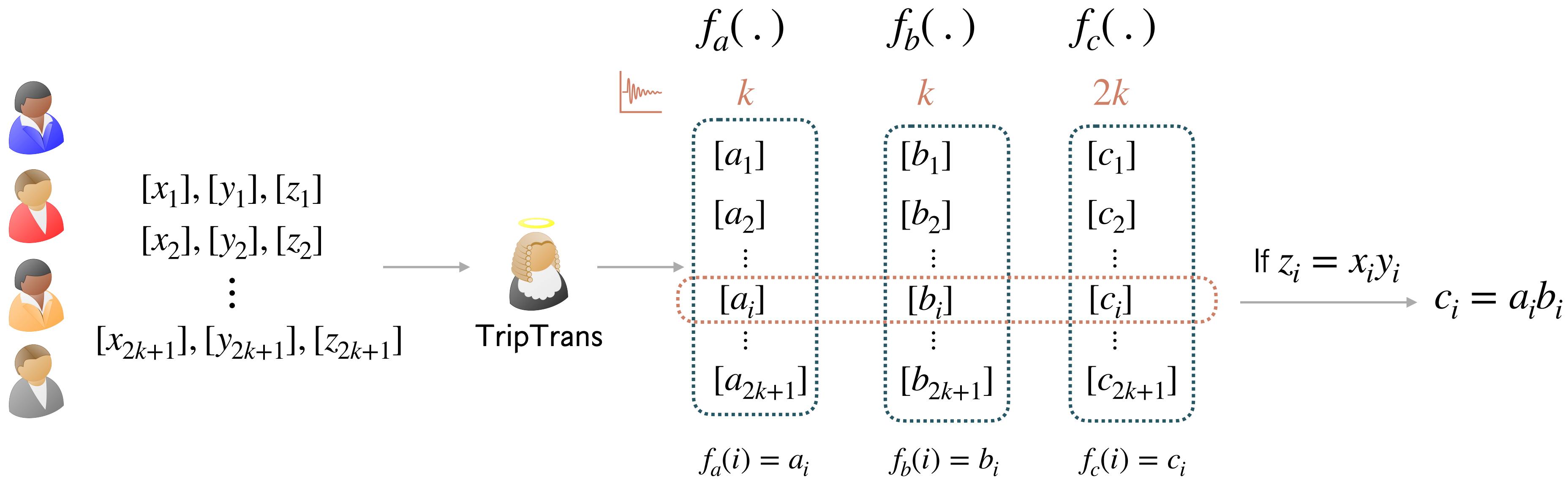
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



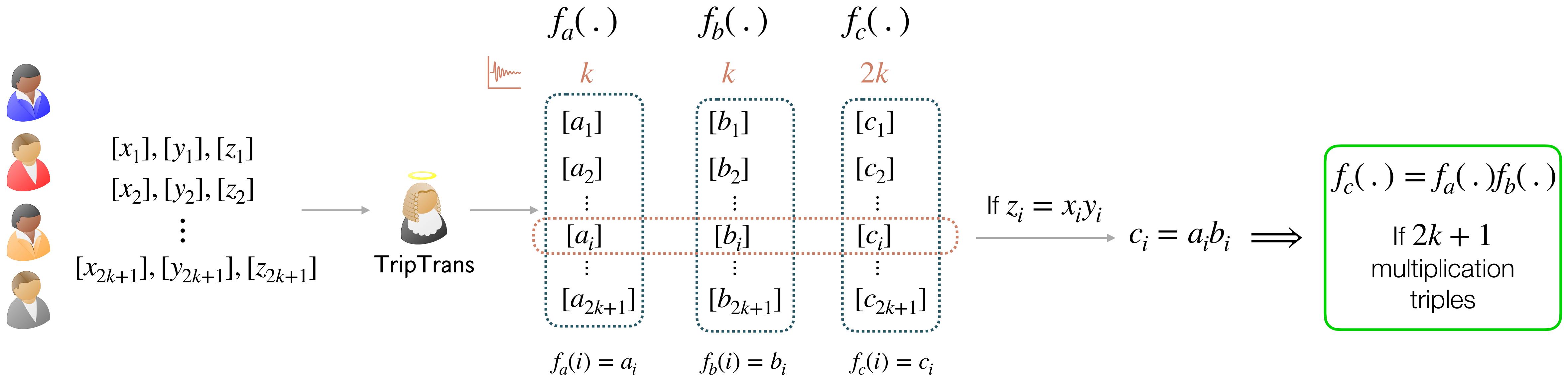
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Perfect HMPC - Triple Transform Functionality



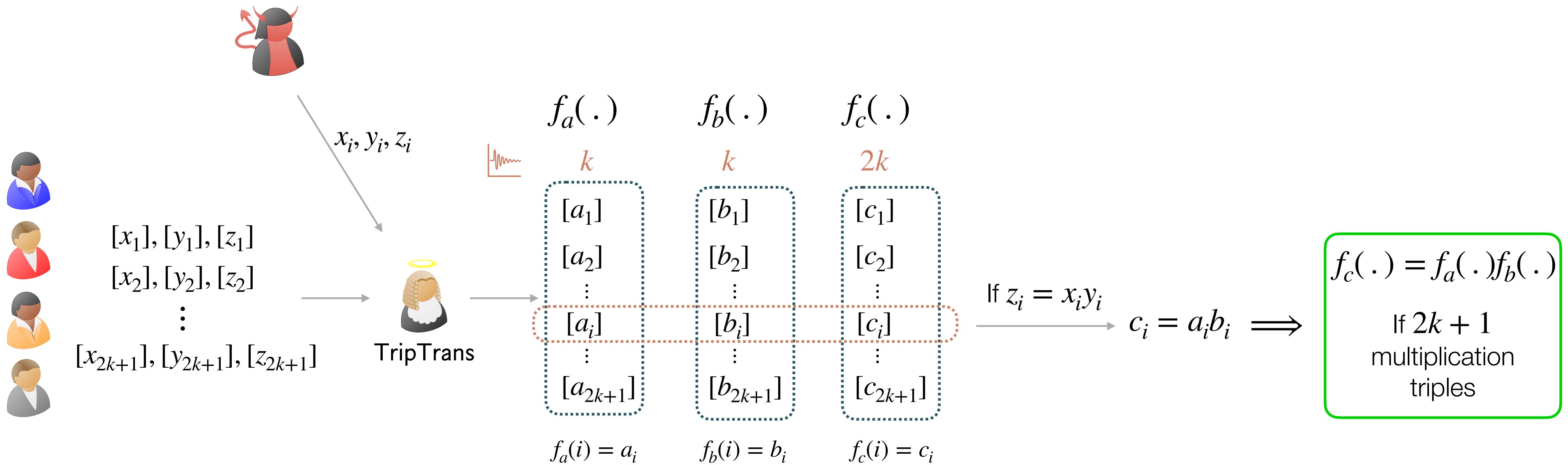
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



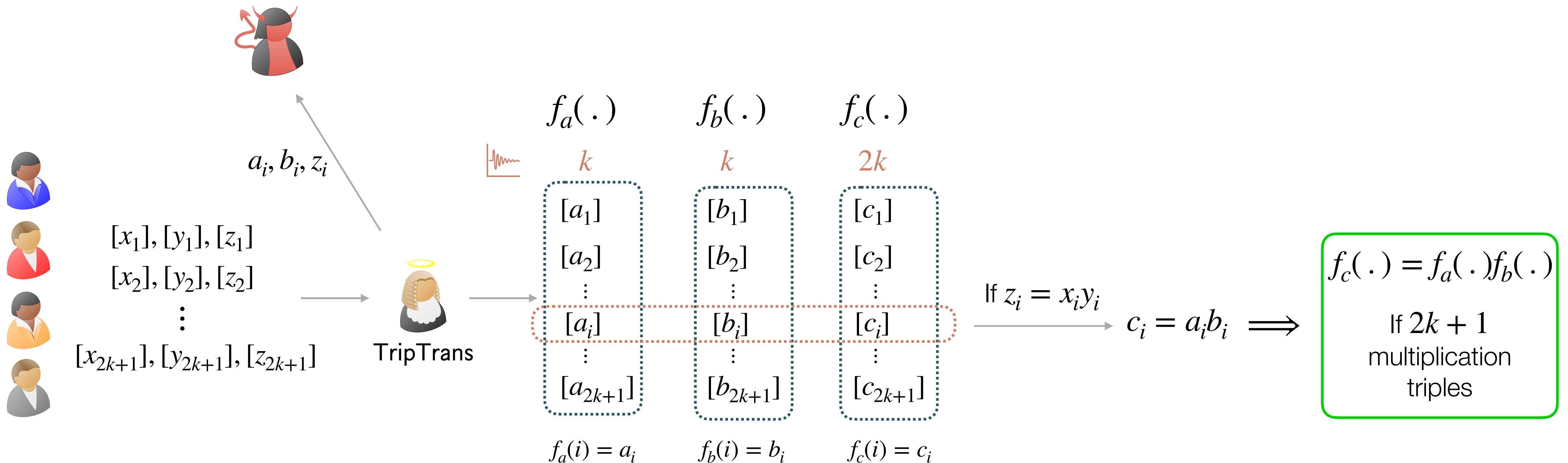
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



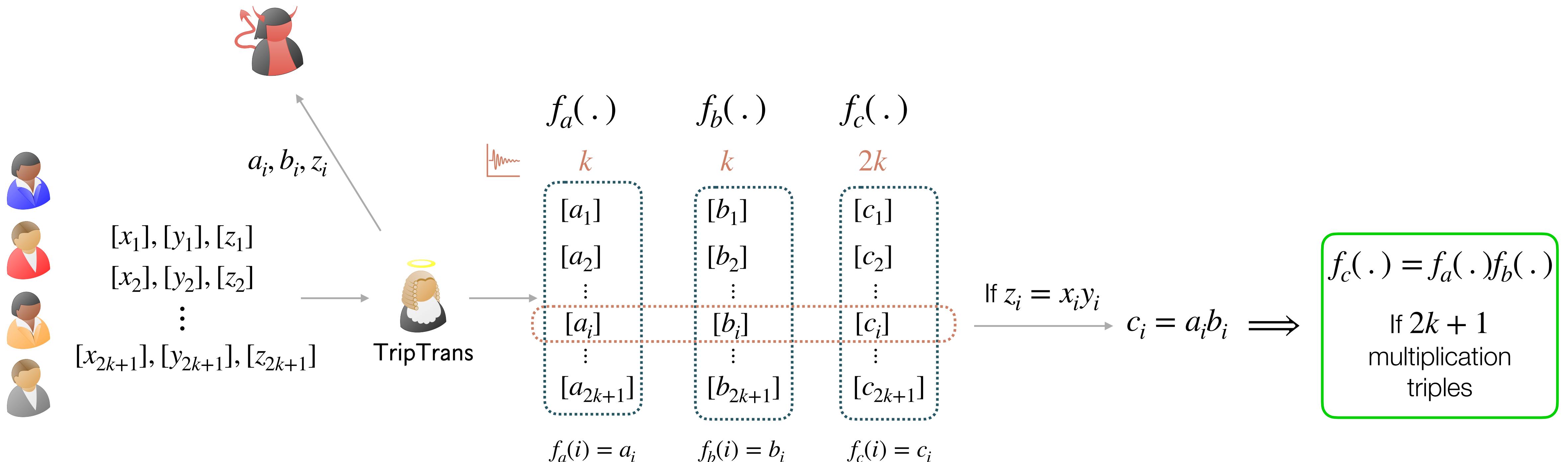
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality

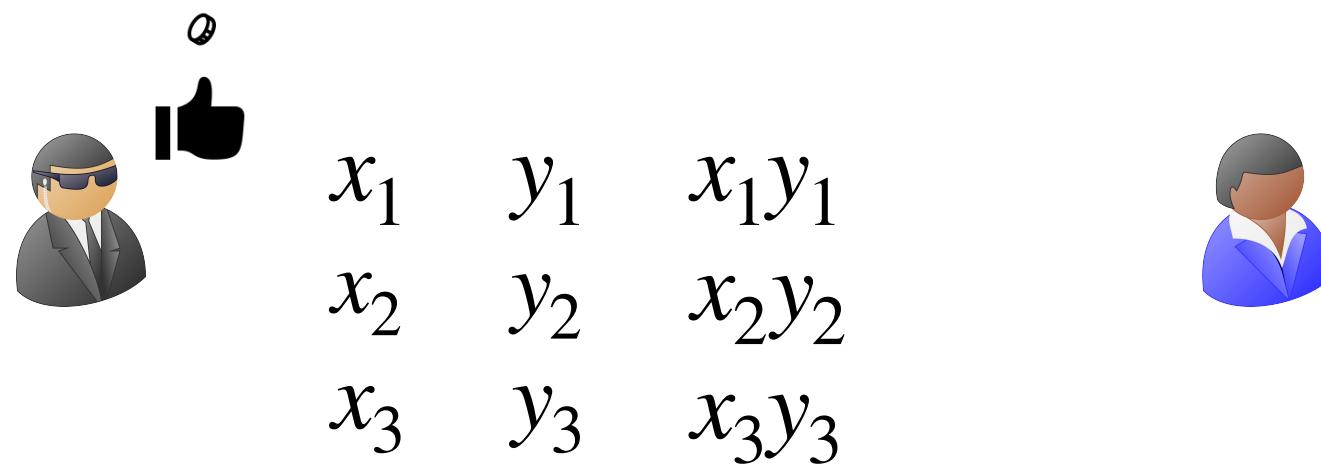


- Random triples \rightarrow correlated random triples
- Completely asynchronous instantiation in [CP17]

Perfect HMPC - Triple Sharing with Party Elimination Protocol



Perfect HMPC - Triple Sharing with Party Elimination Protocol



Perfect HMPC - Triple Sharing with Party Elimination Protocol

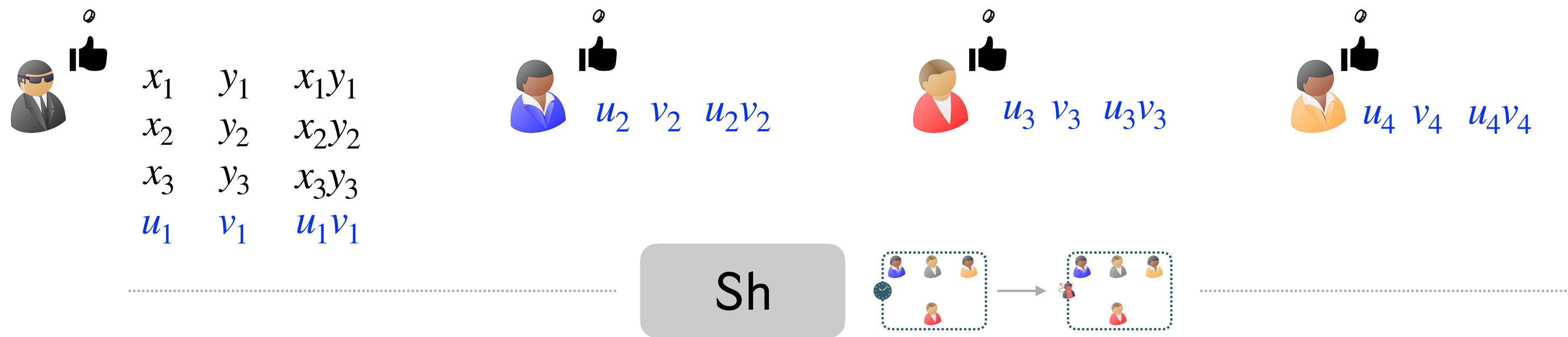
①
 $x_1 \quad y_1 \quad x_1y_1$
 $x_2 \quad y_2 \quad x_2y_2$
 $x_3 \quad y_3 \quad x_3y_3$
 $u_1 \quad v_1 \quad u_1v_1$

②
 $u_2 \quad v_2 \quad u_2v_2$

③
 $u_3 \quad v_3 \quad u_3v_3$

④
 $u_4 \quad v_4 \quad u_4v_4$

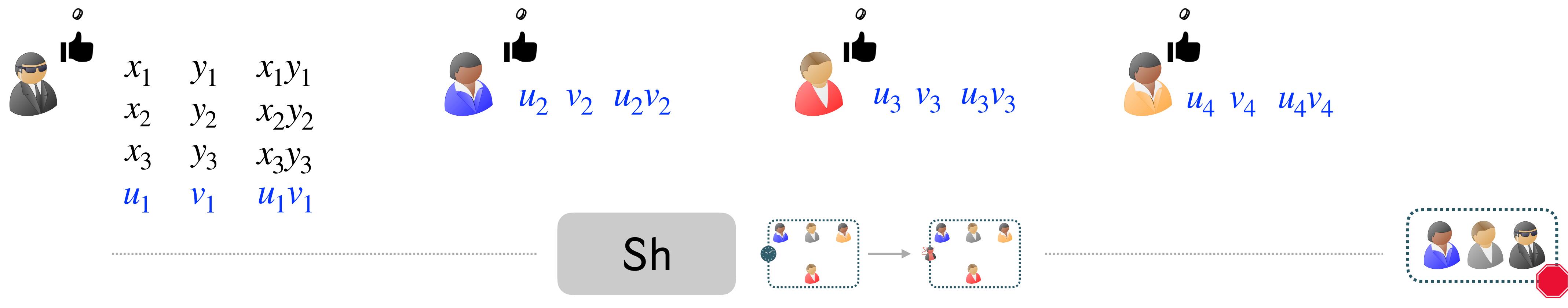
Perfect HMPC - Triple Sharing with Party Elimination Protocol



$[x_1] [y_1] [z_1]$
 $[x_2] [y_2] [z_2]$
 $[x_3] [y_3] [z_3]$

$[u_1] [v_1] [w_1]$
 $[u_2] [v_2] [w_2]$
 $[u_3] [v_3] [w_3]$
 $[u_4] [v_4] [w_4]$

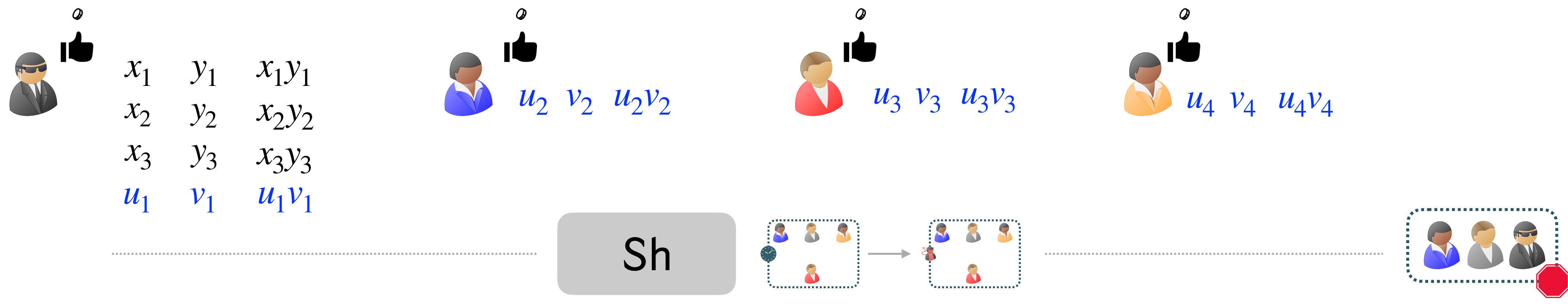
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 $[x_2] [y_2] [z_2]$
 $[x_3] [y_3] [z_3]$

$[u_1] [v_1] [w_1]$
 $[u_2] [v_2] [w_2]$
 $[u_3] [v_3] [w_3]$
 $[u_4] [v_4] [w_4]$

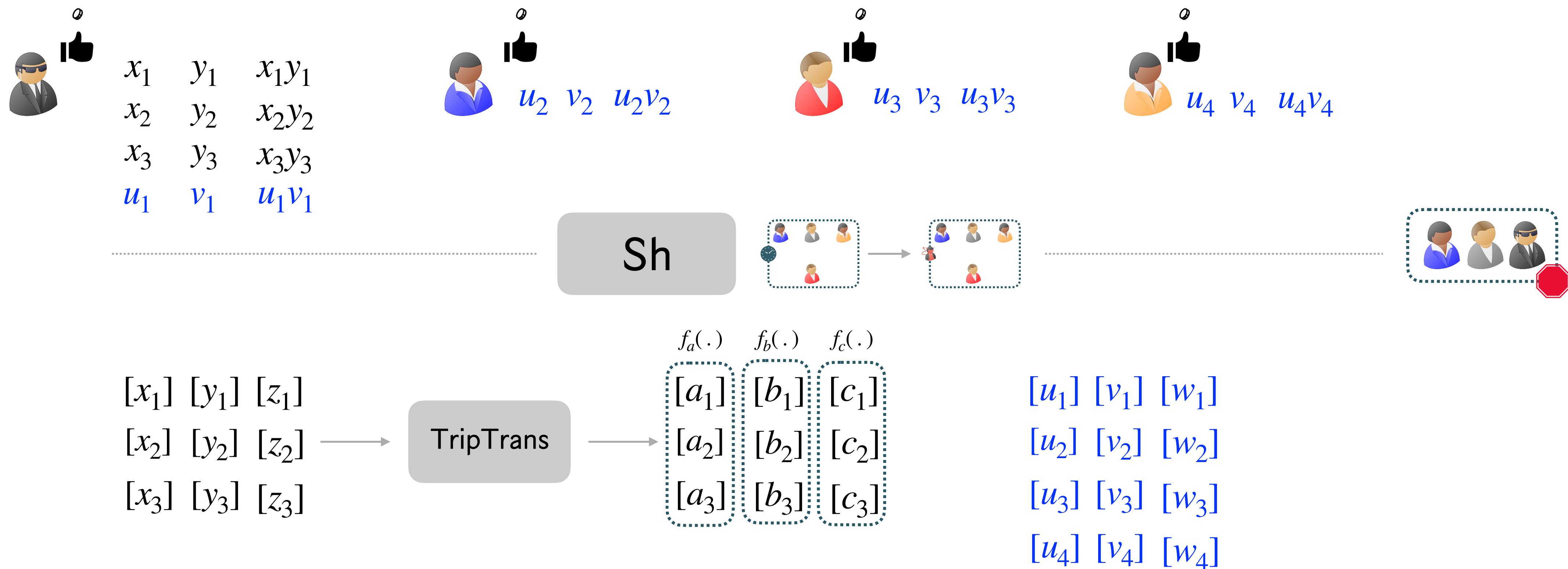
Perfect HMPC - Triple Sharing with Party Elimination Protocol



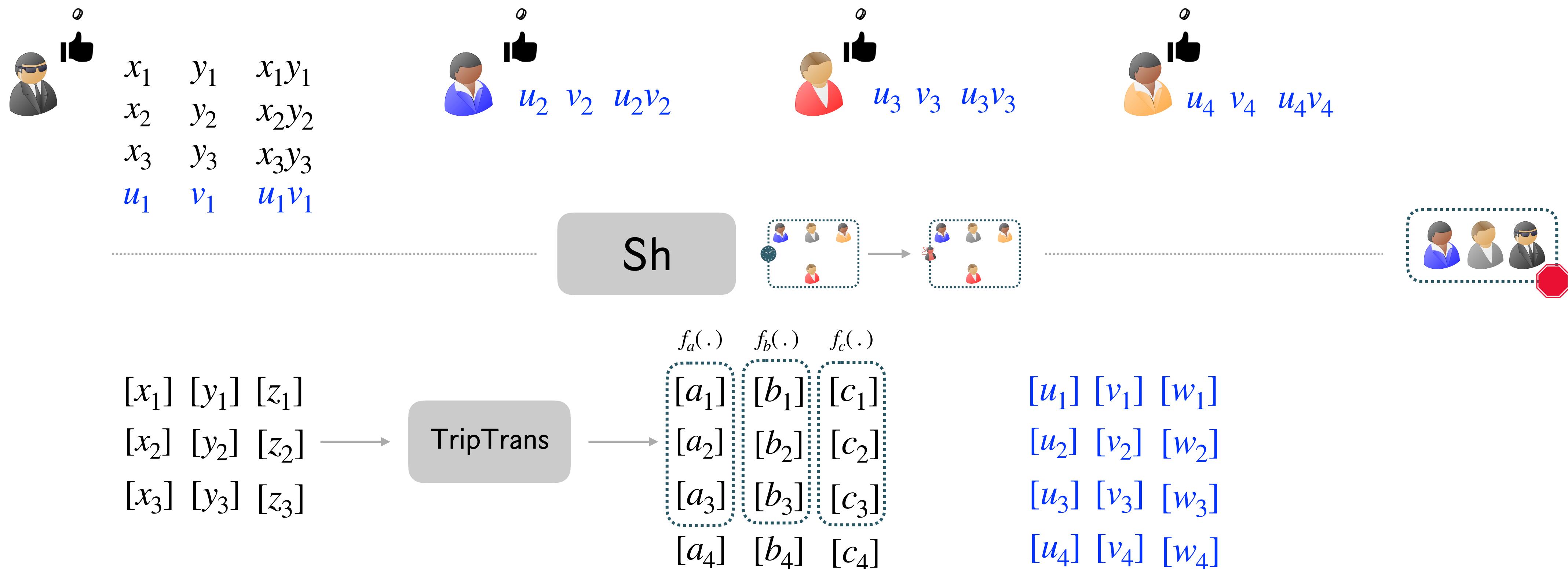
$[x_1] [y_1] [z_1]$ $[a_1] [b_1] [c_1]$ $[u_1] [v_1] [w_1]$
 $[x_2] [y_2] [z_2]$ $[a_2] [b_2] [c_2]$ $[u_2] [v_2] [w_2]$
 $[x_3] [y_3] [z_3]$ $[a_3] [b_3] [c_3]$ $[u_3] [v_3] [w_3]$

 $[u_4] [v_4] [w_4]$

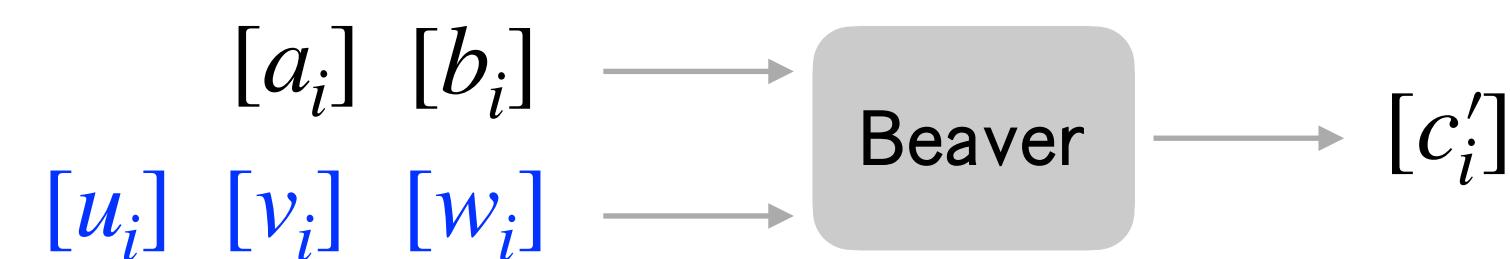
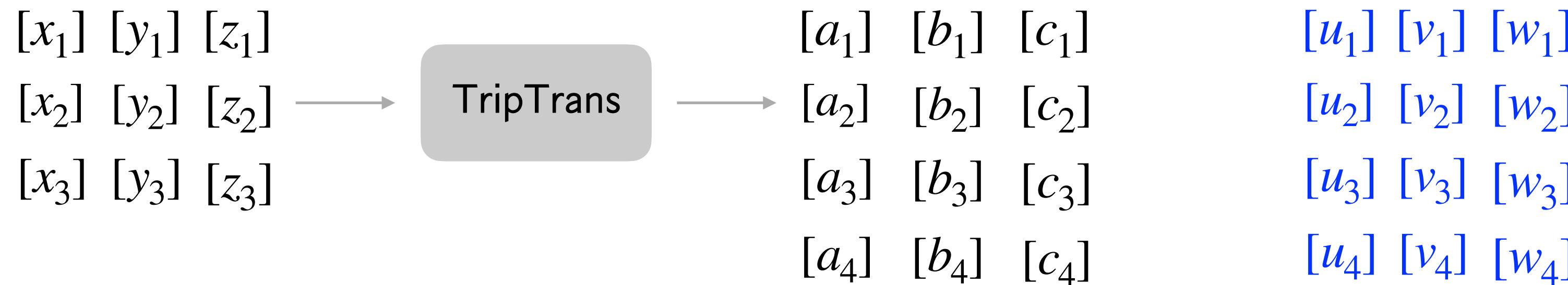
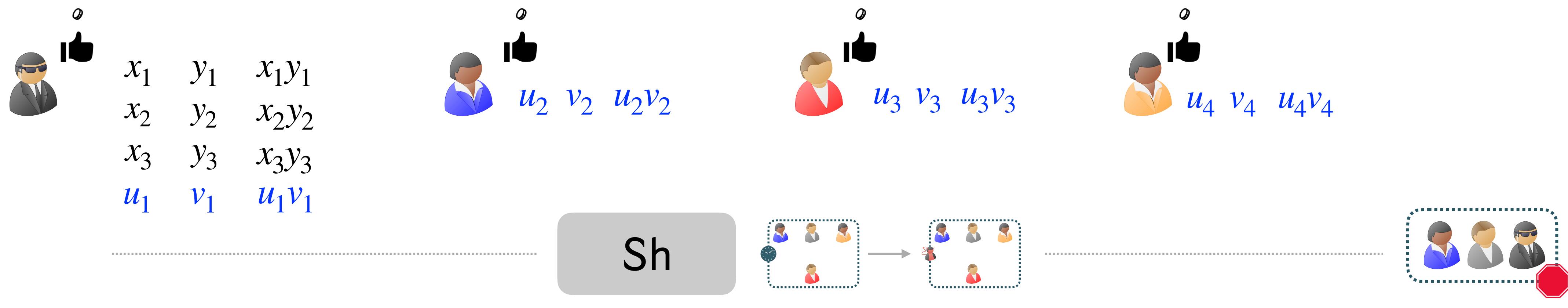
Perfect HMPC - Triple Sharing with Party Elimination Protocol



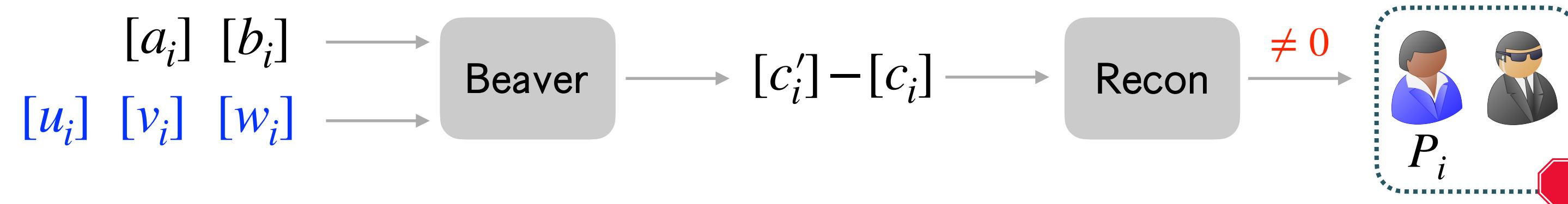
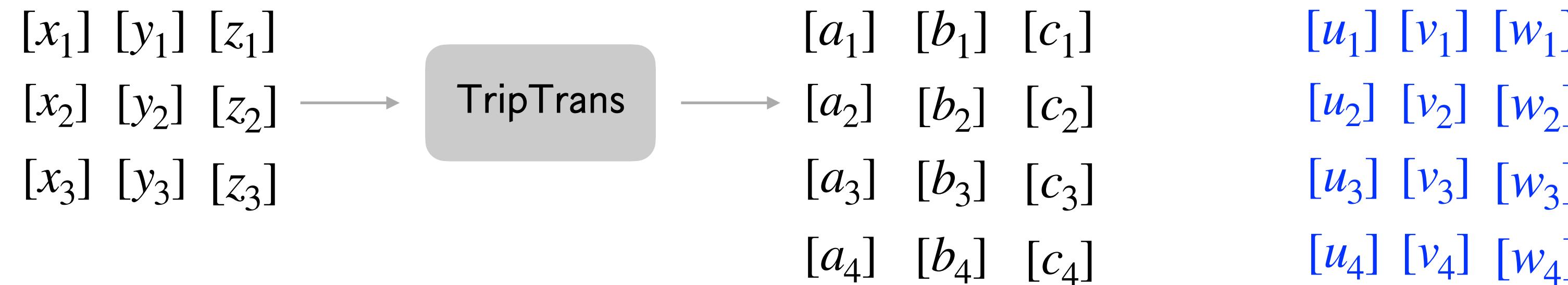
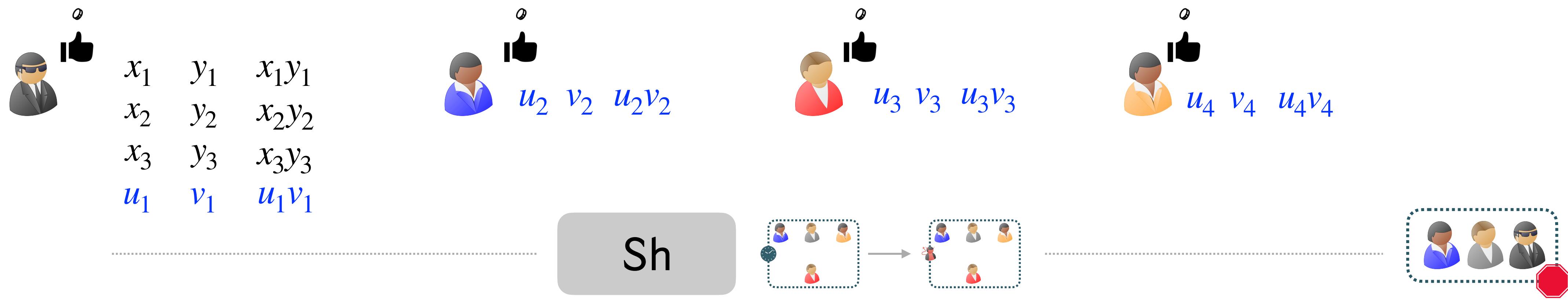
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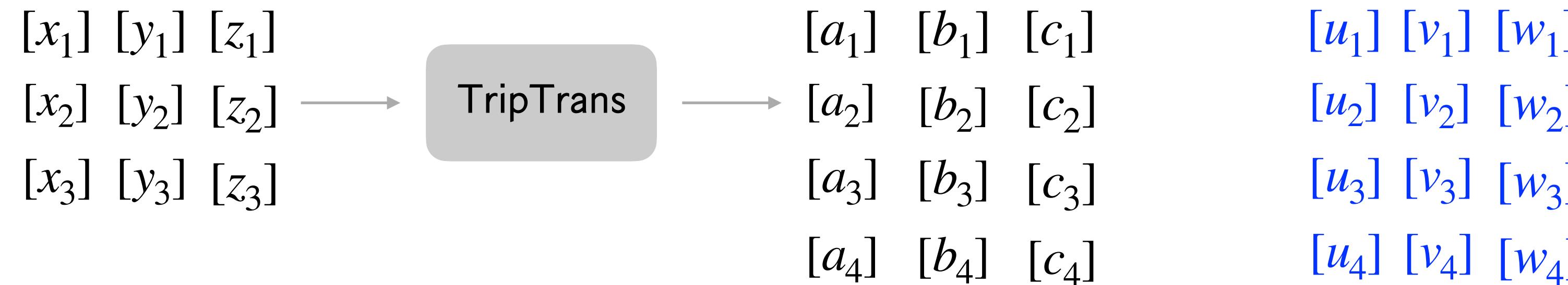
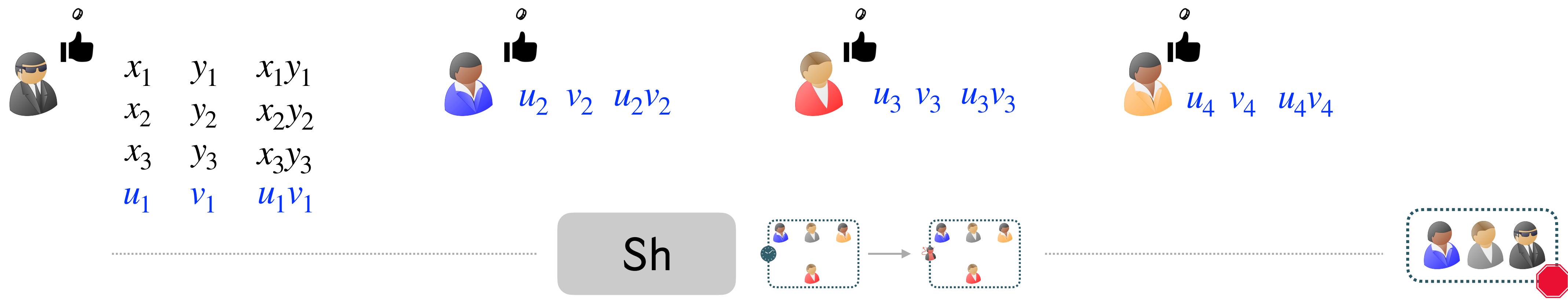
Perfect HMPC - Triple Sharing with Party Elimination Protocol



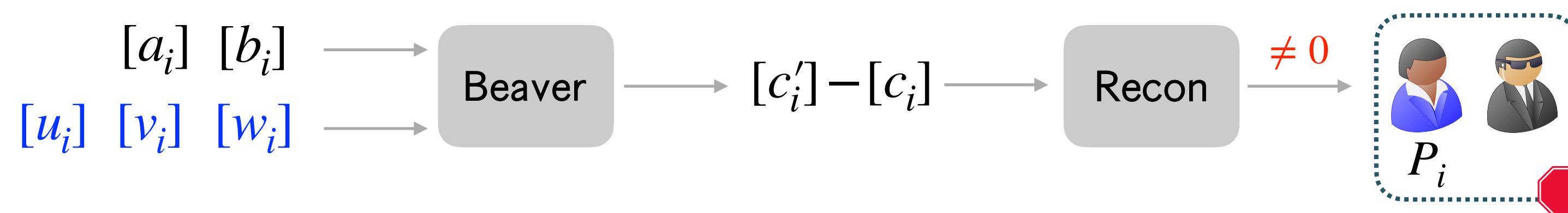
Perfect HMPC - Triple Sharing with Party Elimination Protocol



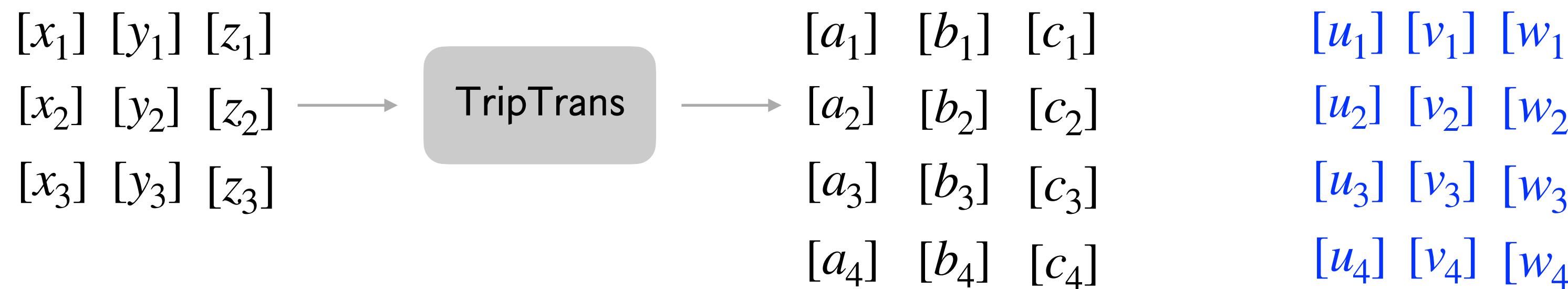
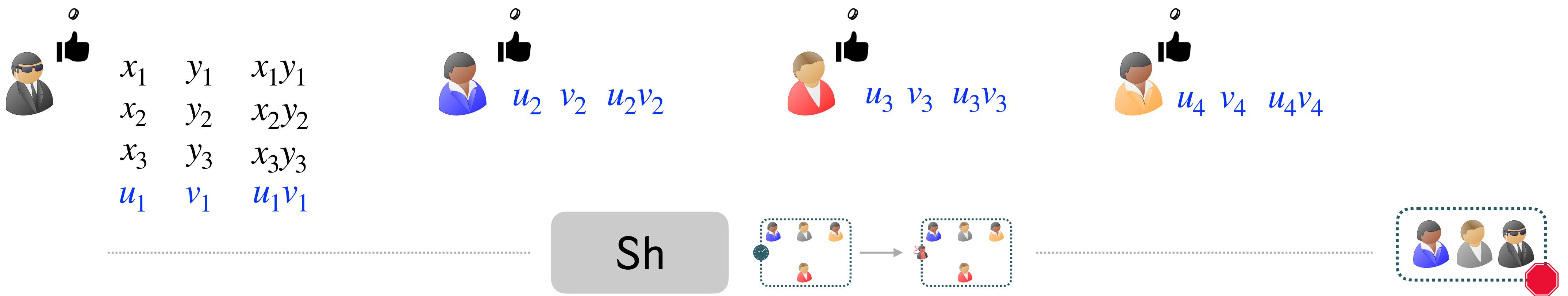
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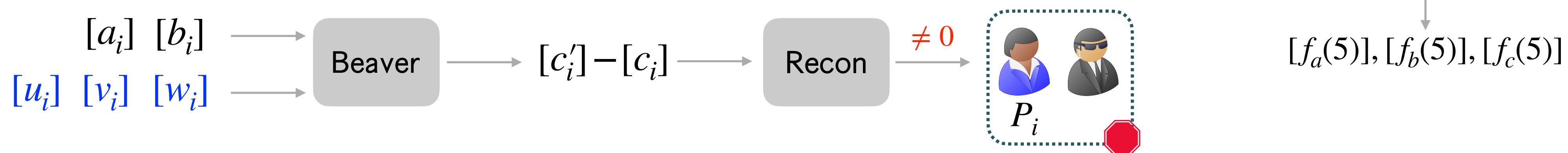
$f_c(\cdot) = f_a(\cdot)f_b(\cdot)$
if all checks hold



Perfect HMPC - Triple Sharing with Party Elimination Protocol



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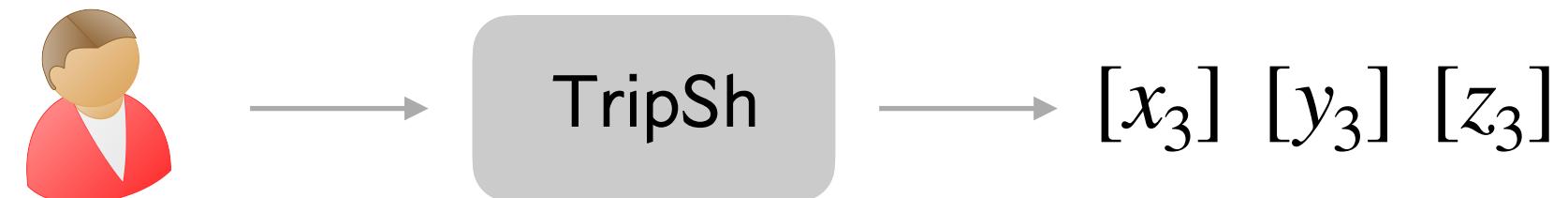
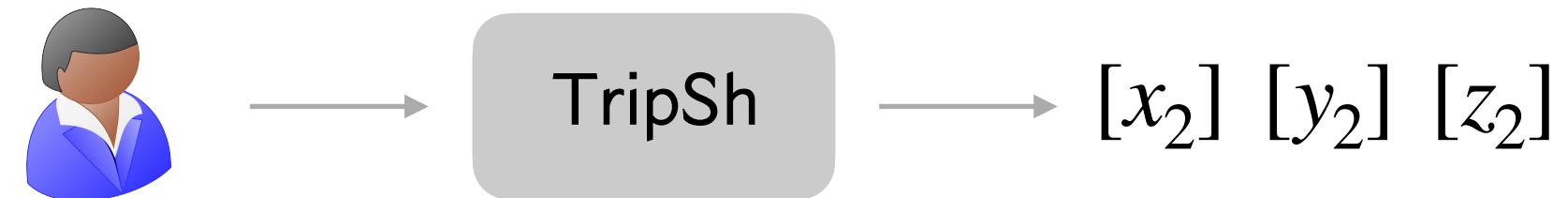
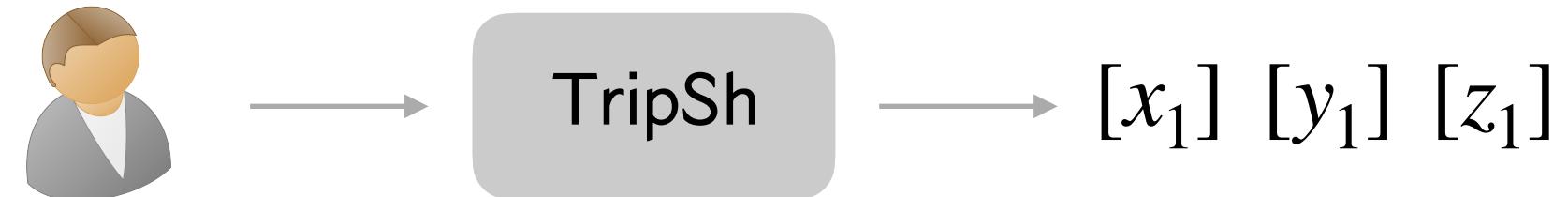


Perfect HMPC - Triple Generation with Party Elimination



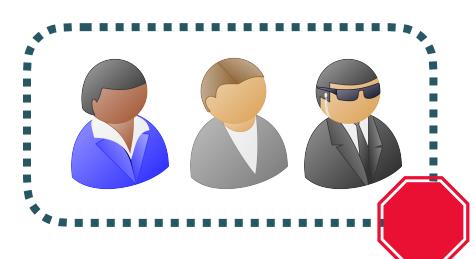
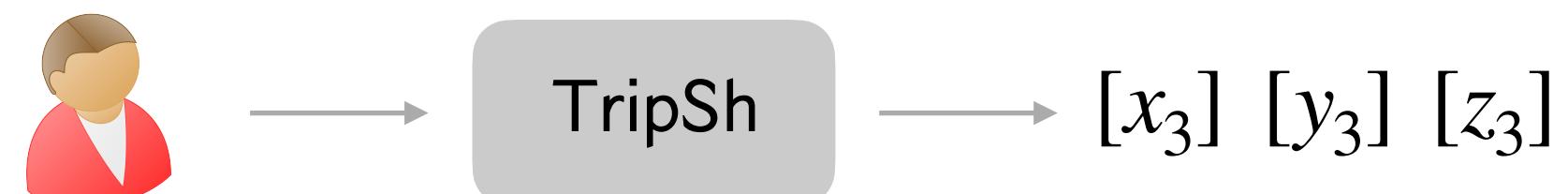
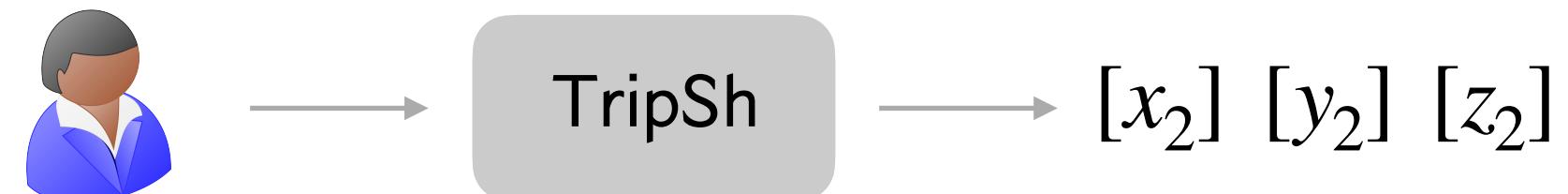
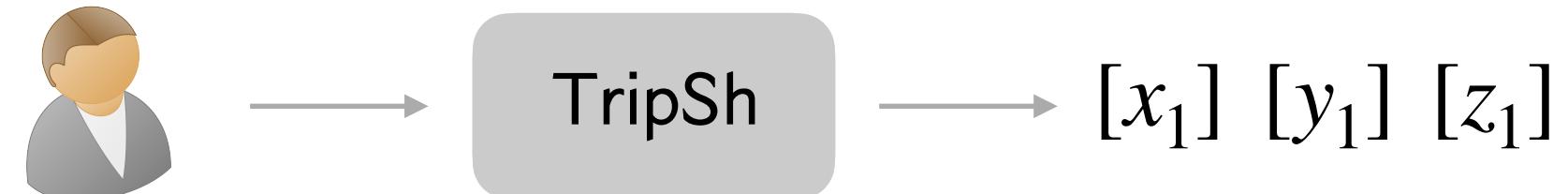
TripGen

Perfect HMPC - Triple Generation with Party Elimination



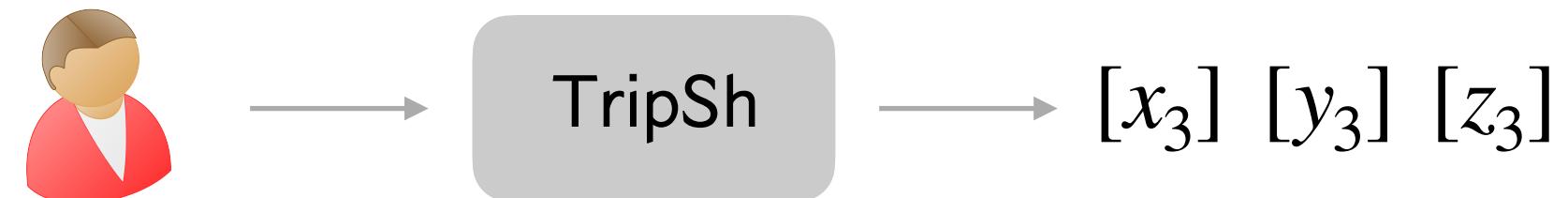
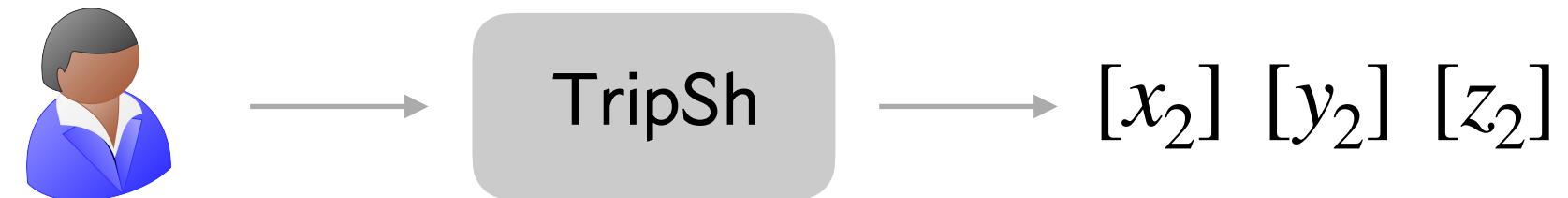
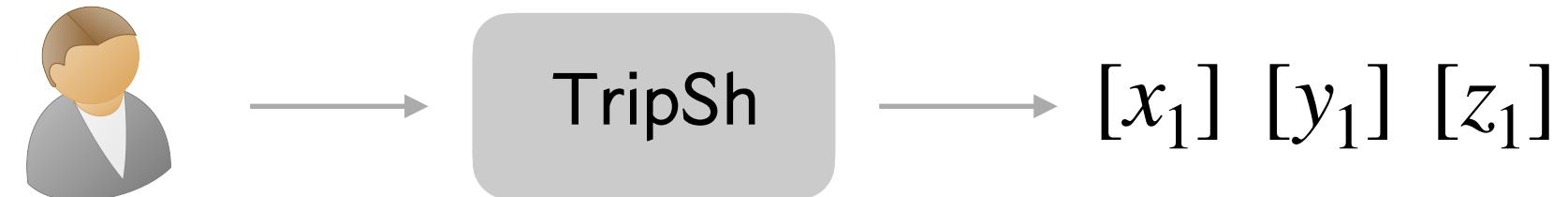
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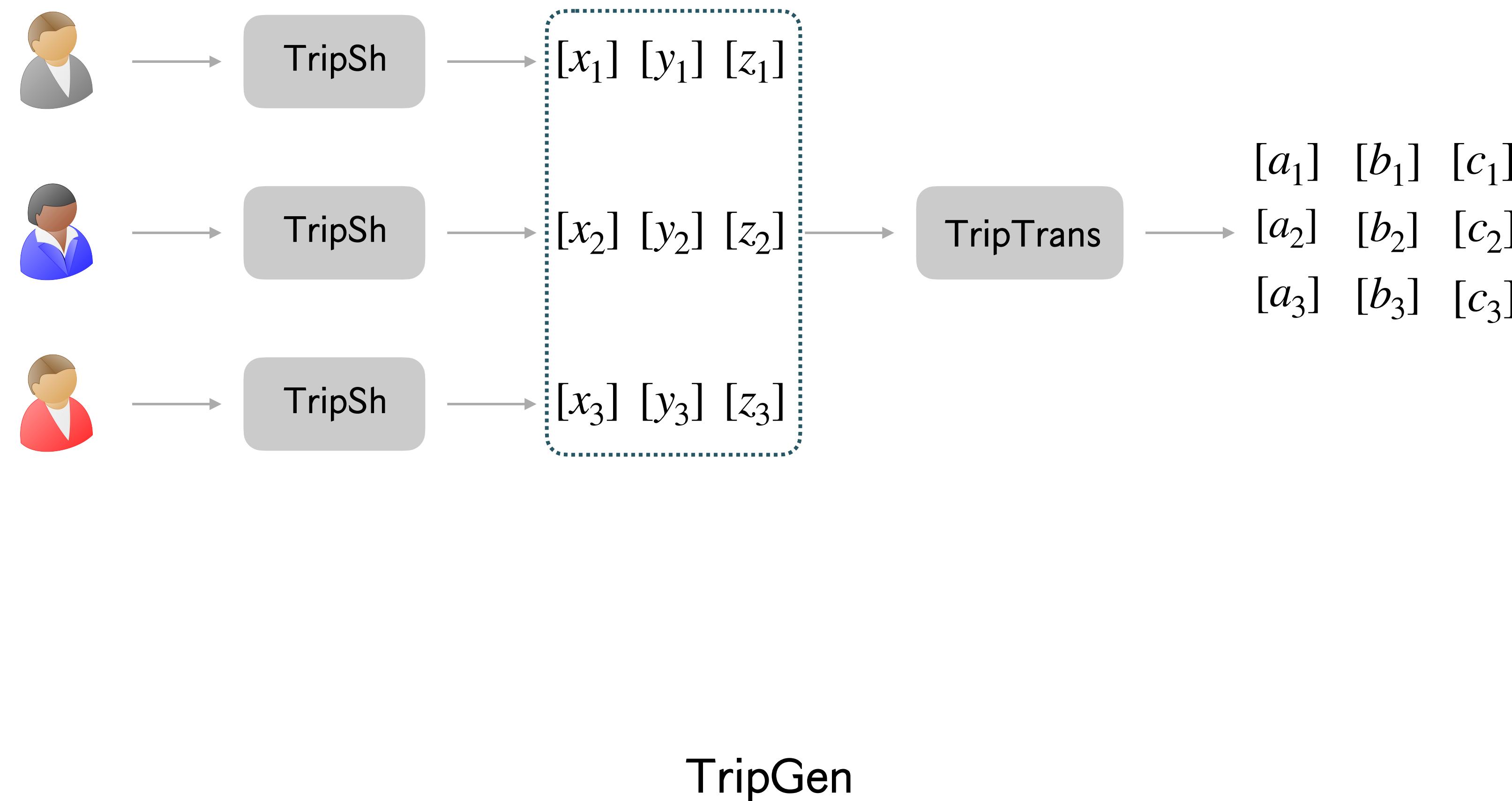
TripGen

Perfect HMPC - Triple Generation with Party Elimination

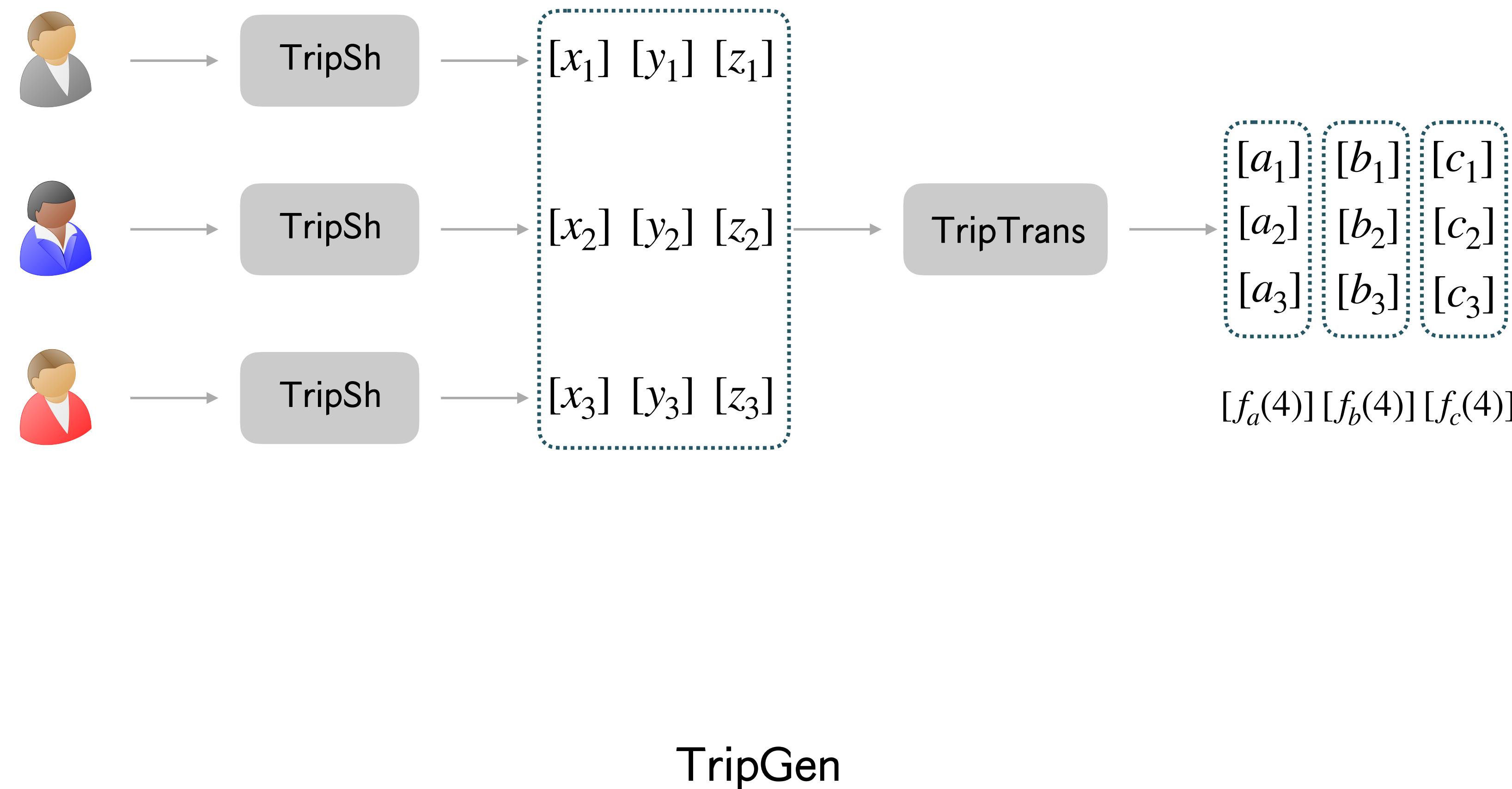


TripGen

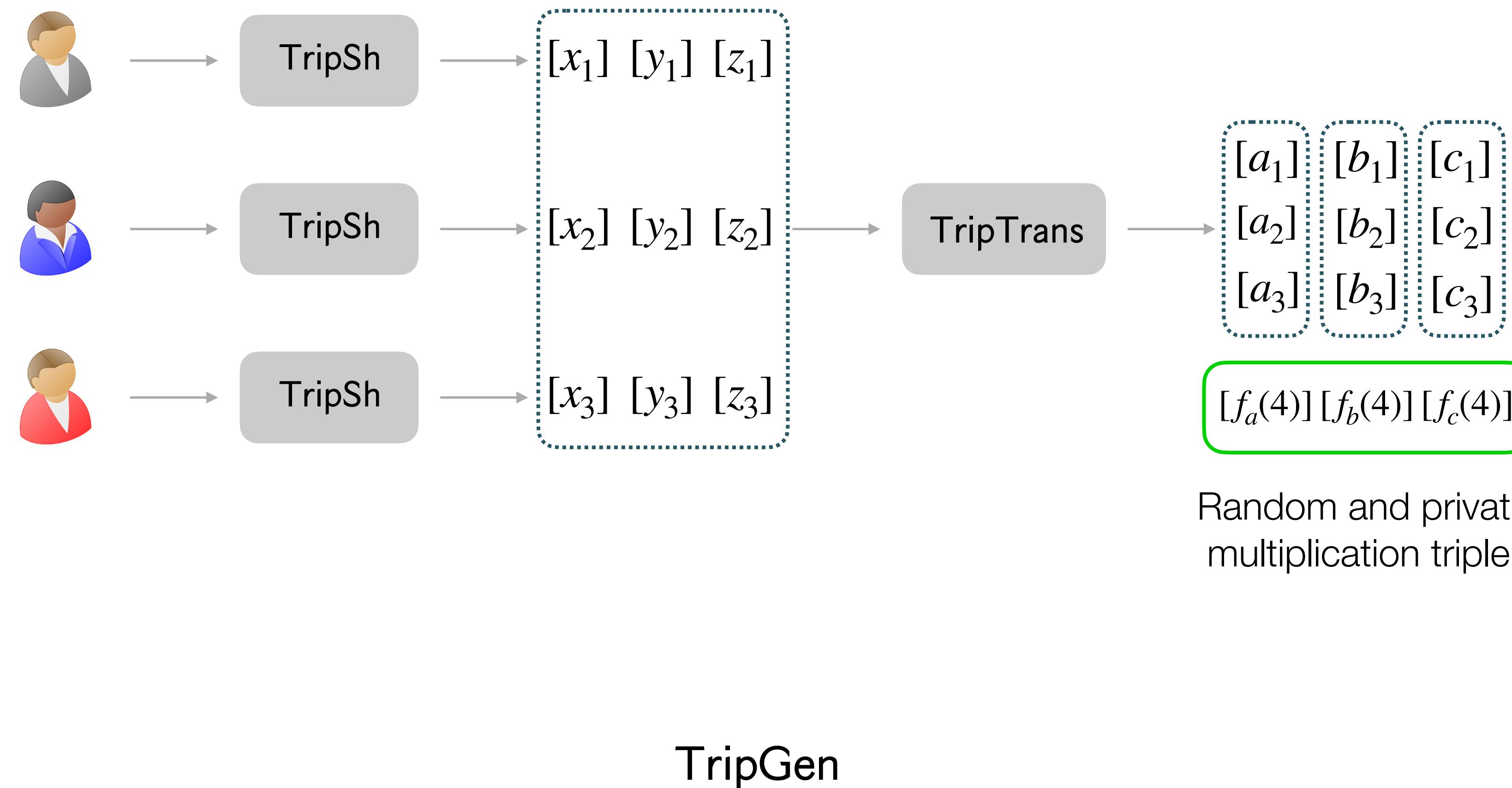
Perfect HMPC - Triple Generation with Party Elimination



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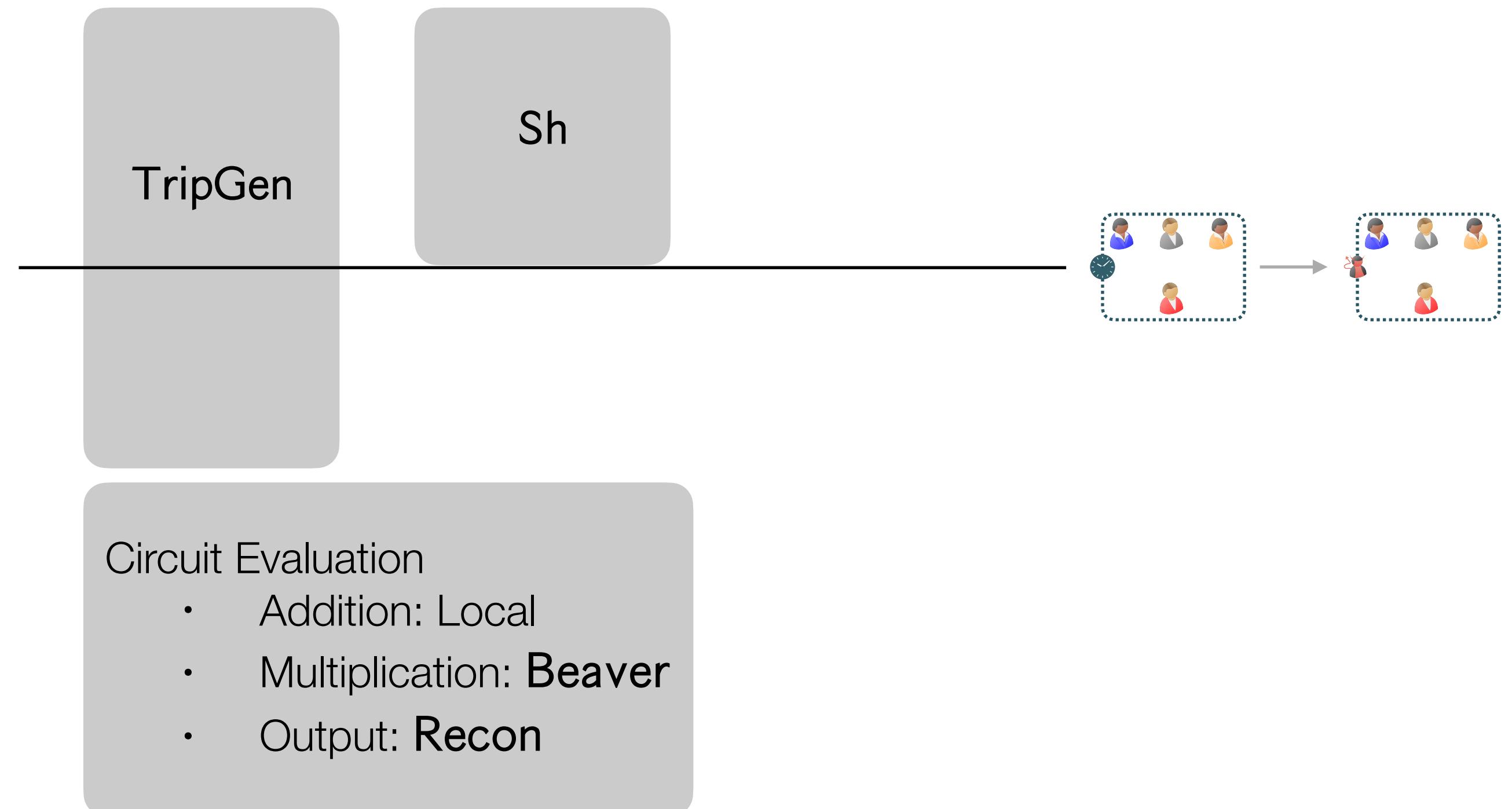


Perfect HMPC - Triple Generation with Party Elimination



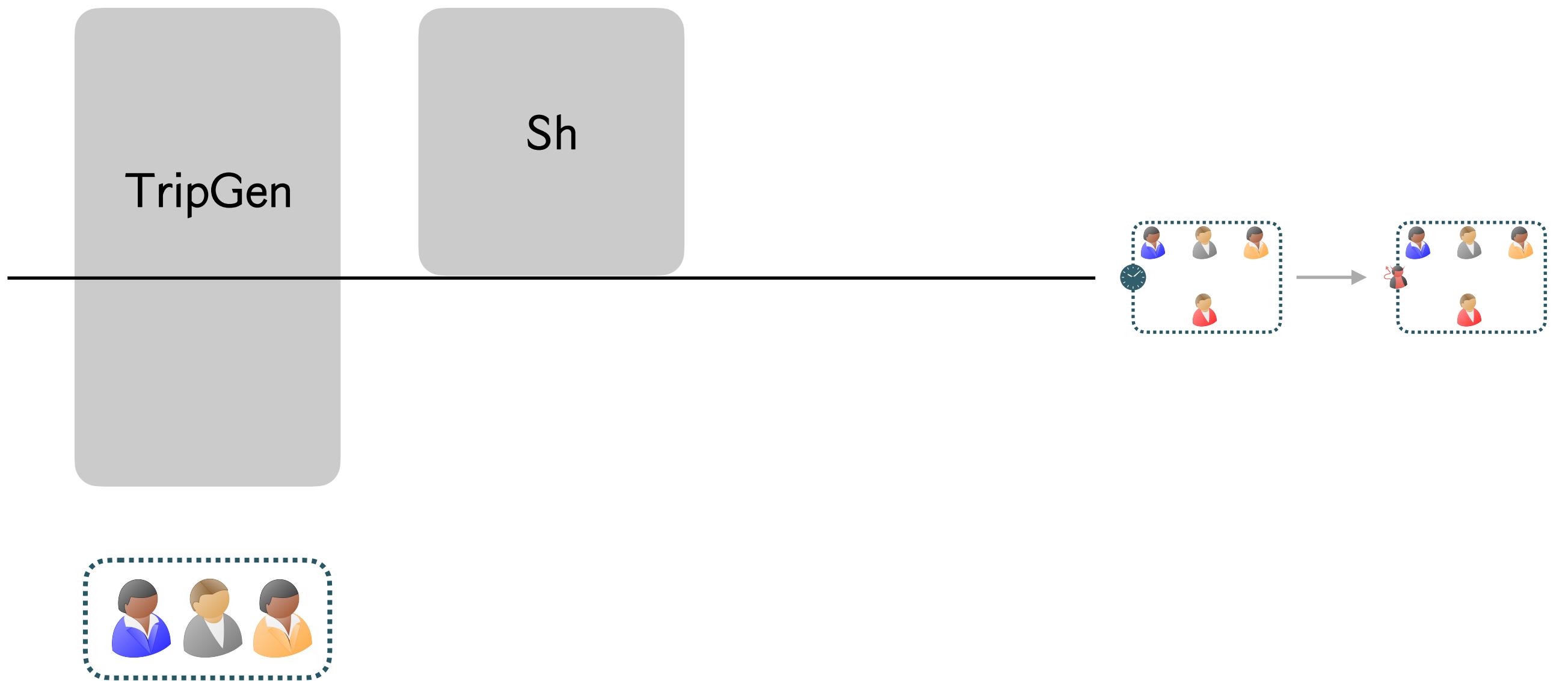
Perfect HMPC

- 3 phases
 - Triple generation phase
 - Input phase
 - Circuit evaluation and output phase



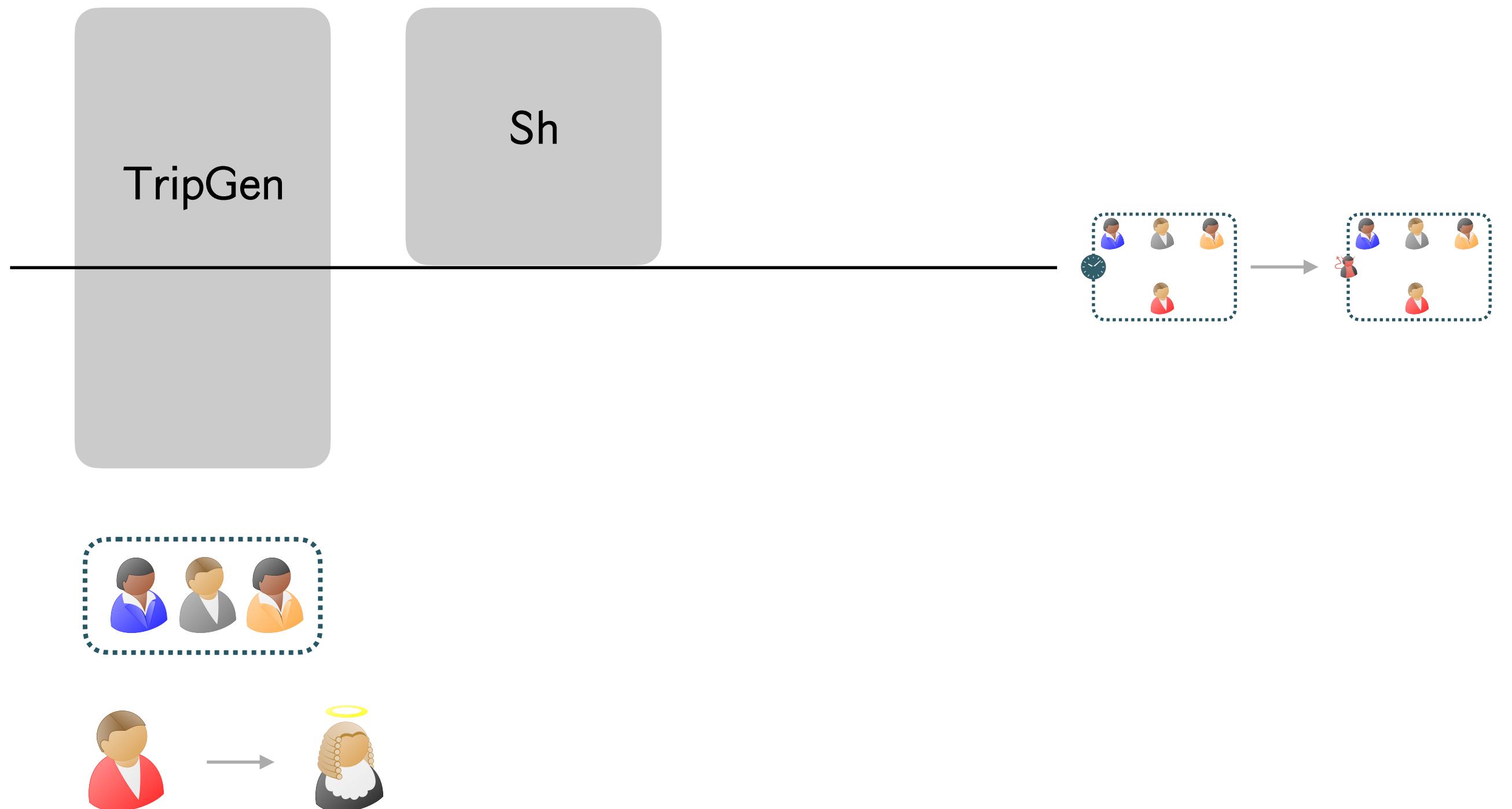
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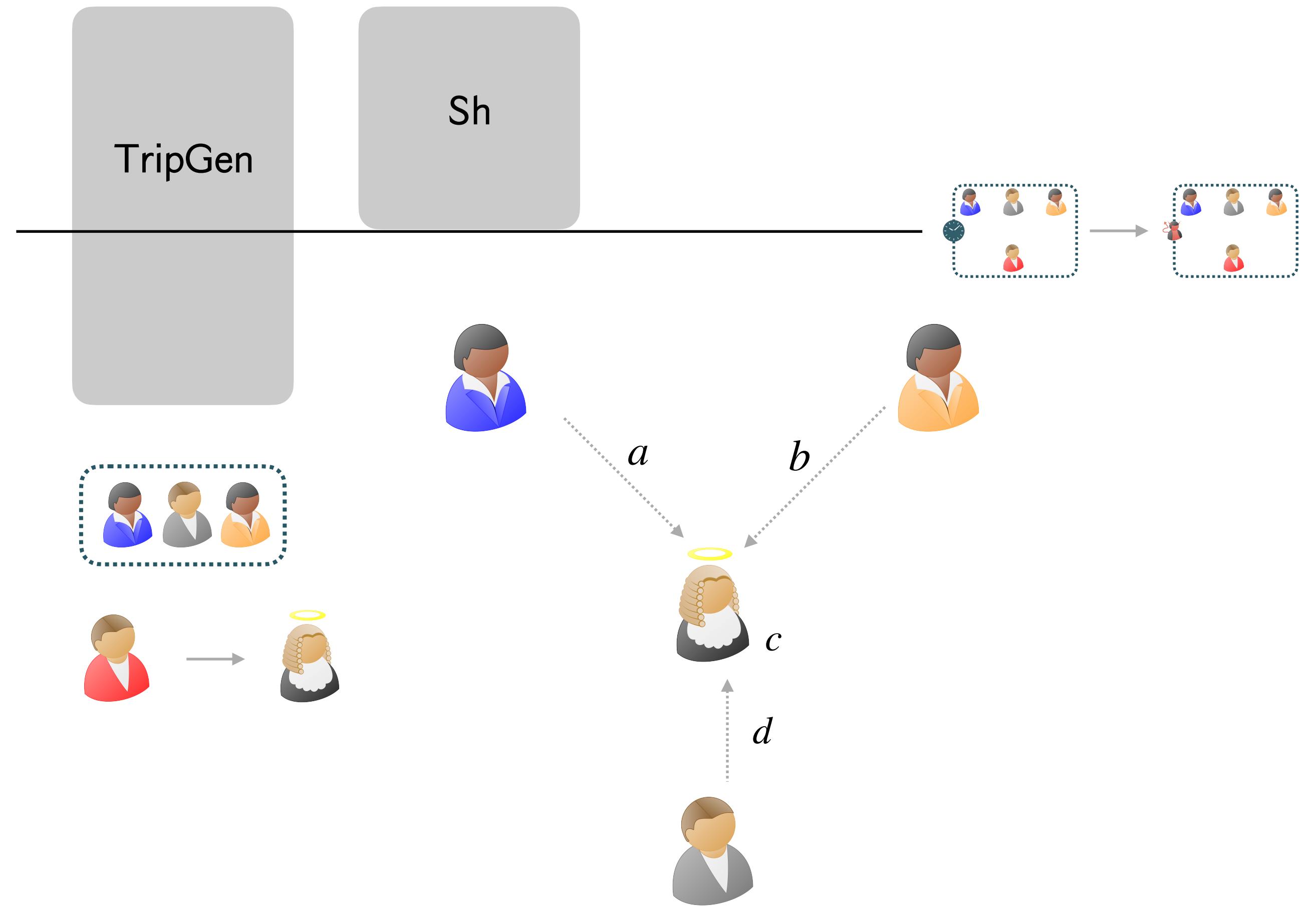
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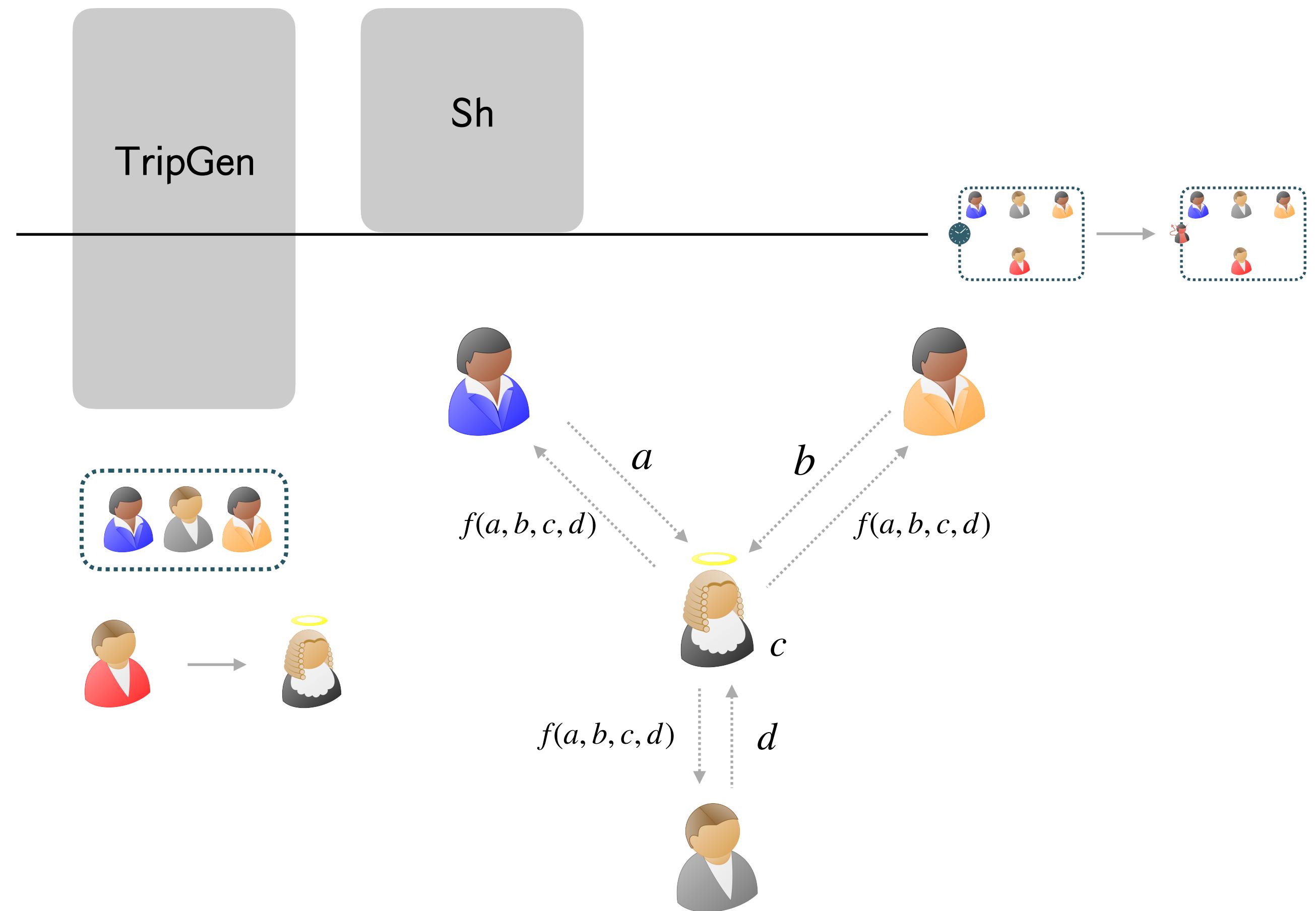
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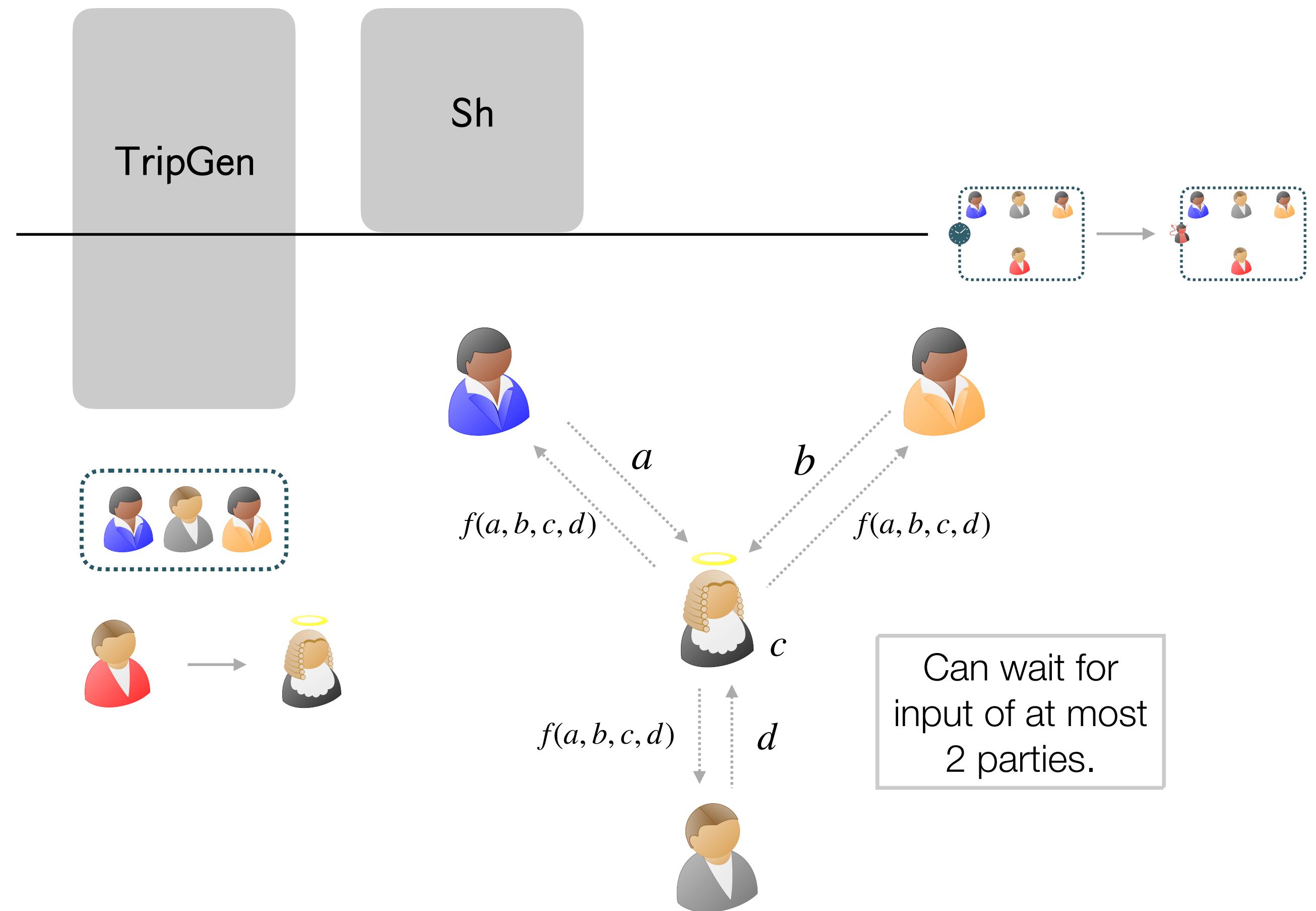
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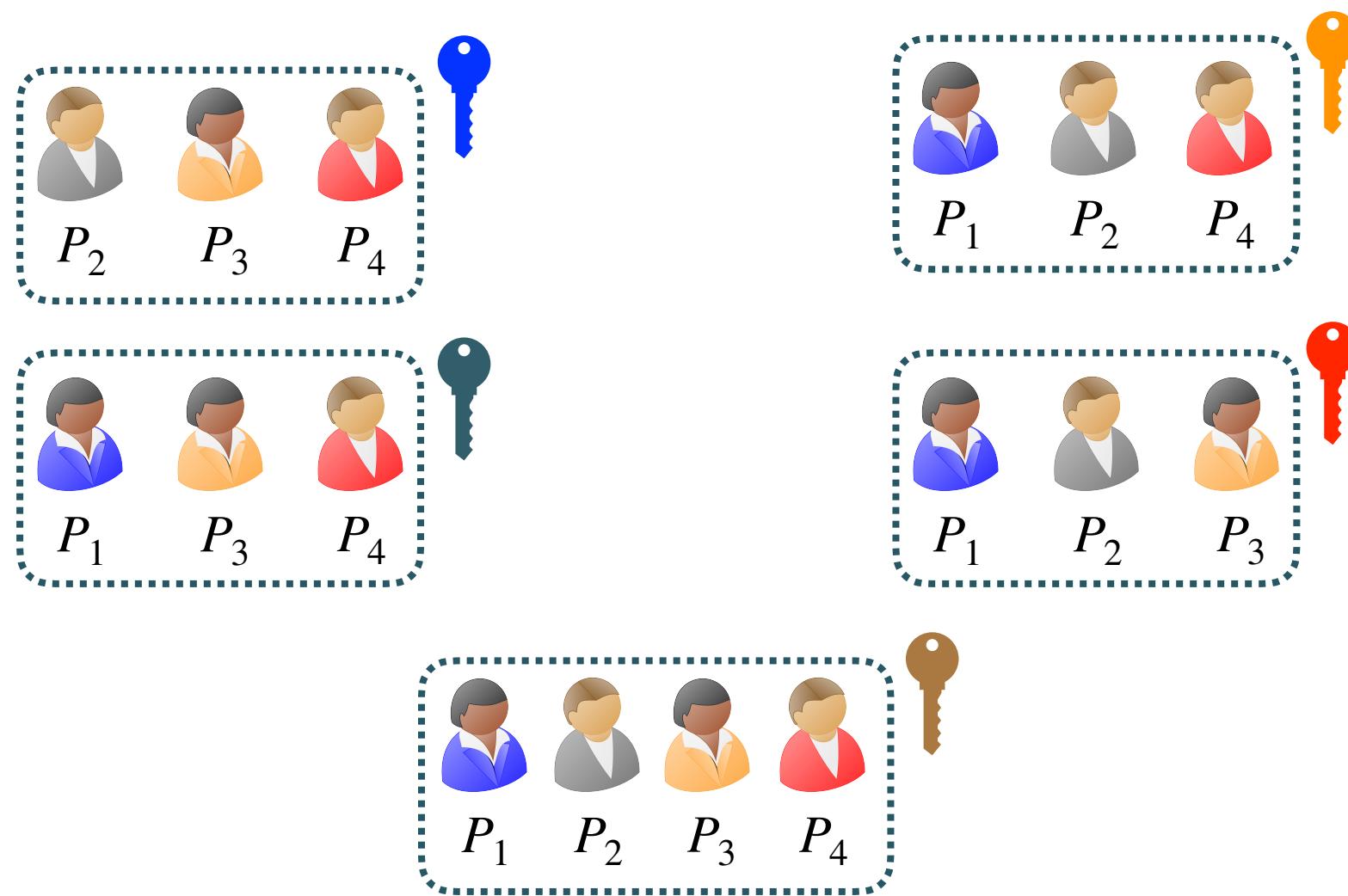


Perfect HMPC

- 3 phases
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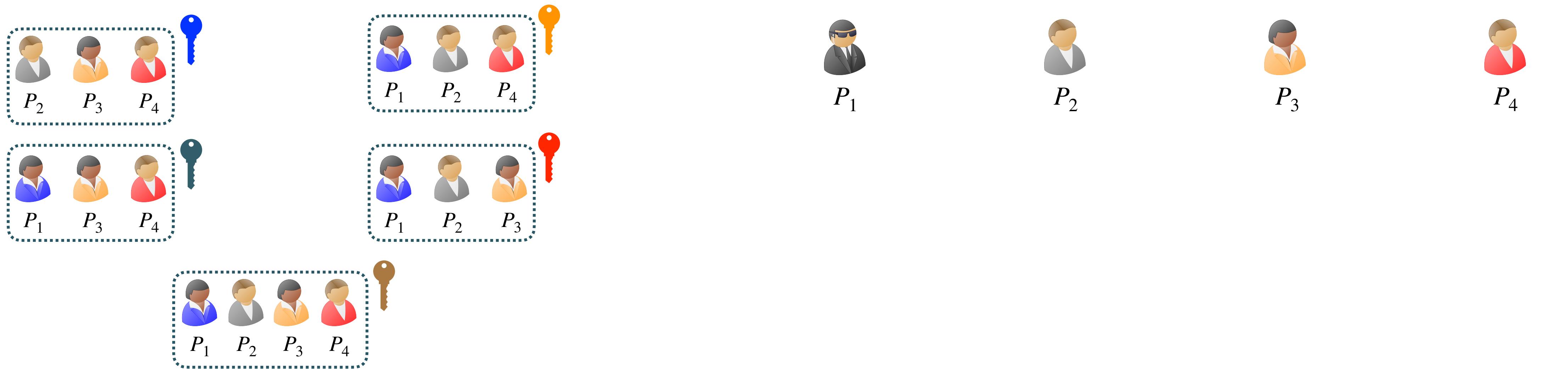


Cryptographically Secure HMP - Secret Sharing and Reconstruction



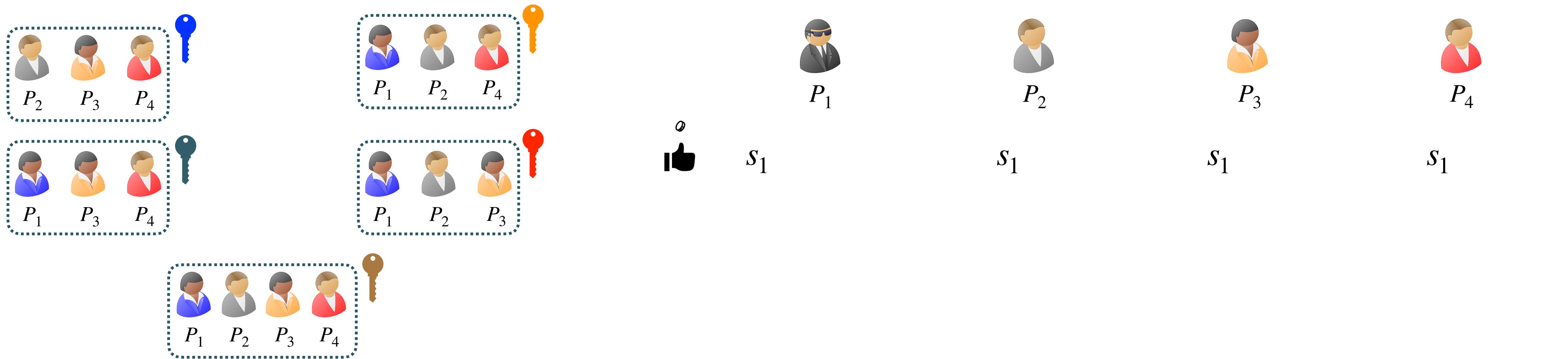
- Assume symmetric-key setup for PRF [AFL+16, CCP+19, MR18]

Cryptographically Secure HMP - Secret Sharing and Reconstruction



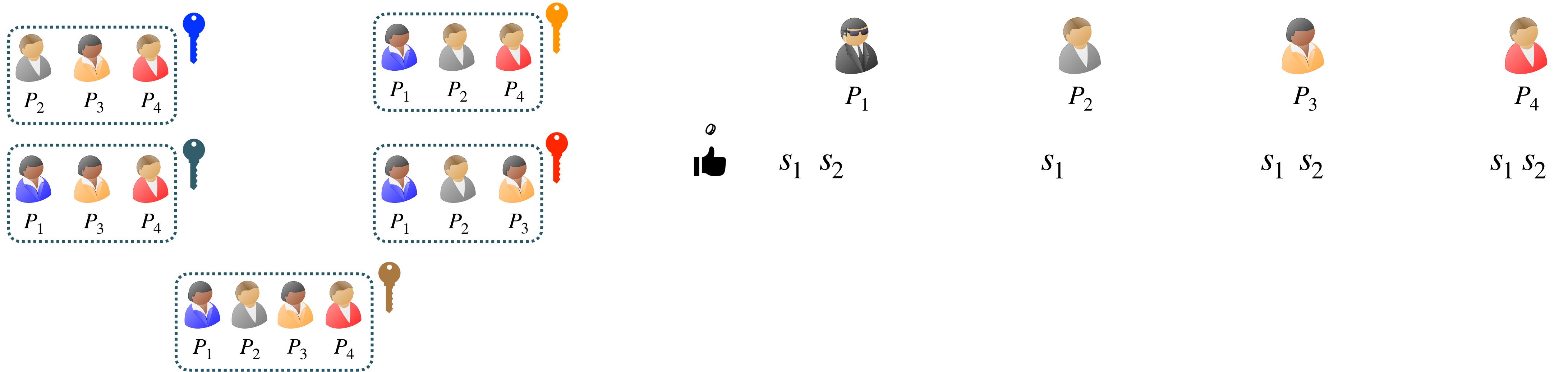
- Assume symmetric-key setup for PRF [AFL+16, CCP+19, MR18]
 - One synchronous round VSS protocol

Cryptographically Secure HMPC - Secret Sharing and Reconstruction



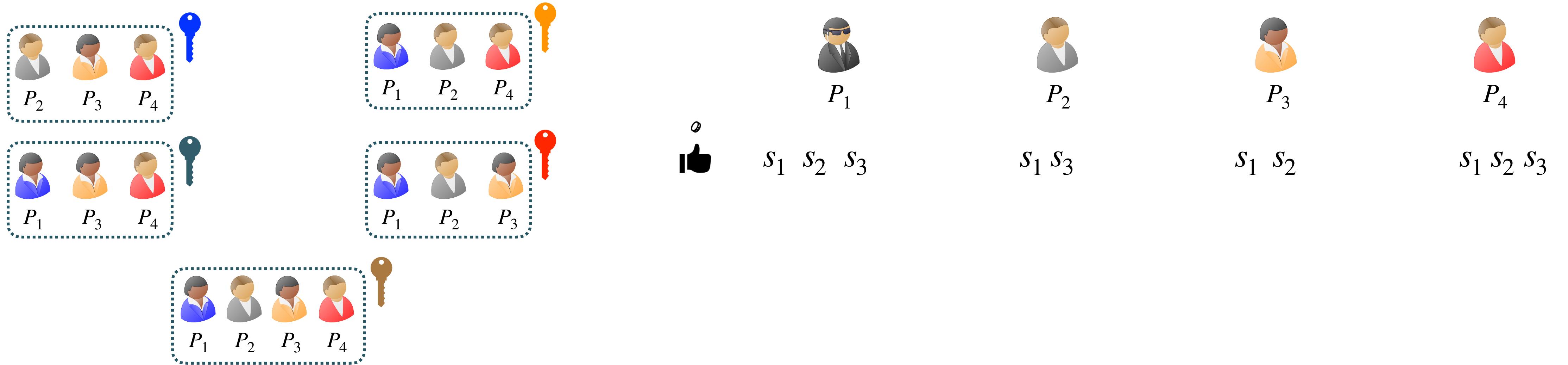
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Cryptographically Secure HMPC - Secret Sharing and Reconstruction



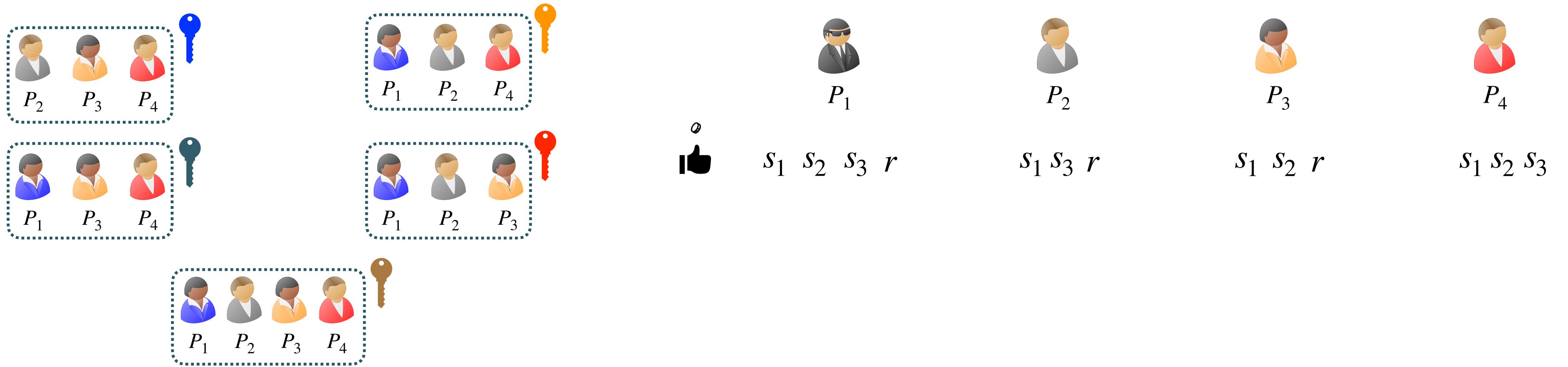
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 - One synchronous round VSS protocol

Cryptographically Secure HMPC - Secret Sharing and Reconstruction



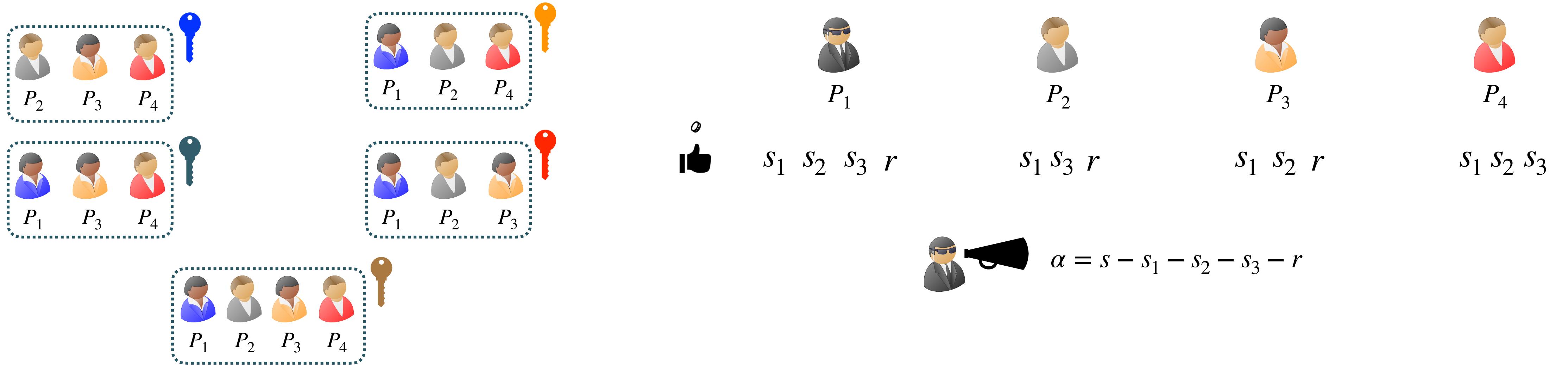
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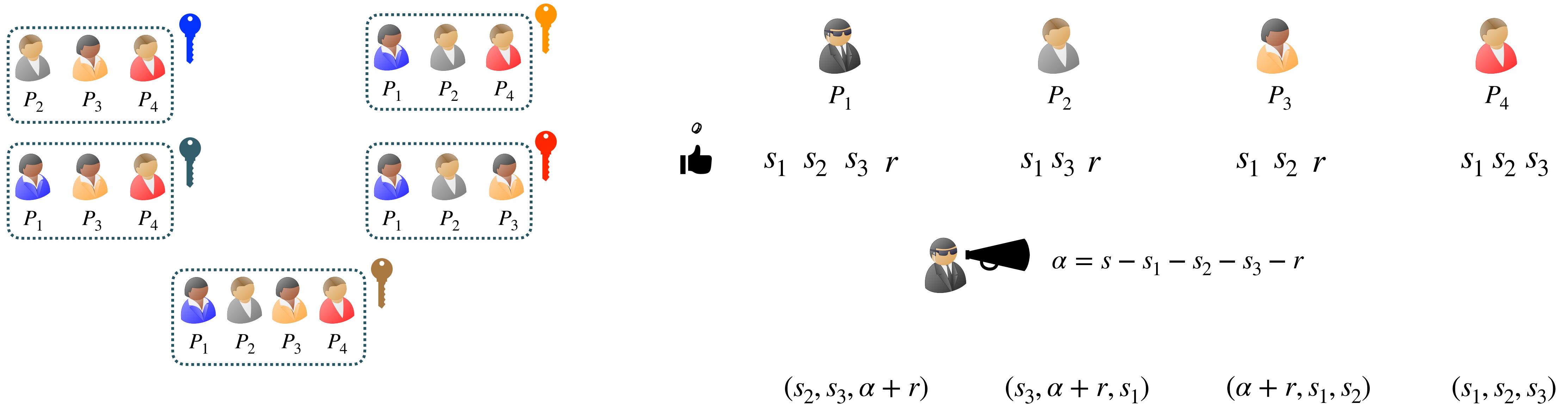
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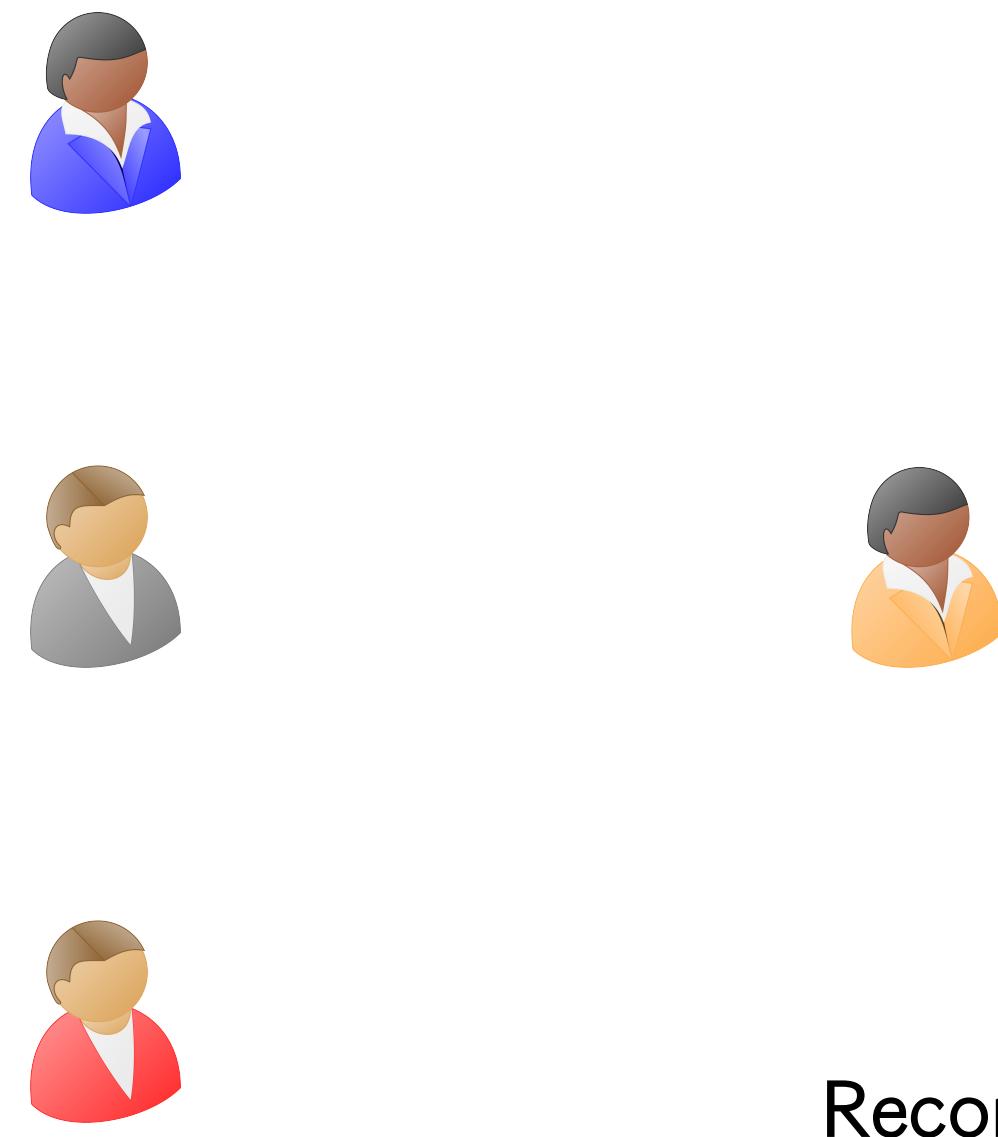
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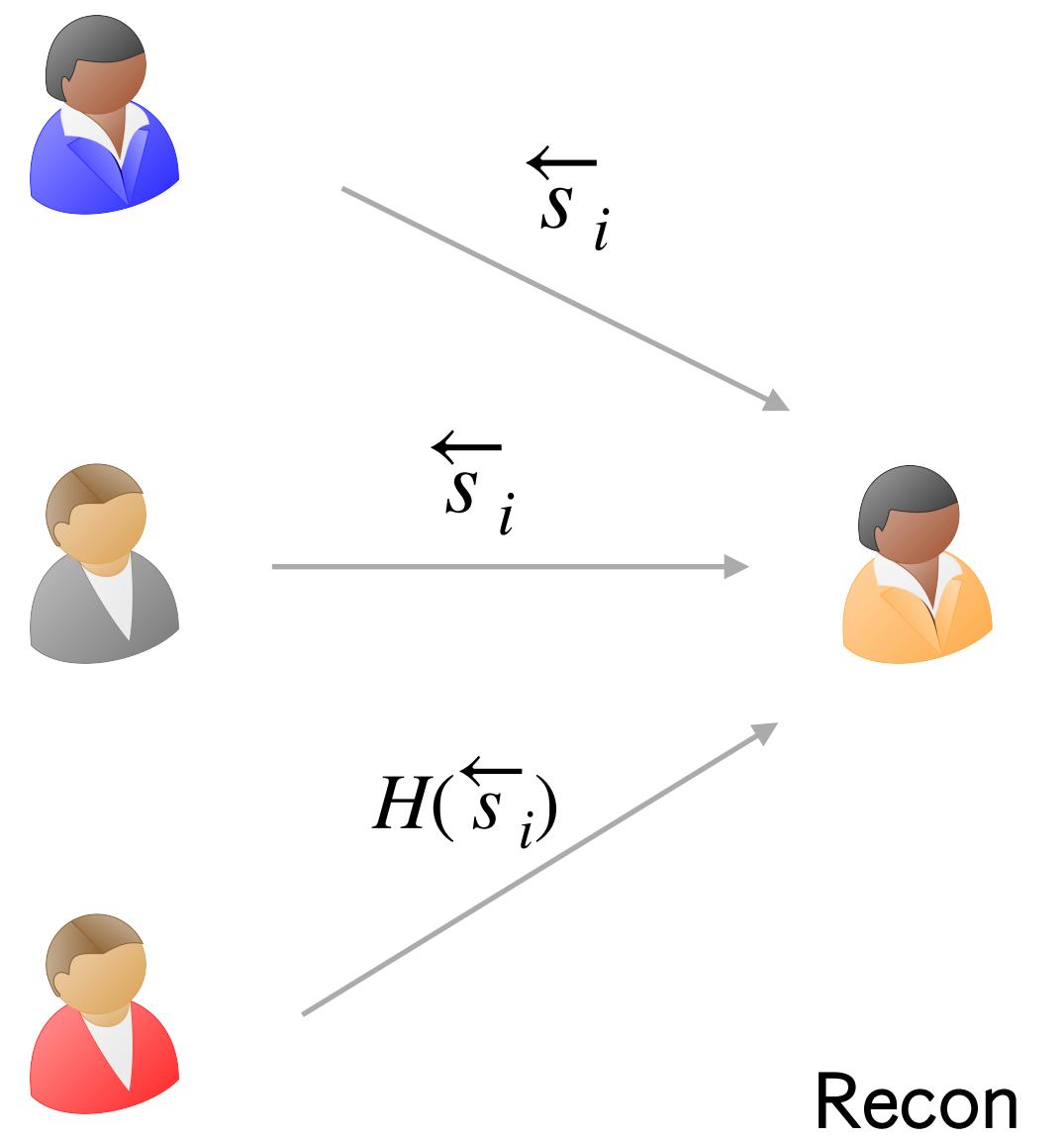
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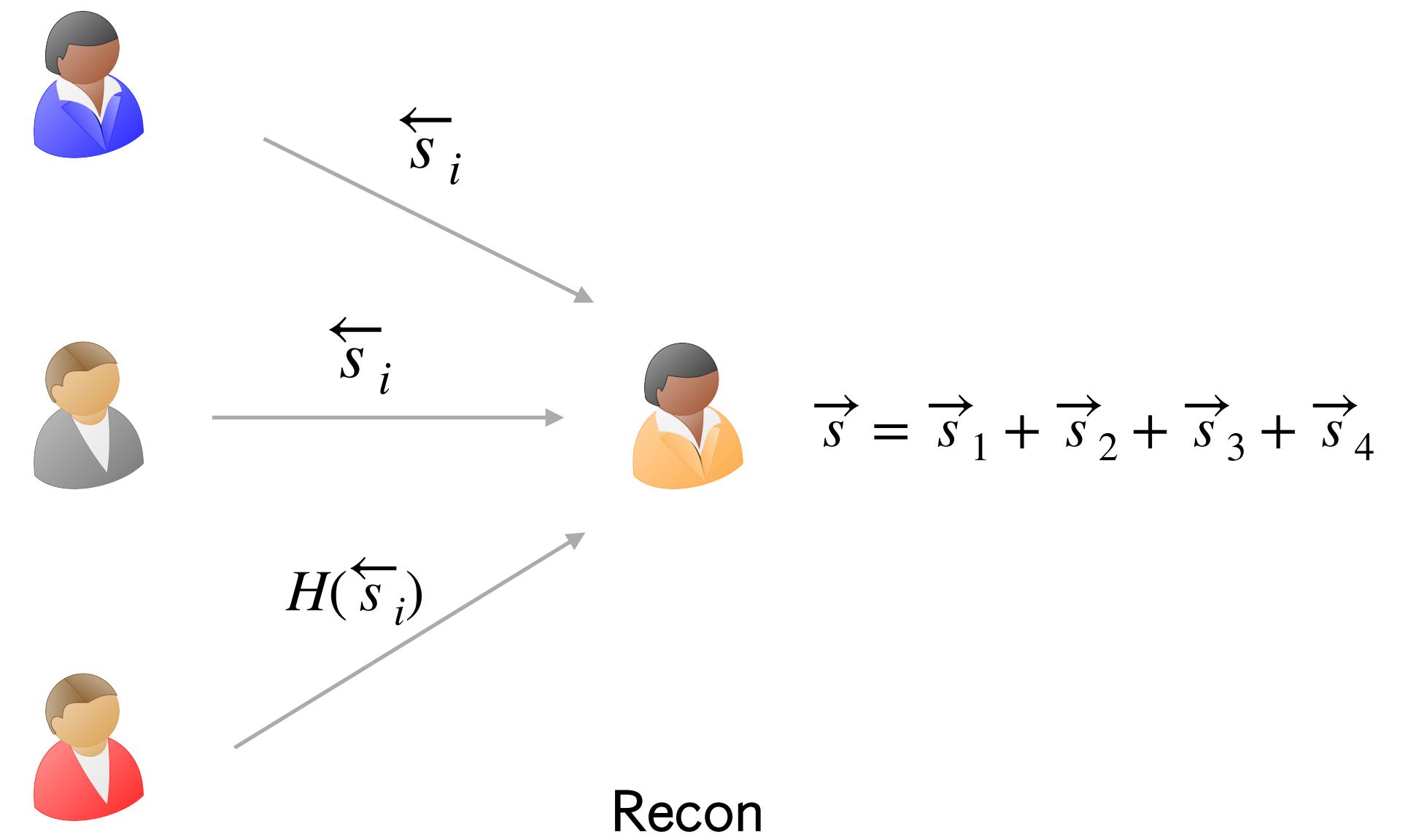
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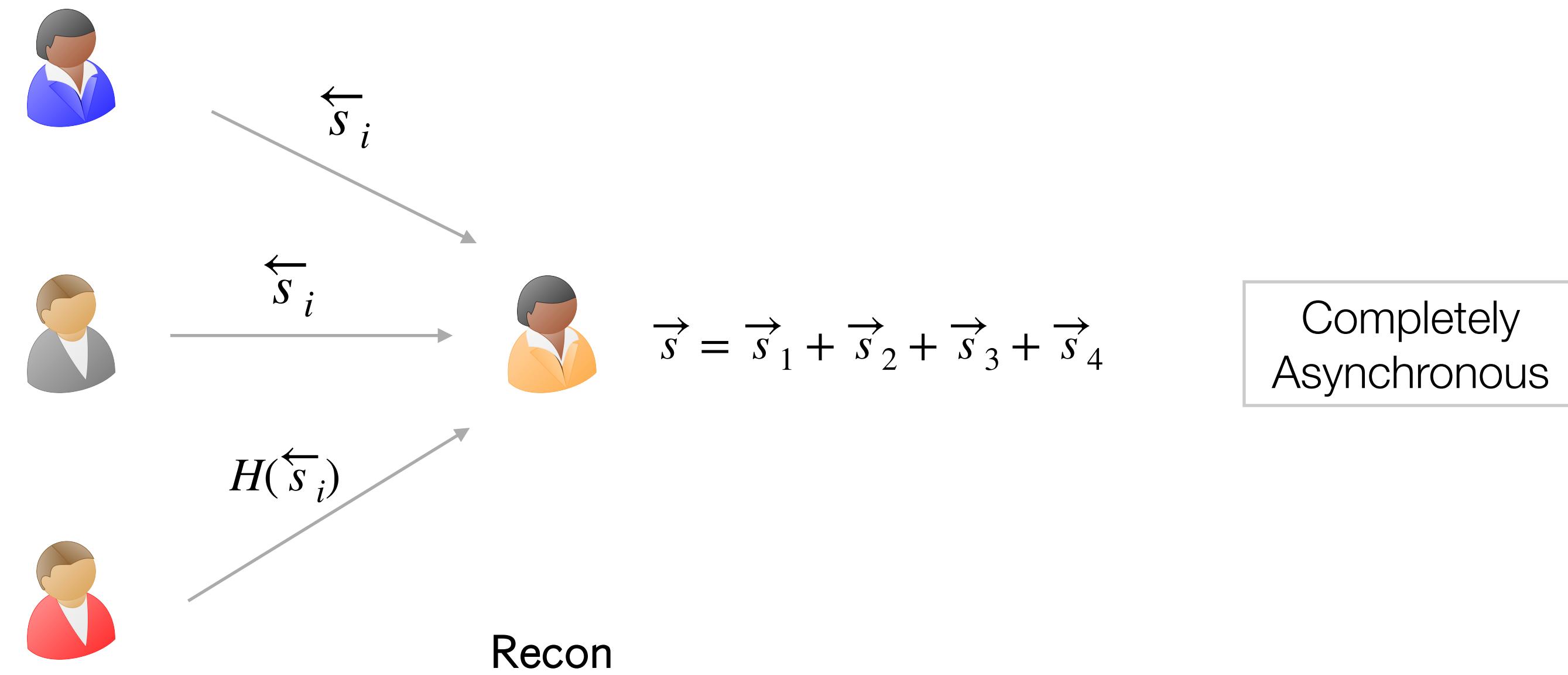
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Cryptographically Secure HMPC - Triple Generation Protocol

- Triple sharing similar to **TripSh**
 - Dealer shares $2l + 1$ triples instead of 3 triples
 - Other parties don't share triples

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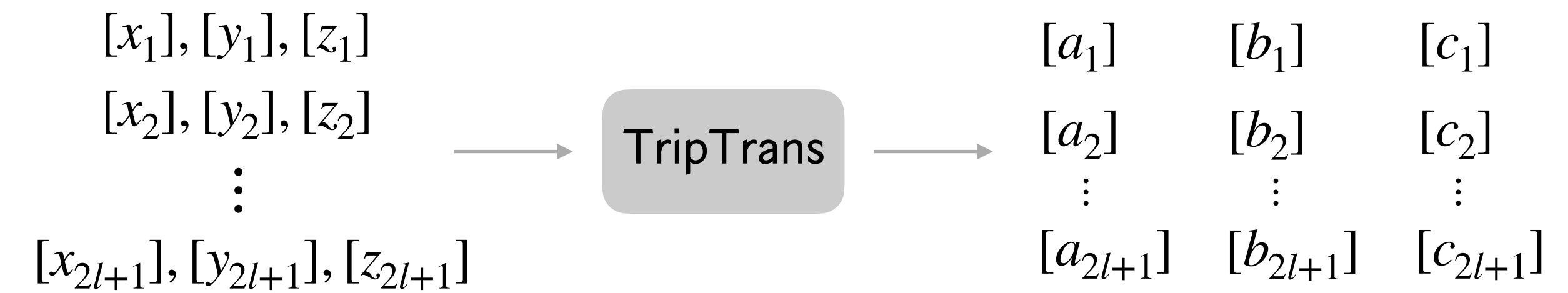
$$[x_1], [y_1], [z_1]$$
$$[x_2], [y_2], [z_2]$$
$$\vdots$$
$$[x_{2l+1}], [y_{2l+1}], [z_{2l+1}]$$

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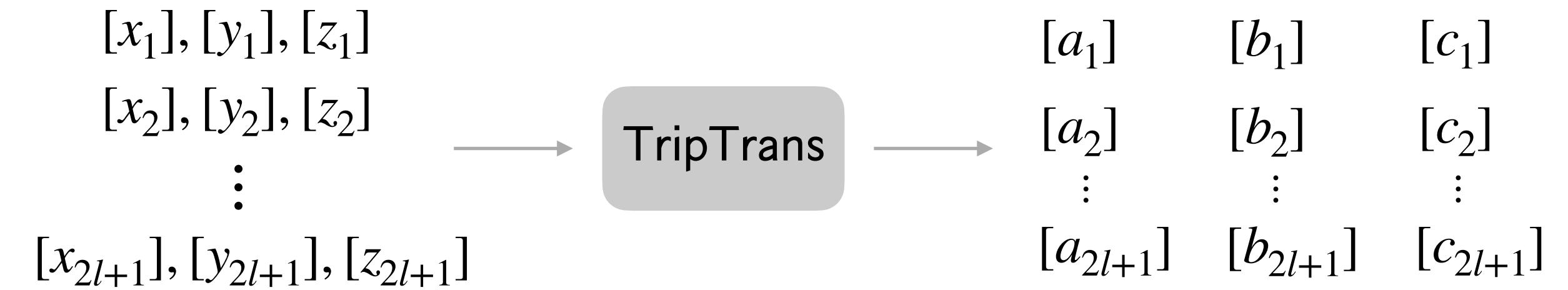


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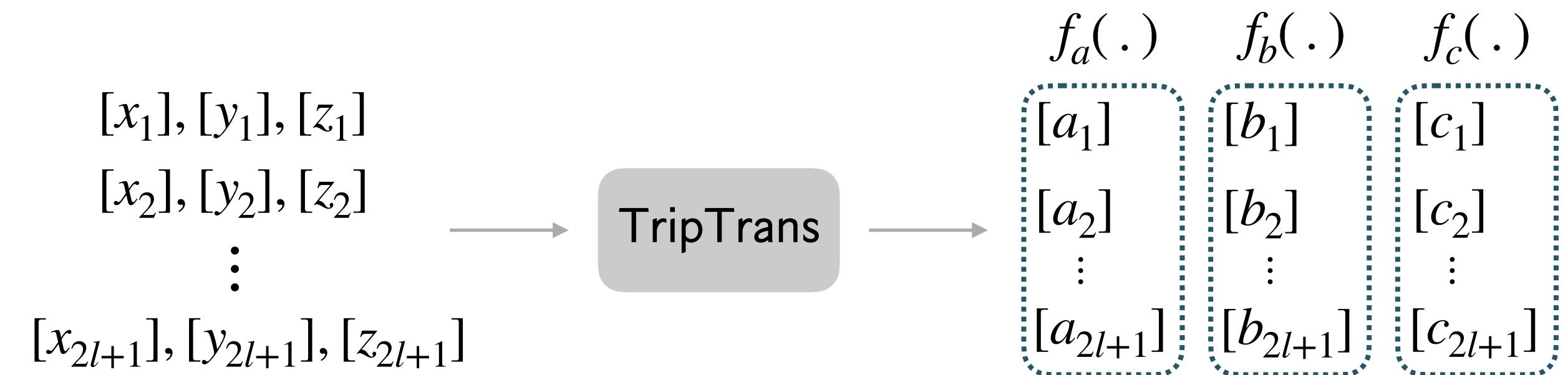
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$r^{\text{!}}$

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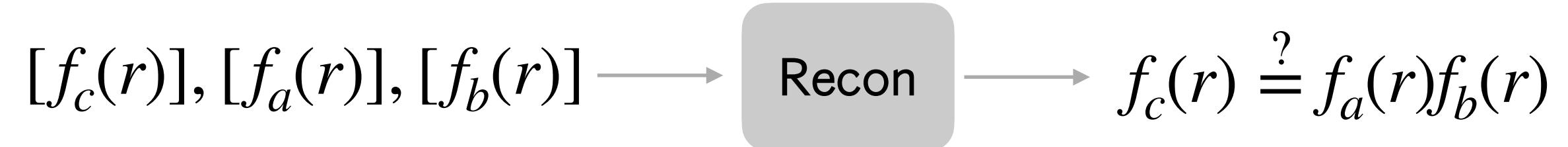
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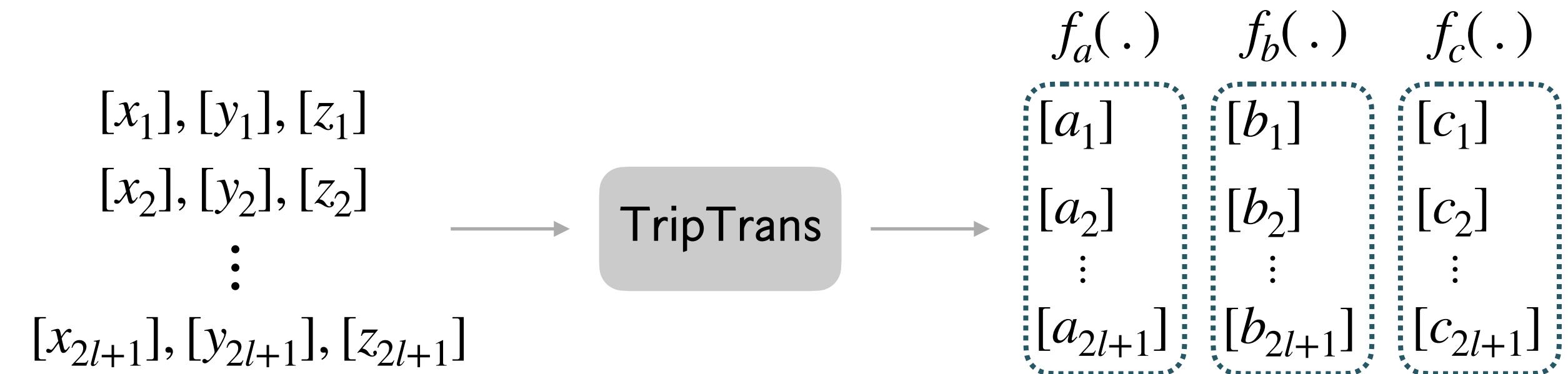
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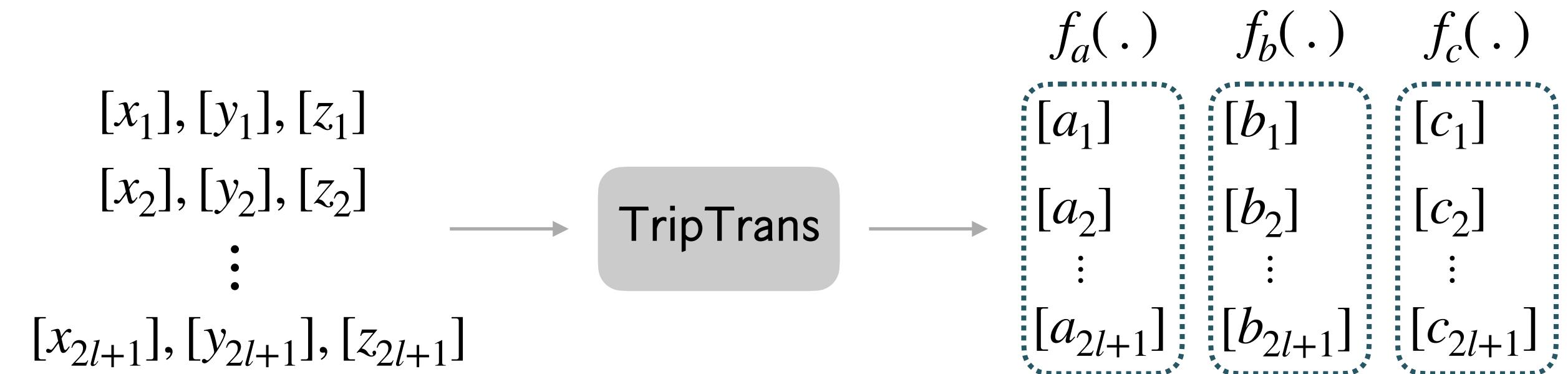
✓ l multiplication triples

✗ l shares of $([0], [0], [0])$

Cryptographically Secure HMPC - Triple Generation Protocol

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$r^{\text{!}}$

- Triple generation similar to

TripGen

- Each instance outputs l triples



✓ l multiplication triples

✗ l shares of $([0], [0], [0])$

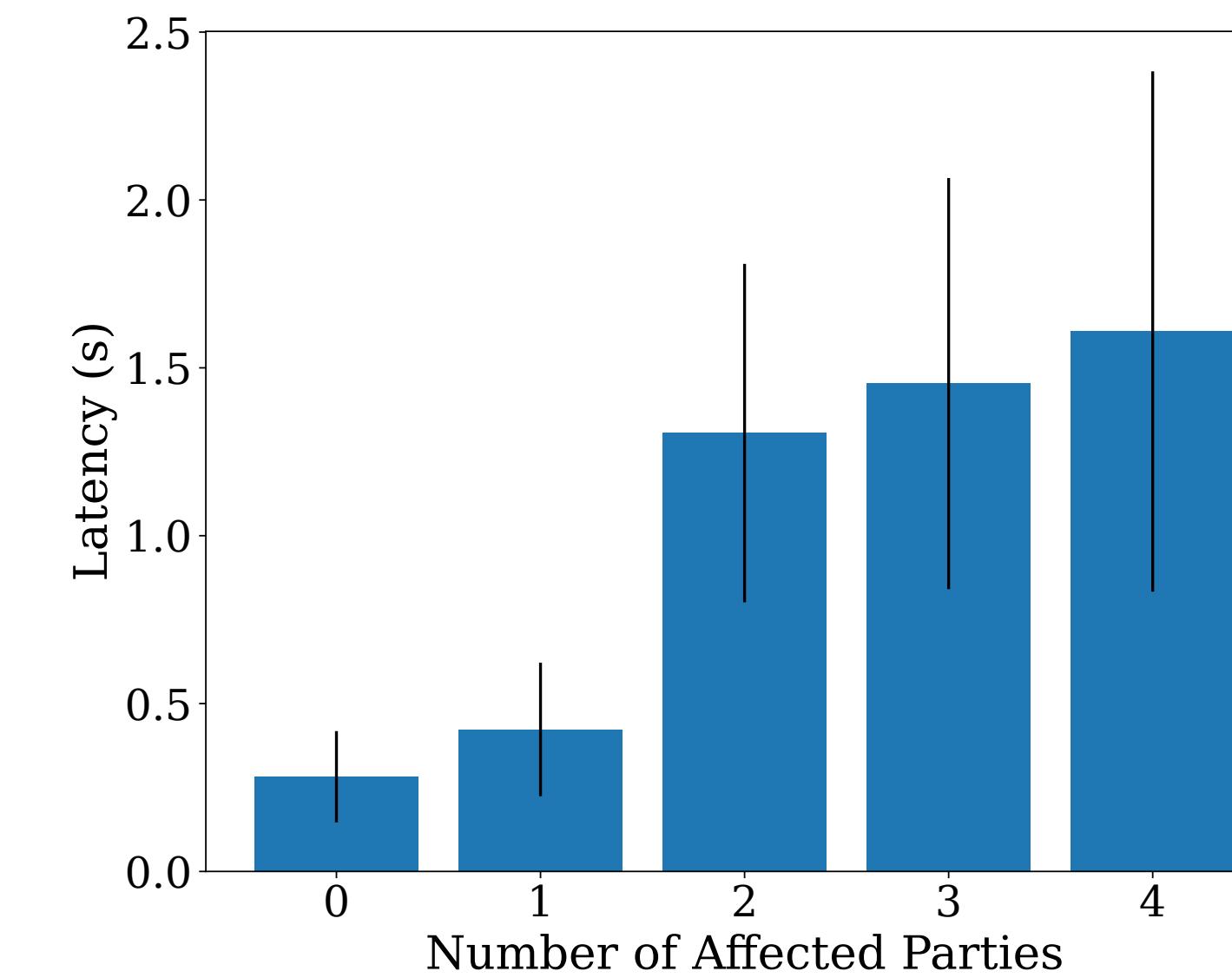
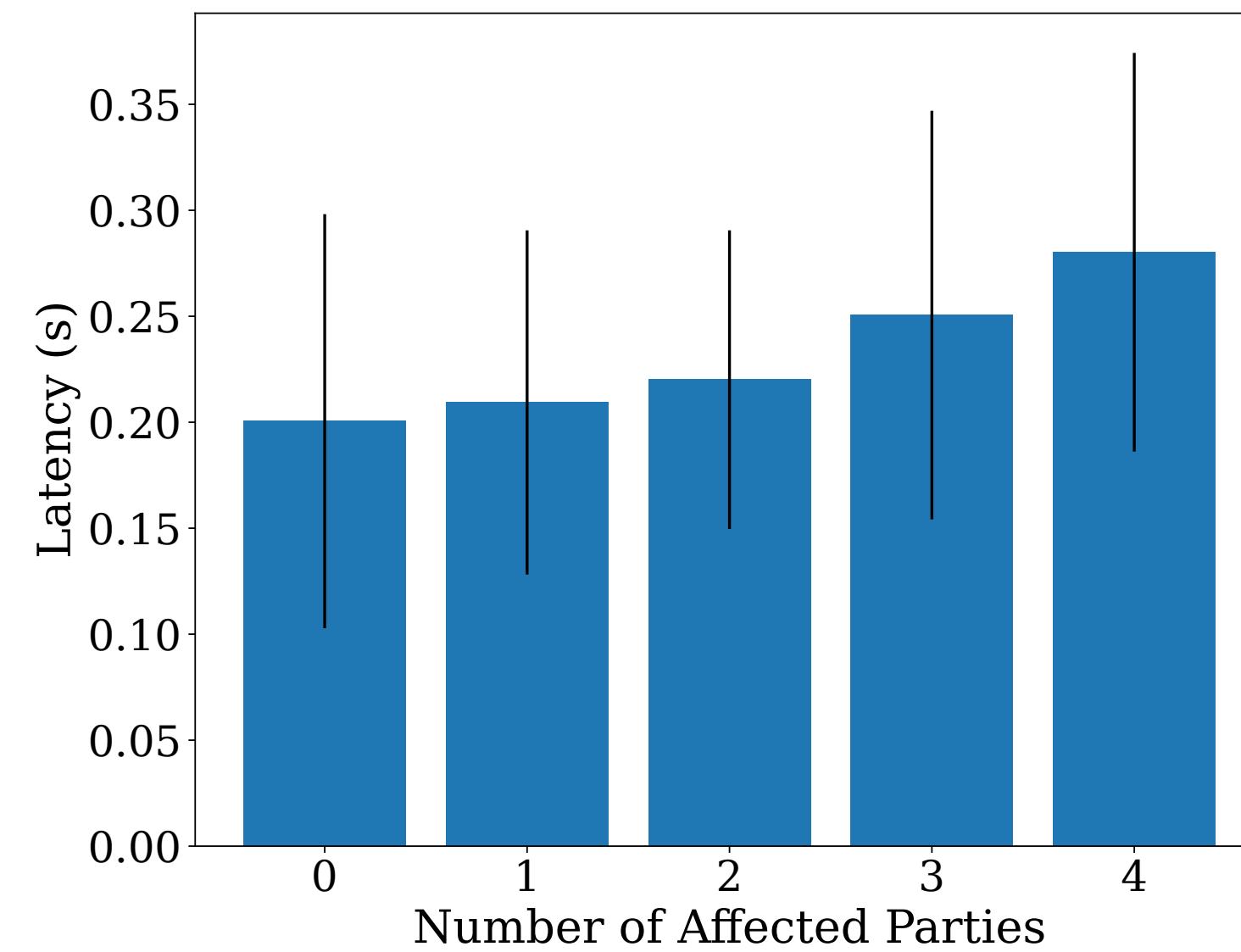
Cryptographically Secure HMPC and AMPC

- Cryptographically secure HMPC
 - Triple generation phase and input phase use 1 synchronous round
 - Circuit evaluation is completely asynchronous
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Cryptographically Secure HMPC and AMPC

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 - Triple generation phase and input phase use 1 synchronous round
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- Cryptographically secure AMPC
 - Similar to Cryptographically secure HMPC
 - No synchronous broadcast \implies ACast and ACS
 - No input provision

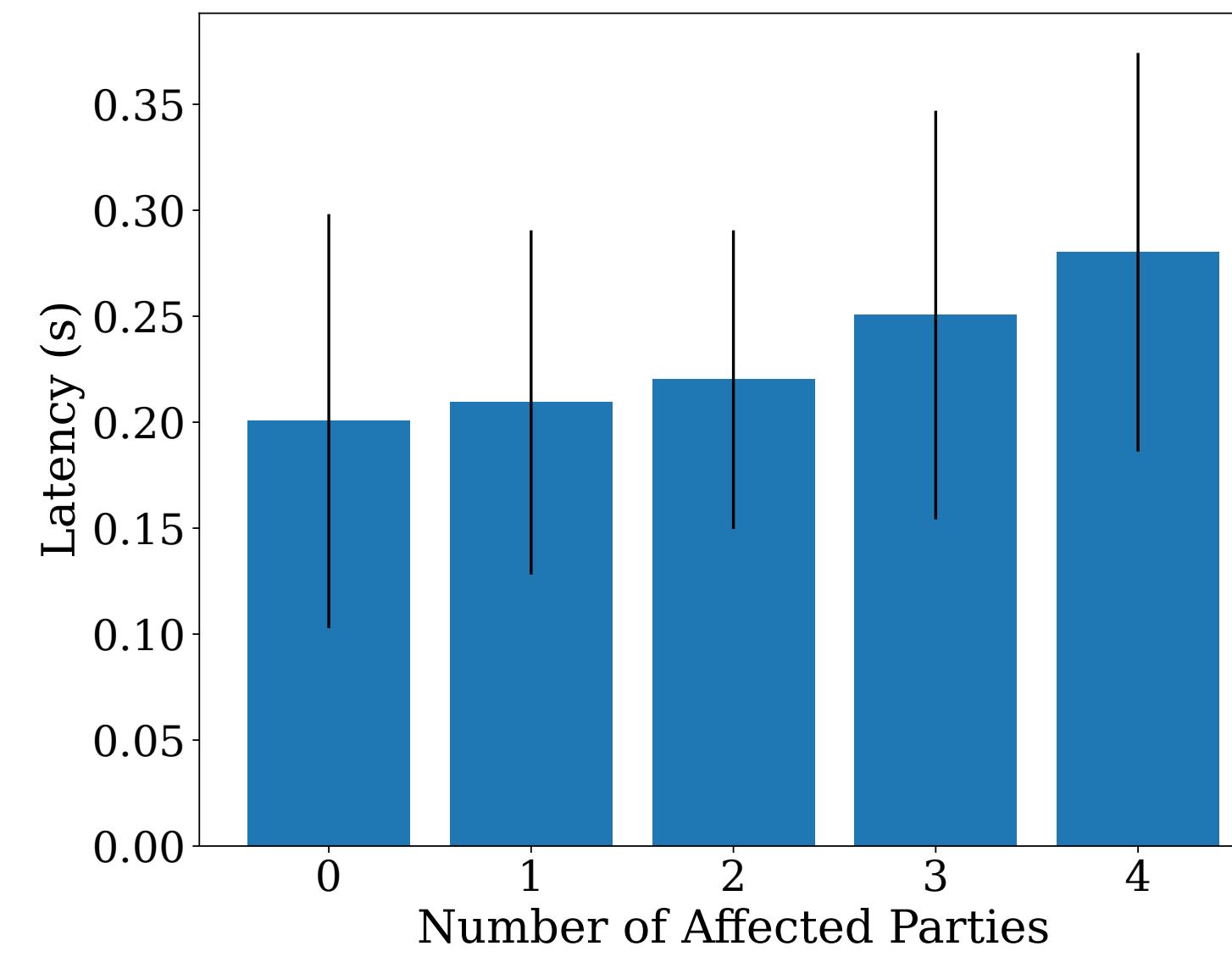
Conclusion



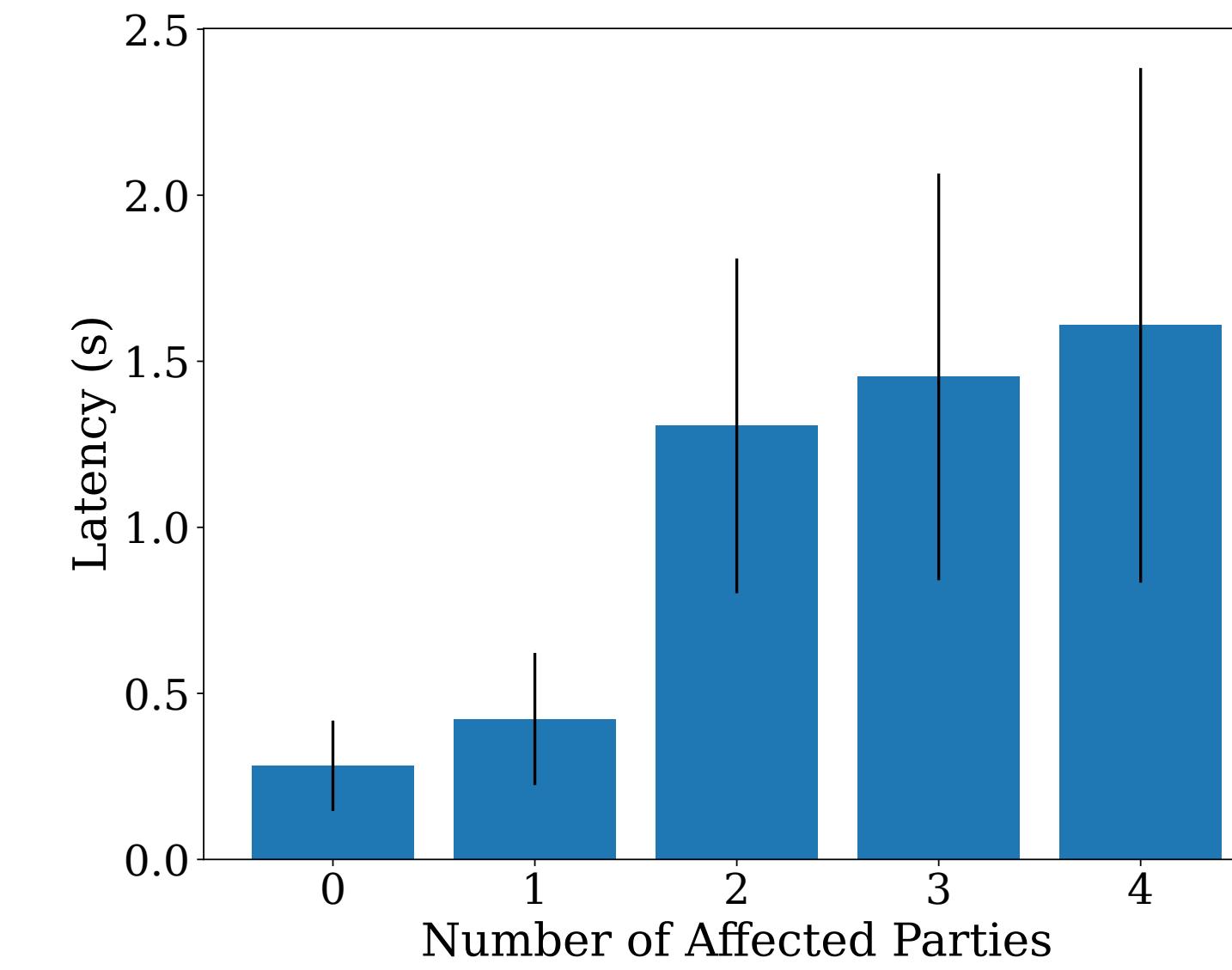
LAN

WAN

Conclusion



LAN



WAN

- Open problems
 - Perfect HMPC protocol for general case
 - Bridging the gap between synchronous and asynchronous MPC protocols

References

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Thank You