Lab Three

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CRAFTING A COMPILER

4.7.A

```
<Start>
< E > < \$ >
<T> plus <E><$>
<F> plus <E><$>
num plus \langle E \rangle \langle \$ \rangle
num plus <T><$>
num plus <T> times <F><$>
num plus <F> times <F><$>
num plus num times <F><$>
num plus num times (\langle E \rangle) \langle \$ \rangle
num plus num times ( <T> plus <E> ) <$>
num plus num times ( <F> plus <E> ) <$>
num plus num times ( num plus \langle E \rangle ) \langle \$ \rangle
num plus num times ( num plus <T> ) <math><$>
num plus num times ( num plus <F> ) <$>
num plus num times ( num plus num ) <$>
num plus num times ( num plus num ) $
```

4.7.B

```
<Start>
< E > < \$ >
<E> $
< T > \$
<T> times <F>$
<T> times num $
<F> times num $
(\langle E \rangle) times num $
(<T> plus <E>) times num $
(< T > plus < T > ) times num $
(< T > plus < F > ) times num $
(< T > plus num) times num $
(< T > plus num) times num $
( <F> plus num ) times num $
((<E>) plus num) times num $
( (<T>) plus num ) times num $
( (<T> times <F> ) plus num ) times num $
((<T> times num) plus num) times num $
((<F> times num) plus num) times num $
( ( num times num ) plus num ) times num $
```

4.7.C

This grammar uses the opposite of PEMDAS to go in opposite operational precedence, and it also favors left-associativity of operators. I say both these things because plus is higher up in the rules of the grammer, and the derivation looks nicer and took fewer steps when done left to right.

5.2c

```
Parser(tokens)
    self.tokens = tokens
    self.currentToken = none

parseStart()
    addBranchNode("start")
    parseValue()
    addLeafNode("$")

parseValue()
    addBranchNode("value")
    if currentToken is num:
        matchAndConsume({num})
    else:
        matchAndConsume({ ( ) )
        parseExpr()
        matchAndConsume({ ( ) ) }
    returnToParent()
```

```
parseExpr()
         addBranchNode("expr")
         if currentToken is plus:
              matchAndConsume({plus})
              parseValue()
              parseValue()
         else:
              matchAndConsume({prod})
              parseValues()
         returnToParent()
    parseValues()
         addBranchNode("values")
         if currentToken is not none:
              parseValue()
              parseValues()
              \# do nothing bc lambda
         returnToParent()
DRAGON
4.2.8
S \rightarrow S S + |S S*|a
string of a a + a *
<S>
<S><S>*
<S><S>+<S>*
a <S>+<S>*
a a + \langle S \rangle *
a a + a *
<S>
<S><S>*
<S> a *
<S><S>+ a *
\langle S \rangle a + a *
a a + a *
```

 \mathbf{C}

A

В



