

# EXPLICIT COMPUTATION WITH COLEMAN INTEGRALS

BU – KEIO WORKSHOP 2019

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## WHY DO WE INTEGRATE THINGS? LOGARITHMS

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## WHY DO WE INTEGRATE THINGS? LOGARITHMS

Take  $\frac{dx}{x}$ , as a differential on the group  $\mathbf{R}^\times$ , this is translation invariant, i.e.  $(a \cdot -)^*(dx/x) = d(ax)/ax = dx/x$ , hence

$$\int_1^t \frac{dx}{x} = \log |t|: \mathbf{R}^\times \rightarrow \mathbf{R}$$

has the property that

$$\int_1^{ab} \frac{dx}{x} = \int_a^{ab} \frac{dx}{x} + \int_1^a \frac{dx}{x} = \int_1^b \frac{dx}{x} + \int_1^a \frac{dx}{x}$$

Integration can define logarithm maps between groups and their tangent spaces.

How do we calculate  $\log |t|$ ? Power series on  $\mathbf{R}_{>0}$  and use the relation  $\log |t| = \frac{1}{2} \log t^2$

## WHY DO WE INTEGRATE THINGS? INTERESTING FUNCTIONS

We have already seen **polylogarithms**, defined recursively by

$$L_1(z) = -\log(1-z), \quad L_k(z) = \int_0^z L_{k-1}(s) \frac{ds}{s} : \mathbb{C} \setminus [1, \infty) \rightarrow \mathbb{C}$$

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These functions can alternatively be described via the power series

$$L_k(z) = \sum_{n=1}^{\infty} \frac{z^n}{n^k}$$

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For instance near a point  $\alpha$ :

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so that

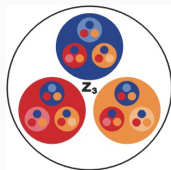
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Bad topology!

$$\int_{\alpha+x} \omega = - \sum \frac{1}{n+1} \left( \frac{-x}{\alpha} \right)^{n+1} + C$$

But we cannot find  $C$ ! There is a different choice in each disk.

## COLEMAN INTEGRATION: MORE PROBLEMS

Now we have functions

$$K\langle t \rangle = \left\{ \sum a_i t^i; a_i \in K, \lim_{i \rightarrow \infty} |a_i| = 0 \right\}$$

and

$$d: T \rightarrow \Omega_T^1$$

and our integral map should send

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So instead we work with a subring of **overconvergent** functions

$$\mathcal{T}^\dagger = \left\{ \sum a_i t^i; a_i \in K, \exists r > 1 \text{ such that } \lim_{i \rightarrow \infty} |a_i| r^i = 0 \right\}.$$

## COLEMAN'S THEOREM

Take  $X/\mathbf{Z}_p$  a genus  $g$  curve, and  $p$  an odd prime.

We pick a lift of the Frobenius map, i.e.  $\phi: X \rightarrow X$  which reduces to the Frobenius on  $X \times \mathbf{F}_p$ , and write  $A^\dagger$  (resp.  $A_{\text{loc}}(X)$ ) for overconvergent (resp. locally analytic) functions on  $X$ .

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### Theorem (Coleman)

There is a  $\mathbf{Q}_p$ -linear map  $\int_b^X: \Omega_{A^\dagger}^1 \otimes \mathbf{Q}_p \rightarrow A_{\text{loc}}(X)$  for which:

$$d \circ \int_b^X = \text{id}: \Omega_{A^\dagger}^1 \otimes \mathbf{Q}_p \rightarrow \Omega_{\text{loc}}^1 \quad \text{"FTC"}$$

$$\int_b^X \circ d = \text{id}: A^\dagger \hookrightarrow A_{\text{loc}}$$

$$\int_b^X \phi^* \omega = \phi^* \int_b^X \omega \quad \text{"Frobenius equivariance"}$$



## COMPUTATION: POLYLOGARITHMS ON $\mathbf{P}^1 \setminus \{0, 1, \infty\}$

Let's revisit the **polylogarithms**

$$L_1(z) = -\log(1-z), \quad L_k(z) = \int_0^z L_{k-1}(s) \frac{ds}{s} : \mathbb{C} \setminus [1, \infty) \rightarrow \mathbb{C}$$

Coleman integration then defines a  $p$ -adic analogue of these functions, with exactly the same definition via iterated integration on  $\mathbf{P}^1 \setminus \{0, 1, \infty\}$ .

(We must choose a branch of the  $p$ -adic logarithm, for simplicity we take the **Iwasawa logarithm** where  $\log_p(p) = 0$ .)

The power series definition still holds near  $z = 0$ , but otherwise we must use Frobenius equivariance to define it.

## COMPUTING POLYLOGARITHMS

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In exactly the same way as:

```
sage: x = RBF(1/3) # Real ball, or do pari(1/3)
sage: (x^2).polylog(4) - 8*x.polylog(4) -
      8*(-x).polylog(4)
[+/- 2.51e-14]
```

## COMPUTATION: GROUP STRUCTURE

If  $X/\mathbf{Q}_p$  is an algebraic group,  $\omega$  is a translation invariant 1-form we have

$$\int_0^{P+Q} \omega = \int_0^P \omega + \int_0^Q \omega \implies \int_0^P \omega = \frac{1}{n} \int_0^{nP} \omega$$

but if  $n = \#\tilde{X}(\mathbf{F}_p)$  then  $nP \in B(0, 1)$  so the integral on the right can be performed locally with only power series.

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but if  $n = \#\tilde{X}(\mathbf{F}_p)$  then  $nP \in B(0, 1)$  so the integral on the right can be performed locally with only power series.

This requires arithmetic in the group, which may be hard. And can only integrate invariant differentials.

## COMPUTATION: $p$ -ADIC COHOMOLOGY

There is an alternate approach via  $p$ -adic cohomology, due to Balakrishnan-Bradshaw-Kedlaya.

Let  $X/\mathbf{Z}_p$  be a smooth curve of good reduction.

Pick a basis  $\omega_1, \dots, \omega_{2g}$  for  $H_{\text{dR}}^1(X)$  and let  $U \subseteq X$  be an affine subspace containing no poles of any  $\omega_i$  and on which we have a lift of Frobenius  $\phi$ .

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If we apply  $\phi^*$  to  $\omega_i$  we may write

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$$\int_{\phi(b)}^{\phi(P)} \omega_i = \int_b^P \phi^*\omega_i = \int_b^P \left( \sum_{j=1}^{2g} M_{ij}\omega_j \right) - \int_b^P df_i$$

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$$\int_{\phi(b)}^{\phi(P)} \omega_i = \int_b^P \left( \sum_{j=1}^{2g} M_{ij} \omega_j \right) - (f_i(P) - f_i(b))$$

$$\Rightarrow \begin{pmatrix} \vdots \\ \int_b^P \omega_i \\ \vdots \end{pmatrix} = (M - I)^{-1} \begin{pmatrix} \vdots \\ f_i(P) - f_i(b) \\ \vdots \end{pmatrix} \text{ if } b = \phi(b), P = \phi(P)$$

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To move outside of  $U$  we have to either work close to the boundary of the removed disks (i.e. in a highly ramified extension). Or use tricks due to the special geometry of the curve (extra automorphisms).

## APPLICATIONS: CHABAUTY'S METHOD

Given  $X/\mathbf{Q}$  a smooth curve and  $p > 2 \cdot \text{genus}(X)$  a prime of good reduction for  $X$  and base point  $b \in X(\mathbf{Q})$ . If

$$\text{rank}(\text{Jac}(X))(\mathbf{Q}) < \text{genus}(X)$$

we can find a differential  $\omega_{\text{ann}} \in H^0(X, \Omega^1)$  such that

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**Note:** We can use either the group theory or  $p$ -adic cohomology method here.

## APPLICATIONS: CHABAUTY-KIM

Minhyong Kim has vastly generalised the above to cases where

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This can be made effective, and computable

### **Theorem (Balakrishnan-Dogra-Muller-Tuitman-Vonk)**

*The (cursed) modular curve  $X_{\text{split}}(13)$  (of genus 3 and jacobian rank 3), has 7 rational points: one cusp and 6 points that correspond to CM elliptic curves whose mod-13 Galois representations land in normalizers of split Cartan subgroups.*

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Their method can also be applied to other interesting curves:

### **Theorem (WIP B.-Bianchi-Triantafillou-Vonk)**

*The modular curve  $X_0(67)^+$  (of genus 2 and jacobian rank 2), has rational points contained in an explicitly computable finite set of 7-adic points.*

## MOTIVATING QUESTION

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For instance Harvey and Minzloff have introduced variants of Kedlaya's algorithm for hyper- and super-elliptic curves that works well when  $p$  is large!

They use interpolation to reduce the work when reducing

$$\phi^* \omega_j \rightsquigarrow \sum M_{ij} \omega_j$$

but its not clear where the functions  $f_i$  went.

Key to their interpolation is the fact that reductions in cohomology are linear in the exponents of  $x, y$ .

## SUPERELLIPTIC CURVES

We can write down a similar recurrence that evaluates the exact forms also, using

$$\left( \sum_{i=0}^N a_i x^i \right) = ((\cdots ((a_N)x + a_{N-1})x + \cdots)x + a_0)$$

### Theorem (B.)

Let

$$C/\mathbf{Z}_{p^n} : y^a = h(x)$$

with  $\gcd(a, \deg(h)) = 1$ ,  $p \nmid a$ , Let  $M$  be the matrix of Frobenius, acting on  $H_{\text{dR}}^1(C)$ , basis  $\{\omega_{i,j} = x^i dx/y^j\}_{i=0,\dots,b-2,j=1,\dots,a}$ , and points  $P, Q \in C(\mathbf{Q}_{p^n})$  known to precision  $p^N$ , if  $p > (aN - 1)b$ , the vector of Coleman integrals  $\left( \int_P^Q \omega_{i,j} \right)_{i,j}$  can be computed in time

$$\tilde{O}\left(g^3 \sqrt{p} n N^{5/2} + N^4 g^4 n^2 \log p\right)$$

to absolute precision  $N - v_p(\det(M - I))$ .