## Szemerédi's regularity lemma

31/6/23

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Recap: Studying Manin's conjecture and (equi)-distribution of

rational points on a Fano variety. Another instance of pseudorandomness emerging with scale is in extremal combinatorics / graph theory. Tao (paraphrased): "The various proofs of Szemerédi's theorem

and related theorems and proofs using measure theory, ergodic theory, graph theory, hypergraph theory, probability theory, information theory, and Fourier analysis share a number of common features and serve as a "Rosetta stone" for connecting these fields, notably they often use dichotomy between randomness

and structure". This time: Give more detail on Szemerédi's regularity lemma (SRL) and its proof, and the reduction to Roth's theorem / Szemerédi's theorem.

Slogan: The vertices of a sufficiently large graph can be partitioned into a fixed number of subsets in a way that the interactions between each behave pseudorandomly.

The statement

There are several variants of SRL, some place more restrictions on the partition (and so appear at first sight stronger). **Definition 1.** Let G = (V, E) be a graph and  $A, B \subseteq V$  two disjoint sets. The density of edges between A and B is

 $d(A,B) = \frac{e(A,B)}{|A||B|}$ We call a pair  $(A, B) \epsilon$ -regular (or uniform) for a given  $\epsilon > 0$ , if for all

 $X \subseteq A, Y \subseteq B$  where  $|X| \ge \epsilon |A|$  and  $|Y| \ge \epsilon |B|$ , we have:  $|d(X,Y) - d(A,B)| \le \epsilon$ 

$$\max\{||V_k| - |V_j||\} \le 1 \text{ or } \left\lfloor \frac{|V|}{k} \right\rfloor \le |V_j| \le \left\lceil \frac{|V|}{k} \right\rceil$$

**Definition 3.** A partition  $V_1, \ldots, V_k$  of the vertices of a graph is said to be  $\epsilon$ -regular if most pairs  $(V_i, V_j)$  are  $\epsilon$ -regular, in the sense that at

**Definition 2.** A partition  $V_1, \ldots, V_k$  of the vertices of a graph is said

to be an equipartition if it is as balanced as possible, i.e.

 $\epsilon \binom{k}{2}$  are not  $\epsilon$ -regular **Theorem 4** (Szemerédi). Let  $\varepsilon > 0$ , and let l be a natural number.

Then there exists an integer L such that every graph G with 
$$|G| \ge l$$
 has an  $\varepsilon$ -uniform equipartition into  $m$  parts for some  $m$  such that  $l \le m \le L$ .

small number of pieces relative to the size of the graph.

upper density contains a 3-term arithmetic progression.

Note that *L* does not depend on *G* or even the size of *G*, so that for large enough graphs we can consider this a partition into a

**Theorem 5** (Roth). A subset of the natural numbers with positive

Szemerédi was motivated by generalizing this result and proved:

**Theorem 6** (Szemerédi). A subset of the natural numbers with positive upper density contains a k-term arithmetic progression for any

Method

any  $n \ge n_0$  and any subset  $A \subseteq \{1, ..., n\}$  satisfying  $|A| \ge \delta n$ , there are distinct elements  $a, b, c \in A$  such that a + c = 2b.

 $B = \{(x, y) : x - y \in A\} \subseteq \{1, \dots, 2n\}^2$ 

**Definition 8** (Corners). A corner in a set  $B \subseteq \{1, ..., n\}^2$  is a triple

**Theorem 7** (Roth'). For every  $\delta > 0$ , there exists an  $n_0$  such that for

## of the form (x, y), (x + h, y), $(x, y + h) \in B$ , 0 < h (anticorner if h < 0). So, given a corner in B we get a 3 term AP (a, b, c) in A from

To prove this we instead let

(x-y,x+h-y,x-y-h).

 $|B| \ge \delta n^2$ , there is a corner in B.

meet at a point of *B*.

So we have at least

Proof of SRL

sets to the equipartition.

**Applications** 

corners). Then construct a tripartite graph where the triangles correspond

Reduce to the weak corners theorem (as above but allowing anti-

There are at least  $\delta n^2$  triangles in this graph from trivials alone so to remove all of them we must remove at least  $\delta n^2$  edges.

(If there are not too many triangles you can remove a small number of edges to remove all triangles.) (this is a special case of the more general Graph removal lemma and

Roth's theorem can be proved directly from the TRL, but the *Corners theorem* is in some sense a stronger version of Roth. We can prove the TRL from the so called triangle counting lemma, from SRL (though other proofs are available).

 We start with a trivial equipartition and may inductively apply this process which must eventually stop (after a number of steps bounded independently of the input graph) at

which point we have proved the lemma.

**Definition 11.** *The* energy *of a partition*  $P = \{V_1, \ldots, V_k\}$  *is* 

 $0 \le q_G(P) = \frac{1}{k^2} \sum_{1 \le i \le k} d_G(V_i, V_j)^2 \le 1$ 

**Lemma 12.** Let G be a graph of order n with an equipartition V =

 $|C_1| = |C_2| = \cdots = |C_k| = c \ge 2^{3k+1}$ .

Suppose that the partition  $\mathcal{P} = (C_i)_{i=0}^k$  is not  $\varepsilon$ -uniform, where  $0 < \varepsilon < \varepsilon$  $\frac{1}{2}$  and  $2^{-k} \leq \varepsilon^5/8$ . Then there is an equitable partition  $\mathcal{P}' = (C_i')_{i=0}^{\ell}$ 

 $q(\mathcal{P}') \ge q(\mathcal{P}) + \frac{\varepsilon^5}{2}$ 

The details of this are quite involved however!

 $\left|C_{ij}\right| \geq \varepsilon \left|C_{i}\right| = \varepsilon c, \left|C_{ji}\right| \geq \varepsilon \left|C_{j}\right| = \varepsilon c$  $\left|d\left(C_{ij},C_{ji}\right)-d\left(C_{i},C_{j}\right)\right|\geq\varepsilon$ We would like to partition simultaneously all possible  $C_{ij}$  into

new sets  $C_h$  so that each  $C_{ij}$  is a union of a bunch of  $C_h$ 's, which would have larger energy, this isn't quite possible, but we can "atomise" each  $C_i$  by making equivalent points not distinguishable by being in different  $C_{ij}$ s. We then pick  $H = 4^k - 2^{k-1}$  different

This new partition can be shown to have an energy of at least  $\epsilon^5/4$ 

 $|c/4^k|$ -subsets, all contained in some atom of  $C_i$ .

## **Theorem 9** (Corners theorem). For every $\delta > 0$ , there exists an $n_0$ such that for any $n \ge n_0$ and any subset $B \subseteq \{1, ..., n\}^2$ satisfying

to (anti or trivial)-corners of *B* and so that all triangles are edge disjoint.

The construction is to have vertices for horizontal, vertical and diagonal lines in  $\{1, \ldots, 2n\}$  and put an edge when two such lines

**Theorem 10** (Triangle Removal Lemma). For all 
$$1 \ge \delta > 0$$
, there exists  $\epsilon > 0$  such that any graph on  $n$  vertices with less than or equal to  $\epsilon n^3$  triangles can be made triangle-free by removing at most  $\delta n^2$  edges.

the *hypergraph removal lemma* (used for the full Szemerédi lemma on 
$$k$$
-term APs).)

Now the above graph must have at least  $\varepsilon n^3$  triangles as the triangle removal lemma does not apply.

 $\epsilon n^3 - \delta n^2$ 

nontrivial triangles, so pick n such that  $\epsilon n > \delta$  and we are done.

modification of the equipartition), without adding too many

To prove this we take, for each 
$$(C_i, C_j)$$
 non-uniform, some sets  $C_{ij} \subset C_i$ ,  $C_{ji} \subset C_j$  witnessing this so that

with  $\ell = k(4^k - 2^{k-1})$  such that

 $\bigcup_{i=0}^k C_i$ 

more.

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