



Computational Structures in Data Science






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Lecture #4: Recursion II and Higher Order Functions

Hackers steal medical data of US Olympic stars
<http://money.cnn.com/2016/09/13/news/wada-hacked-russian-spies/index.html?iid=surge-story-summary>
 This was: September 16, 2016

February 9th, 2018

<http://inst.eecs.berkeley.edu/~cs88>




Administrative issues

- Waitlist should be cleared.
- Based on your feedback: I will record optional lectures going deeper on data, code, algorithms, information, recursion, decision trees, run"time" complexity and other stuff.
- Speaking of recording: ETS will start capturing.

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Computational Concepts today

- More on Recursion
- Runtime (preliminary)
- Higher Order Functions
- Functions as Values
- Functions with functions as argument
- Assignment of function values
- Higher order function patterns
 - Map, Filter, Reduce
- Function factories – create and return functions

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Remember: Sanity Check...

- Recursion is ■ Iteration (i.e., loops)
 - a) more powerful than
 - b) just as powerful as
 - c) less powerful than

YOUR PARTY ENTERS THE TAVERN.

IT GATHER EVERYONE AROUND A TABLE. I HAVE THE ELVES START WHITTILING DICE AND GET OUT SOME PARCHMENT FOR CHARACTER SHEETS.


HEY, NO RECURSING.



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Why Recursion?

- "After Abstraction, Recursion is probably the 2nd biggest idea in this course"
- "It's tremendously useful when the problem is self-similar"
- "It's no more powerful than iteration, but often leads to more concise & better code"
- "It's more 'mathematical'"
- "It embodies the beauty and joy of computing"
- ...

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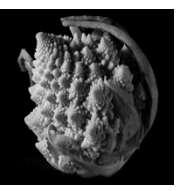
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Why Recursion? More Reasons

- Recursive structures exist (sometimes hidden) in nature and therefore in data!
- It's mentally and sometimes computationally more efficient to process recursive structures using recursion.

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Recursion (unwanted)



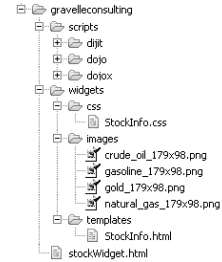
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Example I

List all items on your hard disk



- Files
- Folders contain
 - Files
 - Folders

Recursion!

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List Files in Python

```
def listfiles(directory):
    content = [os.path.join(directory, x) for x in os.listdir(directory)]

    dirs = sorted([x for x in content if os.path.isdir(x)])
    files = sorted([x for x in content if os.path.isfile(x)])

    for d in dirs:
        print d
        listfiles(d)

    for f in files:
        print f
```

Iterative version about twice as much code
and much harder to think about.

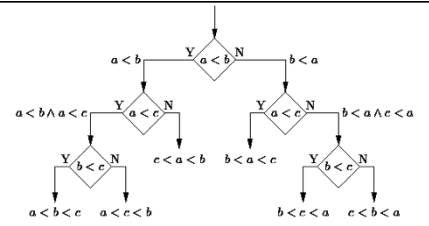
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Example II

Sort the numbers in a list.



Hidden recursive structure: Decision tree!

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Tree Recursion makes Sorting Efficient

Break the problem into multiple smaller sub-problems, and solve them recursively

```
def split(x, s):
    return [i for i in s if i <= x], [i for i in s if i > x]

def qsort(s):
    """Sort a sequence - split it by the first element,
    sort both parts and put them back together."""
    if not s:
        return []
    else:
        pivot = first(s)
        lessor, more = split(pivot, rest(s))
        return qsort(lessor) + [pivot] + qsort(more)

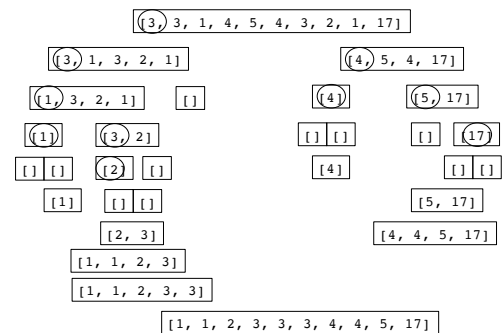
>>> qsort([3,3,1,4,5,4,3,2,1,17])
[1, 1, 2, 3, 3, 3, 4, 4, 5, 17]
```

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QuickSort Example



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More on Recursion

```
def sum_of_squares(n):
    if n < 1:
        return 0
    else:
        return n**2 + sum_of_squares(n-1)
```

- The sum of no numbers is zero
- The sum of 1^2 through n^2 is n^2 plus the sum of 1^2 through $(n-1)^2$

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Recap: Tail Recursion

- All the work happens on the way down the recursion
- On the way back up, just return

```
def sum_up_squares(i, n, accum):
    """Sum the squares from i to n in incr. order"""
    if i > n:
        Base Case
    else:
        Tail Recursive Case

>>> sum_up_squares(1, 3, 0)
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```

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How much ???

- “Time” is required to compute `sum_of_squares(n)`?
 - Recursively?
 - Iteratively?
- “Space” is required to compute `sum_of_squares(n)`?
 - Recursively?
 - Iteratively?
- Count the steps: Recursive is linear, iterative is constant! What about the order of evaluation?
- **Careful:** As taught traditionally, Computer Science has no measurement units convertible to physical time (s) and space (m^3)!

Linear
proportional to n
or for some c

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Recap: Defining Functions

```
def <function name> (<argument list>) :
```

return

expression

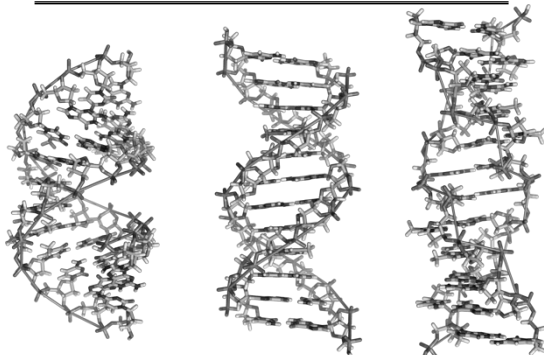
- Generalizes an expression or set of statements to apply to lots of instances of the problem
- A function should *do one thing well*

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Recap: Data or Code?



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Higher Order Functions

- Functions that operate on functions
- A function

```
def odd(x):
    return (x%2==1)

>>> odd(3)
True
```

Why is this
not 'odd'?

- A function that takes a function arg

```
def filter(fun, s):
    return [x for x in s if fun(x)]

>>> filter(odd, [0,1,2,3,4,5,6,7])
[1, 3, 5, 7]
```

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Higher Order Functions (cont)

- A function that returns (makes) a function

```
def leq_maker(c):
    def leq(val):
        return val <= c
    return leq

>>> leq_maker(3)
<function leq_maker.<locals>.leq at 0x1019d8c80>

>>> leq_maker(3)(4)
False

>>> filter(leq_maker(3), [0,1,2,3,4,5,6,7])
[0, 1, 2, 3]
>>>
```

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One more example

- What does this function do?

```
def split_fun(p, s):
    """ Returns <you fill this in>."""
    return [i for i in s if p(i)], [i for i in s if not p(i)]

>>> split_fun(leq_maker(3), [0,1,2,3,4,5,6])
([0, 1, 2, 3], [4, 5, 6])
```

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Three super important HOFs

`map(function_to_apply, list_of_inputs)`
Applies function to each element of the list

`filter(condition, list_of_inputs)`
Returns a list of elements for which the condition is true

`reduce(function, list_of_inputs)`
Reduces the list to a result, given the function

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Recursion with Higher Order Fun

```
def map(f, s):
    if Base Case:
    else:
        Recursive Case

def square(x):
    return x**2

>>> map(square, [2,4,6])
[4, 16, 36]
```

- Divide and conquer

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Using HOF to preserve interface

```
def sum_of_squares(n):
    def sum_upper(i, accum):
        if i > n:
            return accum
        else:
            return sum_upper(i+1, accum + i*i)
    return sum_upper(1,0)
```

- What are the globals and locals in a call to `sum_upper`?
– Try [python tutor](#)
- Lexical (static) nesting of function def within def - vs
- Dynamic nesting of function call within call

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Recap: Quicksort

- Break the problem into multiple smaller sub-problems, and Solve them recursively

```
def split(x, s):
    return [i for i in s if i <= x], [i for i in s if i > x]

def qsort(s):
    """Sort a sequence - split it by the first element,
    sort both parts and put them back together."""
    if not s:
        return []
    else:
        pivot = first(s)
        lessor, more = split(pivot, rest(s))
        return qsort(lessor) + [pivot] + qsort(more)

>>> qsort([3,3,1,4,5,4,3,2,1,17])
[1, 1, 2, 3, 3, 3, 4, 4, 5, 17]
```

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Quicksort with HOF

```
def qsort(s):
    """Sort a sequence - split it by the first element,
    sort both parts and put them back together."""

    if not s:
        return []
    else:
        pivot = first(s)
        lessor, more = split_fun(leg_maker(pivot), rest(s))
        return qsort(lessor) + [pivot] + qsort(more)

>>> qsort([3,3,1,4,5,4,3,2,1,17])
[1, 1, 2, 3, 3, 3, 3, 4, 4, 5, 17]
```

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How much ???

- “Time” is required to compute `quicksort(s)`?
 - “Space” is required?
 - Name of this recursion scheme?
 - Tree recursion
- Logarithmic time
 $c \cdot \log(\text{len}(s))$

Logarithmic to $\text{len}(s)$
 $c \cdot \log(\text{len}(s))$ for some c

Questions?

[illegible]

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