

Computational Geometry
[csci 3250]

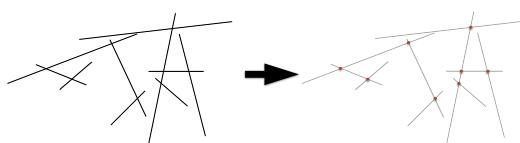
Line segment intersection

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Line segment intersection

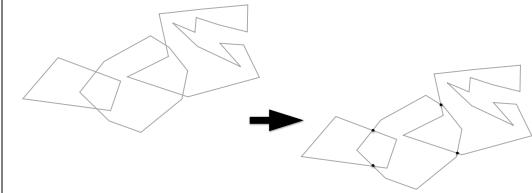
Problem: Given a set of line segments in 2D, find all their pairwise intersections.



2

Line segment intersection

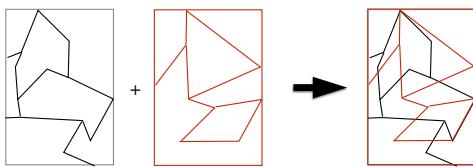
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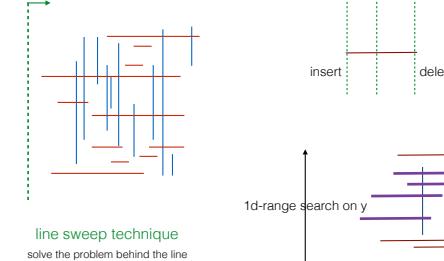
Line segment intersection

Problem: Given a set of line segments in 2D, find all their pairwise intersections.



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Orthogonal line segment intersection



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Orthogonal line segment intersection

- n : size of the input (number of segments)
- k : size of output (number of intersections)

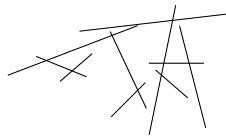
line sweep technique
solve the problem behind the line

Result: The intersections of a set of n orthogonal segments in the plane can be found in $O(n \lg n + k)$ time.

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General line segment intersection

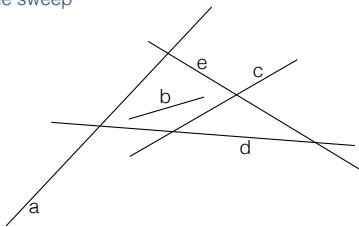
- n : size of the input (number of segments)
- k : size of output (number of intersections)



- Extend sweep line idea
- We'll get an overall bound of $O(n \lg n + k \lg n)$ which improves on the naive $O(n^2)$ when k is small
- The algorithm was developed by Jon Bentley and Thomas Ottmann in 1979
- Simple (once you see it) and practical

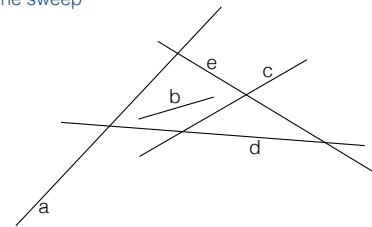
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The sweep



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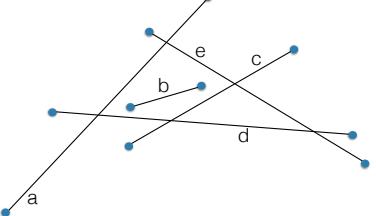
The sweep



- Let X be the set of all x-coords of segments

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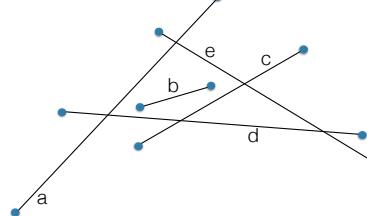
The sweep



- Let X be the set of all x-coords of segments

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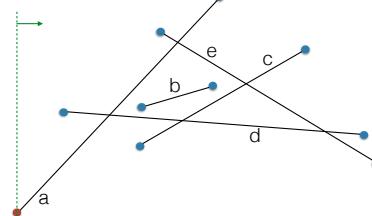
The sweep



- Let X be the set of all x-coords of segments
- Traverse the events in X in order

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The sweep



- Let X be the set of all x-coords of segments
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The sweep

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- Let X be the set of all x-coords of segments
- Traverse the events in X in order

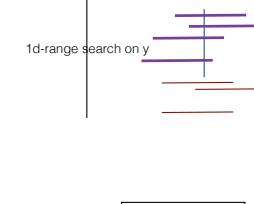
The sweep

- 4 segments are active
- How do we order these segments?
- How do we detect intersections?

How do we order these segments?

a below d a above d

general segments



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Key idea #1

- How do we order the active segments?
- Use above-below order
- Order will flip at intersection point

Key idea #2

- How do we detect intersections?
- Segments that intersect are consecutive in above-below order just before they intersect

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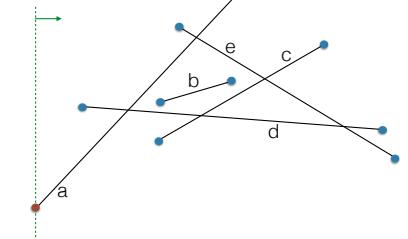
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Bentley-Ottmann sweep

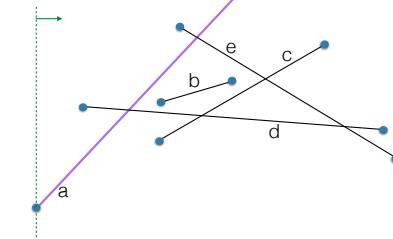
- Let X be the set of all x-coords of segments
- Initialize $AS = \{\}$
- Traverse events in order

Bentley-Ottmann sweep



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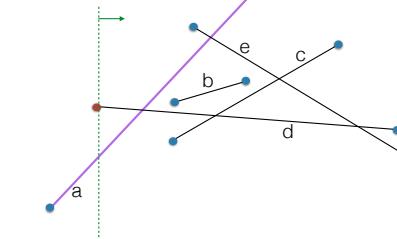
Bentley-Ottmann sweep



- this event is start of segment a:
 - insert a in AS: a

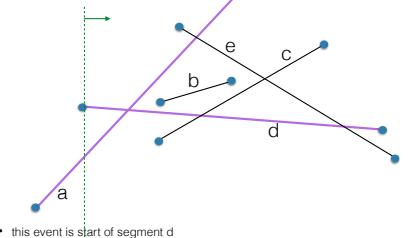
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Bentley-Ottmann sweep



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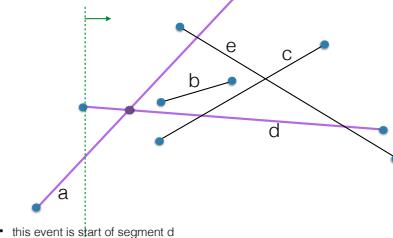
Bentley-Ottmann sweep



- this event is start of segment d
 - insert d in AS: a < d
- check if (d,a) intersect to the right of the line; they do; report point and insert it in the list of future events

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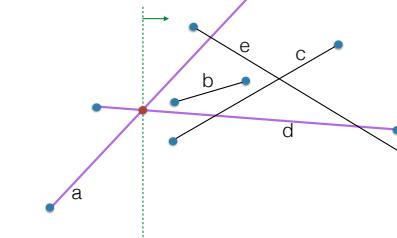
Bentley-Ottmann sweep



- this event is start of segment d
 - insert d in AS, d is above a (a < d)
- check if (d,a) intersect to the right of the line; they do; report point and insert it in the list of future events

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Bentley-Ottmann sweep



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Bentley-Ottmann sweep

- this event is an intersection point of (a,d):
 - flip a and d in AS: a is now above d ($d < a$)

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Bentley-Ottmann sweep

Bentley-Ottmann sweep

- this event is start of segment c:
 - insert c in AS; c is below d ($c < d < a$)
 - check c with its above and below neighbors for intersection to the right of the sweep line; this detects the intersection point of c and d; report it and insert it as future event

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Bentley-Ottmann sweep

- this event is start of segment c:
 - insert c in AS; c is below d ($c < d < a$)
 - check c with its above and below neighbors for intersection to the right of the sweep line; this detects the intersection point of c and d; report it and insert it as future event

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Bentley-Ottmann sweep

- this event is start of segment b:

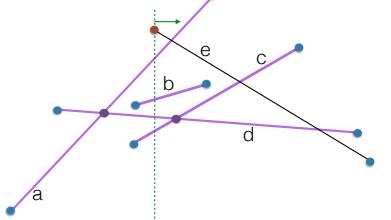
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Bentley-Ottmann sweep

- this event is start of segment b:
 - insert b in AS; $c < d < b < a$
 - check b with its above and below neighbors for intersection to the right of the sweep line; (d,b) don't intersect; (b, a) don't intersect

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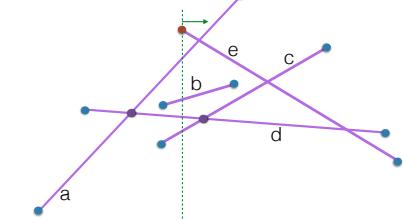
Bentley-Ottmann sweep



- this event is start of segment e:

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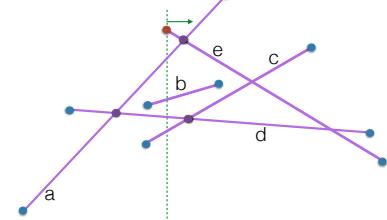
Bentley-Ottmann sweep



- this event is start of segment e:
 - insert e in AS: $c < d < b < a < e$
 - check e with its above and below neighbors for intersection to the right of the sweep line; this detects intersection point of (a,e); report it and insert it as future event

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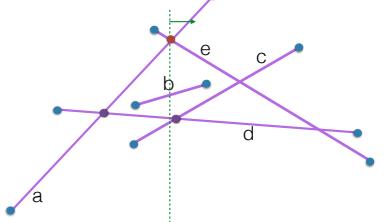
Bentley-Ottmann sweep



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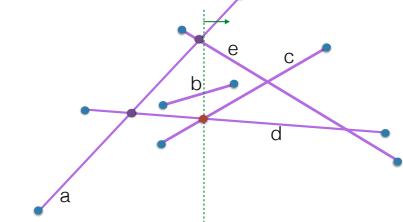
Bentley-Ottmann sweep



- this event is intersection of (a,e):
 - flip a and e: $c < d < b < e < a$
 - check new neighbors (e,b) for intersection to the right of the sweep line; (e,b) don't intersect

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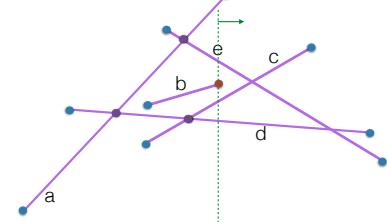
Bentley-Ottmann sweep



- this event is intersection of (c,d):
 - flip c and d: $d < c < b < e < a$
 - check new neighbors (c,b) for intersection to the right of the sweep line; (c,b) don't intersect

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Bentley-Ottmann sweep



- this event is end of segment b:
 - delete b from AS: $d < c < e < a$
 - check new neighbors (c,e) for intersection to the right of the sweep line; this detects the intersection point of (c,e); report it and insert it as future event

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Bentley-Ottmann sweep

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- this event is end of segment b:
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Bentley-Ottmann sweep

-
- this event is end of segment a:
 - delete a from AS: $d < c < e$
 - no new neighbors

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Bentley-Ottmann sweep

-
- this event is end of segment a:
 - delete a from AS: $d < c < e$
 - no new neighbors

46

Bentley-Ottmann sweep

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- this event is the intersection of (c,e):
 - flip c,e in AS: $d < e < c$
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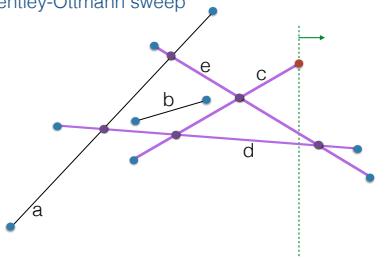
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Bentley-Ottmann sweep

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- this event is the intersection of (c,e):
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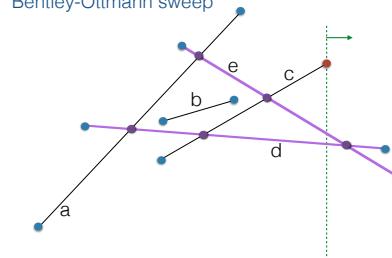
Bentley-Ottmann sweep



- this event is end of segment c:
 - delete c in AS: $d < e$
 - no new neighbors

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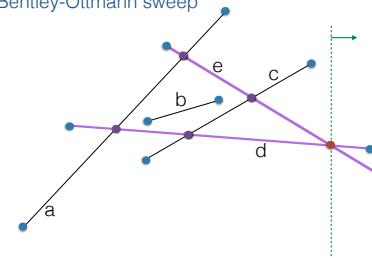
Bentley-Ottmann sweep



- this event is end of segment c:
 - delete c in AS: $d < e$
 - no new neighbors

50

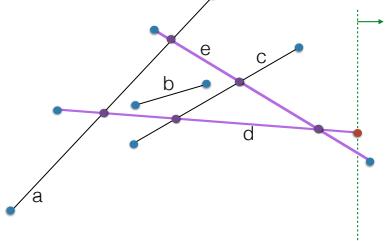
Bentley-Ottmann sweep



- this event is the intersection of (d,e):
 - flip d,e in AS: $e < d$
 - no new neighbors

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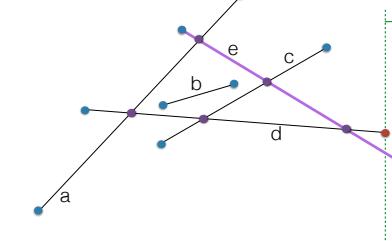
Bentley-Ottmann sweep



- this event is the end of d:
 - delete d in AS: e
 - no new neighbors

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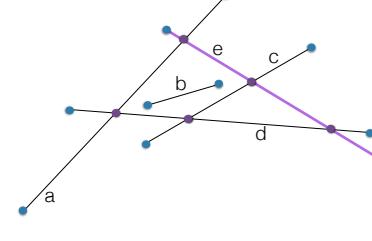
Bentley-Ottmann sweep



- this event is the end of d:
 - delete d in AS: e
 - no new neighbors

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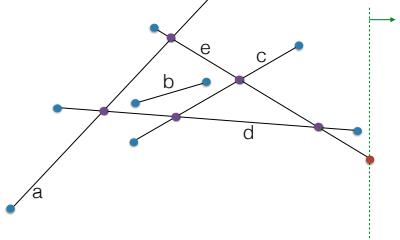
Bentley-Ottmann sweep



- this event is the end of e:
 - delete e in AS:
 - no new neighbors

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Bentley-Ottmann sweep



- this event is the end of e:
 - delete e in AS:
 - no new neighbors

55

Bentley-Ottmann sweep

- Simplifying assumptions
 - no vertical segments
 - no two segments intersect at their endpoints
 - no three (or more) segments have a common intersection
 - all endpoints (of segments) and all intersection points have different x-coordinates
 - no segments overlap
- These assumptions are not realistic for real data.
- But, they don't provide insight into the plane sweep technique, so we omit them
- Real data challenges
 - dealing with degenerate cases
 - dealing with finite precision arithmetic and precision problems

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Bentley-Ottmann sweep

We'll maintain the following invariants during the algorithm:

Bentley-Ottmann sweep

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- Active structure AS:
 - For any position of the sweep line SL, AS contains all active segments (ie segments that start before SL and end after SL)
 - AS is sorted by their y-coordinates of their intersection with SL

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Bentley-Ottmann sweep

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- Active structure AS:
 - For any position of the sweep line SL, AS contains all active segments (ie segments that start before SL and end after SL)
 - AS is sorted by their y-coordinates of their intersection with SL
- Event list EL:
 - For any position of SL, EL contains segment endpoints to the right of SL, and also the intersections to the right of SL of active segments that were/ are neighbors in SL
 - EL is sorted by x-coordinate

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Bentley-Ottmann sweep

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- Active structure AS:
 - For any position of the sweep line SL, AS contains all active segments (ie segments that start before SL and end after SL)
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- Event list EL:
 - For any position of SL, EL contains segment endpoints to the right of SL, and also the intersections to the right of SL of active segments that were/ are neighbors in SL
 - EL is sorted by x-coordinate
- For any position of the sweep line SL, all pairs of intersecting dead segments have been reported.

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Algorithm Bentley-Ottmann (S)

```
//S is a set of n line segments in the plane
• initialize AS = {}
• sort 2n endpoints of all segments in S by x-coord and store them in EventList
• while EventList not empty
    • let e be the next event from EventList; delete it from EL
    //sweep line moves to x=e.x
    • if e is left endpoint of a segment I
        //I becomes active
        • insert I in AS in the right place
        • check if I intersects with I->prev and I->succ in AS to the right of the sweep
          line; if they do, insert their intersection point in the EventList
          //optional: since I.prev and I.succ are not neighbors anymore, we check if they
          intersect and if they do, delete that intersection point from the EventList
    • if e is the right endpoint of a segment
        • ...
    • if e is the intersection of two segments
        • ...
    • end.
```

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Bentley-Ottmann sweep

Questions

- AS
 - What operations do we do on AS?
 - What data structure should we use for AS?
- EL
 - Note that we know a priori the 2n events corresponding to start and end-points of segments, but EL is not static; the events corresponding to intersection points are generated on the fly
 - What operations do we do on EL?
 - What data structure should we use for EL?

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Bentley-Ottmann sweep

Running time

- AS
 - What is the size of AS?
 - $O(n)$
 - How many operations?
 - $O(n+k)$
 - Overall time?
 - $O((n+k)\lg n)$

EventList

- What is the size of EventList?
- $O(n+k)$
- How many operations?
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Bentley-Ottmann sweep

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Bentley-Ottmann sweep runs in $O((n+k)\lg n)$ time.

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