

# MACM 316 Lecture 13

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When we use iterative matrix techniques, we want to know when powers of a matrix become small.

**Def.** We call an  $n \times n$  matrix  $A$  convergent if

$$\lim_{k \rightarrow \infty} (A^k)_{ij} = 0, \text{ for all } i, j$$

**Ex.** Consider  $A = \begin{bmatrix} \frac{1}{2} & 0 \\ 16 & \frac{1}{2} \end{bmatrix}$

$$A^2 = \begin{bmatrix} \frac{1}{4} & 0 \\ 16 & \frac{1}{4} \end{bmatrix}$$

$$A^3 = \begin{bmatrix} \frac{1}{8} & 0 \\ 12 & \frac{1}{8} \end{bmatrix}$$

$$A^4 = \begin{bmatrix} \frac{1}{16} & 0 \\ 8 & \frac{1}{16} \end{bmatrix}$$

$$A^k = \begin{bmatrix} \frac{1}{2^k} & 0 \\ P_k & \frac{1}{2^k} \end{bmatrix}$$

$$\text{where } P_k = \begin{cases} 16 & k = 1 \\ \frac{16}{2^{k-1}} + \frac{1}{2}P_{k-1} & k > 1. \end{cases}$$

Since  $\lim_{k \rightarrow \infty} P_k = 0$ , we also know that  $\lim_{k \rightarrow \infty} P_k = 0$ .  $\therefore A$  is a convergent matrix.

Notice that this convergent matrix has a spectral radius (see Lecture 12 notes, page 5) less than 1.

This generalizes:

**Thm.** The following statements are equivalent:

1.  $A$  is a convergent matrix.
2.  $\rho(A) < 1$
3.  $\lim_{n \rightarrow \infty} A^n x = 0$  for every  $x$
4.  $\lim_{n \rightarrow \infty} \|A^n\| = 0$  for all natural norms  $\|\cdot\|$

Iterative techniques convert the system  $Ax = b$  into an equivalent system of the form  $x = Tx + c$  where  $T$  is a fixed matrix and  $c$  is a vector. An initial vector  $x^{(0)}$  is chosen, and then a sequence of approximate solution vectors is generated:

$$x^{(k)} = Tx^{(k-1)} + c$$

Iterative techniques are rarely used in very small systems (i.e. when  $n^3$  is small). In these cases, iterative techniques may be slower since they require several iterations to obtain the desired accuracy.

**IDEA:** It is possible to “split” the matrix  $A$  :

$$\begin{aligned} Ax &= b \\ [M + (A - M)]x &= b \\ Mx &= b + (M - A)x \\ x &= (I - M^{-1}A)x + M^{-1}b \end{aligned}$$

Iteration becomes

$$x^{(k+1)} = (I - M^{-1}A)x^{(k)} + M^{-1}b$$

We set  $T \cong I - M^{-1}A$  (the amplification matrix) and  $c \cong M^{-1}b$ .

$$x^{(k+1)} = Tx^{(k)} + c$$

How do we choose  $M$ ?

We want:

1.  $M$  easy to “invert”

2.  $M$  “close to  $A$ ” in the sense that  $\rho(T)$  is small.

**Ex.** Let  $M = D = \begin{bmatrix} a_{11} & 0 & 0 \\ 0 & \ddots & 0 \\ 0 & 0 & a_{nn} \end{bmatrix}$

This gives the Jacobi Iterative Method.

In the text’s notation,

$$A = D - L - U$$

Where  $D$  is diagonal,  $L$  is lower triangular, and  $U$  is upper triangular.

$$\begin{aligned} Ax &= b \\ (D - L - U)x &= b \\ Dx &= (L + U)x + b \\ x &= D^{-1}(L + U)x + D^{-1}b \end{aligned}$$

Which results in the iteration

$$x^{(k+1)} = D^{-1}(L + U)x^{(k)} + D^{-1}b$$

Let  $T = D^{-1}(L + U)$  and  $c = D^{-1}b$ .

$$x^{(k+1)} = Tx^{(k)} + c$$

See example in the notes, there are too many matrices to type out in LaTeX. See “Chapter 7.pdf” page 27 (7-36.7)

### 0.0.1 Comments on Jacobi’s Method

$$x^{(k+1)} = D^{-1}(L + U)x^{(k)} + D^{-1}b$$

1. The algorithm requires  $a_{ii} \neq 0$  for  $i = 1, \dots, n$ . If one of the  $a_{ii} = 0$ , and the system is nonsingular, then a reordering of the equations can be performed so that no  $a_{ii} = 0$ .
2. To speed convergence, the equations should be arranged such that  $|a_{ii}|$  is as large as possible.

3. A possible stopping criterion is to iterate until  $\frac{\|x^{(k)} - x^{(k-1)}\|}{\|x^{(k-1)}\|} \leq \epsilon$

If we write out Jacobi's Method

$$x^{(k+1)} = D^{-1}(L + U)x^{(k)} + D^{-1}b$$

we find that

$$x_i^{(k+1)} = \frac{\sum_{j=1; j \neq i}^n (-a_{ij}x_j^{(k)}) + b_i}{a_{ii}}$$

Notice that to compute  $x_i^{(k+1)}$ , the components  $x_i^{(k)}$  are used. But, for  $i > 1$ ,  $x_1^{(k+1)}, x_2^{(k+1)}, \dots, x_n^{(k+1)}$  have already been computed and are likely better approximations to the actual solutions than

$$x_1^{(k)}, x_2^{(k)}, \dots, x_n^{(k)}$$

So it seems reasonable to compute with these most recently computed values.

**i.e.:**

$$x_i^{(k+1)} = \frac{-\sum_{j=1}^{i-1} (a_{ij}x_j^{(k+1)}) - \sum_{j=i+1}^n (a_{ij}x_j^{(k)}) + b_i}{a_{ii}}$$

This is called the Gauss-Seidel iterative technique, and it also has a matrix formulation with  $M \cong (D - L)$  :

$$\begin{aligned} Ax &= b \\ (D - L - U)x &= b \\ (D - L)x &= Ux + b \\ x &= (D - L)^{-1}Ux + (D - L)^{-1}b \end{aligned}$$

$\implies$  iteration becomes

$$x^{(k+1)} = (D - L)^{-1}Ux^{(k)} + (D - L)^{-1}b$$

\*Notice that  $D - L$  is lower triangular. It is invertible  $\iff$  each  $a_{ii} \neq 0$