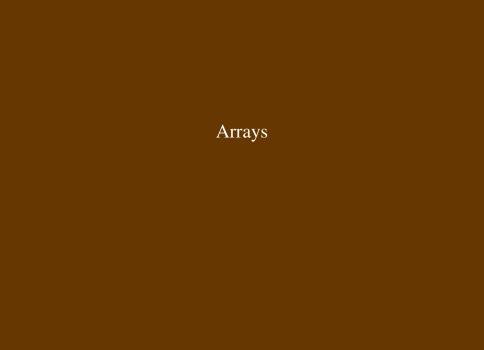


Modern Fortran Programming II

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Outline

- Arrays
- 2 Procedures
- 3 Derived Types
- 4 Object Based Programming
- 5 Exercise
 - Day 1 Exercises

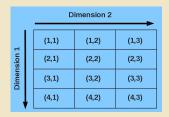


Arrays

- Arrays (or matrices) hold a collection of different values at the same time.
- Individual elements are accessed by subscripting the array.
- A 10 element array is visualized as

	1	2	3	 8	9	10
- 4						

while a 4x3 array as



Each array has a type and each element of holds a value of that type.

Array Declarations

- The dimension attribute declares arrays.
- Usage: dimension (lower_bound:upper_bound)

 Lower bounds of one (1:) can be omitted
- Examples:

```
integer, dimension(1:106) :: atomic_number
real, dimension(3,0:5,-10:10) :: values
character(len=3), dimension(12) :: months
```

• Alternative form for array declaration

```
integer :: days_per_week(7), months_per_year(12)
real :: grid(0:100,-100:0,-50:50)
complex :: psi(100,100)
```

• Another alternative form which can be very confusing for readers

```
integer, dimension(7) :: days_per_week, months_per_year(12)
```

Array Terminology

```
real :: a(0:20), b(3,0:5,-10:10)
   Rank: Number of dimensions.
          a has rank 1 and b has rank 3
 Bounds: upper and lower limits of each dimension of the array.
          a has bounds 0:20 and b has bounds 1:3, 0:5 and -10:10
  Extent: Number of element in each dimension
          a has extent 21 and b has extents 3,6 and 21
    Size: Total number of elements.
          a has size 21 and b has 30
  Shape: The shape of an array is its rank and extent
          a has shape 21 and b has shape (3,6,21)
```

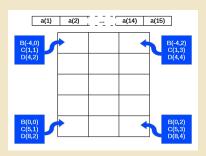
- Arrays are conformable if they share a shape.
- The bounds do not have to be the same

```
c(4:6) = d(1:3)
```

Array Visualization

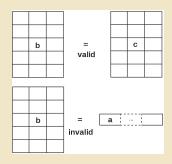
• Define arrays a, b, c and d as follows

```
real,dimension(15) :: a
real,dimension(-4:0,0:2) :: b
real,dimension(5,3) :: c
real,dimension(4:8,2:4) :: d
```



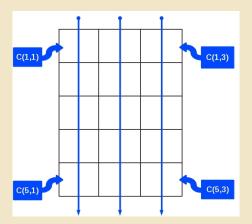
Array Conformance

- Array or sub-arrays must conform with all other objects in an expression
 - ① a scalar conforms to an array of any shape with the same value for every element c = 1.0 is the same as c(:,:) = 1.0
 - 2 two array references must conform in their shape.



Array Element Ordering

• Fortran is a column major form i.e. elements are added to the columns sequentially. This ordering can be changed using the reshape intrinsic.



Array Constructors I

• Used to give arrays or sections of arrays specific values

```
implicit none
integer :: i
integer, dimension(10) :: ints
character(len=5), dimension(3) :: colors
real, dimension(4) :: height
height = (/5.10, 5.4, 6.3, 4.5 /)
colors = (/'red ', 'green', 'blue ' /)
ints = (/ 30, (i = 1, 8), 40 /)
```

• constructors and array sections must conform.

```
ints = (/30, (i = 1, 10), 40/) is invalid
```

- strings should be padded so that character variables have correct length.
- use reshape intrinsic for arrays for higher ranks
- (i = 1, 8) is an implied do.
- You can also specify a stride in the implied do.

```
ints = (/30, (i = 1, 16, 2), 40/)
```

• There should be no space between / and (or)

Array Constructors II

- reshape (source, shape, pad, order) constructs an array with a specified shape shape starting from the elements in a given array source.
- If pad is not included then the size of source has to be at least product (shape).
- If pad is included it has to have the same type as source.
- If order is included, it has to be an integer array with the same shape as shape and the values must be a permutation of (1,2,3,...,N), where N (max value is 7) is the number of elements in shape.

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```

• In Fortran, for a multidimensional array, the first dimension has the fastest index while the last dimension has the slowest index i.e. memory locations are continuous for the last dimension.

Array Constructors III

- The order statement allows the programmer to change this order. The last example above sets the memory location order which is consistent to that in C/C++.
- Arrays can be initialized as follows during variable declaration

```
integer, dimension(4) :: imatrix = (/ 2, 4, 6, 8/)
character(len=*), dimension(3) :: colors = (/'red ', 'green', 'blue '/)}
! All strings must be the same length)
real, dimension(4) :: height = (/5.10, 5.4, 6.3, 4.5/)
integer, dimension(10) :: ints = (/ 30, (i = 1, 8), 40/)
real, dimension(4,3), parameter :: rcell = reshape( (/0.d0, 0.d0, 0.
```

Array Syntax

- Arrays can be treated as a single variable when performing operations
 - set whole array to a constant: a = 0.0
 - 2 can use intrinsic operators between conformable arrays (or sections)

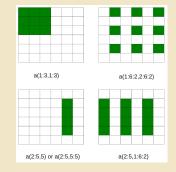
```
b = c * d + b**2
this is equivalent to
```

- 6 elemental intrinsic functions can be used: b = sin(c) + cos(d)
- 4 All operations/functions are applied element by element

Array Sections I

real, dimension(6:6):: a

- a(1:3,1:3) = a(1:6:2,2:6:2) and a(1:3,1:3) = 1.0 are valid
- a(2:5,5) = a(2:5,1:6:2) and a(2:5,1:6:2) = a(1:6:2,2:6:2) are not
- a (2:5, 5) is a 1D section while a (2:5, 1:6:2) is a 2D section



- The general form for specifying sub-arrays or sections is |<bound1>|:|<bound2>||:<stride>|
- The section starts at

 bound1> and ends at or before

 bound2>.
- <stride> is the increment by which the locations are selected, by default stride=1
- <bound1>, <bound2>, <stride> must all be scalar integer expressions.

Array Sections II

```
real, dimension(1:20) :: a
integer :: m,n,k
```

a(:)

```
a (3:9)
a (3:9:1)
a (m:n)
a (m:n:k)
a (15:3:-2)
a (15:3)
a (m:)
a (:n)
a (::2)
a (m:m)
a (m)
are valid sections.
```

```
the whole array elements 3 to 9 in increments of 1 as above elements m through n elements m through n in increments of k elements 15 through 3 in increments of -2 zero size array elements m through 20, default upper bound elements 1, default lower bound through n all elements from lower to upper bound in increments of 2 1 element section array element not a section
```

Array I/O I

```
real, dimension(4,4):: a
```

• Arrays are printed in the order that they appear in memory

```
print *, a
would produce on output
a(1,1),a(2,1),a(3,1),a(4,1),a(1,2),a(2,2),...,a(3,4),a(4,4)
read *, a
```

would read from input and assign array elements in the same order as above

• The order of array I/O can be changed using intrinsic functions such as reshape, transpose Or cshift.

• Example: consider a 3x3 matrix

1	4	7	
2	5	8	
3	6	9	

• The following print statements

```
print *, 'array element = ',a(3,3)
print *, 'array section = ',a(:,2)
print *, 'sub-array = ',a(:3,:2)
print *, 'whole array = ',a
print *, 'array transpose = ',transpose(a)
```

• would produce the following output

```
array element = 9
array section = 4 5 6
sub-array = 1 2 3 4 5 6
whole array = 1 2 3 4 5 6 7 8 9
array transpose = 1 4 7 2 5 8 3 6 9
```

Array Intrinsic Functions I

```
size(x[,n]) The size of x (along the n^{th} dimension, optional)
     shape(x) The shape of x
lbound(x[,n]) The lower bound of x
ubound(x[,n]) The upper bound of x
    minval(x) The minimum of all values of x
    maxval(x) The maximum of all values of x
    minloc(x) The indices of the minimum value of x
    \frac{\text{maxloc}(x)}{\text{maxloc}(x)} The indices of the maximum value of x
   sum(x[,n]) The sum of all elements of x (along the n^{th} dimension, optional)
                sum(x) = \sum_{i,j,k,\dots} x_{i,j,k,\dots}
```

Array Intrinsic Functions II

- product(x[,n]) The product of all elements of x (along the n^{th} dimension, optional) $prod(x) = \prod_{i,j,k,\dots} x_{i,j,k,\dots}$
- transpose(x) Transpose of array x: $x_{i,j} \Rightarrow x_{j,i}$
- $dot_product(x,y)$ Dot Product of arrays x and y: $\sum_i x_i * y_i$
 - matmul(x,y) Matrix Multiplication of arrays x and y which can be 1 or 2 dimensional arrays: $z_{i,j} = \sum_k x_{i,k} * y_{k,j}$
 - conjg(x) Returns the conjugate of x: $a + ib \Rightarrow a ib$
- cshift(ARRAY, SHIFT, dim) perform a circular shift by SHIFT positions to the left on array ARRAY along the dimth dimension

Allocatable Arrays I

Why?

- At compile time we may not know the size an array needs to be
- We may want to change the problem size without recompiling
- The molecular dynamics code was written for 4000 atoms. If you want to run a simulation for 256 and 1024 atoms, do you need to recompile and create two executables?
- Allocatable arrays allow us to set the size at run time.

```
real, allocatable :: force(:,:)
real, dimension(:), allocatable :: vel
```

• We set the size of the array using the allocate statement.

```
allocate(force(natoms, 3))
```

• We may want to change the lower bound for an array

```
allocate(grid(-100,100))
```

Allocatable Arrays II

We may want to use an array once somewhere in the program, say during
initialization. Using allocatable arrays also us to dynamically create the array
when needed and when not in use, free up memory using the deallocate
statement

```
deallocate (force, grid)
```

- Sometimes, we want to check whether an array is allocated or not at a particular part of the code
- Fortran provides an intrinsic function, allocated which returns a scalar logical value reporting the status of an array

```
if ( allocated(grid))deallocate(grid)
if ( .not. allocated(force))allocate(force(natoms,3))
```

Masked Array Assignment: Where Statement

• Masked array assignment is achieved using the where statement

```
where (c < 2)a = b/c
```

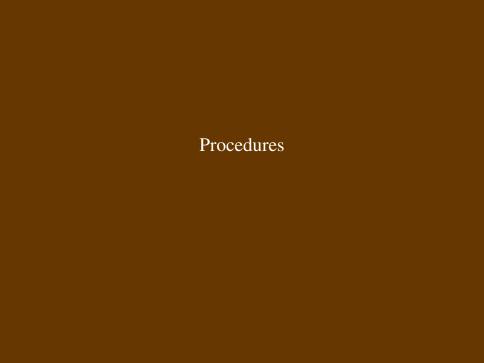
the left hand side of the assignment must be array valued.

the mask (logical expression) and the right hand side of the assignment must all conform

- Fortran 95/2003 introduced the where ... elsewhere ... end where functionality
- where statement cannot be nested

```
! Apply PBC to coordinates
where ( coord(i,:) > boxl(:) )
   coord(i,:) = coord(i,:) - boxl(:)
elsewhere ( coord(i,:) < 0d0 )
   coord(i,:) = coord(i,:) + boxl(:)
end where</pre>
```

```
! Apply PBC to coordinates
do j = 1, 3
   if ( coord(i,j) > boxl(j) ) then
      coord(i,j) = coord(i,j) - boxl(j)
   else if ( coord(i,j) < 0d0 ) then
      coord(i,j) = coord(i,j) + boxl(j)
   endif
end do</pre>
```



Program Units I

- Most programs are hundreds or more lines of code.
- Use similar code in several places.
- A single large program is extremely difficult to debug and maintain.
- Solution is to break up code blocks into procedures

Subroutines: Some out-of-line code that is called exactly where it is coded

Functions: Purpose is to return a result and is called only when the result is

needed

Modules: A module is a program unit that is not executed directly, but

contains data specifications and procedures that may be utilized

by other program units via the use statement.

Program Units II

```
program main
  use module1 ! specify which modules to use
  implicit none ! implicit typing is not recommended
  variable declarations ! declare all variables used in the program
  call routinel(argl,arg2,arg3) ! call subroutine routinel with arguments
  abc = func(arg1,arg2) ! abc is some function of arg1 and arg2
  contains ! internal procedures are listed below
    subroutine routinel (argl, arg2) ! subroutine routinel contents go here
    end subroutine routinel ! all program units must have an end statement
    function func(argl,arg2) ! function funcl contents go here
   end function func
  end program main
```

Program Units III

program md

! Molecular Dynamics code for equilibration of Liquid Argon ! This program simulates the equilibration of Liquid Argon ! potential and velocity verlet algorithm ! This program should be the starting point to learn Modern ! This program is hard coded for 4000 atoms equilibrated at ! 10K with a time step of 0.001 time units and 1000 time steps ! Lets assume that time units is femtoseconds, so total simulation ! time is 1 femtosecond ! Modify this code using the Fortran Concepts learned 1. split code into smaller subunits, modules and/or subroutines a. number of atoms or number of unit cells (you can't do both) c. time step You will need to make use allocatable arrays. If you do not know why, review 4. If you use derived types, can you overload operators? If yes, program it Parallelize the code for OpenMP/OpenACC (can also be done from step 3 or 4) ! There is no time limit for completing this exercise. This exercise is for measuring ! what have you got from the training.

Program Units IV

```
! As an additional exercise, use other potentials such as Morse potential and
 ! All Lennard-Jones Potential parameters are set to 1.
  ! This is code can be used as an introduction to molecular dynamics. There are lot more
 ! concepts in MD that are not covered here.
 ! npartdim : number of unit cells, uniform in all directions. change to nonuniform if you desire
 ! natom : number of atoms
 ! nstep : nummber of simulation time steps
 ! dt : simulation time steps
 ! kb : boltzmann constant, set to 1 for simplicity
 ! mass : mass of Ar atom, set to 1 for simplicity
 ! epsilon, sigma : LJ parameters, set to 1 for simplicity
 ! vel, vel_t0 : nuclear velocities for each step
 ! acc, acc t0 : nuclear acceleration for each step
 ! avtemp : average temperature at current time step
 ! scale : scaling factor to set current temperature to desired temperature
 implicit none
 integer, parameter :: npartdim - 10
 integer, parameter :: natom = 4.d0 * npartdim ** 3
 integer, parameter :: nstep - 1000
 real*8, parameter :: tempK = 10, dt = 1d-3
 integer :: istep
 real *8 :: box1(3), alat
 integer :: n, i, i, k, l
! Can you use derived types for coord, vel. acc and force
 real+8 :: coord t0(natom, 3), coord(natom, 3)
```

Program Units V

```
real *8 :: vel_t0(natom, 3), vel(natom, 3)
real+8 :: acc_t0(natom, 3), acc(natom, 3)
real *8 :: force (natom, 3), pener, mass
real*8 :: vcm(3), r(3), rr, r2, r6, f
real+8 :: avtemp, ke, kb, epsilon, sigma, rcell(3,4), scale
real +8 :: gasdev
alat = 2d0 + (2d0/3d0)
do i - 1, 3
  boxl(i) - npartdim * alat
end do
kb = 1.d0
mass = 1.d0
epsilon - 1.d0
sigma - 1.d0
! Hint: Simplify unit cell creation, maybe in variable declaration
rcell(1,1) = 0d0
rcell(2,1) = 0d0
rcell(3,1) = 0d0
rcell(1,2) = 0.5d0 * alat
rcell(2,2) = 0.5d0 * alat
rcell(3,2) = 0d0
rcell(1,3) = 0d0
rcell(2.3) = 0.5d0 * alat
rcell(3,3) = 0.5d0 * alat
rcell(1,4) = 0.5d0 * alat
rcell(2.4) - 0d0
rcell(3,4) = 0.5d0 * alat
do i - 1, natom
  do 1 - 1, 3
      vel t0(i,i) - 0d0
      acc_t0(i,j) = 0d0
```

Program Units VI

```
end do
end do
n - 1
do i - 1, npartdim
  do | - 1, npartdim
      do k - 1, npartdim
         do 1 - 1, 4
            coord_t0(n,1) = alat * dble(i - 1) + rcell(1,1)
            coord t0(n,2) = alat \star dble(1 - 1) + rcel1(2,1)
            coord_t0(n,3) = alat * dble(k - 1) + rcell(3,1)
         and do
      end do
   end do
end do
open(unit-1,file-'atom.xvz',status-'unknown')
write(1,'(i8)') natom
write(1, *)
do i - 1, natom
   write(1,'(a2,2x,3f12.6)') 'Ar', coord_t0(i,1), coord_t0(i,2), coord_t0(i,3)
end do
close(1)
do i - 1, natom
  do 1 - 1, 3
      vel_t0(i,j) = gasdev()
   end do
end do
```

! Set Linear Momentum to zero

Program Units VII

```
! First get center of mass velocity
vcm = 0d0
do i - 1, natom
  do 1 - 1, 3
      vcm(j) = vcm(j) + vel_t0(i,j)/natom
   end do
end do
do i - 1, natom
  do 1 - 1, 3
      vel_t0(i,j) = vel_t0(i,j) - vcm(j)
and do
ke = 0d0
do i - 1, natom
  do 1 - 1, 3
      ke = ke + mass * vel_t0(i, j) **2
   end do
end do
avtemp = mass \star ke / ( 3d0 \star kb \star ( natom - 1))
print '(a, 2x, 1pe15.8)', 'Initial Average Temperature: ', avtemp
! scale initial velocity to desired temperature
scale = sqrt( tempK / avtemp )
ke = 0d0
do i - 1, natom
  do 1 - 1, 3
      vel_t0(i,j) = vel_t0(i,j) * scale
      ke = ke + mass * vel_t0(i, j) **2
   end do
end do
```

Program Units VIII

```
avtemp = mass \star ke / (3d0 \star kb \star (natom - 1))
print '(a,2x,1pe15.8)', 'Initial Scaled Average Temperature: ', avtemp
! MD Simulation
do istep - 1, nstep
   ! Set coordinates, velocity, acceleration and force at next time step to zero
   do i - 1, natom
      do 1 - 1, 3
        coord(i,i) = 0d0
        vel(i, j) = 0d0
        acc(i,i) = 0d0
   end do
   pener - 0d0
   ! Get new atom positions from Velocity Verlet Algorithm
   ! Hint: OpenMP/OpenACC
   do i - 1, natom
      do 1 - 1, 3
        coord(i,j) = coord_t0(i,j) + vel_t0(i,j) * dt + 0.5d0 * acc_t0(i,j) * dt ** 2
         ! Apply PBC to coordinates
         if (coord(i,j) > boxl(j)) then
            coord(i,j) = coord(i,j) = boxl(j)
        else if (coord(i,j) < 0d0) then
            coord(i, j) = coord(i, j) + boxl(j)
         endif
      end do
   end do
   ! Get force at new atom positions
```

Program Units IX

```
do i - 1, natom - 1
   do j - i + 1, natom
      do k = 1, 3
         r(k) = coord(i,k) = coord(i,k)
         r(k) = r(k) - nint(r(k) / boxl(k)) + boxl(k)
      end do
      rr = r(1) \leftrightarrow 2 + r(2) \leftrightarrow 2 + r(3) \leftrightarrow 2
      r2 = 1.d0 / rr
      r6 - r2 + 3
      pener - pener + 4d0 * r6 * ( r6 - 1.d0 )
      f = 48d0 * r2 * r6 * (r6 - 0.5d0)
      do k = 1, 3
         force(i,k) = force(i,k) + r(k) + f
         force(i,k) = force(i,k) - r(k) + f
      end do
   end do
end do
! Calculate Acceleration and Velocity at current time step
do i - 1, natom
   do 1 - 1, 3
      acc(i,j) = force(i,j) / mass
      vel(i,j) = vel_t0(i,j) + 0.5d0 * (acc(i,j) + acc_t0(i,j)) * dt
   end do
end do
! Set Linear Momentum to zero
! First get center of mass velocity
```

Program Units X

```
vcm = 0d0
do i - 1, natom
   do 1 - 1, 3
      vcm(j) = vcm(j) + vel(i,j)/natom
   end do
end do
do i - 1, natom
  do 1 - 1, 3
      vel(i,j) = vel(i,j) - vcm(j)
   end do
end do
ke = 0d0
do i - 1, natom
   do 1 - 1, 3
      ke = ke + vel(i,i) ** 2
   end do
end do
avtemp = mass \star ke / (3d0 \star kb \star ( natom - 1))
print '(a,2x,18,2x,1pe15.8,1x,1pe15.8)', 'Average Temperature: ', istep, avtemp, pener
scale = sqrt ( tempk/ avtemp )
do i - 1, natom
  do 1 - 1, 3
     acc_t0(i,j) = acc(i,j)
      vel_t0(i,j) = vel(i,j) * scale
   end do
end do
! Write current coordinates to xvz file for visualization
open (unit-1, file-'atom.xyz', position-'append')
```

Program Units XI

```
write(1,'(i8)') natom
    write (1, +)
    do i - 1, natom
        write(1,'(a2,2x,3f12.6)') 'Ar', coord_t0(i,1), coord_t0(i,2), coord_t0(i,3)
     end do
    close(1)
 end do
end program md
double precision function gasdev()
 implicit none
 real +8 :: v1, v2, fac, rsq
 real+8, save :: gset
 logical, save :: available - .false.
 if (available) then
    gasdev - gset
     available - .false.
 6100
    do
        call random_number(v1)
        call random number (v2)
        v1 = 2.d0 \times v1 = 1.d0
        v2 = 2.d0 + v2 = 1.d0
        rsq = v1**2 + v2**2
        if ( rsq > 0.d0 .and. rsq < 1.d0 ) exit
     fac - sqrt(-2.d0 * log(rsq) / rsq)
    gasdev - v1 * fac
    gset - v2 + fac
    available - .true.
 end if
```

end function gasdev

Subroutines I

Call Statement:

- The call statement evaluates its arguments and transfers control to the subroutine
- Upon return, the next statement is executed.

SUBROUTINE Statement:

- The **subroutine** statement declares the procedure and its arguments.
- These are also known as dummy arguments.

• The subroutine's interface is defined by

- The subroutine statement itself
- The declarations of its dummy arguments
- Anything else that the subroutine uses

Subroutines II

Statement Order

- A subroutine statement starts a subroutine
- Any use statements come next
- implicit none comes next, followed by
- rest of the declarations,
- executable statements
- 6 End with a end subroutine statement

Dummy Arguments

- Their names exist only in the procedure and are declared as local variables.
- The dummy arguments are associated with the actual arguments passed to the subroutines.
- The dummy and actual argument lists must match, i.e. the number of arguments must be the same and each argument must match in type and rank.

Subroutines III

```
subroutine verlet(coord, coord_t0, vel, vel_t0, acc, acc_t0, force, pener)
                                                                                     program md
 use precision
 use potential
 use param, only : natom, mass, dt, boxl, pot
  implicit none
                                                                                        real(dp), dimension(:,:), allocatable :: coord_t0, coord
  real(dp), dimension(:,:), intent(in) :: coord_t0, vel_t0, acc_t0
                                                                                        real(dp), dimension(:,:), allocatable :: vel_t0, vel
  real(dp), dimension(:,:), intent(out) :: coord, vel, acc, force
                                                                                        real(dp), dimension(:,:), allocatable :: acc_t0, acc, force
  real(dp), intent(out) :: pener
  integer(ip) :: i, j, k
                                                                                        real(dp) :: pener
  real(dp) :: epot
  real(dp) :: r(3), f(3)
                                                                                        interface
 ! Set coordinates, velocity, acceleration and force at next time step to zero
                                                                                           subroutine verlet(coord, coord_t0, vel_t0, vel, acc_t0, acc, force, pener)
 coord = 0d0 : vel = 0d0 : acc = 0d0 : force = 0d0
                                                                                             use precision
 pener = 0d0
                                                                                             implicit none
 ! Get new atom positions from Velocity Verlet Algorithm
                                                                                             real(dp), dimension(:,:), intent(in) :: coord_t0, vel_t0, acc_t0
 coord = coord_t0 + vel_t0 * dt + 0.5d0 * acc_t0 * dt ** 2
                                                                                             real(dp), dimension(:,:), intent(out) :: coord, vel, acc, force
 do i = 1, natom
                                                                                             real(dp), intent(out) :: pener
                                                                                           end subroutine verlet
    where ( coord(i,:) > boxl(:) )
                                                                                        end interface
    elsewhere ( coord(i,:) < 0d0 )
    end where
  end do
                                                                                        do istep = 1, nstep
 do i = 1, natom - 1
    do i = i + 1, natom
                                                                                           coord = 0d0 : vel = 0d0 : acc = 0d0
                                                                                           force = 0d0 : pener = 0d0
       r = r - nint(r / boxl) * boxl
       select case(pot)
       case ('mp')
                                                                                           call verlet (coord, coord t0, vel t0, vel, acc t0, acc, force, pener)
          call morse ( r. f. epot )
       case default
          call lennard_jones( r, f, epot )
       end select
                                                                                        and do
       pener = pener + epot
       force(i,:) = force(i,:) - f(:)
                                                                                        deallocate(coord_t0,vel_t0,acc_t0,coord,vel,acc,force)
    end do
 end do
                                                                                      end program md
 acc = force / mass
```

vel = vel t0 + 0.5d0 * (acc + acc t0) * dt

end subroutine verlet

Internal Procedures

- Internal procedures appear just before the last end statement and are preceded by the contains statement.
- Internal procedures can be either subroutines or functions which can be accessed only by the program, subroutine or module in which it is present
- Internal procedures have declaration of variables passed on from the parent program unit
- If an internal procedure declares a variable which has the same name as a variable from the parent program unit then this supersedes the variable from the outer scope for the length of the procedure.

Functions

- functions operate on the same principle as subroutines
- The only difference is that function returns a value and does not involve the call statement

```
module potential
  use precision
  implicit none
  real(dp) :: r2, r6, d2, d
  real(dp), parameter :: de = 0.176d0, a = 1.4d0, re = 1d0
  real(dp) :: exparre
  subroutine lennard jones (r.f.p)
   ! V = 4 * epsilon * [ (sigma/r)**12 - (sigma/r)**6 ]
    implicit none
    real(dp), dimension(:), intent(in) :: r
    real(dp), dimension(:), intent(out) :: f
    real(dp), intent(out) :: p
    r2 = 1.d0 / dot product(r,r)
    f = dvdr_1j(r2, r6) * r
   p = pot_1j(r2, r6)
  end subroutine lennard_jones
  subroutine morse(r,f,p)
    implicit none
    real(dp), dimension(:), intent(in) :: r
    real(dp), dimension(:), intent(out) :: f
```

```
d2 = dot_product(r,r)
    d = sart(d2)
   exparre = exp(-a * (d - re))
   f = dvdr mp(exparre) * r
   p = pot mp(exparre)
 end subroutine morse
  function pot li(r2, r6)
    implicit none
    real(dp), intent(in) :: r2, r6
    real(dp) :: pot li
   pot 1i = 4d0 * r6 * ( r6 - 1.d0 )
 end function pot_lj
  function pot mp(exparre)
    implicit none
    real(dp), intent(in) :: exparre
    real(dp) :: pot mp
   pot mp = de \star ( 1d0 - exparre ) \star \star 2
 end function pot_mp
 function dvdr li(r2,r6)
   implicit none
    real(dp), intent(in) :: r2, r6
   real(dp) :: dvdr_lj
   dvdr_1j = 48d0 * r2 * r6 * (r6 - 0.5d0)
 end function dvdr_lj
 function dvdr_mp(exparre)
    implicit none
   real(dp), intent(in) :: exparre
    real(dp) :: dvdr_mp
   dvdr_mp = 2d0 * de * a * (1d0 - exparre) * exparre
 end function dvdr_mp
end module potential
```

real(dp), intent(out) :: p

Array-valued Functions

function can also return arrays

```
module potential
 use precision
 implicit none
 real(dp) :: r2, r6, d2, d
 real(dp), parameter :: de = 0.176d0, a = 1.4d0, re = 1d0
 real(dp) :: exparre
containe
 subroutine lennard_jones(r,f,p)
         - 48 * r**(-8) * [ r**(-6) - 0.5 ] * i for epsilon-sigma-1
    implicit none
    real(dp), dimension(:), intent(in) :: r
    real(dp), dimension(:), intent(out) :: f
    real(dp), intent(out) :: p
    r2 = 1.d0 / dot_product(r,r)
    r6 = r2 \leftrightarrow 3
    f - dvdr_lj(r2, r6, r)
    p = pot 11(r2, r6)
 end subroutine lennard_jones
 subroutine morse(r,f,p)
    implicit none
    real(dp), dimension(:), intent(in) :: r
    real(dp), dimension(:), intent(out) :: f
    real(dp), intent(out) :: p
```

```
d2 - dot product(r,r)
    d = sqrt(d2)
    exparre - exp(-a + (d - re))
    f - dvdr_mp(exparre,r)
    p - pot mp(exparre)
  end subroutine morse
  function pot 11(r2, r6)
    implicit none
    real(dp), intent(in) :: r2, r6
    real(dp) :: pot_lj
    pot_1 = 4d0 \times r6 \times (r6 - 1.d0)
  end function pot_li
  function pot mp(exparre)
    implicit none
    real(dp), intent(in) :: exparre
    real(dp) :: pot mp
    pot_mp = de * (1d0 - exparre) **2
  end function pot mp
  function dvdr_li(r2,r6,r)
    implicit none
    real(dp), intent(in) :: r2, r6, r
    real(dp), dimension(size(r)) :: dvdr_li
    dvdr 11 = 48d0 + r2 + r6 + (r6 - 0.5d0) + r
  end function dvdr_lj
  function dvdr_mp(exparre,r)
    implicit none
    real(dp), intent(in) :: exparre, r
    real(dp), dimension(size(r)) :: dvdr_mp
    dvdr mp = 2d0 * de * a * (1d0 - exparre) * exparre * r
  end function dvdr_mp
end module potential
```

Recursive Procedures

- In Fortran 90, recursion is supported as a feature
 - recursive procedures call themselves
 - Precursive procedures must be declared explicitly
 - Tecursive function declarations must contain a result keyword, and
 - one type of declaration refers to both the function name and the result variable.

```
program fact
 implicit none
 integer :: i
 print *, 'enter integer whose factorial you want to calculate
 read *, i
 print '(i5,a,i20)', i, '! = ', factorial(i)
contains
 recursive function factorial(i) result(i fact)
   integer, intent(in) :: i
   integer :: i_fact
   if (i > 0) then
      i_fact = i * factorial(i - 1)
   else
      i_fact = 1
   end if
 end function factorial
end program fact
```

```
[apacheco@qb4 Exercise] ./factorial
enter integer whose factorial you want to calculate
10
10! = 3628800
[apacheco@qb4 Exercise] ./fact1
Enter an integer < 15
10
10! = 3628800
```

- The interface statement is the first statement in an interface block.
- The interface block is a powerful structure that was introduced in FORTRAN 90.
- When used, it gives a calling procedure the full knowledge of the types and characteristics of the dummy arguments that are used inside of the procedure that it references.
- This can be a very good thing as it provides a way to execute some safety checks when compiling the program.
- Because the main program knows what argument types should be sent to the referenced procedure, it can check to see whether or not this is the case.
- If not, the compiler will return an error message when you attempt to compile the program.

```
subroutine verlet (coord, coord_t0, vel, vel_t0, acc, acc_t0, force,
        pener)
 use precision
 use param, only : natom, mass, dt, boxl, pot
 implicit none
 real(dp), dimension(:,:), intent(in) :: coord t0, vel t0, acc t0
 real(dp), dimension(:,:), intent(out) :: coord, vel, acc, force
 real(dp), intent(out) :: pener
 integer(ip) :: i
 interface
     subroutine get pot force(coord, force, pener)
      use precision
      implicit none
       real(dp), dimension(:,:), intent(in) :: coord
      real(dp), dimension(:,:), intent(out) :: force
       real(dp), intent(out) :: pener
     end subroutine get_pot_force
 end interface
 ! Set coordinates, velocity, acceleration and force at next time
 coord - 0d0 : vel - 0d0 : acc - 0d0
 ! Get new atom positions from Velocity Verlet Algorithm
 coord - coord t0 + vel t0 * dt + 0.5d0 * acc t0 * dt ** 2
 do i - 1, natom
    ! Apply PBC to coordinates
     where ( coord(i,:) > boxl(:) )
       coord(i,:) = coord(i,:) - boxl(:)
     elsewhere ( coord(i,:) < 0d0 )
       coord(i,:) = coord(i,:) + boxl(:)
    end where
 and do
 ! Get Potential and force at new atom positions
 call get pot force(coord, force, pener)
 ! Calculate Acceleration and Velocity at current time step
```

```
acc - force / mass
  vel - vel_t0 + 0.5d0 * ( acc + acc_t0 ) * dt
end subroutine verlet
subroutine get pot force(coord, force, pener)
  use precision
  use potential
  use param, only : natom, boxl
  implicit none
  real(dp), dimension(:,:), intent(in) :: coord
  real(dp), dimension(:,:), intent(out) :: force
  real(dp), intent(out) :: pener
  integer(ip) :: i, i
  real(dp) :: epot
 real(dp) :: r(3), f(3)
  force - 0d0
  do i - 1, natom - 1
    do 1 - i + 1, natom
       r = r - nint(r / boxl) * boxl
       select case(pot)
        case ('mp')
           call morse ( r, f, epot )
        case default
           call lennard iones ( r. f. epot )
        end select
       pener - pener + epot
        force(i,:) = force(i,:) + f(:)
        force(1,:) = force(1,:) - f(:)
     and do
  end do
end subroutine get pot force
```

```
subroutine verlet (coord, coord_t0, vel, vel_t0, acc, acc_t0, force,
        pener)
 use precision
 use param, only : natom, mass, dt, boxl, pot
 implicit none
 real(dp), dimension(:,:), intent(in) :: coord t0, vel t0, acc t0
 real(dp), dimension(:,:), intent(out) :: coord, vel, acc, force
 real(dp), intent(out) :: pener
 integer(ip) :: i
 ! Set coordinates, velocity, acceleration and force at next time
 coord - 0d0 ; vel - 0d0 ; acc - 0d0
 ! Get new atom positions from Velocity Verlet Algorithm
 coord = coord_t0 + vel_t0 * dt + 0.5d0 * acc_t0 * dt ** 2
 do i - 1, natom
     where ( coord(i,:) > boxl(:) )
       coord(i,:) = coord(i,:) - boxl(:)
     elsewhere ( coord(i,:) < 0d0 )
       coord(i,:) = coord(i,:) + boxl(:)
     end where
 end do
 ! Get Potential and force at new atom positions
 call get_pot_force(coord, force, pener)
 ! Calculate Acceleration and Velocity at current time step
 acc = force / mass
 vel = vel t0 + 0.5d0 * (acc + acc t0) * dt
contains
```

```
subroutine get_pot_force(coord, force, pener)
    use potential
    implicit none
    real(dp), dimension(:,:), intent(in) :: coord
    real(dp), dimension(:,:), intent(out) :: force
    real(dp), intent(out) :: pener
    integer(ip) :: i, j
    real(dp) :: epot
    real(dp) :: r(3), f(3)
    pener - 0d0
    force - 0d0
    do i = 1, natom = 1
       do i - i + 1, natom
          r(:) = coord(i,:) - coord(i,:)
          r - r - nint(r / boxl) + boxl
          select case (pot)
          case ('mp')
             call morse ( r. f. epot )
          case default
             call lennard_jones ( r, f, epot )
          end select
          pener - pener + epot
          force(i,:) = force(i,:) + f(:)
          force(i,:) = force(i,:) - f(:)
       end do
    and do
  end subroutine get_pot_force
end subroutine verlet
```

Here since subroutine get_pot_force is an internal procedure, no interface is required since it
is already implicit and all variable declarations are carried over from subroutine verlet

Argument Association

• Recall from MD code example the invocation

```
call linearmom(vel_t0)
```

• and the subroutine declaration

```
subroutine linearmom(vel)
```

- vel_t0 is an actual argument and is associated with the dummy argument vel
- In subroutine linearmom, the name vel is an alias for vel_t0
- If the value of a dummy argument changes, then so does the value of the actual argument
- The actual and dummy arguments must correspond in type, kind and rank.

Local Objects

- In subroutine linearmom,
 - i and vom are local objects.
- Local Objects
 - are created each time a procedure is invoked
 - are destroyed when the procedure completes
 - do not retain their values between calls
 - do not exist in the programs memory between calls

Example

```
subroutine linearmom(vel)
 use precision
 use param, only : natom
 implicit none
 real(dp), dimension(:,:), intent(inout) :: vel
 integer(ip) :: i
 real(dp) :: vcm(3)
 ! First get center of mass velocity
 vcm = 0d0
 do i - 1, 3
    vcm(i) - sum(vel(:,i))
 end do
 vcm - vcm / real(natom, dp)
 ! Now remove center of mass velocity from all atoms
 do i - 1, natom
    vel(i,:) = vel(i,:) - vcm(:)
 end do
end subroutine linearmom
```

Optional & Keyword Arguments I

- Optional Arguments
 - allow defaults to be used for missing arguments
 - make some procedures easier to use
- once an argument has been omitted all subsequent arguments must be keyword arguments
- the present intrinsic can be used to check for missing arguments
- if used with external procedures then the interface must be explicit within the procedure in which it is invoked.

```
subroutine get_temp(vel,boltz)
use pracision
use param, only: natom, avtemp, mass, kb
implicit none
real(dp), dimension(:,:), intent(in):: vel
real(dp), optional:: boltz
integer(lp):: i
real(dp):: ke

if (present(boltz)) kb - boltz
ke - 0:00
do i - 1, natom
ke - ke + dot_product(vel(i,:),vel(i,:))
end do
avtemp - mass * ke / ( 3:00 * kb * real( natom - 1, dp))
end subroutine get temp
```

```
subroutine initialize(coord_t0, vel_t0, acc_t0)
...
Interface
subsputtine linearmom(vel)
subsputtine linearmom(vel)
subsputtine linearmom
subroutine linearmom
subroutine get_temp(vel, boltz)
use precision
implicit none
real(dp), dimension(:,:), intent(inout) :: vel
end subroutine get_temp(vel, boltz)
use precision
implicit none
real(dp), dimension(:,:), intent(in) :: vel
real(dp), optional :: boltz
end subroutine get_temp
end interface
...
call get_temp(vel_t0)
```

Optional & Keyword Arguments II

- Keyword Arguments
 - allow arguments to be specified in any order
 - makes it easy to add an extra argument no need to modify any calls
 - helps improve readability of the program
 - are used when a procedure has optional arguments
- once a keyword is used, all subsequent arguments must be keyword arguments
- if used with external procedures then the interface must be explicit within the procedure in which it is invoked.

```
subroutine initialize(coord, vel, acc)

...

real(dp),dimension(:,:), intent(out) :: coord, vel
, acc
end subroutine initialize

program md
...

call initialize(coord_t0, vel_t0, acc_t0)
end program md
...
end program md
```

Optional & Keyword Arguments III

• subroutine initialize can be invoked using

using the positional argument invocationusing keyword arguments

program md
...
 interface
 subroutine initialize(coord, vel, acc)
 use precision
 implicit none
 real(dp), dimension(:,:), intent(out) :: coord, vel, acc
 end subroutine initialize
 end interface
...
! All three calls give the same result.
 call initialize(coord_t0, vel_t0, acc_t0)
 call initialize(coord_coord_t0, acc=acc_t0, vel=vel_t0)
 call initialize(coord_t0, acc=acc_t0, vel=vel_t0)

Dummy Array Arguments

- There are two main types of dummy array argument:
 - explicit-shape: all bounds specified

```
real, dimension(4,4), intent(in):: explicit_shape
```

The actual argument that becomes associated with an explicit shape dummy must conform in size and shape

- 2 assumed-shape: no bounds specified, all inherited from the actual argument
 - real, dimension(:,:), intent(out):: assumed_shape
 - An explicit interface must be provided
- 3 assumed-size: final dimension is specified by *
 - real :: assumed_size(dim1, dim2, *)

Commomly used in FORTRAN, use assumed-shape arrays in Modern Fortran.

• dummy arguments cannot be (unallocated) allocatable arrays.

Explicit-shape Arrays

```
program md
  use precision
  use param
  implicit none
  integer(ip) :: n, i, j, k, 1
  real(dp), dimension(:,:), allocatable :: coord_t0, vel_t0, acc_t0
  real(dp), dimension(:,:), allocatable :: coord, vel, acc, force
 allocate(coord(natom, 3), coord_t0(natom, 3))
  allocate (vel (natom, 3), vel t0 (natom, 3))
  allocate (acc (natom, 3), acc t0 (natom, 3))
  allocate (force (natom, 3))
  ! Initialize coordinates and random velocities
 call initialize (coord t0, vel t0, acc t0)
end program md
subroutine initialize (coord t0, vel t0, acc t0)
  use precision
  use param, only : natom, npartdim, alat, rcell
  implicit none
  real(dp), dimension(natom, 3) :: coord_t0, vel_t0, acc_t0
  integer(ip) :: n, i, j, k, 1
  coord_t0 = 0d0 ; vel_t0 = 0d0 ; acc_t0 = 0d0
end subroutine initialize
```

Assumed-Shape Arrays

```
program md
 use precision
 use param
 implicit none
 integer(ip) :: n, i, j, k, l
 real(dp), dimension(:,:), allocatable :: coord t0, vel t0, acc t0
 real(dp), dimension(:,:), allocatable :: coord, vel, acc, force
 interface
     subroutine initialize(coord_t0, vel_t0, acc_t0)
      use precision
       implicit none
       real(dp), dimension(:,:), intent(out) :: coord_t0, vel_t0, acc_t0
     end subroutine initialize
 end interface
 allocate (coord (natom, 3), coord_t0 (natom, 3))
 allocate(vel(natom, 3), vel t0(natom, 3))
 allocate (acc(natom, 3), acc_t0(natom, 3))
 allocate (force (natom, 3))
 call initialize(coord_t0, vel_t0, acc_t0)
end program md
subroutine initialize(coord_t0, vel_t0, acc_t0)
 use precision
 use param, only : natom, npartdim, alat, rcell
 implicit none
 real(dp), dimension(:,:), intent(out) :: coord_t0, vel_t0, acc_t0
 integer(ip) :: n, i, i, k, 1
 ! Set initial coordinates, velocity and acceleration to zero
 coord t0 - 0d0 ; vel t0 - 0d0 ; acc t0 - 0d0
end subroutine initialize
```

Automatic Arrays

- Automatic Arrays: Arrays which depend on dummy arguments
 - their size is determined by dummy argumentsthey cannot have the save attribute or be initialized.
- The size intrinsic or dummy arguments can be used to declare automatic arrays.

```
program main
  implicit none
 integer :: i, j
  real, dimension(5,6) :: a
 call routine(a,i,j)
  contains
    subroutine routine(c,m,n)
      integer :: m, n
      real, dimension(:,:), intent(inout) :: c ! assumed shape array
      real :: b1(m,n)
                                                  ! automatic array
      real, dimension(size(c,1),size(c,2)) :: b2 ! automatic array
    end subroutine routine
end program main
```

Save Attribute and Arrays

- Declaring a variable (or array) as save gives it a static storage memory.
- i.e information about variables is retained in memory between procedure calls.

```
subroutine something(iarg1)
implicit none
integer, intent(in) :: iarg1
real,dimension(:,:),allocatable,save :: a
real, dimension(:,:),allocatable :: b
...
if (.not.allocated(a))allocate(a(i,j))
allocate(b(j,i))
...
deallocate(b)
end subroutine something
```

- Array a is saved when something exits.
- Array b is not saved and needs to be allocated every time in something and deallocated, to free up memory, before something exits.

Intent

- intent attribute was introduced in Fortran 90 and is recommended as it
 - allows compilers to check for coding errors
 - a facilitates efficient compilation and optimization
- Declare if a parameter is

• A variable declared as intent(in) in a procedure cannot be changed during the execution of the procedure (see point 1 above)

Modules I

- Modules were introduced in Fortran 90 and have a wide range of applications.
- Modules allow the user to write object based code.
- A module is a program unit whose functionality can be exploited by other programs which attaches to it via the use statement.
- A module can contain the following
 - global object declaration: replaces Fortran 77 **COMMON** and **INCLUDE** statements
 - interface declaration: all external procedures using assumed shape arrrays, intent and keyword/optional arguments must have an explicit interface
 - 9 procedure declaration: include procedures such as subroutines or functions in modules. Since modules already contain explicit interface, an interface statement is not required

Modules II

```
module precision
implicit none
save
integer, parameter :: ip - selected_int_kind(15)
integer, parameter :: dp - selected_real_kind(15)
end module precision
module param
use precision
implicit none
integer(ib) :: neartdim. natom, nate, istep
```

```
real(dp):: tempK, dt, boxl(3), alat, mass
real(dp):: avtemp, ke, kk, epsilon, sigma, scale
real(dp).dimension(j,4):: rcell - reshape( // &
0.50+00, 0.50+00, 0.00+00, &
0.50+00, 0.50+00, 0.00+00, &
0.50+00, 0.50+00, 0.50+00, &
0.50+00, 0.00+00, 0.50+00 /), (/ 3, 4 /))
character(len-2):: pot
end module param
```

- within a module, functions and subroutines are called module procedures.
- module procedures can contain internal procedures
- module objects that retain their values should be given a save attribute
- modules can be used by procedures and other modules, see module precision.
- modules can be compiled separately. They should be compiled before the program unit that uses them.

Observe that in my examples with all code in single file, the modules appear before the main program and subroutines.

Modules III

Visibility of module procedures

- By default, all module procedures are public i.e. they can accessed by program units that use the module using the use statement
- To restrict the visibility of the module procedure only to the module, use the private statement
- In the module potential, all functions which calculate forces can be declared as private as follows

```
module potential
use precision
implicit none
real(dp) :: r2, r6, d2, d
real(dp), parameter :: de = 0.176d0, a = 1.4d0, re = 1d0
real(dp) :: exparre
public :: lennard_jones, morse, pot_lj, pot_mp
private :: dvdr_lj, dvdr_mp

contains
...
end module potential
```

Program Units in the MD code can directly call
 lennard_jones, morse, pot_lj and pot_mp but cannot access dvdr_lj and
 dvdr_mp

Modules IV

Using Modules

 The use statement names a module whole public definitions are to be made accessible.

To use all variables from module param in program md:

```
program md
  use param
  ...
end program md
```

module entities can be renamed

To rename pot and dt to more user readable variables:

```
use param, pot => potential, dt => timestep
```

- It's good programming practice to use only those variables from modules that are neccessary to avoid name conflicts and overwrite variables.
- For this, use the use <modulename>, only statement

```
subroutine verlet(coord,force,pener)
  use param,only : dp,npart,boxl,tstep
...
end subroutine verlet
```

Compiling Modules I

- Consider the MD code containing a main program md.f90, modules precision.f90, param.f90 and potential.f90 and subroutines initialize.f90, verlet.f90, linearmom.f90 and get_temp.f90.
- In general, the code can be compiled as

```
ifort -o md md.f90 precision.f90 param.f90 potential.f90 initialize.f90 \
    verlet.f90 linearmom.f90 get_temp.f90
```

- Most compilers are restrictive in the order of compilation.
- The order in which the sub programs should be compiled is
 - Modules that do not use any other modules.
 - Modules that use one or more of the modules already compiled.
 - Repeat the above step until all modules are compiled and all dependencies are resolved.
 - Main program followed by all subroutines and functions (if any).
- In the MD code, the module precision does not depend on any other modules and should be compiled first
- The modules param and potential only depend on precision and can be compiled in any order

Compiling Modules II

• The main program and subroutines can then be compiled

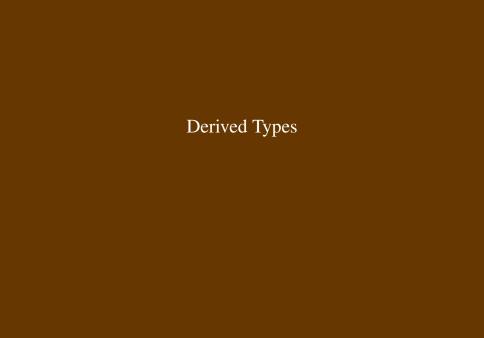
```
ifort -o md md.f90 precision.f90 param.f90 potential.f90 initialize.f90 \
verlet.f90 linearmom.f90 get_temp.f90
```

• modules are designed to be compiled independently of the main program and create a .mod files which need to be linked to the main executable.

• The main program can now be compiled as

```
ifort -o md md.f90 initialize.f90 verlet.f90 linearmom.f90 get_temp.f90 \
   -I{path to directory containing the .mod files}
```

• The Makefile tutorial will cover this aspect in more detail.



Derived Types I

- Defined by user (also called structures)
- Can include different intrinsic types and other derived types
- Components are accessed using the percent operator (%)
- Only assignment operator (=) is defined for derived types
- Can (re)define operators see operator overloading
- Derived type definitions should be placed in a module.
- Previously defined type can be used as components of other derived types.

```
type line_type
real:: x1, y1, x2, y2
end type line_type

type(line_type) :: a, b

type vector_type
    type(line_type) :: line ! defines x1,y1,x2,y2
    integer :: direction ! O=nodirection, l=(x1,y1)->(x2,y2)
end type vector_type

type(vector_type) :: c, d
```

Derived Types II

- values can be assigned to derived types in two ways
 - 1 component by component individual component may be selected using the % operator
 - as an object the whole object may be selected and assigned to using a constructor

```
a%x1 = 0.0; a%x2 = 0.5; a%y1 = 0.0; a%y2 = 0.5
c%direction = 0
c%line%x1 = 0.0; c%line%x2 = 1.0
c%line%y1 = -1.0; c%line%y2 = 0.0
b = line_type(0.0, 0.0, 0.5, 0.5)
d%line = line_type(0.0, -1.0, 1.0, 0.0)}
d = vector_type( d%line, 1 )
l or
d = vector_type( line_type(0.0, -1.0, 1.0, 0.0), 1)
```

Derived Types III

• Assignment between two objects of the same derived type is intrinsically defined In the previous example: a = b is allowed but a = c is not.

```
coord_t0(n) %x = alat * real(i - 1, dp) + rcell(1,1)
coord_t0(n) %y = alat * real(j - 1, dp) + rcell(2,1)
coord_t0(n) %z = alat * real(k - 1, dp) + rcell(3,1)
OR
    x = alat * real(i - 1, dp) + rcell(1,1)
y = alat * real(j - 1, dp) + rcell(2,1)
z = alat * real(k - 1, dp) + rcell(3,1)
coord_t0(n) = dynamics(x, y, z)
```

- I/O on Derived Types
 - Can do normal I/O on derived types
 print *, a will produce the result 1.00.51.5
 print *, c will produce the result 2.00.00.00.0
- Arrays and Derived Types
 - Can define derived type objects which contain non-allocatable arrays and arrays of derived type objects
- Derived Type Valued Functions

Derived Types IV

• Functions can return results of an arbitrary defined type.

Private Derived Types

• A derived type can be wholly private or some of its components hidden

```
module data

type :: position
    real, private :: x, y, z
    end type position
    type, private :: acceleration
    real, private :: x, y, z
    end type acceleration
    contains

...
end module data
```

 Program units that use data have position exported but not it's components x, y, z and the derived type acceleration

Generic Procedures I

- In Fortran, most intrinsic functions are generic in that their type is determined by their argument(s)
- For example, the abs (x) intrinsic function comprises of
 - ① cabs: called when x is complex
 - 2 abs: called when x is real
 - 3 iabs: called when x is integer
- These sets of functions are called *overload sets*
- Fortran users may define their own overload sets in an interface block

```
interface clear
  module procedure clear_real, clear_type, clear_typelD
end interface
```

• The generic name clear is associated with specific names

```
clear_real, clear_type, clear_type1D
```

Generic Procedures II

```
module dynamic data
  type dynamics
     real(dp) :: x, y, z
  end type dynamics
  interface dot product
    module procedure dprod
  end interface dot product
  interface clear
     module procedure clear real, clear type,
            clear type1D
  end interface
contains
  function dprod(a,b) result(c)
    type(dynamics).intent(in) :: a.b
    real(dp) :: c
    c = a x * b x + a v * b v + a z * b z
  end function dprod
  subroutine clear real(a)
    real(dp), dimension(:,:), intent(out) :: a
    a = 0.00
  end subroutine clear real
  subroutine clear_type(a)
    type(dynamics), dimension(:), intent(out) ::
    a%x = 0d0; a%y = 0d0; a%z = 0d0
  end subroutine clear type
  subroutine clear typelD(a)
    type(dynamics),intent(out) :: a
    a%x = 0d0; a%y = 0d0; a%z = 0d0
  end subroutine clear type1D
end module dynamic data
```

```
program md
  use dynamic data
  type (dynamics), dimension (:), allocatable :: coord, coord
         0.vel.force
  allocate (coord (npart), coord0 (npart), vel (npart), force (
         npart))
     do i=1.npart
        v2t = v2t + dot product(vel(i), vel(i))
     enddo
end program md
subroutine setup(coord, vel, coord0)
  type (dynamics) :: vt
  call clear (coord)
  call clear (coord0)
  call clear (vel)
  call clear(vt)
end subroutine setup
```

Generic Procedures III

- The dot_product intrinsic function is overloaded to inleude derived types
- The procedure clear is overloaded to set all components of derived types and all elements of 2D real arrays to zero.

Operator Overloading I

- Intrinsic operators such as +, -, * and / can be overloaded to apply to all types of data
- Recall, for derived types only the assignment (=) operator is defined
- In the MD code, coord_t(i) = coord_t0(i) is well defined, but vel_t(i) = vel_t(i) * scalef is not
- Operator overloading as follows
 - specify the generic operator symbol in an interface operator statement
 - specify the overload set in a generic interface
 - Second declare the module procedures (functions) which define how the operations are implemented.
 - 4 these functions must have one or two non-optional arguments with intent (in) which correspond to monadic or dyadic operators

Operator Overloading II

```
module dynamic data
  type dynamics
     real(dp) :: x, y, z
  end type dynamics
  interface operator (*)
    module procedure scale tr, scale rt
  end interface operator (*)
  interface operator (+)
    module procedure add
 end interface operator (+)
contains
 type (dynamics) function scale tr(a,b) result(c)
    type (dynamics), intent(in)::a
    real(dp),intent(in) :: b
    type (dynamics) :: c
    c%x = a%x * b
   c%v = a%v * b
```

```
c%z = a%z * b
  end function scale tr
  type (dynamics) function scale rt(b,a) result(c)
    type (dynamics), intent(in)::a
    real(dp),intent(in) :: b
    type(dynamics) :: c
    c%x = b * a%x
   c%v = b * a%v
    c%z = b * a%z
  end function scale rt
  type (dynamics) function add(a,b) result(c)
    type (dynamics), intent(in) :: a,b
   type(dynamics) :: c
    c%x = a%x + b%x
   c%v = a%v + b%v
    c_{2} = a_{2} + b_{2}
  end function add
end module dynamic data
```

• The following operations are now defined for derived types a, b, c and scalar r

```
c = a * r

c = r * a

c = a + b
```

Operator Overloading III

 If operator overloading is not defined, the above operations would have to be executed as follows whereever needed

```
C%x = a%x * r
c%y = a%y * r
c%z = a%z * r

c%x = r * a%x
c%y = r * a%y
c%z = r * a%z

c%x = a%x + b%x
c%y = a%y + b%y
c%z = a%z + b%z
```

Object Based Programming

OOP Concepts

• Fortran 90 has some Object Oriented facilities such as

- data abstraction: user defined types (covered)
- a data hiding private and public attributes (covered)
- encapsulation modules and data hiding facilities (covered)
- inheritance and extensibility super-types, operator overloading and generic procedures (covered)
- opplymorphism user can program his/her own polymorphism by generic overloading
- 6 resuability modules

Pointers I

- In Fortran, a pointer variable or simply a pointer is best thought of as a "free-floating" name that may be associated with or "aliased to" some object.
- The object may already have one or more other names or it may be an unnamed object.
- The object represent data (a variable, for example) or be a procedure.
- A pointer is any variable that has been given the pointer attribute.
- A variable with the pointer attribute may be used like any ordinary variable.

Pointers II

• Each pointer is in one of the following three states:

undefined condition of each **pointer** at the beginning of a **program**, unless it has been initialized null not an alias of any data object

associated it is an alias of some target data object

• pointer objects must be declared with the pointer attribute

```
real, pointer :: p
```

 Any variable aliased or "pointed to" by a pointer must be given the target attribute

```
real, target :: r
```

• To make p an alias to r, use the pointer assignment statement

```
p => r
```

Pointers III

The variable declared as a pointer may be a simple variable as above, an array
or a structure

```
real, dimension(:), pointer :: v
```

• pointer v declared above can now be aliased to a 1D array of reals or a row or column of a multi-dimensional array

```
real, dimension(100,100), target :: a
v => a(5,:)
```

• pointer variables can be used as any other variables

```
For example, print *, v and print *, a(5,:) are equivalent
```

```
v = 0.0 is the same as a (5, :) = 0.0
```

• pointer variables can also be an alias to another pointer variable

• Consider the following example

```
real, target :: r
real, pointer :: p1, p2
r = 4.7
p1 => r
p2 => r
print *, r, p1, p2
r = 7.4
print *, r, p1, p2
```

- The output on the screen will be
- Changing the value of r to 7.4 causes the value of both p1 and p2 to change to 7.4

• Consider the following example

```
real, target :: r1, r2
real, pointer :: p1, p2
r1 = 4.7; r2 = 7.4
p1 => r1; p2 => r2
print *, r1, r2, p1, p2
p1 = p2
print *, r1, r2, p1, p2
```

• The output on the screen will be

```
4.7 7.4 4.7 7.4
4.7 4.7 4.7 4.7
```

- The assignment statement p2 = p1 has the same effect of r2 = r1 since p1 is an alias to r1 and p2 is an alias to r2
- The allocate statement can be used to create space for a value and cause a
 pointer to refer to that space.

Pointers V

allocate (p1) creates a space for one real number and makes p1 an alias to that space.

- No real value is stored in that space so it is neccessary to assign a value to p1
- p1 = 4.7 assigns a value 4.7 to that allocated space
- Before a value is assigned to p1, it must either be associated with an unnamed target using the allocate statement or be aliased with a target using the pointer assignment statement.
- deallocate statement dissociates the pointer from any target and nullifies it deallocate (p1)

Pointer Intrinsic Functions

null intrinsic

- pointer variables are undefined unless they are initialized
- **pointer** variable must not be reference to produce a value when it is undefined.
- It is sometime desirable to have a **pointer** variable in a state of not pointing to anything
- The null intrinsic function nullifies a pointer assignment so that it is in a state of not pointing to anything

```
p1 => null()
```

- If the target of p1 and p2 are the same, then nullifying p1 does not nullify p2
- If p1 is null and p2 is pointing to p1, then p2 is also nullified.

associated intrinsic

 The associated intrinsic function queries whether a pointer varibale is pointing to, or is an alias for another object.

```
associated (p1, r1) and associated (p2, r2) are true, but associated (p1, r2) and associated (p2, r1) are false
```

Extended Data Types I

• Recall the derived type example which has as a component another derived type

```
type, public :: line_type
  real :: x1, y1, x2, y2
end type line_type
type, public :: vector_type
type(line_type) :: line !position of center of sphere
  integer :: direction ! 0=no direction, 1=(x1,y1)->(x2,y2)
end type vector_type
```

• An object, c, of type vector_type is referenced as c%line%x1, c%line%y1, c%line%x2, c%line%y2 and c%direction which can be cumbersome.

Extended Data Types II

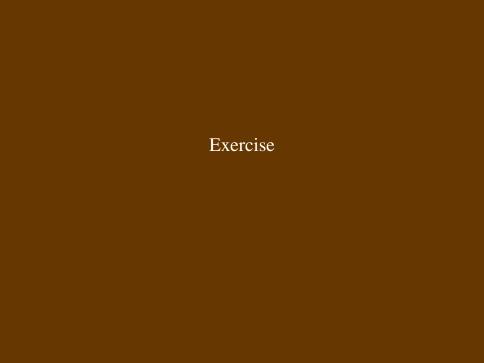
• In Fortran, it is possible to extend the base type line_type to other types such as vector_type and painted_line_type as follows

```
type, public, extends(line_type) :: vector_type
  integer :: direction
end type vector_type
type, public, extends(line_type) :: painted_line_type
  integer :: r, g, b ! rgb values
end type painted_line_type
```

- An object,c of type vector_type inherits the components of the type line_type and has components x1, y1, x2, y2 and direction and is referenced as c%x1, c%y1, c%x1, c%y2 and c%direction
- Similarly, an object, d of type painted_line_type is referenced as d%x1, d%y2, d%x2, d%y2, d%r, d%g and d%b
- The three derived types constitute a class; the name of the class is the name of the base type line_type

References

- Fortran 95/2003 Explained, Michael Metcalf
- Modern Fortran Explaned, Michael Metcalf
- Guide to Fortran 2003 Programming, Walter S. Brainerd
- Introduction to Programming with Fortran: with coverage of Fortran 90, 95, 2003 and 77, I. D. Chivers
- Fortran 90 course at University of Liverpool,
 http://www.liv.ac.uk/HPC/F90page.html
- Introduction to Modern Fortran, University of Cambridge, http://www.ucs.cam.ac.uk/docs/course-notes/unix-courses/Fortran
- Scientific Programming in Fortran 2003: A tutorial Including Object-Oriented Programming, Katherine Holcomb, University of Virginia.



Hands-On Exercise: Molecular Dynamics

- Molecular Dynamics code for melting of solid Argon using Lennard-Jones Potential.
- Your goal is to rewrite the code using Modern Fortran concepts that you have grasped.
- This exercise is more of a "What concepts have I learned of Modern Fortran?", so there are multiple correct solutions
- Code can be obtained from http://www.hpc.lsu.edu/training/archive/tutorials.php:
- md-orig.f90 is the original code that you should begin working on (this is the same code that was shown in todays slides)
- There is no "correct solution", however there are multiple solutions md-v{1-5}.f90 based on various concepts presented.
- It's entirely up to you to decide which solution you want to arrive at.
- Compare the results of your edited code with that of md-v0.out. If the results are not the same, debug your code.

Calculate pi by Numerical Integration I

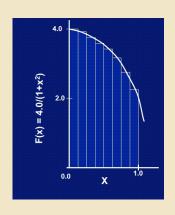
We know that

$$\int_0^1 \frac{4.0}{(1+x^2)} \, dx = \pi$$

 So numerically, we can approxiate pi as the sum of a number of rectangles

$$\sum_{i=0}^{N} F(x_i) \Delta x \approx \pi$$

Meadows et al, A "hands-on" introduction to OpenMP, SC09



Calculate pi by Numerical Integration II

Algorithm 1 Pseudo Code for Calculating Pi

```
\begin{array}{l} \mathbf{program} \ \mathsf{CALCULATE\_PI} \\ step \leftarrow 1/n \\ sum \leftarrow 0 \\ \mathbf{do} \ i \leftarrow 0 \cdots n \\ x \leftarrow (i+0.5) * step; sum \leftarrow sum + 4/(1+x^2) \\ \mathbf{end} \ \mathbf{do} \\ pi \leftarrow sum * step \\ \mathbf{end} \ \mathbf{program} \end{array}
```

 SAXPY is a common operation in computations with vector processors included as part of the BLAS routines

$$y \leftarrow \alpha x + y$$

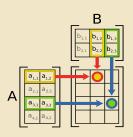
• Write a SAXPY code to multiply a vector with a scalar.

Algorithm 2 Pseudo Code for SAXPY

```
program SAXPY n \leftarrow some large number x(1:n) \leftarrow some number say, 1 y(1:n) \leftarrow some other number say, 2 a \leftarrow some other number ,say, 3 do i \leftarrow 1 \cdots n y_i \leftarrow y_i + a * x_i end do end program SAXPY
```

Matrix Multiplication I

- Most Computational code involve matrix operations such as matrix multiplication.
- Consider a matrix C which is a product of two matrices A and B: Element i,j of C is the dot product of the i^{th} row of A and j^{th} column of B
- Write a MATMUL code to multiple two matrices.



Matrix Multiplication II

Algorithm 3 Pseudo Code for MATMUL

```
\begin{aligned} & \textbf{program MATMUL} \\ & m, n \leftarrow \text{some large number} \leq 1000 \\ & \text{Define } a_{mn}, b_{nm}, c_{mm} \\ & a_{ij} \leftarrow i+j; b_{ij} \leftarrow i-j; c_{ij} \leftarrow 0 \\ & \textbf{do } i \leftarrow 1 \cdots m \\ & \textbf{do } j \leftarrow 1 \cdots m \\ & c_{i,j} \leftarrow \sum_{k=1}^n a_{i,k} * b_{k,j} \\ & \textbf{end do} \\ & \textbf{end do} \\ & \textbf{end program MATMUL} \end{aligned}
```