

Tutorial 8B: Type inference

Exercise 1: In the lambda-calculus, the **values** or **answers** of a computation are the **normal forms** of reduction: those λ -terms without redexes.

a) Convince yourself that if a term is in normal form, it has the shape

$$\lambda x_1. \dots \lambda x_n. x N_1 \dots N_k$$

where each N_i is in normal form. (Either n or k may be zero, that is, a term might just look like $x N_1 \dots N_k$ or like $\lambda x_1. \lambda x_2. x$, for instance.)

b) Using this characterization, write down all the closed normal forms (i.e. normal forms with no free variables) of the following types:

- $o \rightarrow o$
- $o \rightarrow o \rightarrow o$
- $(o \rightarrow o) \rightarrow o \rightarrow o$
- $(o \rightarrow o) \rightarrow o$

Exercise 2: Recall the following standard terms in the lambda-calculus:

$$\text{true} \triangleq \lambda x. \lambda y. x$$

$$\text{false} \triangleq \lambda x. \lambda y. y$$

$$\text{iszero} \triangleq \lambda n. n (\lambda f. \text{false}) \text{true}$$

$$\text{add} \triangleq \lambda m. \lambda n. \lambda f. \lambda x. m f (n f x)$$

Find the most general types of the following terms, or show that none exists.

- a) the numerals 2, 1, and 0
- b) the terms true and false
- c) add
- d) add 0 1
- e) iszero
- f) iszero 2
- g) $\lambda x. x x$
- h) $\lambda x. \lambda f. x (f x)$
- i) $\lambda f. \lambda x. \lambda y. f (x y) (y x)$
- j) $\lambda f. f (\lambda x. f (\lambda y. x))$