

# Lecture 11: Selecting prior information

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## Information entropy

# Prequel to the principle of maximum entropy

- You have a discrete random variable  $X$ .
- You know what values it takes, say  $x_1, \dots, x_N$ .
- You also have some information about it, e.g., the expectation of  $X$  is 0.5, the variance 0.1, etc.
- What probability distribution do you assign to  $X$ ?

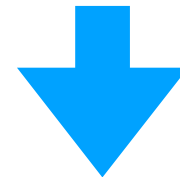
# Prequel to the principle of maximum entropy



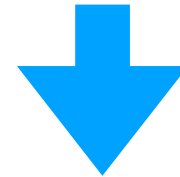
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*The knowledge of average values does give a reason for preferring some possibilities to others, but we would like [...] to assign a probability distribution which is as uniform as it can be while agreeing with the available information.*

— E. T. Jaynes



The uniform is the most “uncertain” distribution.



We need to assign the distribution that has the maximum uncertainty while being consistent with the data.

# Measure of uncertainty

- You can think of the probability mass function of  $X$  as a vector  $p = (p_1, \dots, p_N)$ .
- We are looking for a function  $\mathbb{H}(p_1, \dots, p_N)$  that tells how much uncertainty there is in this probability distribution.
- In 1948, Claude Shannon posed and answer this problem in the paper “A Mathematical Theory of Communication.”
- The function he came up with is called “information entropy.”

# What did Shannon do?

- He assumed that  $\mathbb{H}(p_1, \dots, p_N)$  is just a real number.
- He posed some obvious axioms for  $\mathbb{H}(p_1, \dots, p_N)$ , e.g., it should be continuous, it should be maximized when given the uniform distribution.
- Then he did a little bit of math and proved that:

$$\mathbb{H}(p_1, \dots, p_N) = - \sum_{i=1}^N p_i \log p_i \quad (\text{information entropy})$$

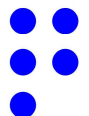


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# Notational convention for information entropy

$X$  takes values in  $\{x_1, x_2, \dots\}$

$$H[p(X)] := - \sum_x p(x) \log p(x) = - \mathbb{E}[\log p(X)].$$



# Information entropy of a distribution with two outcomes

$$X = \begin{cases} 0, & p_0 \\ 1, & p_1 = 1 - p_0 \end{cases}$$

$$\begin{aligned} H[p(X)] &= - \sum_x p(x) \log p(x) \\ &= -p_0 \log p_0 - p_1 \log p_1 \\ &= -p_0 \log p_0 - (1-p_0) \log (1-p_0) \end{aligned}$$

