

Skip Lists

- provide a clever compromise to efficiently support search/update operations
- a skip list for map M consists of a series of lists $\{S_0, S_1, \dots, S_n\}$
 - each list S_i for map M stores a subset of the items from the map sorted by inc. keys
 - each list also has 2 sentinel keys denoted $-\infty$ and $+\infty$, where $-\infty$ smaller than every possible key and $+\infty$ larger than every possible key
- skip lists also satisfy:
 - 1.) List S_0 contains every item of map M
 - 2.) For $i=1, \dots, n-1$, list S_i contains randomly generated subset of items of list S_{i-1}
 - 3.) List S_n contains only $-\infty$ and $+\infty$
- it is customary to view/think of skip lists with list S_0 at bottom and lists S_1, \dots, S_n above it
- we refer to n as the height of skip list S
- intuitively, the lists are set up so that S_{i+1} contains more or less alt. items of S_i

Search/Update in Skip Lists

- Searching: suppose we are given search key k ;
 - we begin skip search by setting position at top left position of S
 - then perform following steps:
 - 1.) if $S.\text{below}(p)$ is None, search terminates -- we are at bottom and have located the item in S with the largest key \leq search key k . Otherwise, drop down to next lower level in current tower by setting $p = S.\text{below}(p)$
 - 2.) Starting at position p , move p forward until it is at rightmost position on present level such that $\text{key}(p) \leq k$. Called scan forward step
 - 3.) Return to step 1

Algo SkipSearch(k)

$p = \text{start}$

while $\text{below}(p) \neq \text{None}$

$p = \text{below}(p)$

while $k > \text{key}(\text{next}(p))$

$p = \text{next}(p)$

return p