

Quadratic-Time Algo

def prefix-average(S):

 $n = \text{len}(S)$   $\longrightarrow$  executes in constant time,  $O(1)$  $A = [0] * n$  // create list of  $n$  zeros  $\rightarrow$  const # of primitive elements, thus  $O(1)$ for  $j$  in range( $n$ ):total = 0 // computing  $S[0] + \dots + S[j]$ for  $i$  in range( $j+1$ ):total +=  $S[i]$  $A[j] = \text{total} / (j+1)$  // record average

return A

 $\rightarrow$  there are 2 nested for-loops;

outer loop {

- outer loop (controlled by  $j$ ) is executed  $n$  times for  $j = 0, \dots, n-1$
- thus, total = 0 AND  $A[j] = \text{total} / (j+1)$  executed  $n$  times each
- implies these statements contribute # of primitive operations proportional to  $n$ , thus  $O(n)$  time

inner loop {

- body of inner loop (controlled by  $i$ ) is executed  $j+1$  times
- thus, total +=  $S[i]$  is executed  $1+2+3+\dots+n$  times
- $1+2+3+\dots+n = \frac{n(n+1)}{2}$  }  $O(n^2)$

 $\rightarrow$  The running time is thus  $O(n^2)$ .

def prefix-average1(S):

 $n = \text{len}(S)$  $A = [0] * n$ for  $j$  in range( $n$ ): $A[j] = \text{sum}(S[0:j+1]) / (j+1)$ 

return A

 $\rightarrow$  replaced inner loop using single expr. sum to compute partial sum $\rightarrow$   $\text{sum}(S[0:j+1])$  is a function call and takes  $O(j+1)$  time $\rightarrow$  the slice also takes  $O(j+1)$  time $\rightarrow$  thus, this implementation is also  $O(n^2)$