

## Searches

- attempt to locate a particular key in a BST by viewing it as a decision tree
- in the algorithm, if  $k$  occurs in a subtree rooted at  $p$ , a call to `TreeSearch` results in the position at which the key is found

Algorithm `TreeSearch(T, p, k)`:

if  $k == p.key()$  then

return  $p$

else if  $k < p.key()$  and  $T.left(p)$  is not None

return `TreeSearch(T, T.left(p), k)`

else if  $k > p.key()$  and  $T.right(p)$  is not None

return `TreeSearch(T, T.right(p), k)`

return  $p$

## Insertion

- the map command `M[k] = v` begins with a search for  $k$ 
  - if found, reassign
  - otherwise, insert a new node for the item into underlying tree  $T$  in place of empty subtree that was reached at end of failed search

Algorithm `TreeInsert(T, k, v)`:

$p = \text{TreeSearch}(T, T.root(), k)$

if  $k = p.key$  then

set  $p$  value to  $v$

else if  $k < p.key()$  then

add node with item  $(k, v)$  as left child of  $p$

else

add node with item  $(k, v)$  as right child of  $p$