

Exo-1 for  $c \in \mathbb{R}^d$ ,  $b \in \mathbb{R}^n$ ,  $A \in \mathbb{R}^{n \times d}$ , we want to solve

$$\min_x c^T x \quad \text{st} \quad \begin{cases} Ax = b & (\text{P}) \\ x \geq 0 & \end{cases} \quad \text{and} \quad \max_y b^T y \quad \text{st} \quad A^T y \leq c \quad (\text{D})$$

(1) The dual problem can be written as

$$L(\lambda, \gamma) = \min_x L(x, \lambda, \gamma) = \min_x [c^T x + \gamma^T (Ax - b) - \lambda^T x]$$

however, we have  $\frac{\partial L(x, \lambda, \gamma)}{\partial x} = c^T + \gamma^T A - \lambda^T$

$$\frac{\partial L(x, \lambda, \gamma)}{\partial x} = 0 \Leftrightarrow c^T + \gamma^T A - \lambda^T = 0$$

As  $L$  is linear in  $x$ , we have

$$L(\lambda, \gamma) = \begin{cases} -\gamma^T b & \text{if } A^T \gamma - \lambda + c = 0 \\ -\infty & \text{if } A^T \gamma - \lambda + c \neq 0 \end{cases}$$

① we want to maximize  $\gamma$  with  $A^T \gamma - \lambda + c = 0$   
 i.e.  $\max_{\gamma} -b^T \gamma$  st  $A^T \gamma + c \leq 0$

(2) The dual problem can be written as

$$L(\lambda) = \min_y L(\lambda, y) = \min_y [-b^T y + \lambda^T (A^T y - c)]$$

$$= \min_y (b^T + \lambda^T A^T) y - \lambda^T c$$

$$y = \begin{cases} -\lambda^T c & \text{if } -b^T + \lambda^T A^T = 0 \\ -\infty & \text{otherwise} \end{cases} \quad \text{with } \lambda \geq 0$$

$L(\lambda, y)$  linear in  $y$

② we want to minimize

$$+ \lambda^T c \quad \text{st} \quad \begin{cases} b^T + \lambda^T A^T = 0 \\ \lambda \geq 0 \end{cases}$$

$$\boxed{\min_{\lambda} + c^T \lambda \quad \text{st} \quad \begin{cases} b^T + A^T \lambda = 0 \\ \lambda \geq 0 \end{cases}}$$

(3) The dual problem can be written as

$$\begin{aligned}
 L(\lambda_1, \lambda_2, \gamma) &= \min_{x, y} c^T x - b^T y + \gamma^T (Ax - b) - \lambda_1^T x + \lambda_2^T (A^T y - c) \\
 &= \min_{x, y} (c^T + \gamma^T A - \lambda_1^T) x + (\lambda_2^T A^T - b^T) y - \gamma^T b - \lambda_1^T c \\
 &\stackrel{\gamma}{=} \begin{cases} -\gamma^T b - \lambda_1^T c & \text{if } c^T + \gamma^T A = \lambda_1^T \text{ and } b^T = \lambda_2^T A^T \\ -\infty & \text{otherwise} \end{cases} \\
 &\text{L linear in } x \text{ and } y
 \end{aligned}$$

(i) we want to maximize

$$\max_{\gamma, \lambda_1, \lambda_2} -\gamma^T b - \lambda_1^T c \quad \text{st} \quad \begin{cases} c^T + \gamma^T A = \lambda_1^T \\ \lambda_1 \geq 0 \quad \lambda_2^T \leq 0 \\ A\lambda_2 = b \end{cases}$$

The condition on  $\lambda_1$  can be forgotten as  $\lambda_1$  does not appear in the min; we just have to keep  $A^T \gamma + c \geq 0$ . Finally, maximizing on  $\gamma$  is the same as maximizing on  $-\gamma$ . We can thus write that our problem is equivalent to

$$\begin{array}{ll} \min_{\gamma, \lambda_2} & b^T \gamma - c^T \lambda_2 \quad \text{st} \quad \begin{cases} c \leq A^T \gamma \quad \lambda_2 \geq 0 \\ A\lambda_2 = b \end{cases} \end{array}$$

(i) The problem is self-dual.

(4) The problem (Self-Dual) is separable in  $x$  and  $y$ . Indeed, as the problem is bounded and feasible, we can separate the conditions in (Self-Dual), giving

$$(\text{Self-Dual}) \Leftrightarrow \min_x c^T x + \min_y -b^T y \quad \text{st} \quad \begin{cases} Ax = b \\ x \geq 0 \\ A^T y \leq c \end{cases}$$

; which shows that

(Self-Dual) is equivalent to solving (P) and (D)

As the problem is feasible, we have shown during the course that strong duality holds for (D). we thus have

$$\max_y b^T y = \max_{\lambda} w^T + c^T \lambda \text{ s.t. } \begin{cases} -b + A\lambda = 0 \\ \lambda \geq 0 \end{cases}$$

which concludes

③ the optimal value for (Self-Dual) is 0

Ex 2 For  $A \in \mathbb{R}^{n \times d}$ ,  $b \in \mathbb{R}^n$ , we want to solve  $\min_x \|Ax - b\|_2^2 + \|x\|_1$

(1) let  $x, y \in \mathbb{R}^d$ , we want to find  $f^*(y) = \sup_z (y^T z - \|z\|_1)$

we have

$$y^T z - \sum_{i=1}^d |z_i| = \sum_{i=1}^d x_i y_i - \sum_{i=1}^d |z_i|$$

\* Case n°1 :  $\exists i_0 \in [1, d] : y_{i_0} > 1$ . Then we take  $x_i = \begin{cases} n & \text{if } i = i_0 \\ 0 & \text{otherwise} \end{cases}$  and we have

$$y^T z - \sum_{i=1}^d |z_i| = y_{i_0} x_{i_0} - z_{i_0} = n(y_{i_0} - 1) \xrightarrow{n \rightarrow +\infty} +\infty$$

\* Case n°2 :  $\exists i_0 \in [1, d] : y_{i_0} < -1$ . We take  $x_i = \begin{cases} -n & \text{if } i = i_0 \\ 0 & \text{otherwise} \end{cases}$  and we have

$$y^T z - \sum_{i=1}^d |z_i| = -ny_{i_0} - n = -n(y_{i_0} + 1) \xrightarrow{n \rightarrow +\infty} +\infty$$

\* Case n°3 :  $\forall i \in [1, d], |y_i| \leq 1$  then

$$\sum_{i=1}^d x_i y_i - \sum_{i=1}^d |z_i| = \sum_{i=1}^d \underbrace{x_i y_i}_{\leq |z_i|} - |z_i| \leq \sum_{i=1}^d 0 = 0. \text{ As } 0 \text{ is}$$

reached by choosing  $x=0$ , we have  $f^*(y)=0$  in this case

Q.  $f^*(y) = \begin{cases} +\infty & \text{if } \|y\|_\infty > 1 \\ 0 & \text{if } \|y\|_\infty \leq 1 \end{cases}$

(2) Let's compute the dual of (RLS)  
we have

$$(RLS) \Leftrightarrow \inf_{x,y} \|y\|_2^2 + \|x\|_1 \text{ st } y = Ax - b.$$

Now

$$\begin{aligned} L(\gamma) &= \inf_{x,y} \|y\|_2^2 + \|x\|_1 + \gamma^T(y - Ax + b) \\ &= \inf_x (\|x\|_1 - \gamma^T A x) + \inf_y (\|y\|_2^2 + \gamma^T y) + \gamma^T b \\ &= \underbrace{\sup_x (\gamma^T A x - \|x\|_1)}_{= 0 \text{ if } \|\gamma^T A\|_\infty \leq 1; +\infty \text{ otherwise}} + \inf_y (\|y\|_2^2 + \gamma^T y) + \gamma^T b. \end{aligned}$$

The function  $g: y \mapsto \|y\|_2^2 + \gamma^T y$  is convex and differentiable

$$Dg(y) = 2y + \gamma = 0 \Leftrightarrow y = -\frac{1}{2}\gamma$$

we thus have

$$\begin{aligned} L(\gamma) &= \begin{cases} \frac{1}{4} \|\gamma\|_2^2 + \left(-\frac{1}{2}\right) \|\gamma\|_2^2 + \gamma^T b & \text{if } \|\gamma^T A\|_\infty \leq 1 \\ +\infty & \text{if } \|\gamma^T A\|_\infty > 1 \end{cases} \\ &= \begin{cases} -\frac{1}{4} \|\gamma\|_2^2 + \gamma^T b & \text{if } \|A^T \gamma\|_\infty \leq 1 \\ +\infty & \text{otherwise} \end{cases}. \end{aligned}$$

The dual problem is thus

$$\boxed{\max_{\gamma} -\frac{1}{4} \|\gamma\|_2^2 + \gamma^T b \text{ if } \|A^T \gamma\|_\infty \leq 1}$$

### Eo3

(1) we consider  $\min_{w, z} \frac{1}{n} \mathbf{1}^T z + \frac{1}{2} \|w\|_2^2$  st  $\begin{cases} z_i > 1 - y_i w^T x_i \\ z_i \geq 0 \end{cases} \forall i$

$$\begin{aligned} \min_w \frac{1}{n} \sum_{i=1}^n L(w, x_i, y_i) + \frac{1}{2} \|w\|_2^2 &= \mathcal{L} \min_{w, z} \frac{1}{n} \sum_{i=1}^n \max\{0, 1 - y_i w^T x_i\} + \frac{1}{2} \|w\|_2^2 \\ &= \mathcal{L} \min_{w, z} \frac{1}{n} \sum_{i=1}^n z_i + \frac{1}{2} \|w\|_2^2 \text{ st } z_i = \max\{0, 1 - y_i w^T x_i\} \\ &= \mathcal{L} \min_{w, z} \frac{1}{n} \mathbf{1}^T z + \frac{1}{2} \|w\|_2^2 \text{ st } z_i = \max\{0, 1 - y_i w^T x_i\} \end{aligned}$$

As we have  $z_i \geq 0 \quad \forall i \in \{1, n\}$ , if we consider  $a$  such that

$\exists i_0 : a_{i_0} > \max\{0, 1 - y_{i_0} w^T x_{i_0}\}$ , then we will have.

$$\frac{1}{n} \mathbf{1}^T a < \frac{1}{n} \mathbf{1}^T b \quad \text{where } \begin{cases} b \text{ respects the conditions} \\ b = \begin{cases} a_i & \text{for } i \neq i_0 \\ \max\{0, 1 - y_i w^T x_i\} & \text{if } i = i_0 \end{cases} \end{cases}$$

we thus have

$$\min_w \frac{1}{n} \sum L(w, x_i, y_i) + \frac{1}{2} \|w\|_2^2 = \mathcal{L} \min_{w, z} \frac{1}{n} \mathbf{1}^T z + \frac{1}{2} \|w\|_2^2 \text{ st } \begin{cases} z \geq 0 \\ z_i = 1 - y_i w^T x_i \end{cases}$$

(1) Sep 2 is equivalent to Sep 1, we just have a  $c$  multiplicative added.

(2) we consider

$$\begin{aligned} L((\lambda_i)_{i \in \{1, n\}}, \pi) &= \min_{w, z} \frac{1}{n} \mathbf{1}^T z + \frac{1}{2} \|w\|_2^2 + \sum_{i=1}^n \lambda_i (1 - y_i w^T x_i - z_i) - \pi^T z \\ &= \min_{w, z} \left[ \frac{1}{n} \mathbf{1}^T - \lambda - \pi \right]^T z + \min_w \left[ \frac{1}{2} \|w\|_2^2 - \sum_i \lambda_i y_i w^T x_i \right] \\ &+ \sum_i \lambda_i \\ &= \begin{cases} \min_w \frac{1}{2} \|w\|_2^2 - \sum_{i=1}^n \lambda_i y_i w^T x_i + \sum_i \lambda_i & \text{if } \frac{1}{n} \mathbf{1}^T - \lambda - \pi = 0 \\ -\infty & \text{otherwise} \end{cases} \end{aligned}$$

we consider  $g: w \mapsto \frac{1}{2} \|w\|_2^2 - \sum_i \lambda_i y_i x_i^T w$   
 $g$  is differentiable and convex and we have.

$$\nabla g(w) = w - \sum_i \lambda_i y_i x_i = 0 \Leftrightarrow w = \sum_i \lambda_i y_i x_i.$$

Consequently, we have

$$\begin{aligned} g(w_{\text{min}}) &= \frac{1}{2} \left( \sum_i \lambda_i y_i x_i \right) \left( \sum_j \lambda_j y_j x_j \right) - \sum_i \lambda_i y_i \left( \sum_j \lambda_j y_j x_j \right)^T x_i \\ &= \frac{1}{2} \sum_{1 \leq i, j \leq n} \lambda_i y_i \lambda_j y_j x_i^T x_j - \sum_{1 \leq i, j \leq n} \lambda_i \lambda_j y_i y_j x_i^T x_j \\ &= -\frac{1}{2} \sum_{1 \leq i, j \leq n} \lambda_i \lambda_j y_i y_j x_i^T x_j. \end{aligned}$$

(Q) The dual problem of (Sep 2) is

$$\max_{\lambda, \pi} \sum_{i=1}^n \lambda_i - \frac{1}{2} \sum_{1 \leq i, j \leq n} \lambda_i \lambda_j y_i y_j x_i^T x_j \quad \text{st} \quad \begin{cases} \frac{1}{n\tau} (1 - \lambda - \pi) = 0 \\ \lambda \geq 0 \\ \pi \leq 0 \end{cases}$$

i.e.

$$\max_{\lambda} \sum_{1 \leq i \leq n} \lambda_i - \frac{1}{2} \sum_{1 \leq i, j \leq n} \lambda_i \lambda_j y_i y_j x_i^T x_j \quad \text{st} \quad 0 \leq \lambda \leq \frac{1}{n\tau} \leq 1$$

**Evo4** we consider (K)  $\min_x c^T x$  st  $\sup_{a \in S} a^T x \leq b$  with  $S = \{a : C^T a \leq d\}$   
 As shown by the hint, we will first focus on  $\sup_a a^T x$  st  $C^T a \leq d$  (Q)  
 In this case, as we suppose the problem feasible, we have  $p^* = d^*$  (strong duality) as shown during class.

The dual<sup>dualization</sup> of (Q) is [as  $\sup_a a^T x = \inf_a -a^T x$ ]

$$g(\lambda) = \inf_a (-a^T x + \lambda^T C^T a - \lambda^T d) = \inf_a \underset{\text{linear in } a}{(-x^T + \lambda^T C^T) a - \lambda^T d}$$

$$= \begin{cases} -\lambda^T d & \text{if } C\lambda = x \\ -\infty & \text{otherwise} \end{cases}$$

dual problem: maximize  $-d^T \lambda$  st  $C\lambda = x$

Q: dual problem: minimize  $d^T z$  st  $Cz = x$   $\begin{matrix} \lambda \geq 0 \\ z \geq 0 \end{matrix}$   
 we have

$$(K) \Leftrightarrow \min_x c^T x \text{ st } \inf_z d^T z \leq b$$

Let  $x \in \mathbb{R}^n$  such that  $\exists z \in \mathbb{R}^m$ :  
 and  $\begin{cases} Cz = x \\ z \geq 0 \end{cases}$

$$\left\{ \begin{array}{l} d^T z \leq b \\ d^T z = \sup_{a \in S} a^T z \end{array} \right. \quad \begin{array}{l} z = x \\ z \geq 0 \end{array} \quad D_x := \{ \text{feasible points} \}$$

$$d^T z = \inf_{t \in D_x} d^T t \Leftrightarrow d^T z = \sup_{a \in S} a^T z \Leftrightarrow d^T z = \sup_{a \in S} (C^T a)^T z$$

$$\Leftrightarrow d^T z = d^T z \text{ which is true.}$$

Consequently, we have

$$(K) \Leftrightarrow \min_z c^T z \text{ st } \begin{array}{l} d^T z \leq b \\ C^T z = x \\ z \geq 0. \end{array}$$

# Exo 5

1. The Lagrangian can be written

$$L(x, \lambda, \gamma) = c^T x + \lambda^T (Ax - b) - \gamma^T x + x^T \text{diag}(\gamma) x \\ = x^T \text{diag}(\gamma) x + (c^T + \lambda^T A - \gamma^T) x - \lambda^T b$$

let's minimize this over  $x$ .

$$\frac{\partial L(x, \lambda, \gamma)}{\partial x} = 2x^T \text{diag}(\gamma) + c^T + \lambda^T A - \gamma^T = 0 \Leftrightarrow 2\text{diag}(\gamma)x = \gamma - A\lambda - c$$

Thus, we take  $y_i = \begin{pmatrix} y_i \\ 1 \end{pmatrix}$ ; we have

$$\gamma^T \text{diag}(\gamma) y_i = (y_i - y_{n+1}) \begin{pmatrix} y_i & y_{n+1} \\ 1 & y_{n+1} \end{pmatrix} = \sum_i y_i \cdot y_i^2$$

$$L(\lambda, \gamma) = \frac{1}{4} (\gamma^T - \lambda^T A - c^T) \text{diag}\left(\frac{1}{\gamma}\right) (\gamma - A^T \lambda - c) \quad \text{or} \quad (c^T + \lambda^T A - \gamma^T) \text{diag}\left(\frac{1}{\gamma}\right) (\gamma - A^T \lambda - c) \\ - \lambda^T b$$

$$\bar{\rho} = -\frac{1}{4} \sum_i \frac{(c_i + (A^T \lambda)_i - \gamma_i)^2}{\gamma_i} - \lambda^T b$$

$$y = \gamma - A^T \lambda - c$$

Dual problem:

$$\sup_{\lambda, \gamma} -\frac{1}{4} \sum_i \frac{(c_i + (A^T \lambda)_i - \gamma_i)^2}{\gamma_i} - \lambda^T b \quad \text{st} \quad \gamma \geq 0$$

$$\Leftrightarrow \sup_{\lambda} \sup_{\gamma \geq 0} - \frac{(c + (A^T \lambda)_i - \gamma_i)^2}{\gamma} - \lambda^T b \quad \text{st} \quad \lambda \geq 0$$

$$\Leftrightarrow \boxed{\sup_{\lambda} -5\lambda + \sum_{i=1}^n \min \{ 0, c_i + A_i^T \lambda \} \quad \text{st} \quad \lambda \geq 0}$$

2. The Lagrangian of the LP relaxation is

$$L(x, \lambda_1, \lambda_2, \lambda_3) = c^T x + \lambda_1^T (Ax - b) + \lambda_2^T (x - 1) - \lambda_3^T x \\ = (c^T + \lambda_1^T A + \lambda_2^T - \lambda_3^T) x - \lambda_1^T b - \lambda_2^T 1 - \lambda_3^T 1$$

$$L(\lambda_1, \lambda_2, \lambda_3) = \begin{cases} -\lambda_1^T b & \text{if } c^T + \lambda_1^T A + \lambda_2^T - \lambda_3^T = 0 \\ -\infty & \text{otherwise} \end{cases}$$

The dual problem is thus  
 $\sup -5\lambda_1 - 11\lambda_2 \text{ st } A^T\lambda + c + \lambda_2 - \lambda_3 = 0$ , which is equivalent  
 $\lambda_{1,2,3} \geq 0$

to the Lagrange relaxation problem.

- (c) The lower bounds for both problems are the same