

Much Ado About ☒Z

Alessandro Warth
Viewpoints Research

STEPS

... toward the reinvention of
programming



The STEPS Project

- **Goal** - To create a highly useful end-user system including:
 - operating system
 - programming environment
 - “applications”
 - graphics, sound, ...

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- **Goal** - To create a highly useful end-user system including:
 - operating system
 - programming environment
 - “applications”
 - graphics, sound, ...
- 
- personal
computing*

A photograph of a man with dark, wavy hair and a prominent mustache, wearing a maroon hoodie. He is smiling and looking towards the right side of the frame. In his hands, he holds a large, dark-colored book or volume. A white speech bubble originates from his mouth, containing the text "... in under 20,000 LOC!" in a black, sans-serif font.

... in
under
20,000
LOC!



... in
under
20,000
LOC!



Squeak 3.0

Welcome to... Squeak 3.0

Squeak is a work in progress based on Smalltalk-80, with which it is still reasonably compatible. Every Squeak release includes all source code for the Squeak system, as well as all source code for its Virtual Machine (VM, or interpreter, also written in Smalltalk).

Browser openBrowser
 [Blue items in this window are active text. If an item contains a URL, it will require internet access and may take a while to load].

Not only is all source code included, and changeable at will, it is also completely open and free. The Squeak system image runs bit-identically across all platforms, and VMs are available for just about every computer and operating system available. The history of the Squeak project can be read at
<http://st.cs.uiuc.edu/Smalltalk/Squeak/docs/OSPLA.Squeak.html>

The Squeak license and most other relevant information can be found on the Squeak Home Page,
<http://www.Squeak.org>

Morphic
 This release of Squeak uses the Morphic user interface. Squeak also includes an MVC architecture for building graphical projects (see the world menu 'open...').

Game Project



Process Browser

Method Finder

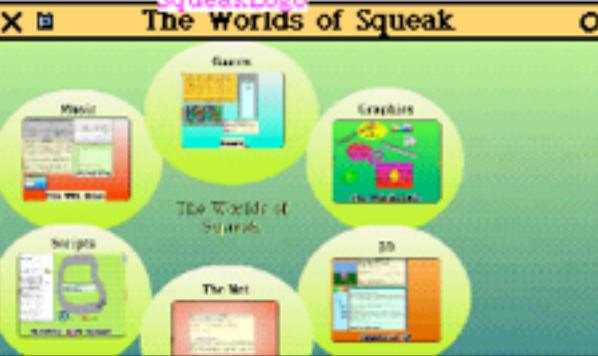
```
#(1 2 3 4). #(2 3). true
#(1 2 3 4) includesAllOf: #(2 3) --> true
#(1 2 3 4) includesAnyOf: #(2 3) --> true
#(1 2 3 4) windowReqNewLabel: #(2 3) --> true
#(1 2 3 4) ~= #(2 3) --> true
#(1 2 3 4) === #(2 3) --> true
```

Type a fragment of a selector in the top pane. Accessor methods are listed in the bottom pane.

Or, use an example to find a method in the system. Enter the receiver, args, and answer in the top pane with per-line separators. 3. 4. 7

SqueakLogo

The Worlds of Squeak



Senders of add:afterIndex: [4]
OrderedCollection hierarchy

Collections-Sequenceable

ProtoObject	Object	Collection	SequenceableCollection	OrderedCollection	GraphicSymbol	SortedCollection
accessing	copying	adding	removing	enumerating	private	testing
add:	adding	add:after:	remove:	next:	next:	test:
add:after:	add:at:	add:before:	remove:at:	next:butLast:	next:butFirst:	test:at:
addAll:	addAll:	addAllFirst:	removeAll:	next:allButLast:	next:allButFirst:	test:allBut:
addAllFirst:	addAllLast:	addFirst:	removeAllBut:	next:allButFirst:	next:allButLast:	test:allButButFirst:
addFirst:	addLast:	addLast:	removeLast:	next:butFirst:	next:butLast:	test:butFirst:
addLast:	remove:	remove:	removeFirst:	next:butLast:	next:butFirst:	test:butLast:

as an element of the receiver. Put it in Answer newObject."

firstIndex + index.

Squeak
~200 thousand LOC

open...
 dismiss this menu
 browser
 package browser
 method finder
 workspace
 file list
 file...
 transcript
 inner world
 simple change sort
 dual change sort
 email reader
 web browser
 IRC chat
 mvc project
 morphic project

T o o l s

Why?

- “Put people in charge of their own software destinies”
 - Can’t understand 40,000,000 LOC (an entire library!)
 - Can “own” 20,000 LOC (one 400-page book)

3 C 006 - 04766

STEREO

JOHN LENNON PLASTIC ONO BAND

POWER TO THE PEOPLE



APPLE



Why? (cont'd)

- Didactic value!
- Curriculum for univ. students to learn about powerful ideas, building complex systems...
- May even be useful at high-school level



The Path to 20K LOC

- Experimenting w/ new...
 - abstractions
 - PLs
 - DSLs

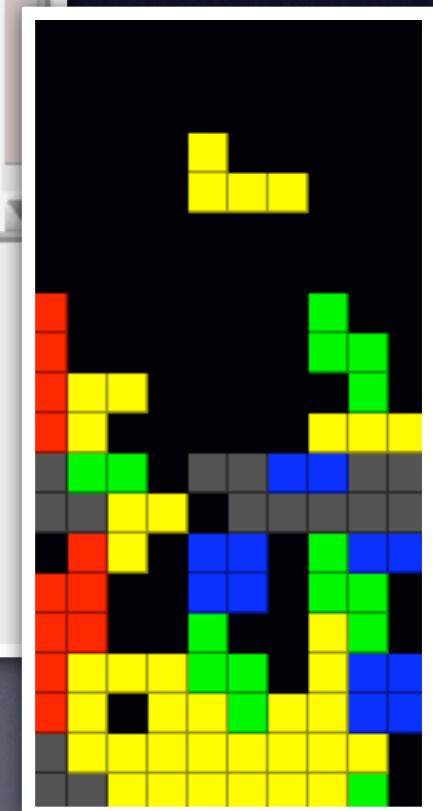
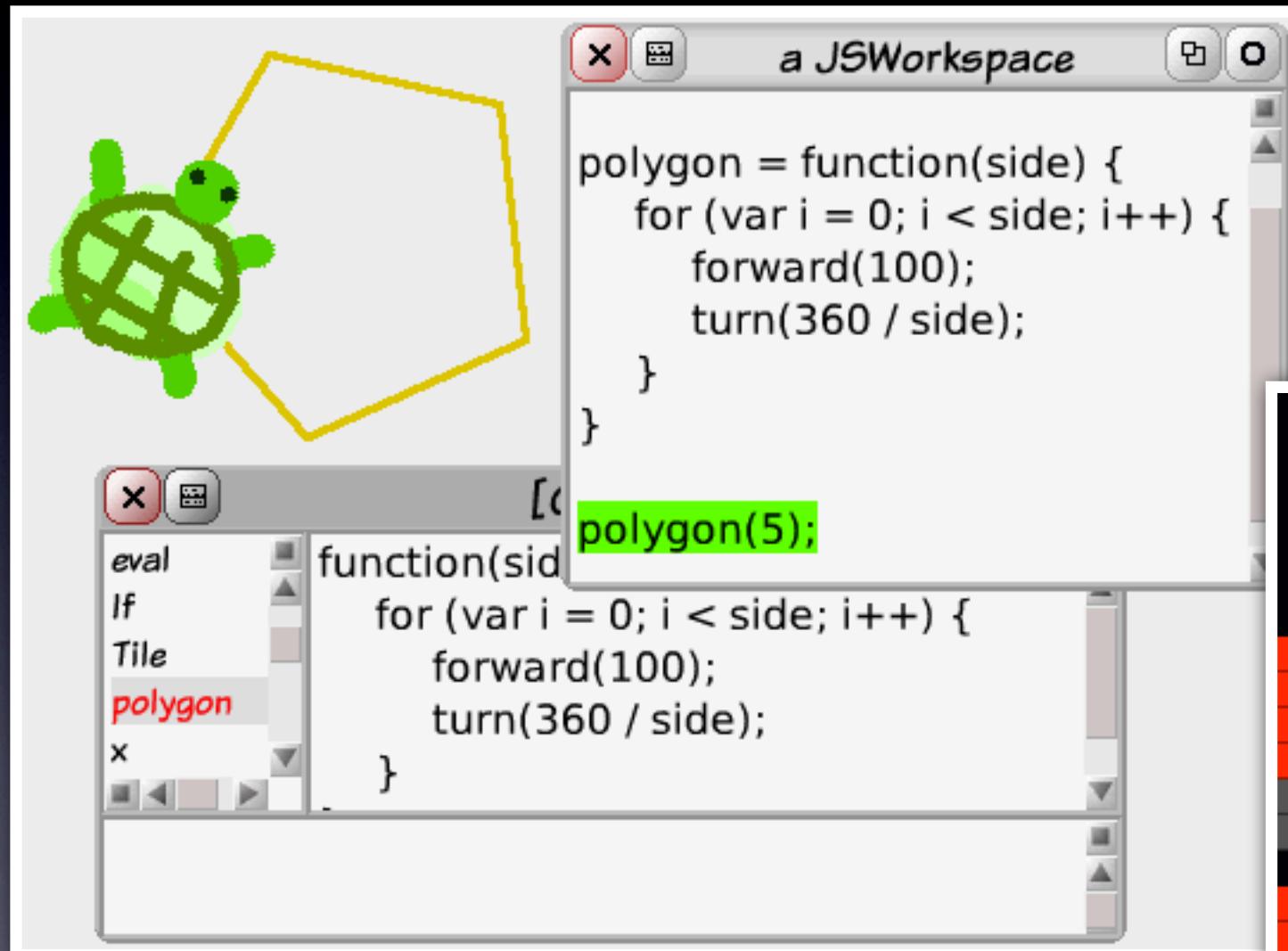


OMeta

Experimenting w/ Programming Languages

JavaScript (OMeta/Squeak)

~350 LOC



Sun's Lively Kernel (OMeta/COLA)

~300 LOC

Array

Body

Canvas

CheapMenuMorph

Color

ColorPickerMorph

Date

DropShadowCanvas

Element

Function

HandMorph

HandleMorph

InputEvent

Morph

MouseHandlerForDragging

Number

Object

PasteUpMorph

Pen

Point

PrimCanvas

PrimTextBox

PrimTextLine

Rectangle

Shape

StepHandler

String

TextMorph

WorldMorph

WorldState

collect

copyWithoutAll

includes

join

pop

push

sort

splice

toString

unshift

The interface features a sidebar on the left containing a list of kernel objects and their associated methods. On the right, there is a graphical workspace displaying several geometric shapes: a red rectangle, a green oval, a cyan triangle, a large blue circle, and a yellow star. Below these shapes, the word "Unbordered" is written. A code editor window is open at the bottom right, showing a snippet of JavaScript-like pseudocode for drawing concentric polygons.

```
P = new Pen();
P.setPenColor(Color.red);
for(var i=1; i<=50; i++)
  { P.go(2*i); P.turn(60); };
P.drawLines();
```

Toylog (OMeta/Squeak)

- Get children interested in logic!
- Front-end to Prolog, runs on Squeak
- ~70 LOC

Homer is Bart's father.

Marge is Bart's mother.

x is y's parent if x is y's father or
or x is y's mother.

Homer is not bart's parent.

Marge is bart's parent.

x is Bart's parent?

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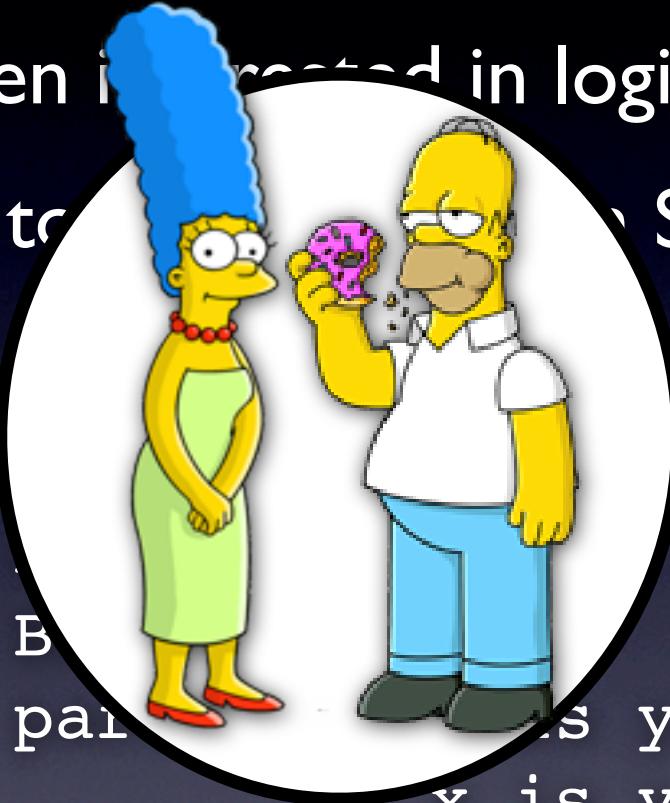
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x is Bart's parent?

Toylog (OMeta/Squeak)

- Get children interpreted in logic!
- Front-end to Squeak
- ~70 LOC



Homer is Bart's father or
Marge is Bart's mother.
 $x \text{ is } y\text{'s parent} \rightarrow x \text{ is } y\text{'s father or } x \text{ is } y\text{'s mother.}$

Homer is not bart's parent.

Marge is bart's parent.

→ $x \text{ is Bart's parent?}$

Prolog - OMeta/JS 2.0 Workspace

important info: (nevermind) ▾

go to project: (nevermind) ▾

previous versions of this project

Instructions [+/-](#)

Play Area [+/-](#)

Source [+/-](#)

```
ometa PrologTranslator <: Parser {  
    variable = spaces firstAndRest(`upper, `letterOrDigit):name      -> new Var(name.join('')),  
    symbol   = spaces firstAndRest(`lower, `letterOrDigit):name      -> new Sym(name.join('')),  
    clause   = symbol:sym "(" listOf(`expr, ','):args ")"          -> new Clause(sym, args),  
    expr     = clause | variable | symbol,  
    clauses  = listOf(#clause, ','),  
    rule     = clause:head ":-" clauses:body           "."          -> new Rule(head, body)  
                  | clause:head                         "."          -> new Rule(head, []),  
    prog     = (rule:r & clause -> r)*:rs clause:q "." spaces end -> {rules: rs, query: q}  
}
```

```
translateCode = function(x) {  
    var prog = PrologTranslator.matchAll(x, "prog")  
    solve(prog.query, prog.rules)  
}
```

```
nat(z).  
nat(s(X)) :- nat(X).  
nat(X).
```

~90 LOC

print it (ctrl+p)

do it (ctrl+d)

save it (ctrl + s)

**Portable
Programming
Language
Prototypes!**



A photograph of two men in a dimly lit bar. One man on the left, wearing a blue and white striped shirt, is playing a guitar. The other man on the right, wearing a grey t-shirt, is also playing a guitar. A speech bubble originates from the man in the grey t-shirt.

Forget Guitar
Hero... I could be the next
Dan Ingalls!

OMeta/JS



Undo

- An important feature in most applications
- Not just about fixing mistakes:
enables *exploration* w/o fear
- learn by trying things out
(errors not a big deal)
- tool for experimenting w/ different choices

SURE..

You're right in liking

MEAT



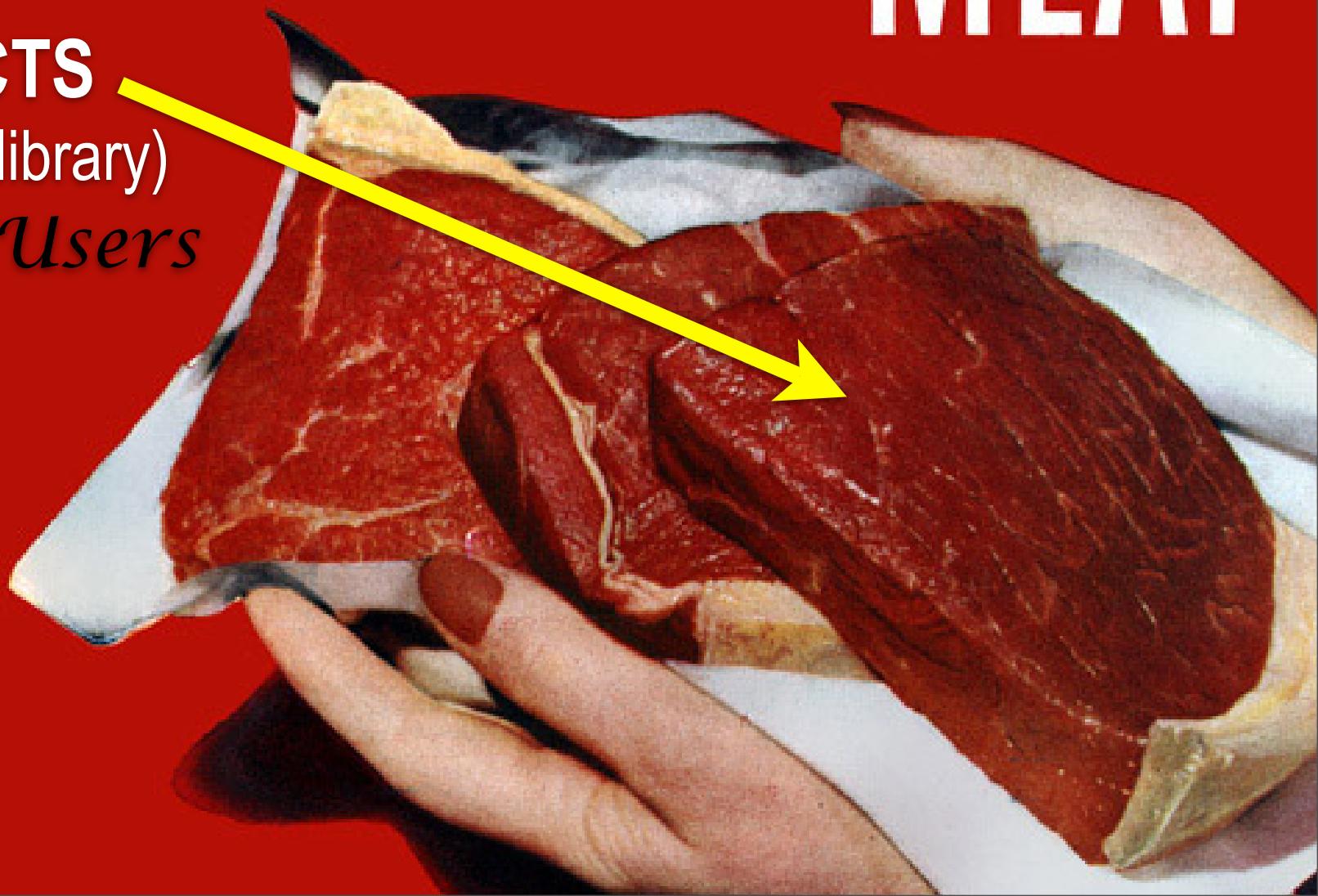
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You're right in liking MEAT

UOBJECTS

(framework / library)

Undo for Users



SURE..

You're right in liking MEAT

UOBJECTS

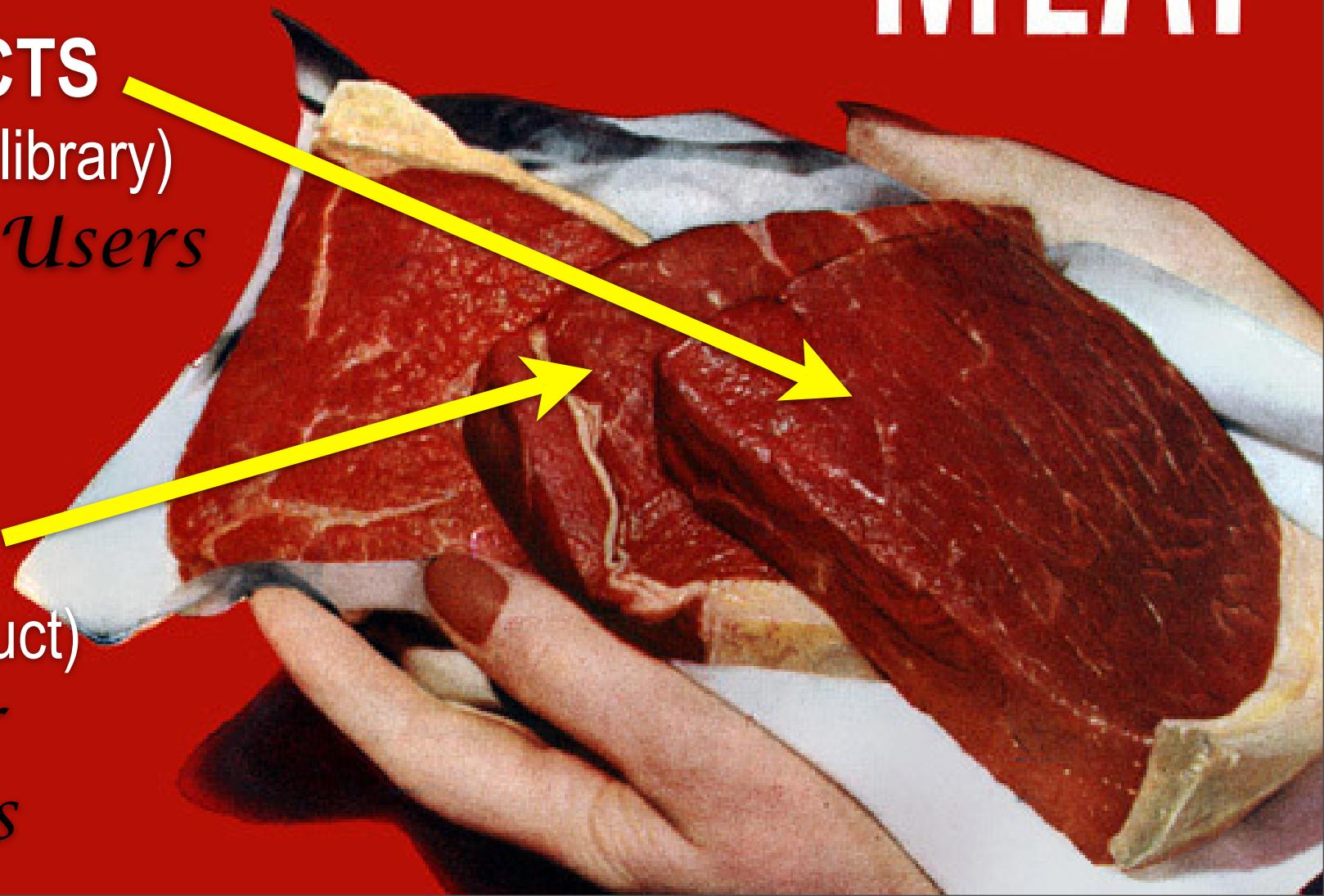
(framework / library)

Undo for Users

WORLDS

(language construct)

*Undo for
Programs*





SURE..

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MEAT

UOBJECTS

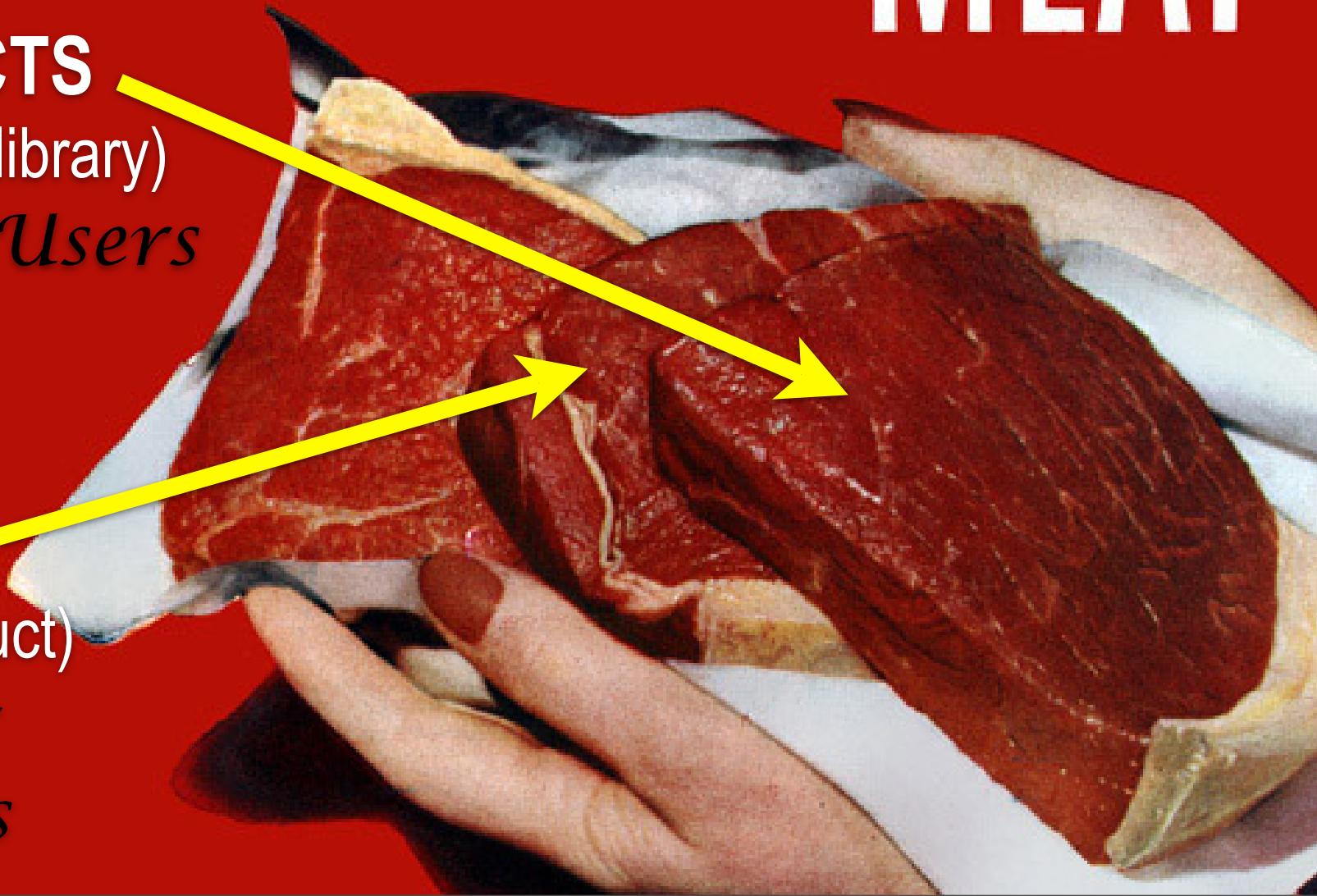
(framework / library)

Undo for Users

WORLDS

(language construct)

*Undo for
Programs*



Part I

UObjects:
Undo for Users

The Status Quo

- Most apps support *linear undo*
- ... which is implemented using:
 - *command* design pattern
 - *memento* design pattern

Command

do = ...

undo = do⁻¹

Command

$\text{do} = \dots$

$\text{undo} = \text{do}^{-1}$

may throw
away info

... need to
keep it for
undo

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must include
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must be
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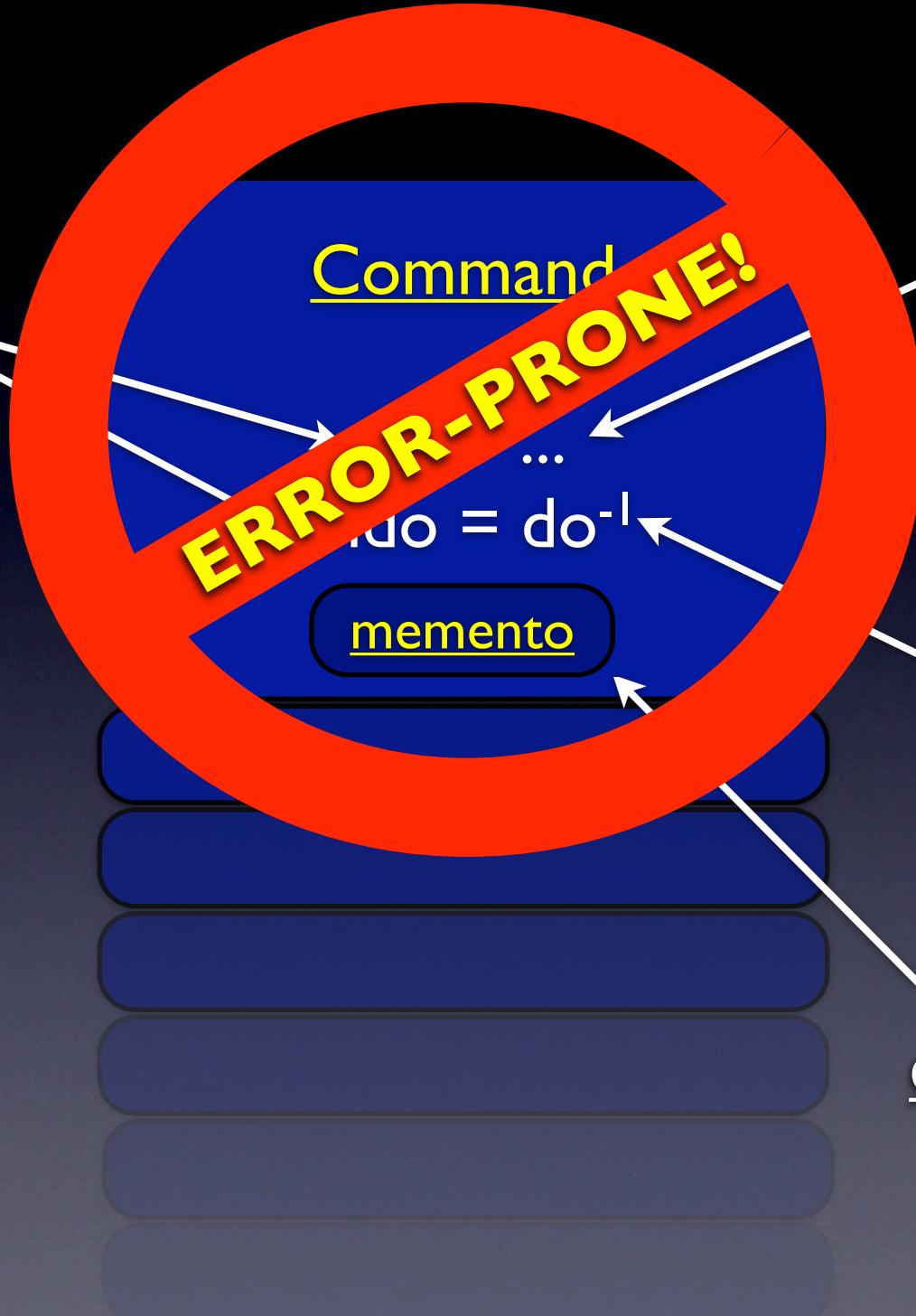
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Proposed Approach

- Why not generate memento on the fly?
 - i.e., record original values of all variables modified
 - (which may belong to multiple objects)
- Undo writes old values back into object(s)
- No need for error-prone idiom

Programming Model

- **UObject** — Undoable Object
 - operations: {#at, #at:put:, ...}
 - may only be modified inside...
- **UTransaction** — Undoable Transaction
 - may modify any no. of **UObjects**
 - operations: {#undo}

```
t1 := UTransaction eval: [  
    obj1 foo:'new'.  
    obj2 bar: 1234.  
    obj1 foo:'newer'.  
].
```

obj1
foo is 'old'

obj2
bar is 5

```
t2 := t1 undo.  "undo"  
t2 undo.        "redo"
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→ t1 := UTransaction eval: [
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t1



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t2 := t1 undo.  "undo"  
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t1
obj1's foo was 'old'
obj2's bar was 5

obj1
foo is 'new' → obj2
bar is 1234

```
t1 := UTransaction eval: [  
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t1 *
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t2

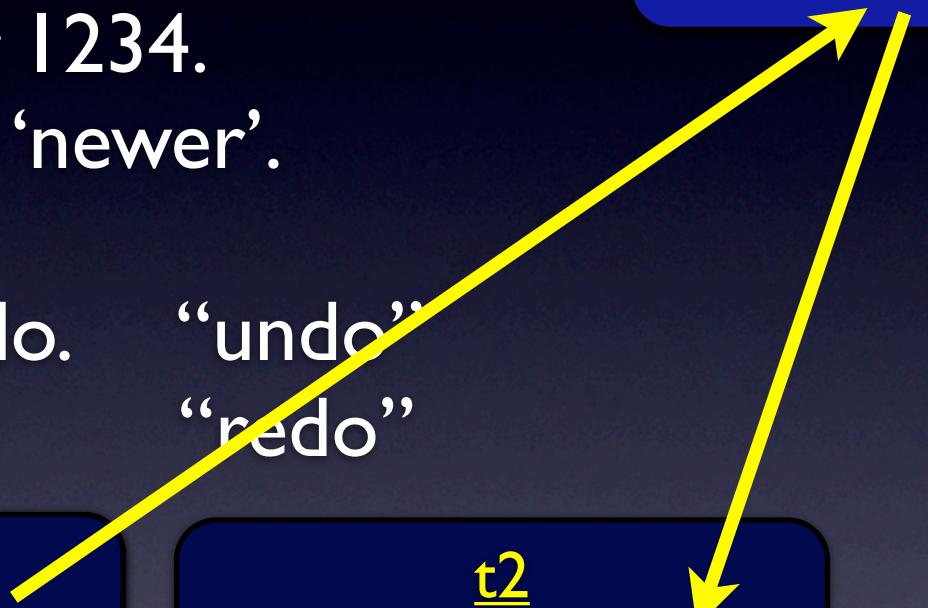
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t3

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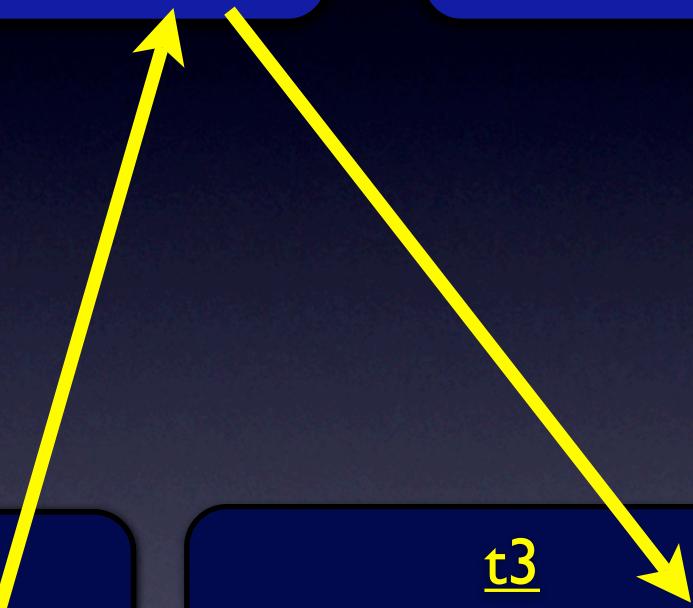
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t3

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t3
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t2
obj1's foo was 'newer'
obj2's bar was 1234

obj1
foo is 'newer'
obj2
bar is 5

t3
obj1's foo was 'old'

```
t1 := UTransaction eval: [  
    obj1 foo:'new'.  
    obj2 bar: 1234.  
    obj1 foo:'newer'.  
].
```

```
t2 := t1 undo.  "undo"  
→ t2 undo.      "redo"
```

t1
obj1's foo was 'old'
obj2's bar was 5

t2
obj1's foo was 'newer'
obj2's bar was 1234

t3
obj1's foo was 'old'
obj2's bar was 5

obj1
foo is 'newer'

obj2
bar is 5



```
t1 := UTransaction eval: [  
    obj1 foo: 'new'.  
    obj2 bar: 1234.  
    obj1 foo: 'newer'.  
].
```

```
t2 := t1 undo.  "undo"  
→ t2 undo.      "redo"
```

t1
obj1's foo was 'old'
obj2's bar was 5

t2
obj1's foo was 'newer'
obj2's bar was 1234

t3
obj1's foo was 'old'
obj2's bar was 5

obj1
foo is 'newer'

obj2
bar is 1234

```
t1 := UTransaction eval: [  
    obj1 foo:'new'.  
    obj2 bar: 1234.  
    obj1 foo:'newer'.  
].
```

obj1
foo is 'newer'

obj2
bar is 1234

```
t2 := t1 undo.  "undo"  
→ t2 undo.      "redo"
```

t1
obj1's foo was 'old'
obj2's bar was 5

t2
obj1's foo was 'newer'
obj2's bar was 1234

t3
obj1's foo was 'old'
obj2's bar was 5

```
t1 := UTransaction eval: [  
    obj1 foo:'new'.  
    obj2 bar: 1234.  
    obj1 foo:'newer'.  
].
```

```
t2 := t1 undo.  "undo"  
t2 undo.        "redo"
```

t1
obj1's foo was 'old'
obj2's bar was 5

t2
obj1's foo was 'newer'
obj2's bar was 1234

obj1
foo is 'newer'

obj2
bar is 1234

Trouble w/ Linear Undo

action_{i+3}



action_{i+2}

action_{i+1}

action_i

action_{i-1}

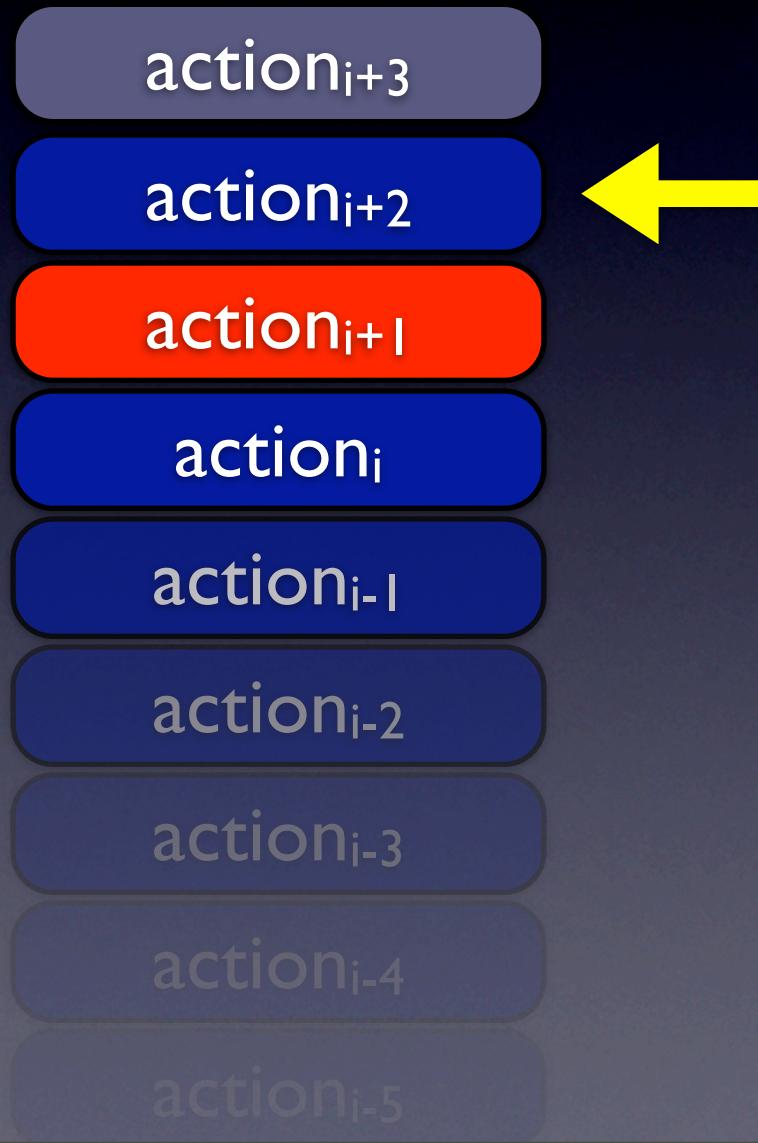
action_{i-2}

action_{i-3}

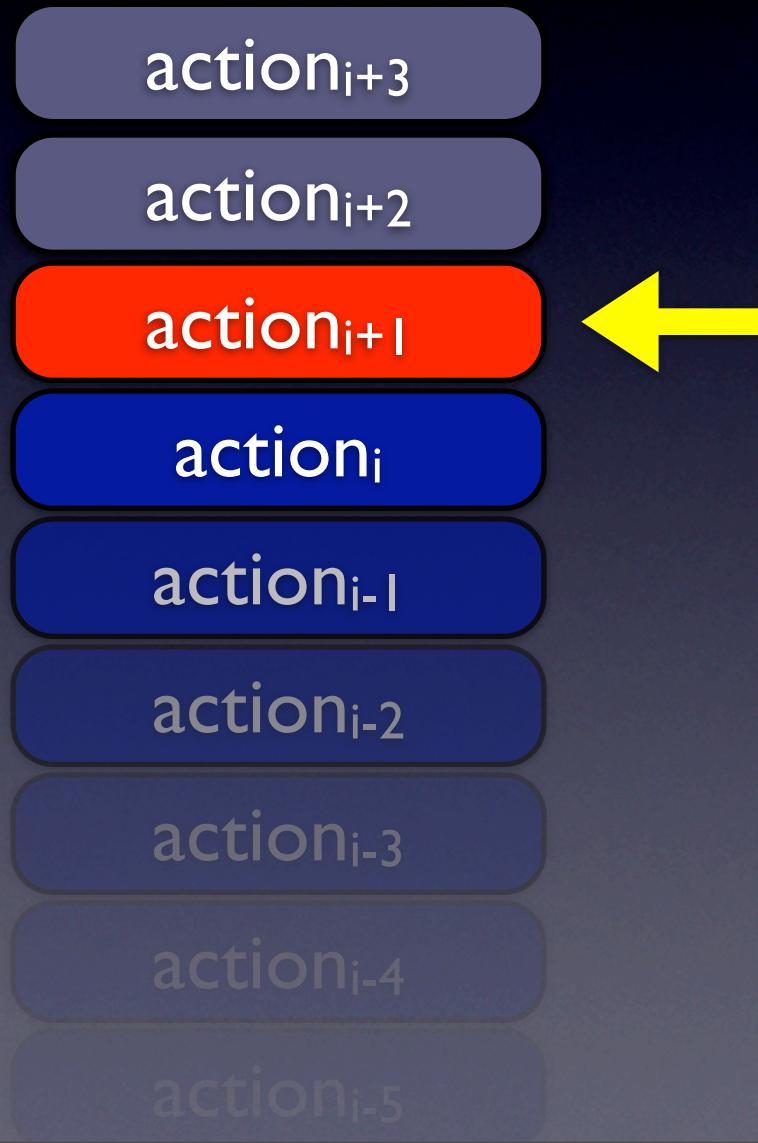
action_{i-4}

action_{i-5}

Trouble w/ Linear Undo



Trouble w/ Linear Undo



Trouble w/ Linear Undo

action_{i+3}

action_{i+2}

action_{i+1}

action_i

action_{i-1}

action_{i-2}

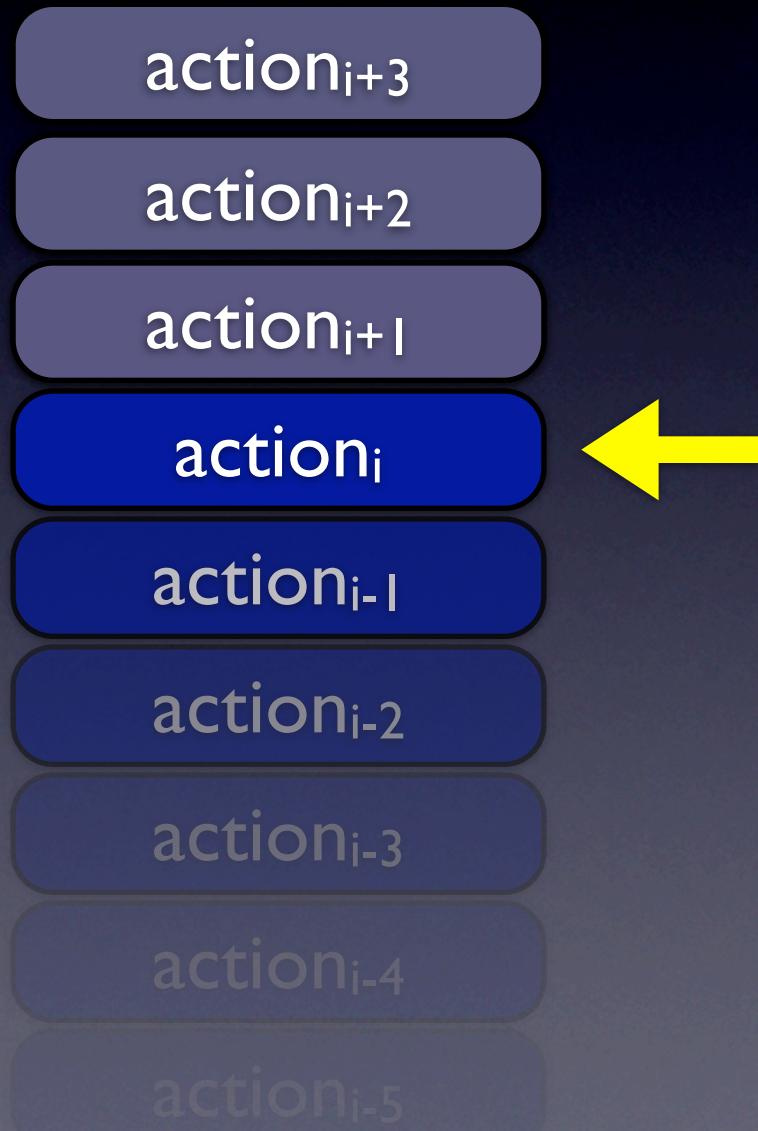
action_{i-3}

action_{i-4}

action_{i-5}



Trouble w/ Linear Undo



Problem:
can't redo
action_{i+2} and
action_{i+3} w/o
redoing
action_{i+1}

Want Selective Undo

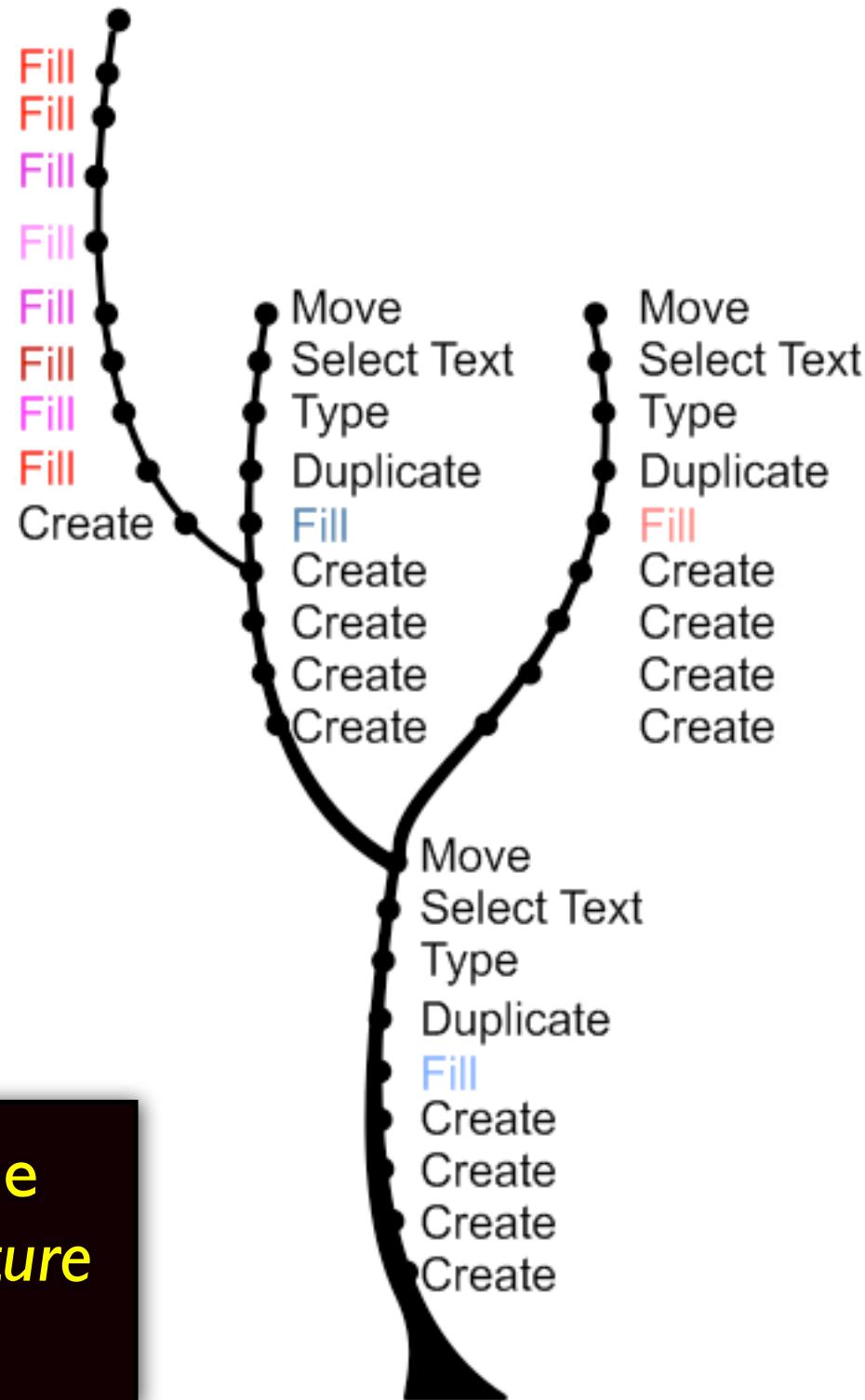
- Undo a command without first undoing commands that were issued afterwards
- **BUT** some commands are based on effects of earlier commands
 - gets tricky!

Selective Undo (Sort Of)

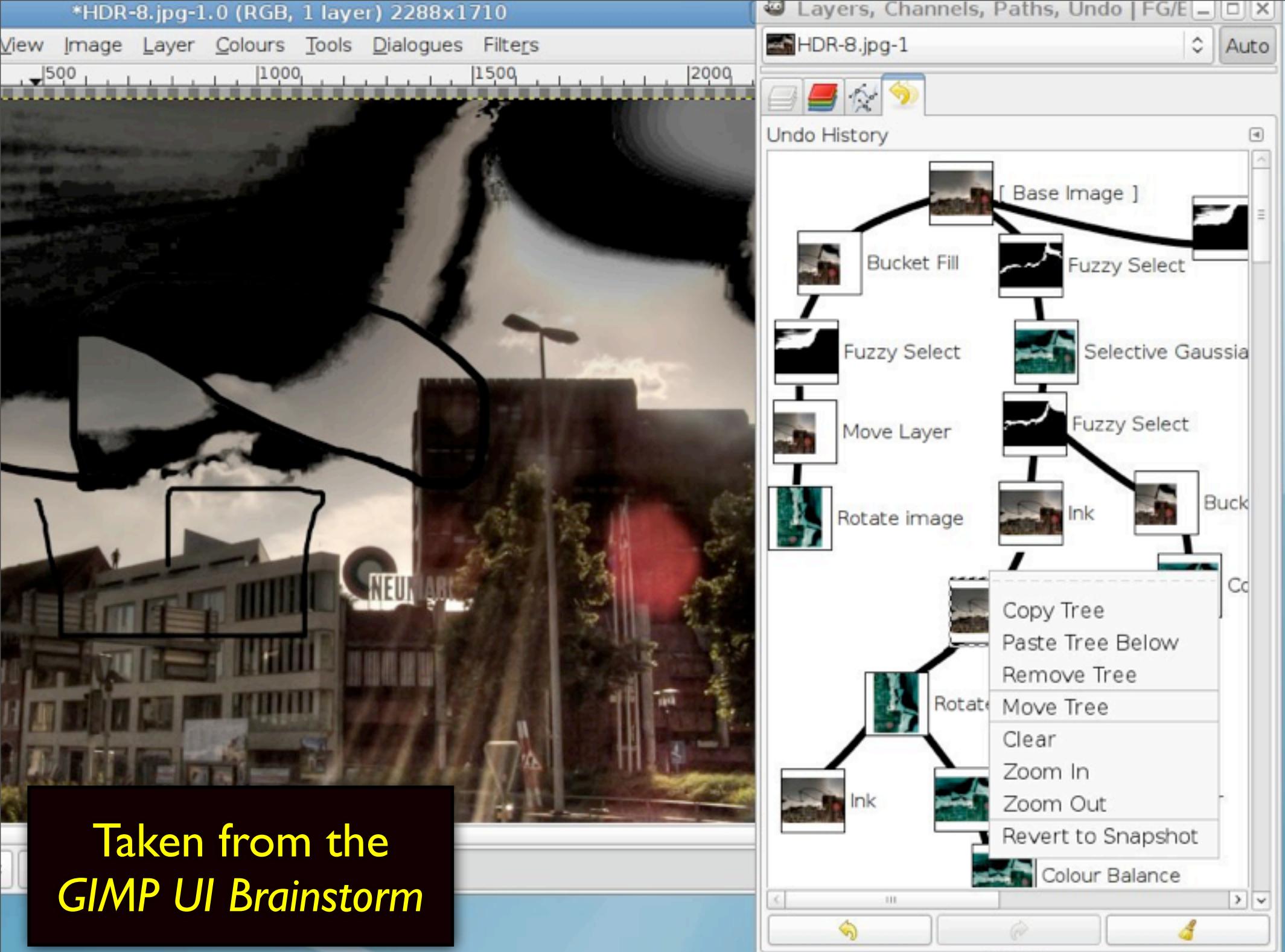
- Can undo action_{i+1} directly (no stack required)
- ... but UTransaction's undo is *transitive*
 - *Undoing a transaction t will transitively undo all later transactions that modified one or more objects modified by t*
- Still stack-like, but *only related operations are undo'ed transitively*
 - A kind of “selective undo” that makes sense

(Too Big a Hammer?)

- It may be!
- ... **but we could write the program so that different aspects of an object are stored in different “sub-objects”**
 - keeps mechanism easy to understand
- **Another option:** take into account what properties of what objects were modified
- No clear winner yet



Taken from the
CorelDRAW Feature
Request Site



Part II

Worlds:

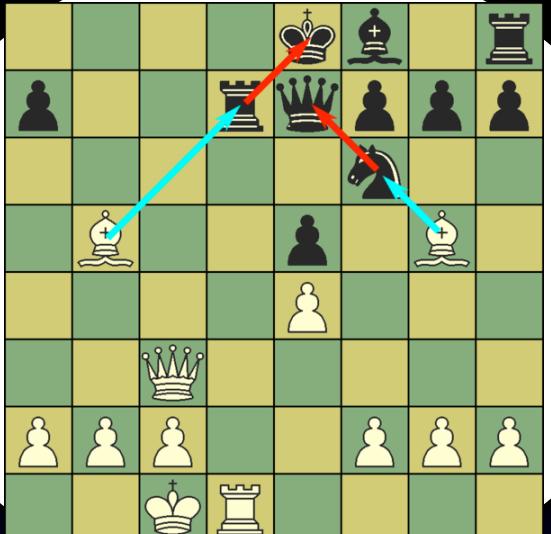
Undo for Programs





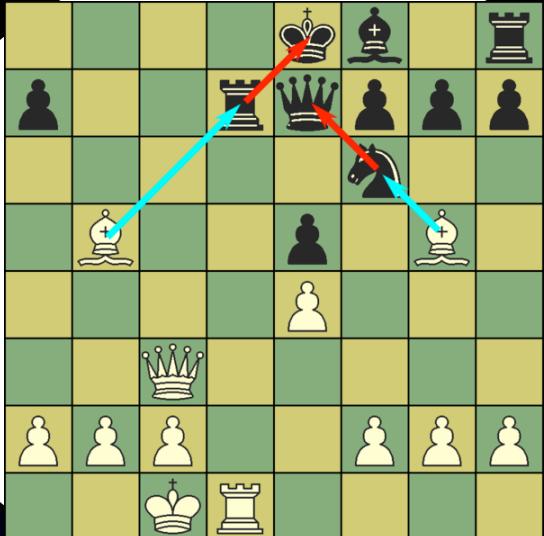
what if...
I take his knight with
my bishop?





what if...
I take his knight with
my bishop?





what if...
I take his knight with
my bishop?



Part II

Worlds:
Undo for Programs

Part II

What
if...? *Worlds:*
 for Programs

I'm Talking About...

- Programming language support for
 - “thought experiments”, a.k.a.,
 - “possible worlds reasoning”
- How? By enabling programmers to **control the scope of side effects**.

About Side Effects

- Not *all* side effects!
- Only **changes to the program store**, e.g.,
 - global, local, instance, and class variables
 - arrays
 - ...

Worlds

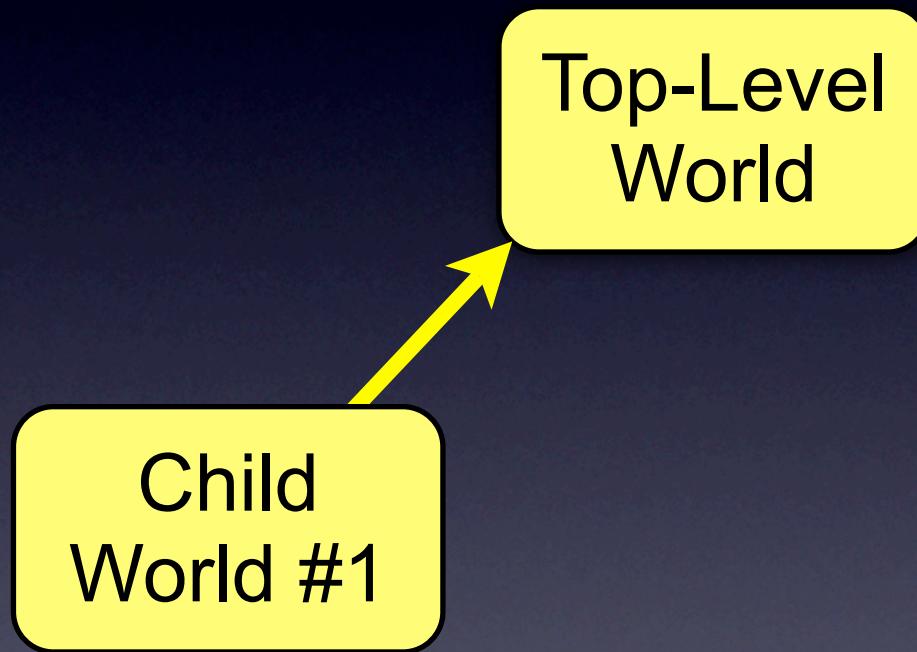
- A simple and expressive model for controlling the scope of side effects
- Worlds: new kind of *first-class store*
 - allows multiple versions of the program store to co-exist
 - organized hierarchically
 - Worlds/Squeak and Worlds/JS

The Programming Model

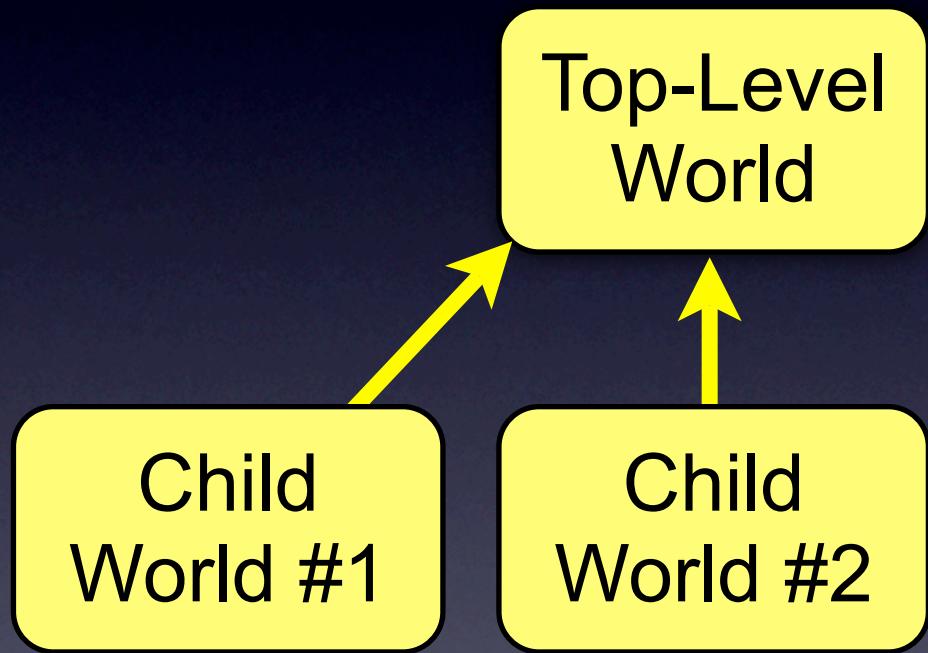
The Programming Model

Top-Level
World

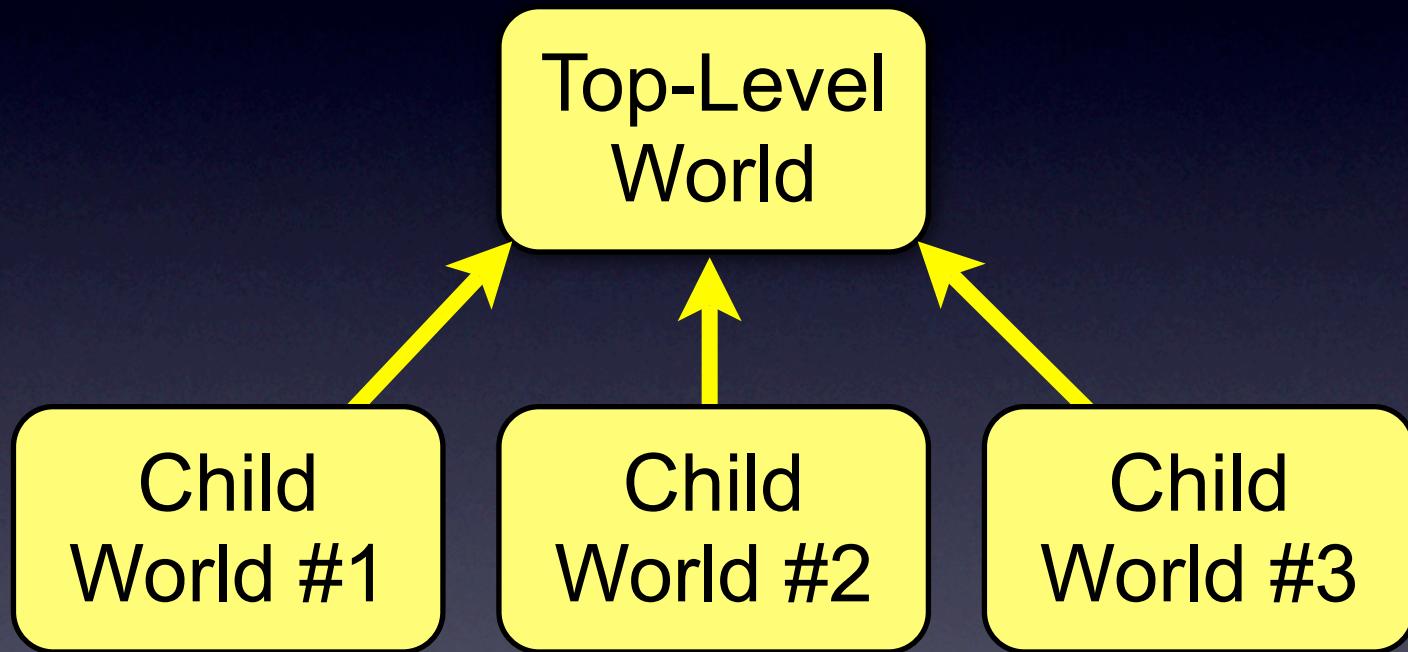
The Programming Model



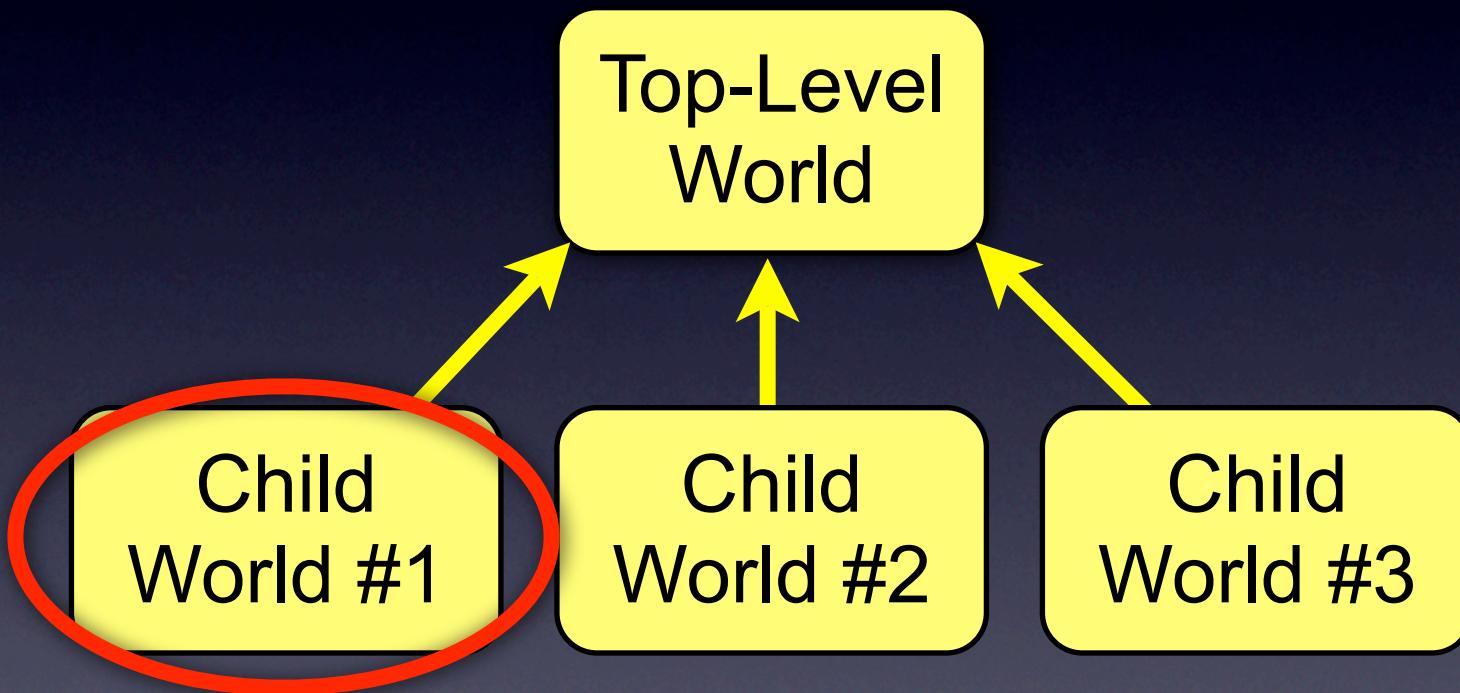
The Programming Model



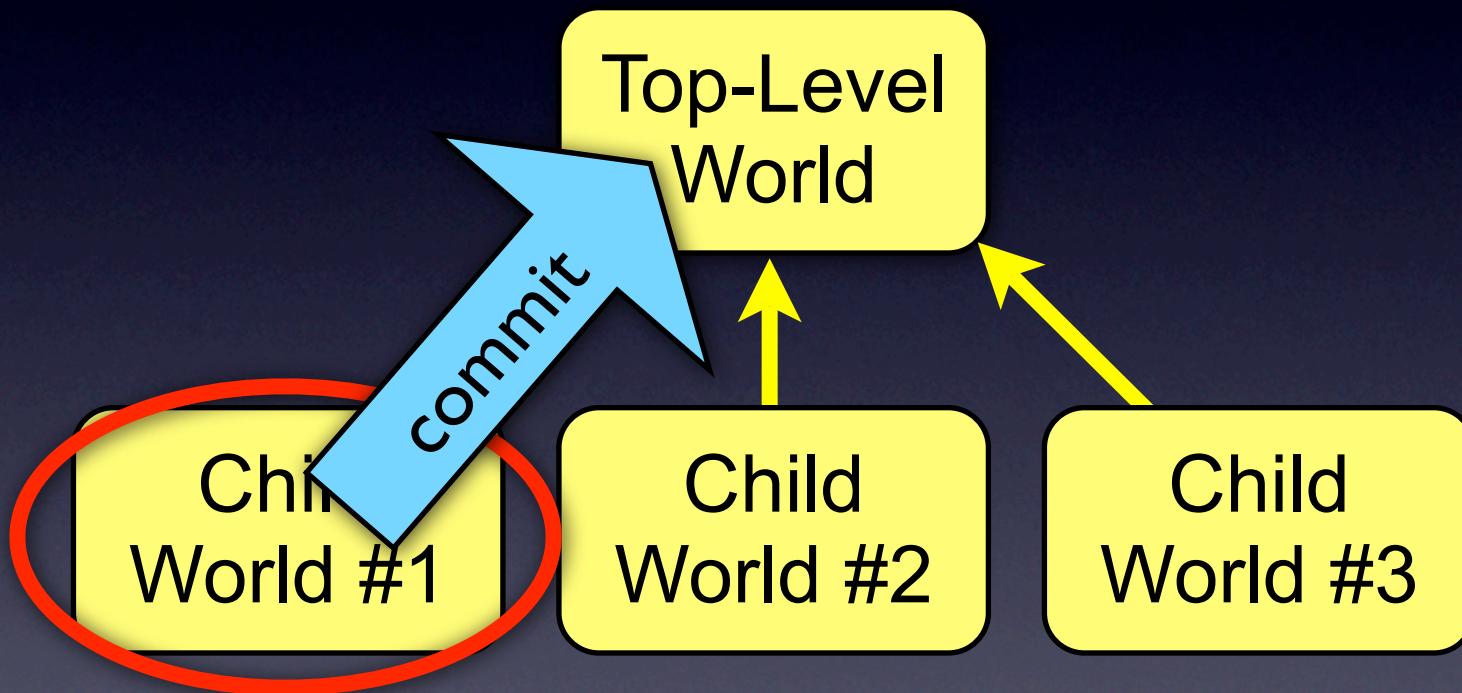
The Programming Model



The Programming Model

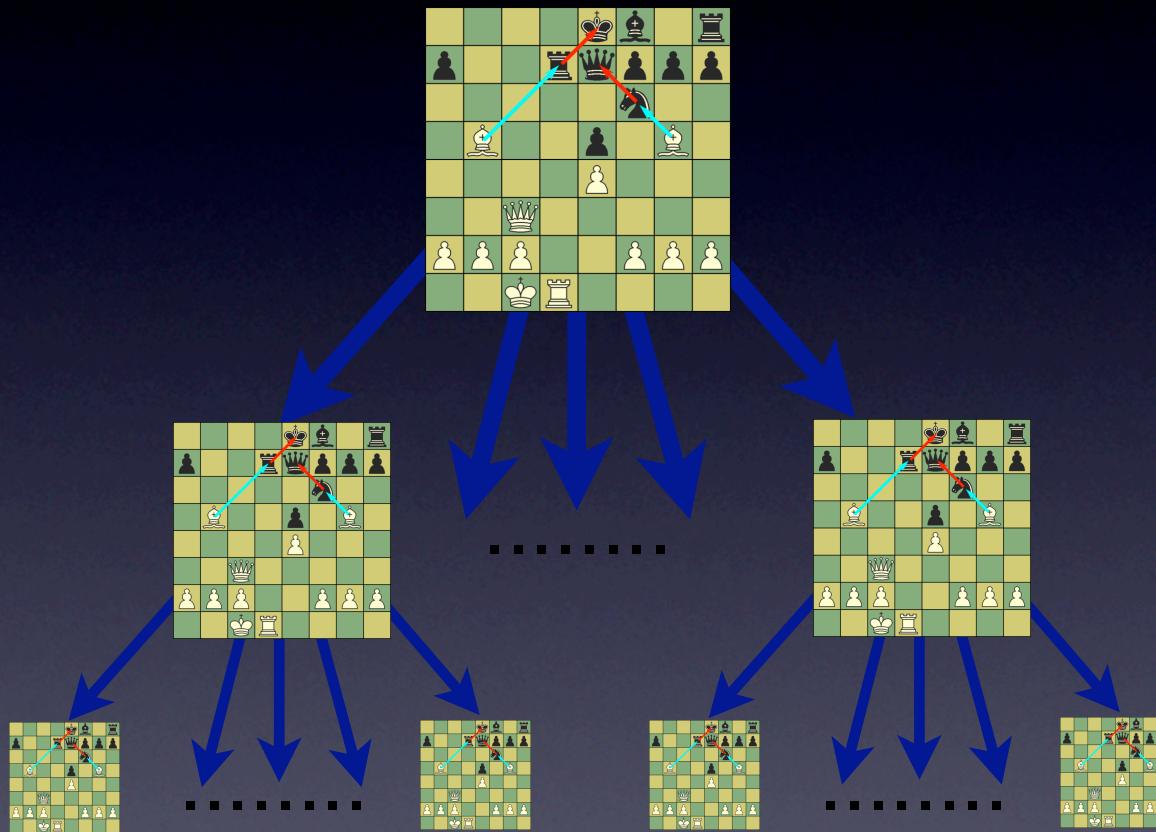


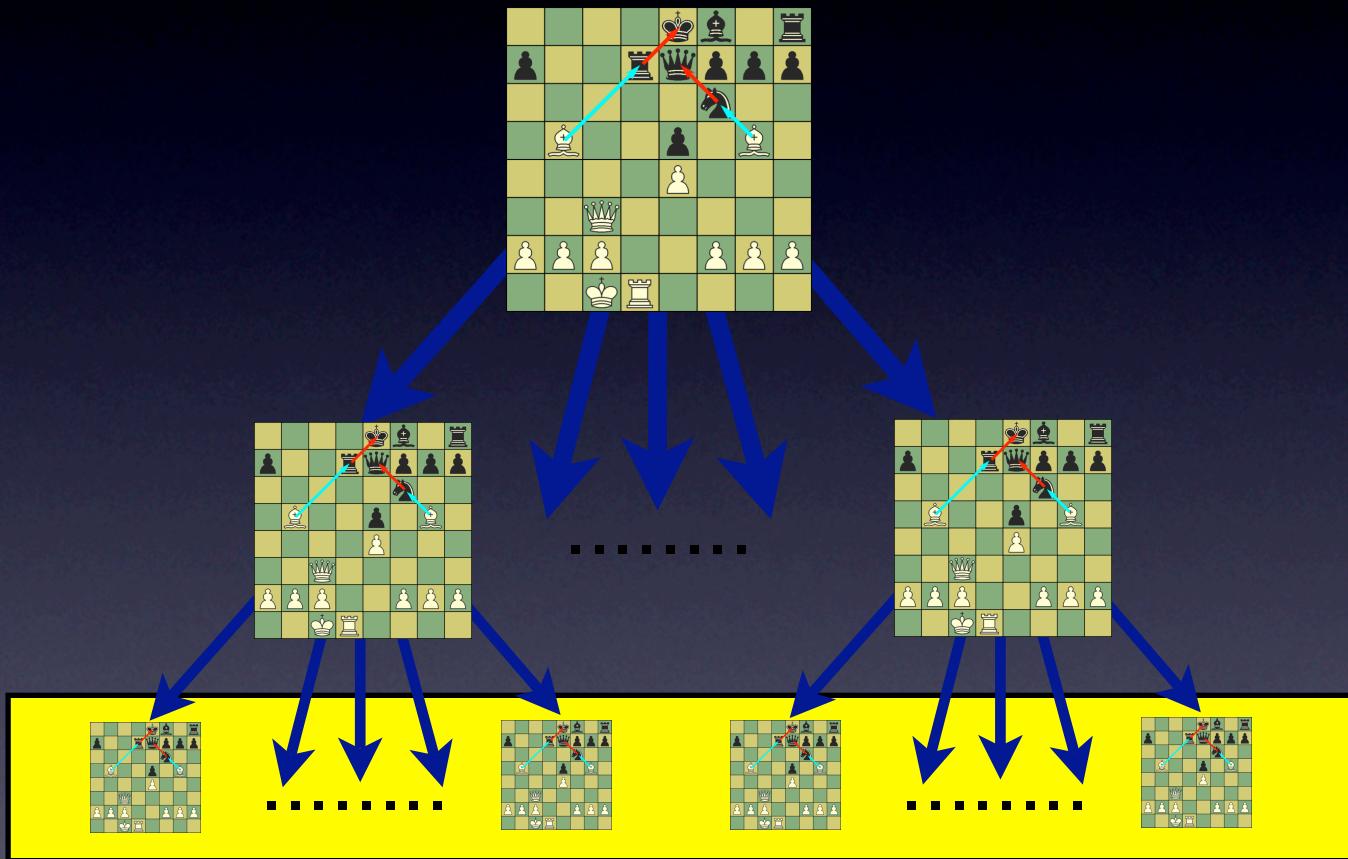
The Programming Model

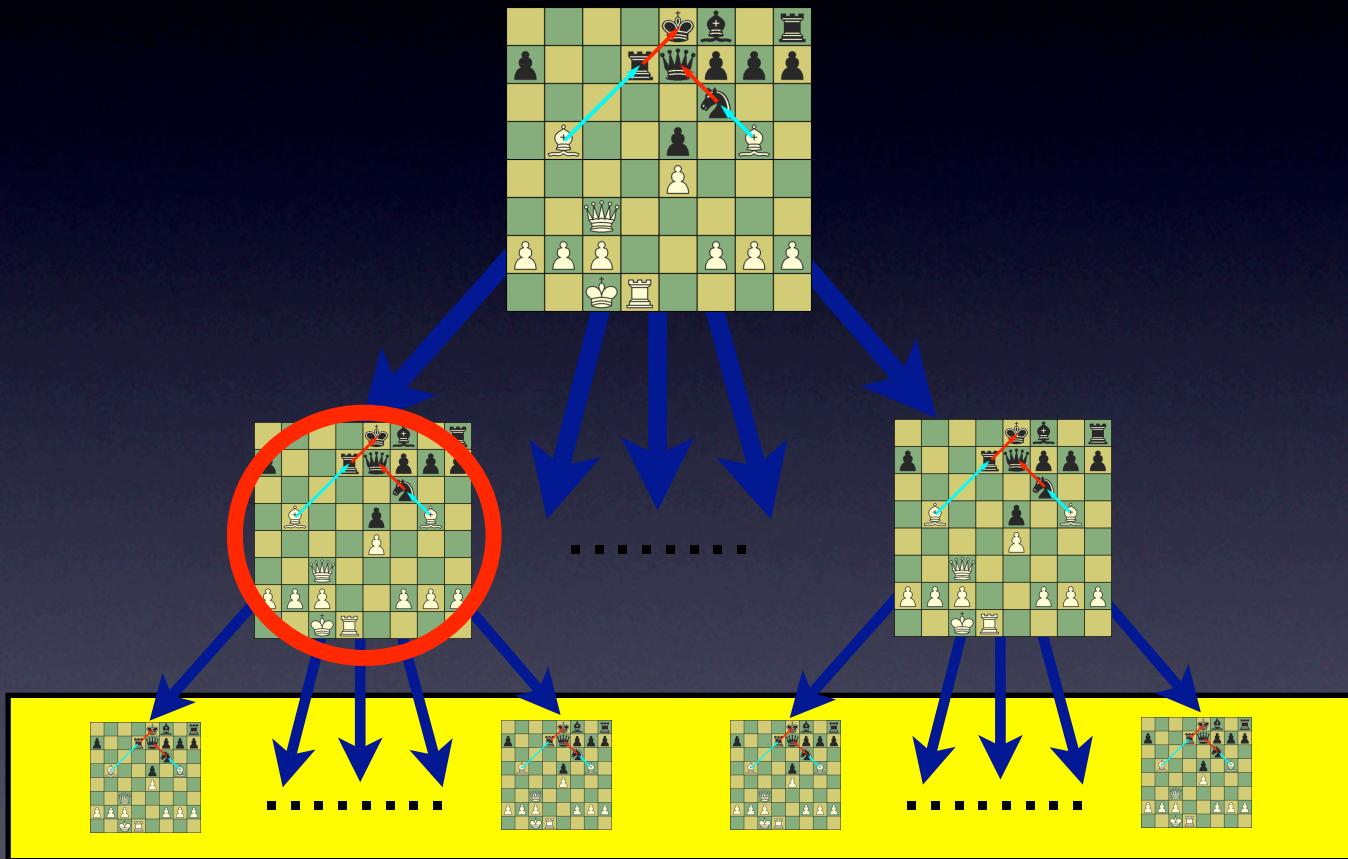


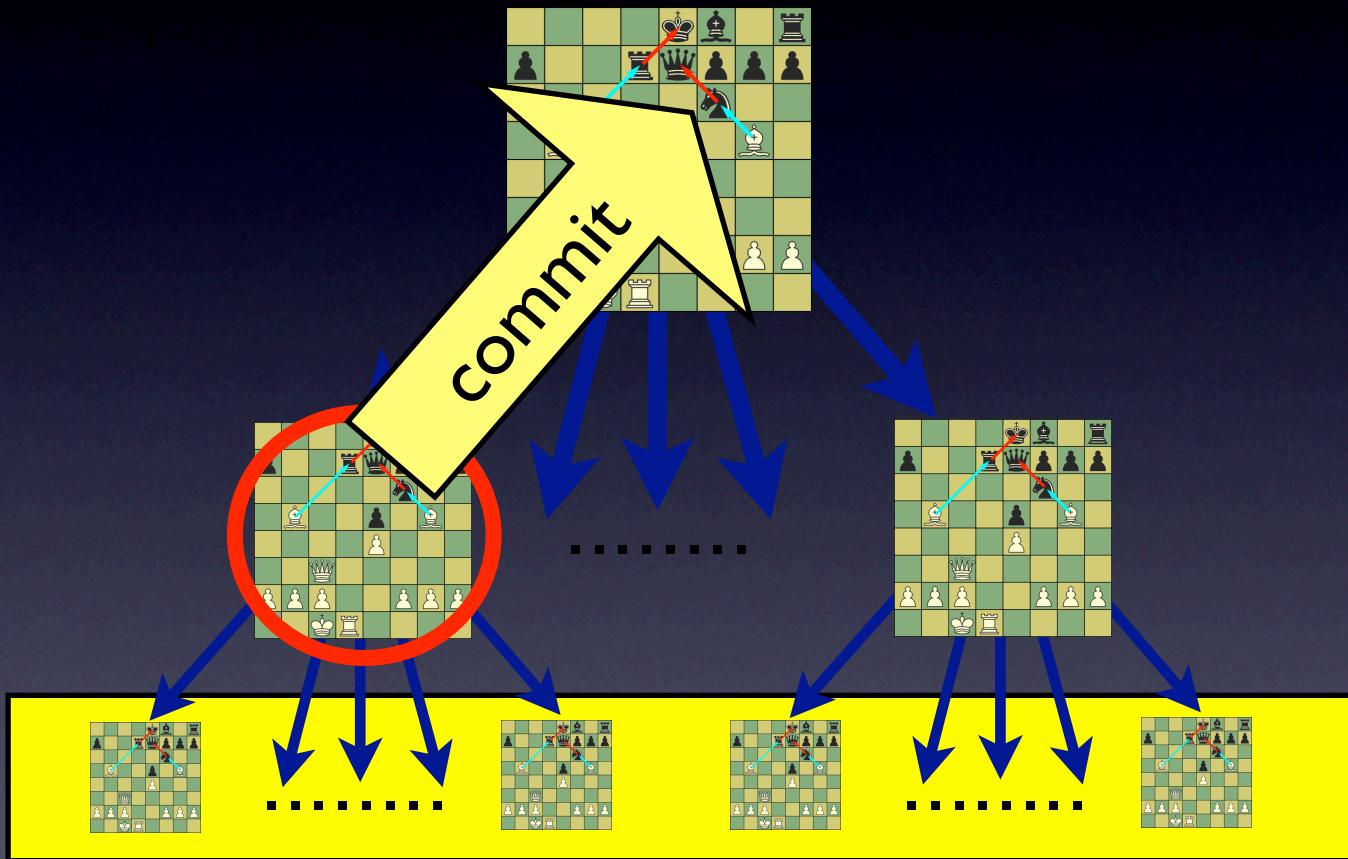














Worlds/Squeak

- `thisWorld`
- `w sprout`
- `w commit`
- `w eval: [...]`

Exception Handling

```
[  
  xs do: [:x |  
    x update  
  ]  
] on: Exception do: [  
  ...  
]
```

Exception Handling

```
[  
  ← save state of  
  xs do: [:x |  
    x update  
  ]  
 ] on: Exception do: [  
  ...  
 ]
```

Exception Handling

```
[           save state of  
  xs do: [:x |  
    x update  
]  
] on: Exception do: [  
  ...  
]
```

Diagram illustrating the flow of exception handling:

- A yellow arrow points from the opening bracket [at the top left towards the xs do: block.
- A yellow arrow points from the closing bracket] at the bottom left towards the ellipsis ... in the on: Exception do: block.
- The text "save state of collection's elements" is positioned to the right of the first arrow.
- The text "restore state of collection's elements" is positioned to the right of the second arrow.

Exception Handling

```
[  
  xs do: [:x |  
    x update  
  ]  
] on: Exception do: [  
  ...  
]
```

Exception Handling

```
[
```

```
thisWorld sprout eval: [
```

```
  xs do: [:x |
```

```
    x update
```

```
].
```

```
thisWorld commit
```

```
]
```

```
] on: Exception do: [
```

```
]
```

Exception Handling

```
[
```

```
thisWorld sprout eval: [
```

```
  xs do: [:x |  
    x update
```

```
].
```

```
thisWorld commit
```

```
]
```

```
] on: Exception do: [  
  
]
```

Sandboxing

```
sandbox = thisWorld.sprout();
in sandbox {
    eval(untrustedCode);
}
```

Sandboxing

```
disableDangerousStuff = function() {  
    alert = null;  
    Object.prototype.forbiddenMethod = null;  
    ...  
}
```

```
sandbox = thisWorld.sprout();  
in sandbox {  
    disableDangerousStuff();  
    eval(untrustedCode);  
}
```

Extension Methods in JS

```
Number.prototype.fact = function() {  
    if (this == 0)  
        return 1;  
    else  
        return this * (this - 1).fact();  
};  
  
print(5.fact());
```

[^] Extension Methods in JS *scoped*

```
myModule = thisWorld.sprout();
in myModule {
    Number.prototype.fact = function() { ... };
}

in myModule {
    print(5.fact());
}
```

Back to OMeta

rhyme = fee fie foe fum
I fiddle dee dee

Back to OMeta

```
rhyme = fee fie foe fum  
        | fiddle dee dee
```

```
fee    = ... -> ...  
fie    = ... -> ...  
foe    = ... -> ...  
fum    = ... -> ...
```

Back to OMeta

```
rhyme = fee fie foe fum  
        | fiddle dee dee
```

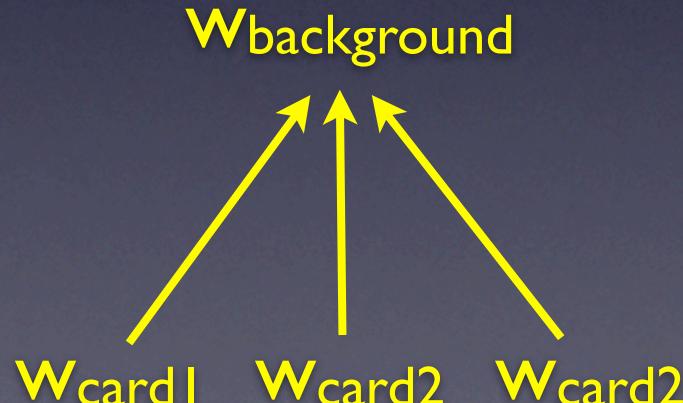
fee	=	...	->	...
fie	=	...	->	...
foe	=	...	->	...
fum	=	...	->	...

Case Study I

- Variant of OMeta/JS in which backtracking rolls back the side effects of rules' semantic actions
- OMeta implemented in JS, and Worlds/JS is a proper superset of JS
 - Re-implemented OR, kleene-*^{*}, etc. using worlds
 - very, very difficult to do w/o (something like) Worlds

Case Study II

- Hypercard-like system implemented w/ Worlds (w/ Ted Kaehler and Yoshiki Ohshima)
- All backgrounds and cards in a stack are really *just one card*, viewed through different worlds



$W_{background}$ contains the default state of the card, which is shared by all cards



W_{cardi} overrides the state of the card, as it appears in $W_{background}$

Future Work

Future Work

- Invariants!
 - register inter- *and* intra-object inv's dynamically
 - modify objects in transactions
 - all relevant invariants checked at end of transaction
 - only commit transaction if all inv's hold



A group of three people are gathered around a computer monitor, examining financial data. The monitor displays a spreadsheet titled "Spreadsheet: Q1 Budget-2.xls". A yellow sticky note overlaid on the screen reads "Budget Updates due Friday". The spreadsheet shows various budget items and analysis, including a section on "ROI Analysis" and "Budget Justification". To the left of the monitor, a web browser window is open, showing a chart titled "Economic Trends" for the "United States". The chart displays employment percentages from 1990 to 2000.

Future Work (cont'd)

- Mechanisms for synchronizing distributed, decentralized systems like
 - TeaTime [Reed '78]
 - Virtual Time / Time Warp [Jefferson '85]
- ... rely on support for speculative execution
- (May be able to do even better w/ Worlds)

Future Work (cont'd)

- Worlds: a model for programming multi-core architectures?
 - e.g., choosing among optimizations
 - will need efficient, HW-based impl.

Part III

Churrasco!

Worlds vs. UObjects

support for spec. execution,
possible worlds reasoning

transitive undo

very general, b/c it
works on every object in
the system

only affects objects
designed to work with it

very dangerous, b/c it
works on every object in
the system

only affects objects
designed to work with it

pro +

con -

Questions?

For more info...
<http://tinlizzie.org/~awarth>