CSC 225 - Assignment 3 Analysis Report Alex Holland

public static int CountInversions(int[] A){

```
int invCount = 0;
                                         //Track the number of inversions in the array.
                                         //Temporary array value holder.
 int temp;
 for (int i = 0; i < A.length - 1; i++){
                                         //1. Loop while i is one less than the size of the array.
    if (A[i] > A[i + 1]) {
                                         //2. Determine if element 1 and element2 should be swapped.
                                         //3. Hold the element 1 value in temp.
       temp = A[i];
       A[i] = A[i + 1];
                                         //4. Set element 1 to be element 2.
       A[i + 1] = temp;
                                         //5. Set element 2 to be temp which holds element 1.
       invCount++;
                                         //6. Increase the inversion count.
       if (i >= 1){
                                         //7. If two elements are swapped go back two indexes and
         i -= 2;
                                         // determine If the new index and next element should also be
                                         // swapped.
       }
    }
 return invCount;
                                         //Return the number of inversion in the array.
}
```

The loop executes A.length - 2 times.

The implemented algorithm is O(n + k) on an array with n elements and k inversions, resulting in a O(n) algorithm when $k \in O(n)$.

Counting operations of worst case time running time T(n)

Primitive Operations:

- Assignments (A)
- Comparisons (C)
- Array indexing (I)
- Add, subtract (S)

CountInversions(int[]A):

Input: An array of elements.

Output: The number of counted inversions.

```
invCount \leftarrow 0
                                     1A
temp \leftarrow 0
                                     1A
for i \leftarrow 0 to A. length -2 do
                                     1A + (n - 1)(1C + 1A + 1S) + 1C
  if A[i] > A[i + 1] then
                                     (n-1)(1C+2I+1S)
    swap(i, i + 1)
                                     (n-1)(4A + 4I + 3S)
  end
                                     (n-1)(1C)
  if i \geq 1
    i \leftarrow i - 2
                                     (n-1)(1A+1S)
  end
end
```

return invCount

$$(n-1)(1A)$$

$$T(n) = 1 + 1 + 1 + 3(n - 1) + 1 + 3(n - 1) + 11(n - 1) + (n - 1) + 2(n - 1) + 1(n - 1)$$

 $T(n) = 4 + 3n - 3 + 3n - 3 + 11n - 11 + n - 1 + 2n - 2 + n - 1$
 $T(n) = 21n - 17$

We can see that $k \geq n$.