9f. Reductions

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INTRODUCTION

Reductions are a powerful technique for **operations that take collections as input and return a single element**. They naturally arise in a wide range of contexts, such as when computing summary statistics (e.g., averages, variances, or maxima of collections).

The reduction process works by iteratively applying an operation to pairs of elements, accumulating the results at each step until the final output is obtained. For example, to compute $\boxed{\text{sum}(x)}$ for a vector \boxed{x} , a reduction would start by adding the first two elements, then add the third element to the total, and continue this cumulative process until all elements have been summed.

The method is particularly convenient when we need to transform a vector's elements prior to aggregating the result. By operating on scalars, **reductions sidestep the need to materialize intermediate outputs**, thus reducing memory allocations. This means that, if for example you have to compute $\boxed{\text{sum}(\log .(x))}$, a reduction would avoid the creation of the intermediate vector $\boxed{\log .(x)}$.

INTUITION

Reductions are typically implemented using a for-loop, with an operator applied to pairs of elements and the resulting output updated in each iteration. A classic example is the summation of all numeric elements in a vector. This involves applying the addition operator + to pairs of elements, iteratively updating the accumulated sum. This process is demonstrated below.

```
x = rand(100)

foo(x) = sum(x)

julia> foo(x)
48.447
```

```
x = rand(100)

function foo(x)
    output = 0.

for i in eachindex(x)
        output = output + x[i]
    end

    return output
end

julia> foo(x)
48.447
```

```
x = rand(100)

function foo(x)
    output = 0.

for i in eachindex(x)
    output += x[i]
    end

    return output
end

julia> foo(x)
48.447
```

In reductions, it's common to see implementations like the last tab, which are based on <u>update operators</u>. Recall that they turn an expression like x = x + a into x + a.

IMPLEMENTING REDUCTIONS

Reductions are implemented by applying either a <u>binary operator</u> or a two-argument function during each iteration. An example of a binary operator is +, which we used above. However, we could've also used its two-argument function form, replacing output = output + \times [i] with output = +(output, \times [i]).

The use of two-argument functions expands the scope of reductions, enabling us to compute more complex operations. For instance, it allows us to compute the maximum value of a vector x. This requires the function x, where x where x and x are x and x and x and x are x are x and x are x are x are x and x are x and x are x and x are x are x are x are x are x and x are x are x and x are x are x and x are x are x are x and x are x are x are x are x are x and x are x are x are x are x are x and x are x are x and x are x are x and x are x are x are x and x are x and x are x are x are x are x are x and x are x are x are x and x are x are x are x and x are x and x are x and x are x are x are x and x are x are x are x and x are x are x and x are x are x and x are x are x are x and x are x are x and x are x are

Formally, a reduction requires the binary operation to satisfy **two mathematical properties**:

• **Associativity**: the way in which operations are grouped must not change the result. For example, addition is associative because (a + b) + c = a + (b + c).

• **Existence of an identity element**: there exists an element that, when combined with any other element through a binary operation, leaves that element unchanged. For example, the identity element of addition is 0 because a + 0 = a.

The following list indicates the identity elements of each operation.

•	,	
Sum	0	

Operation Identity Element

Product 1

Maximum -Inf

Minimum Inf

The relevance of identity elements lies in that they constitute the initial values of the iterative process. Based on these identity elements, we next implement reductions for several operations. The examples make use of the function fool to show the desired outcome, while fool provides the same output via a reduction.

```
x = rand(100)

foo1(x) = sum(x)

function foo2(x)
    output = 0.

for i in eachindex(x)
    output += x[i]
    end

    return output
end
```

```
x = rand(100)

fool(x) = prod(x)

function foo2(x)
    output = 1.

for i in eachindex(x)
    output *= x[i]
    end

    return output
end
```

```
x = rand(100)
foo1(x) = maximum(x)

function foo2(x)
   output = -Inf

for i in eachindex(x)
    output = max(output, x[i])
  end

return output
end
```

```
x = rand(100)

foo1(x) = minimum(x)

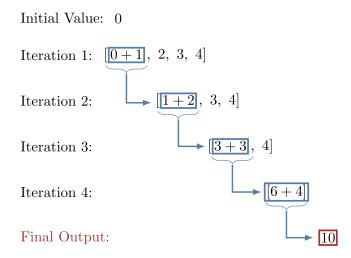
function foo2(x)
  output = Inf

for i in eachindex(x)
     output = min(output, x[i])
  end

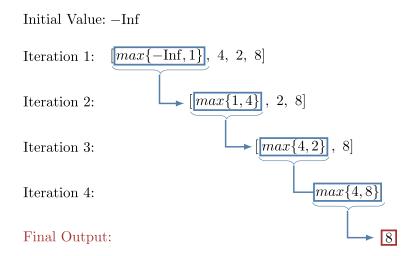
return output
end
```

We can also visually illustrate how reductions operate in these examples, as we do below.

REDUCTION 1: sum of [1,2,3,4]



REDUCTION 2: maximum of [1,4,2,8]



AVOIDING MEMORY ALLOCATIONS VIA REDUCTIONS

One of the primary advantages of reductions is that they avoid the memory allocation of intermediate results.

To illustrate this, consider the operation $\boxed{\text{sum}(\log_{\cdot}(x))}$ for a vector \boxed{x} . Its computation involves transforming \boxed{x} into $\boxed{\log_{\cdot}(x)}$, and then summing the transformed elements. By default, broadcasting creates a new vector to store the result of $\boxed{\log_{\cdot}(x)}$, thereby allocating memory. However, in many cases like this one, we're only interested in the final sum, not the intermediate result. Therefore, an approach that bypasses the allocation of $\boxed{\log_{\cdot}(x)}$ is beneficial. Reductions accomplish this by defining a scalar $\boxed{\text{output}}$, which is iteratively updated by summing the transformed values of \boxed{x} .

```
x = rand(100)

foo1(x) = sum(log.(x))

function foo2(x)
    output = 0.

    for i in eachindex(x)
        output += log(x[i])
    end

    return output
end

julia> @btime foo1($x)
    315.584 ns (1 allocation: 896 bytes)
    julia> @btime foo2($x)
    296.119 ns (0 allocations: 0 bytes)
```

```
x = rand(100)

foo1(x) = prod(log.(x))

function foo2(x)
    output = 1.

    for i in eachindex(x)
        output *= log(x[i])
    end

    return output
end

julia> @btime foo1($x)
    311.840 ns (1 allocation: 896 bytes)
julia> @btime foo2($x)
    296.061 ns (0 allocations: 0 bytes)
```

```
x = rand(100)
foo1(x) = maximum(log.(x))
function foo2(x)
    output = -Inf

for i in eachindex(x)
    output = max(output, log(x[i]))
    end

    return output
end

julia> @btime foo1($x)
    482.602 ns (1 allocation: 896 bytes)
julia> @btime foo2($x)
    374.961 ns (0 allocations: 0 bytes)
```

```
x = rand(100)

foo1(x) = minimum(log.(x))

function foo2(x)
    output = Inf

for i in eachindex(x)
    output = min(output, log(x[i]))
    end

    return output
end

julia> @btime foo1($x)
    487.156 ns (1 allocation: 896 bytes)

julia> @btime foo2($x)
    368.502 ns (0 allocations: 0 bytes)
```

REDUCTIONS VIA BUILT-IN FUNCTIONS

The previous examples implemented reductions through explicit for-loops. Unfortunately, this approach can compromise readability due to the verbosity of for-loops. To address the issue, Julia offers several streamlined alternatives to implement reductions.

In particular, several common operations come with additional methods that transform a vector and then apply a reduction. This is the case of functions like [sum], [prod], [maximum], and [minimum]. Their syntax is given by [foo(<transforming function>, x)], where [foo] is one of the functions mentioned and [x] is the vector to be transformed. For instance, the following examples consider reductions for the transformed vector $[2 \cdot x]$.

```
x = rand(100)

foo(x) = prod(log, x) #same output as prod(log.(x))

julia> @btime foo($x)
    294.763 ns (0 allocations: 0 bytes)
```

```
x = rand(100)

foo(x) = maximum(log, x) #same output as maximum(log.(x))

julia> @btime foo($x)

579.940 ns (0 allocations: 0 bytes)
```

```
x = rand(100)

foo(x) = minimum(log, x) #same output as minimum(log.(x))

julia> @btime foo($x)

577.516 ns (0 allocations: 0 bytes)
```

To keep matters simple, we've used the built-in function \log for transforming x. However, the approach supports any form of function, including anonymous ones.

```
x = rand(100)

foo(x) = sum(a -> 2 * a, x)  #same output as sum(2 .* x)

julia> @btime foo($x)
  6.493 ns (0 allocations: 0 bytes)
```

```
x = rand(100)

foo(x) = prod(a -> 2 * a, x)  #same output as prod(2 .* x)

julia> @btime foo($x)
  6.741 ns (0 allocations: 0 bytes)
```

Finally, all these functions accept transforming functions that require multiple arguments. To incorporate this possibility, it's necessary to call the multiple variables using $\boxed{\mathtt{zip}}$, and referring to each variable through indexes. We illustrate this below, where the transforming function is $\boxed{\mathtt{x} \cdot \mathtt{*} \cdot \mathtt{y}}$.

```
x = rand(100); y = rand(100)

foo(x,y) = sum(a -> a[1] * a[2], zip(x,y)) #same output as sum(x .* y)

julia> @btime foo($x)
    29.127 ns (0 allocations: 0 bytes)
```

```
x = rand(100); y = rand(100)

foo(x,y) = prod(a -> a[1] * a[2], zip(x,y)) #same output as prod(x .* y)

julia> @btime foo($x)

48.031 ns (0 allocations: 0 bytes)
```

THE "REDUCE" AND "MAPREDUCE" FUNCTIONS

Beyond the specific functions we've discussed, reductions can be applied to any operation that meets the requirements for a valid reduction. In Julia, this is implemented through the functions reduce and mapreduce, where reduce applies the reduction directly, while mapreduce transforms the collection's elements prior to reducing it.

It's worth remarking that reductions with <code>sum</code>, <code>prod</code>, <code>max</code>, and <code>min</code> should still be done via the dedicated functions. The reason is that they've been optimized for their respective tasks. In this context, our primary use case of <code>reduce</code> and <code>mapreduce</code> is for other types of reductions not covered or when packages provide their own implementations of these functions. ²

FUNCTION "REDUCE"

The function reduce uses the syntax reduce(<function>, x), where <function> is a two-argument function. The following example demonstrates its use.

```
x = rand(100)
foo(x) = reduce(+, x) #same output as sum(x)

julia> @btime foo($x)
6.168 ns (0 allocations: 0 bytes)
```

```
x = rand(100)

foo(x) = reduce(*, x)  #same output as prod(x)

julia> @btime foo($x)
6.176 ns (0 allocations: 0 bytes)
```

```
x = rand(100)
foo(x) = reduce(min, x) #same output as minimum(x)

julia> @btime foo($x)
    167.440 ns (0 allocations: 0 bytes)
```

Note all the examples presented could've been implemented as <u>we did previously</u>, where we directly applied <code>sum</code>, <code>prod</code>, <code>maximum</code> and <code>minimum</code>.

FUNCTION "MAPREDUCE"

The function mapreduce combines the functions map and reduce: before applying the reduction, mapreduce transforms vectors via the function map. Recall that map(foo,x) transforms each element of the collection x by applying foo element-wise. Thus, mapreduce(<transformation>, < reduction>, x) first transforms x's elements through map, and then applies a reduction to the resulting output.

To illustrate its use, we make use of a log transformation.

```
x = rand(100)

foo(x) = mapreduce(log, *, x) #same output as prod(log.(x))

julia> @btime foo($x)

294.618 ns (0 allocations: 0 bytes)
```

```
x = rand(100)

foo(x) = mapreduce(log, min, x) #same output as minimum(log.(x))

julia> @btime foo($x)

577.505 ns (0 allocations: 0 bytes)
```

Just like with reduce, note that the examples could've been implemented directly as we did previously, through the functions sum, prod, maximum, and minimum.

mapreduce can also be used with anonymous functions and transformations requiring multiple arguments. Below, we illustrate both, whose implementations are the same as with sum, prod, maximum, and minimum.

```
x = rand(100); y = rand(100)

foo(x,y) = mapreduce(a -> a[1] * a[2], *, zip(x,y)) #same output as prod(x .* y)

julia> @btime foo($x)

48.221 ns (0 allocations: 0 bytes)
```

```
x = rand(100); y = rand(100)

foo(x,y) = mapreduce(a -> a[1] * a[2], max, zip(x,y)) #same output as maximum(x .* y)

julia> @btime foo($x)

175.634 ns (0 allocations: 0 bytes)
```

REDUCE OR MAPREDUCE?

reduce can be considered as a special case of mapreduce, where the latter transforms x through the identity function, identity(x). Likewise, mapreduce(<transformation>, <operator>, x) produces the same result as reduce(<operator>, map(<transformation>, x)). However, mapreduce is more efficient, since it avoids the allocation of the transformed vector. This is demonstrated below, where we compute x through a reduction.

```
x = rand(100)

foo(x) = mapreduce(a -> 2 * a, +, x)

julia> @btime foo($x)
6.372 ns (0 allocations: 0 bytes)
```

FOOTNOTES

^{1.} In the section <u>Lazy Operations</u>, we'll explore an alternative with broadcasting that doesn't materialize intermediate results either.

^{2.} For instance, the package Folds provides a parallelized version of both map and mapreduce, enabling the utilization of all available CPU cores. Its syntax is identical to Julia's built-in functions.