Lecture 2b: optimal transport with entropic regularization

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https://www.math-econ-code.org/regularized-optimal-transport

Learning objectives

- ► Entropic regularization
- ► The log-sum-exp trick
- ► The Iterated Proportional Fitting Procedure (IPFP)

References for block 9

- ► [OTME], Ch. 7.3
- Peyré, Cuturi, Computational Optimal Transport, Ch. 4.

Entropic regularization of the optimal transport problem

► Consider the problem

$$\max_{\pi \in \mathcal{M}(p,q)} \sum_{ij} \pi_{ij} \Phi_{ij} - \sigma \sum_{ij} \pi_{ij} \ln \pi_{ij}$$

where $\sigma > 0$. The problem coincides with the optimal assignment problem when $\sigma = 0$. When $\sigma \to +\infty$, the solution to this problem approaches the independent coupling, $\pi_{ij} = p_i q_i$.

▶ Later on, we will provide microfoundations for this problem, and connect it with a number of important methods in economics (BLP, gravity model, Choo-Siow...). For now, let's just view this as an extension of the optimal transport problem.

Dual of the regularized problem

Let's compute the dual by the minimax approach. We have

$$\max_{\pi \geq 0} \min_{u,v} \sum_{ij} \pi_{ij} \left(\Phi_{ij} - u_i - v_j - \sigma \ln \pi_{ij} \right) + \sum_i u_i p_i + \sum_j v_j q_j$$

thus

$$\min_{u,v} \sum_{i} u_i p_i + \sum_{i} v_j q_j + \max_{\pi \geq 0} \sum_{ii} \pi_{ij} \left(\Phi_{ij} - u_i - v_j - \sigma \ln \pi_{ij} \right)$$

▶ By FOC in the inner problem, one has $\Phi_{ii} - u_i - v_i - \sigma \ln \pi_{ii} - \sigma = 0$, thus

$$\pi_{ij} = \exp\left(\frac{\Phi_{ij} - u_i - v_j - \sigma}{\sigma}\right)$$

and $\pi_{ij} \left(\Phi_{ij} - u_i - v_j - \sigma \ln \pi_{ij} \right) = \sigma \pi_{ij}$, thus the dual problem is

$$\min_{u,v} \sum_{i} u_i p_i + \sum_{i} v_j q_j + \sigma \sum_{i} \exp\left(\frac{\Phi_{ij} - u_i - v_j - \sigma}{\sigma}\right).$$

After replacing v_i by $v_i + \sigma$, the dual is

$$\min_{u,v} \sum_{i} u_{i} p_{i} + \sum_{j} v_{j} q_{j} + \sigma \sum_{j} \exp \left(\frac{\Phi_{ij} - u_{i} - v_{j}}{\sigma} \right) - \sigma. \tag{V1}$$
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Another expression of the dual

► Claim: the problem is equivalent to

$$\min_{u,v} \sum_{i} u_{i} p_{i} + \sum_{j} v_{j} q_{j} + \sigma \log \sum_{i,j} \exp \left(\frac{\Phi_{ij} - u_{i} - v_{j}}{\sigma} \right)$$
 (V2)

▶ Indeed, let us go back to the minimax expression

$$\min_{u,v} \sum_{i} u_i p_i + \sum_{j} v_j q_j + \max_{\pi \geq 0} \sum_{ij} \pi_{ij} \left(\Phi_{ij} - u_i - v_j - \sigma \ln \pi_{ij} \right)$$

we see that the solution π has automatically $\sum_{ij}\pi_{ij}=1$; thus we can incorporate the constraint into

$$\min_{u,v} \sum_{i} u_i p_i + \sum_{j} v_j q_j + \max_{\pi \geq 0: \sum_{ij} \pi_{ij} = 1} \sum_{ij} \pi_{ij} \left(\Phi_{ij} - u_i - v_j - \sigma \ln \pi_{ij} \right)$$

which yields (V2).

Another expression of the dual (ctd)

Expression (V2) is interesting because, taking any $\hat{\pi} \in M(p, q)$, (V2) reexpresses as

$$\max_{u,v} \sum_{ij} \hat{\pi}_{ij} \left(\frac{\Phi_{ij} - u_i - v_j}{\sigma} \right) - \log \sum_{ij} \exp \left(\frac{\Phi_{ij} - u_i - v_j}{\sigma} \right)$$

therefore if the parameter is $\theta = (u, v)$, observations are ij pairs, and the likelihood of ij is

$$\pi_{ij}^{ heta} = rac{\exp\left(rac{\Phi_{ij} - u_i - v_j}{\sigma}
ight)}{\sum_{ij} \exp\left(rac{\Phi_{ij} - u_i - v_j}{\sigma}
ight)}$$

▶ Hence, (V2) will coincide with the maximum likelihood in this model.

A third expression of the dual problem

Consider

$$\min_{u,v} \sum_{i} u_i p_i + \sum_{j} v_j q_j$$

$$s.t. \sum_{i,j} \exp\left(\frac{\Phi_{ij} - u_i - v_j}{\sigma}\right) = 1$$

▶ It is easy to see that the solutions of this problem coincide with (V2). Indeed, the Lagrange multiplier is forced to be one. In other words,

$$\min_{u,v} \sum_{i} u_{i} p_{i} + \sum_{j} v_{j} q_{j}$$

$$s.t. \sigma \log \sum_{i,j} \exp\left(\frac{\Phi_{ij} - u_{i} - v_{j}}{\sigma}\right) = 0$$
(V3)

Small-temperature limit and the log-sum-exp trick

▶ Recall that when $\sigma \rightarrow 0$, one has

$$\sigma \log \left(e^{a/\sigma} + e^{b/\sigma} \right) o \max \left(a, b \right)$$

▶ Indeed, letting $m = \max(a, b)$,

$$\sigma \log \left(e^{a/\sigma} + e^{b/\sigma} \right) = m + \sigma \log \left(\exp \left(\frac{a - m}{\sigma} \right) + \exp \left(\frac{b - m}{\sigma} \right) \right), \tag{1}$$

and the argument of the logarithm lies between 1 and 2.

▶ This simple remark is actually a useful numerical recipe called the log-sum-exp trick: when σ is small, using (1) to compute $\sigma \log \left(e^{a/\sigma} + e^{b/\sigma} \right)$ ensures the exponentials won't blow up.

A third expression of the dual problem

▶ Back to the third expression, with $\sigma \rightarrow 0$, one has

$$\min_{u,v} \sum_{i} u_{i}p_{i} + \sum_{j} v_{j}q_{j}$$

$$s.t. \max_{ij} (\Phi_{ij} - u_{i} - v_{j}) = 0$$
(V3)

► This is exactly equivalent with the classical Monge-Kantorovich expression

$$\min_{u,v} \sum_{i} u_i p_i + \sum_{j} v_j q_j$$

$$s.t. \Phi_{ij} - u_i - v_j \le 0$$
(V3)

Computation

- ▶ We can compute $\min F(x)$ by two methods:
 - ▶ Either by gradient descent: $x(t+1) = x_t \epsilon_t \nabla F(x_t)$. (Steepest descent has $\epsilon_t = 1/|\nabla F(x_t)|$.)
 - ▶ Or by coordinate descent: $x_i(t+1) = \arg\min_{x_i} F(x_i, x_{-i}(t))$.
- ▶ Why do these methods converge? Let's provide some justification. We will decrease x_t by εd_t , were d_t is normalized by $|d_t|_p := \left(\sum_{i=1}^n d_t^i\right)^{1/p} = 1$. At first order, we have

$$F(x_t - \epsilon d_t) = F(x_t) - \epsilon d_t^{\mathsf{T}} \nabla F(x_t) + O(\epsilon^1).$$

- ▶ We need to maximize $d_t^T \nabla F(x_t)$ over $|d_t|_p = 1$.
 - ▶ For p = 2, we get $d_t = \nabla F(x_t) / |\nabla F(x_t)|$
 - ► For p = 1, we get $d_t = sign(\partial F(x_t)/\partial x^i)$ if $|\partial F(x_t)/\partial x^i| = \max_j |\partial F(x_t)/\partial x^j|$, 0 otherwise.

► Here, gradient descent is

$$u_{i}\left(t+1
ight)=u_{i}\left(t
ight)-\epsilonrac{\partial F}{\partial u_{i}}\left(u\left(t
ight),v\left(t
ight)
ight), ext{ and}$$
 $v_{j}\left(t+1
ight)=v_{j}\left(t
ight)-\epsilonrac{\partial F}{\partial v_{j}}\left(u\left(t
ight),v\left(t
ight)
ight)$

while coordinate descent is

$$\frac{\partial F}{\partial u_{i}}\left(u_{i}\left(t+1\right),u_{-i}\left(t\right),v\left(t\right)\right)=0,\text{ and }\frac{\partial F}{\partial v_{i}}\left(u\left(t\right),v_{j}\left(t+1\right),v_{-j}\left(t\right)\right)=0.$$

Gradient descent

Gradient of objective function in V1:

$$\left(p_i - \sum_j \exp\left(\frac{\Phi_{ij} - u_i - v_j}{\sigma}\right), q_j - \sum_i \exp\left(\frac{\Phi_{ij} - u_i - v_j}{\sigma}\right)\right)$$

► Gradient of objective function in V2:

$$\left(p_i - \frac{\sum_{j} \exp\left(\frac{\Phi_{ij} - u_i - v_j}{\sigma}\right)}{\sum_{ij} \exp\left(\frac{\Phi_{ij} - u_i - v_j}{\sigma}\right)}, q_j - \frac{\sum_{i} \exp\left(\frac{\Phi_{ij} - u_i - v_j}{\sigma}\right)}{\sum_{ij} \exp\left(\frac{\Phi_{ij} - u_i - v_j}{\sigma}\right)}\right)$$

Coordinate descent on objective function in V1:

$$\begin{split} & p_{i} = \sum_{j} \exp \left(\frac{\Phi_{ij} - u_{i}\left(t+1\right) - v_{j}\left(t\right)}{\sigma} \right), \\ & q_{j} = \sum_{i} \exp \left(\frac{\Phi_{ij} - u_{i}\left(t\right) - v_{j}\left(t+1\right)}{\sigma} \right) \end{split}$$

that is

$$\left\{ \begin{array}{l} u_{i}\left(t+1\right) = \sigma \log \left(\frac{1}{p_{i}} \sum_{j} \exp \left(\frac{\Phi_{ij} - v_{j}\left(t\right)}{\sigma}\right)\right) \\ v_{j}\left(t+1\right) = \sigma \log \left(\frac{1}{q_{j}} \sum_{i} \exp \left(\frac{\Phi_{ij} - u_{i}\left(t\right)}{\sigma}\right)\right) \end{array} \right.$$

this is called the Iterated Fitting Proportional Procedure (IPFP), or Sinkhorn's algorithm.

► Coordinate descent on objective function in V2 does not yield a closed-form expression.

IPFP, linear version

Letting $a_i = \exp\left(-u_i/\sigma\right)$ and $b_j = \exp\left(-v_j/\sigma\right)$ and $K_{ij} = \exp\left(\Phi_{ij}/\sigma\right)$, one has $\pi_{ij} = a_ib_jK_{ij}$, and the procedure reexpresses as

$$\left\{ \begin{array}{l} a_{i}\left(t+1\right)=p_{i}/\left(\mathit{Kb}\left(t\right)\right)_{i} \text{ and } \\ b_{j}\left(t+1\right)=q_{j}/\left(\mathit{K}^{\mathsf{T}}a\left(t\right)\right)_{j}. \end{array} \right.$$

The log-sum-exp trick

▶ The previous program is extremely fast, partly due to the fact that it involves linear algebra operations. However, it breaks down when σ is small; this is best seen taking a log transform and returning to $u^k = -\sigma \log a^k$ and $v^k = -\sigma \log b^k$, that is

$$\begin{cases} u_i^k = \mu_i + \sigma \log \sum_j \exp\left(\frac{\Phi_{ij} - v_j^{k-1}}{\sigma}\right) \\ v_j^k = \zeta_j + \sigma \log \sum_i \exp\left(\frac{\Phi_{ij} - u_i^k}{\sigma}\right) \end{cases}$$

where $\mu_i = -\sigma \log p_i$ and $\zeta_i = -\sigma \log q_i$.

▶ One sees what may go wrong: if $\Phi_{ij} - v_j^{k-1}$ is positive in the exponential in the first sum, then the exponential blows up due to the small σ at the denominator. However, the "log-sum-exp trick" can be used in order to avoid this issue.

The log-sum-exp trick (ctd)

► Consider

$$\left\{ \begin{array}{l} \tilde{\mathbf{v}}_{i}^{k} = \max_{j} \left\{ \Phi_{ij} - \mathbf{v}_{j}^{k} \right\} \\ \tilde{\mathbf{u}}_{j}^{k} = \max_{i} \left\{ \Phi_{ij} - \mathbf{u}_{i}^{k} \right\} \end{array} \right.$$

(the indexing is not a typo: \tilde{v} is indexed by i and \tilde{u} by j).

One has

$$\begin{cases} u_i^k = \mu_i + \tilde{v}_i^{k-1} + \sigma \log \sum_j \exp\left(\frac{\Phi_{ij} - v_j^{k-1} - \tilde{v}_i^k}{\sigma}\right) \\ v_j^k = \zeta_j + \tilde{u}_j^k + \sigma \log \sum_i \exp\left(\frac{\Phi_{ij} - u_i^k - \tilde{u}_j^k}{\sigma}\right) \end{cases}$$

and now the arguments of the exponentials are always nonpositive, ensuring the exponentials don't blow up.