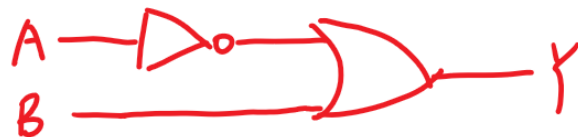


A.1.1  
(1.)

		B	
		0	1
A	0	1	1
	1	0	1

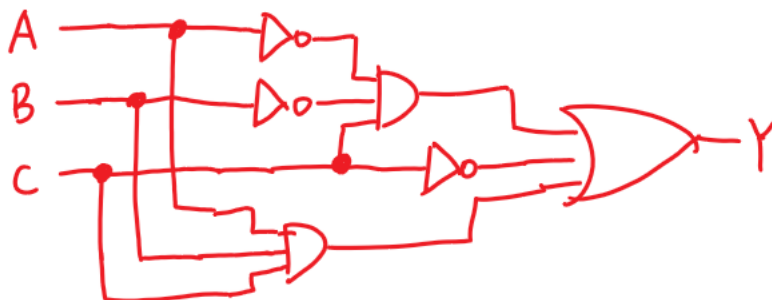
$$Y = \bar{A} + B$$



(2.)

		BC			
		00	01	11	10
A	0	1	1		1
	1	1		1	1

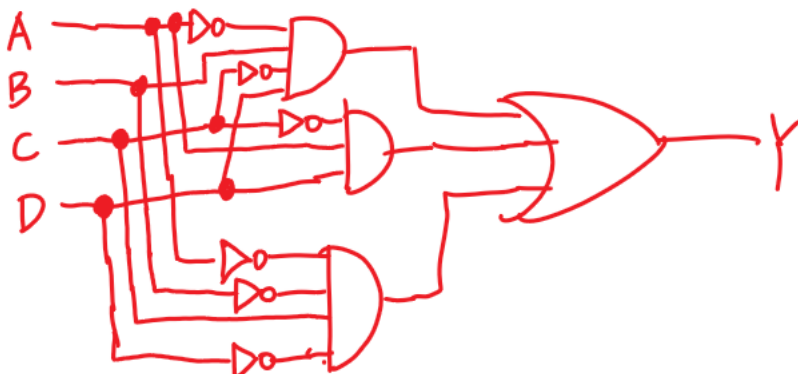
$$Y = \bar{C} + \bar{A}\bar{B}C + ABC$$



(3.)

		CD			
		00	01	11	10
AB	00	0	0	0	1
	01	0	1	0	0
	11	0	1	0	0
	10	0	1	0	0

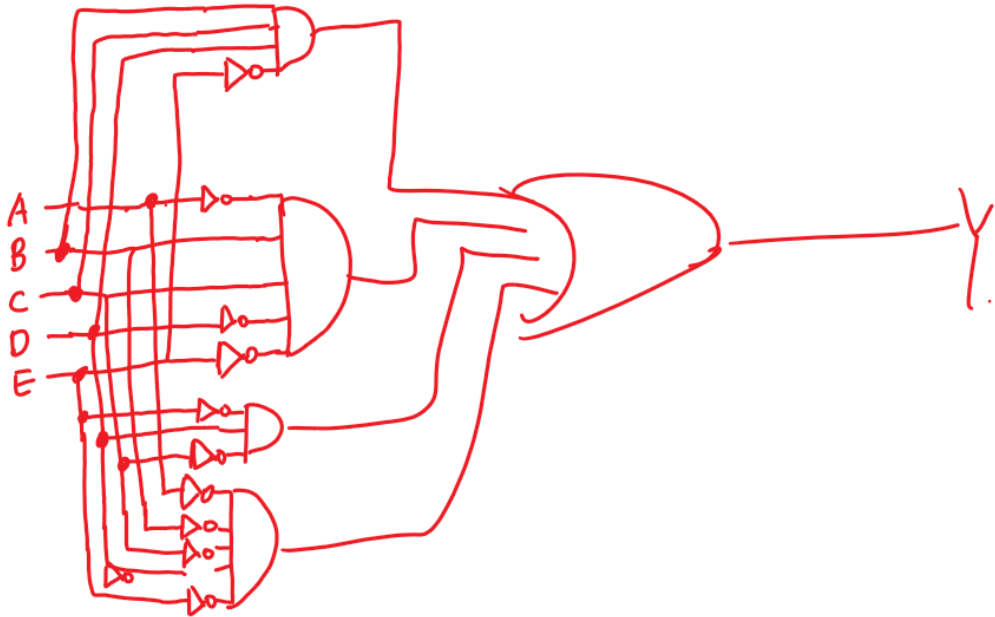
$$Y = \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}\bar{C}D + A\bar{C}\bar{D}$$



(4)

		CDE							
		000	001	011	010	110	111	101	100
AB	00	1	0	0	1	0	0	0	0
	01	0	0	0	1	1	0	0	1
	11	0	0	0	1	1	0	0	0
	10	0	0	0	1	0	0	0	0

$$Y = \bar{A}\bar{B}\bar{C}\bar{D}\bar{E} + \bar{C}D\bar{E} + BCD\bar{E} + \bar{A}BC\bar{D}\bar{E}$$



A.2.1(a)

when  $N = 3$

$$Y = A \oplus B \oplus C$$

A	B	C	$A \oplus B$	$A \oplus B \oplus C$
0	0	0	0	0
0	0	1	0	1
0	1	0	1	1
0	1	1	1	0
1	0	0	1	1
1	0	1	1	0
1	1	0	0	0
1	1	1	0	1

When  $N = 3$ , the output will be TRUE when there are exact 1 input is TRUE or all of the three inputs are TRUE

(b) When  $N = 4$

A	B	C	D	$A \oplus B$	$A \oplus B \oplus C$	$A \oplus B \oplus C \oplus D$
0	0	0	0	0	0	0
0	0	0	1	0	0	1
0	0	1	0	0	1	1
0	0	1	1	0	1	0
0	1	0	0	1	1	1
0	1	0	1	1	1	0
0	1	1	0	1	0	0
0	1	1	1	1	0	1
1	0	0	0	1	1	1
1	0	0	1	1	1	0
1	0	1	0	1	0	0
1	0	1	1	1	0	1
1	1	0	0	0	0	0
1	1	0	1	0	0	1
1	1	1	0	0	1	1
1	1	1	1	0	1	0

When  $N = 3$ , the output will be TRUE when there are exact 1 input is FALSE or TRUE.

A2.2 (i)

a	b	c	$b \oplus c$	$a(b \oplus c)$	ab	ac	$ab \oplus ac$
0	0	0	0	0	0	0	0
0	0	1	1	0	0	0	0
0	1	0	1	0	0	0	0
0	1	1	0	0	0	0	0
1	0	0	0	0	0	0	0
1	0	1	1	1	0	1	1
1	1	0	1	1	1	0	1
1	1	1	0	0	1	1	0

prove by truth table,  $a(b \oplus c) = ab \oplus ac$

(2)

a	b	c	$a \oplus b$	$(a \oplus b) \oplus c$	$b \oplus c$	$a \oplus (b \oplus c)$
0	0	0	0	0	0	0
0	0	1	0	1	1	1
0	1	0	1	1	1	1
0	1	1	1	0	0	0
1	0	0	1	1	0	1
1	0	1	1	0	1	0
1	1	0	0	0	1	0
1	1	1	0	1	0	1

prove by truth table  
 $(a \oplus b) \oplus c = a \oplus (b \oplus c)$

A2.3

a	b	c	$a \oplus b \oplus c$	$a \odot b$	$a \odot b \odot c$
0	0	0	0	1	0
0	0	1	1	1	1
0	1	0	1	0	1
0	1	1	0	0	0
1	0	0	1	0	1
1	0	1	0	0	0
1	1	0	0	1	0
1	1	1	1	1	1

prove by truth table,  $a \oplus b \oplus c = a \odot b \odot c$

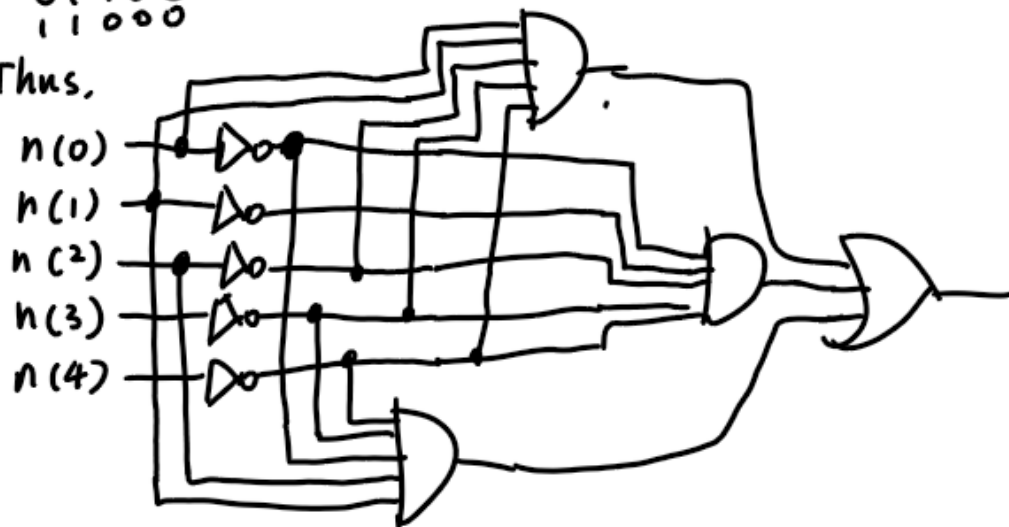
### A3.3

	n(0)	n(1)	n(2)	n(3)	n(4)	chkdiv3	chkdiv4	chkdiv12(d)
0	0	0	0	0	0	1	1	1
1	0	0	0	0	0	1	0	0
2	0	0	0	0	1	0	0	0
3	0	0	0	0	1	1	1	0
4	0	0	0	1	0	0	0	1
5	0	0	0	1	0	1	0	0
6	0	0	0	1	1	0	1	0
7	0	0	0	1	1	1	0	0
8	0	1	0	0	0	0	0	1
9	0	1	0	0	0	1	1	0
10	0	1	0	1	0	0	0	0
11	0	1	0	1	1	1	0	0
12	0	1	1	0	0	0	1	1
13	0	1	1	0	1	1	0	0
14	0	1	1	1	1	0	0	0
15	0	1	1	1	1	1	1	0
16	1	0	0	0	0	0	0	1
17	1	0	0	0	0	1	0	0
18	1	0	0	1	0	0	1	0
19	1	0	0	1	1	1	0	0
20	1	0	1	0	0	0	0	1
21	1	0	1	0	1	1	1	0
22	1	0	1	1	0	0	0	0
23	1	0	1	1	1	1	0	0
24	1	1	0	0	0	0	1	1
25	1	1	0	0	1	1	0	0
26	1	1	0	1	0	0	0	0
27	1	1	0	1	1	1	0	0
28	1	1	1	0	0	0	0	1
29	1	1	1	0	1	1	0	0
30	1	1	1	1	0	1	0	0
31	1	1	1	1	1	0	0	0

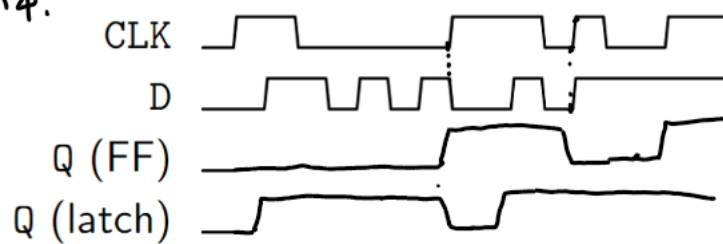
From the truth table ,

00000  
01100  
11000

Thus,

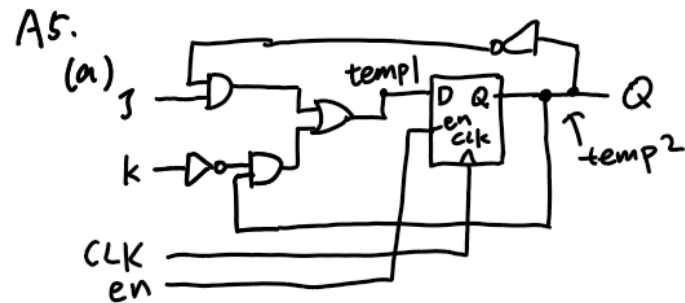


### A4.

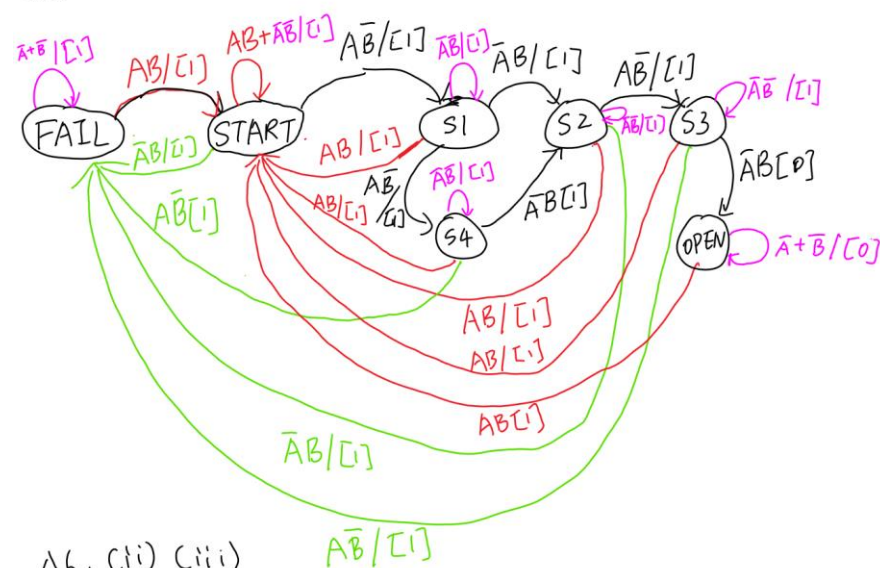


A3.4

After I decompose these two approaches into two-inputs gates, the first approach using divisible by 12 is equal to both divisible by 3 and 4 has **46** two-inputs gates. If directly write  $\text{chkdiv12}$  by its truth table, there are **14** two-inputs gates. Thus, the first approach is more physically larger.



A6. (i)



A6. (ii) (iii)

STATE REPRESENT	Current State	A	B	Next State	Output
0000 FAIL	0000	0	0	0000	1
0001 START	0000	0	1	0000	1
0010 S1	0000	1	0	0000	1
0011 S2	0000	1	1	0001	1
0100 S3	0001	0	0	0001	1
0101 S4	0001	0	1	0000	1
0110 OPEN	0001	1	0	0010	1
	0001	1	1	0001	1
	0010	0	0	0010	1
	0010	0	1	0011	1
	0010	1	0	0101	1
	0010	1	1	0001	1
	0011	0	0	0011	1
	0011	0	1	0000	1
	0011	1	0	0100	1
	0011	1	1	0001	1
	0100	0	0	0100	1
	0100	0	1	0110	0
	0100	1	0	0000	1
	0100	1	1	0001	1
	0101	0	0	0000	1
	0101	0	1	0011	1
	0101	1	0	0000	1
	0101	1	1	0001	1
	0110	0	0	0110	0
	0110	0	1	0110	0
	0110	1	0	0110	0
	0110	1	1	0001	1