

Foundations of Machine Learning — Homework Assignment 1

Anirudhan J Rajagopalan
N18824115
ajr619

October 11, 2015

A. PAC Learning

1

Algorithm A: Given a sample S , the algorithm returns the tightest interval I' containing all points labeled with 1 in sample S .

Error intervals: We define error regions as the intervals that lie between $[a, b]$ but outside I' .

Proof of PAC learnability :

Let I be the target concept — $[a, b]$.

Let $\epsilon > 0$ be the error.

Let m be the number of samples in S .

Let I_s be the tightest interval formed by the the points labelled 1 from sample S .

Let $\Pr[I_s]$ denote the probability mass of the interval defined by I_s .

$$\Pr[I_s] > \epsilon$$

Lets assume that the interval I_s (denoted by $[a', b']$) has error ϵ . The error can be found in intervals $[a, a']$ and $(b', b]$ denoted by I_1 and I_2 respectively. If we assume that the error is equally distributed across the two regions, we can denote the error to be $\epsilon/2$ for each of I_1 and I_2 . Each point in the error region has a probability of $(1 - \frac{\epsilon}{2})$. So Probability of the error in sample being greater than ϵ can be written as

$$\begin{aligned} \Pr_{S \sim \mathcal{D}^m} [R(I_S) > \epsilon] &\leq \sum_{i=1}^2 \Pr_{S \sim \mathcal{D}^m} [I \cap I_i = \emptyset] \\ &\leq 2(1 - \epsilon/2)^m \\ &\leq 2e^{-m\epsilon/2} \end{aligned}$$

Equating the RHS to δ gives us the sample complexity.

$$\begin{aligned} 2e^{-m\epsilon/2} &= \delta \\ \frac{2}{\delta} &= e^{m\epsilon/2} \\ \log \frac{2}{\delta} &= \frac{m\epsilon}{2} \\ m &= \frac{2}{\epsilon} \log \frac{2}{\delta} \end{aligned}$$

2

Algorithm A: Given a sample \mathcal{S} for target concept C containing p closed intervals,

1. If there are p separate sequence of positively labeled points in the training data, then return the union of the p tightest intervals containing the positive points.
2. Otherwise, return $(p - i)$ tightest intervals, each containing a sequence of positive labels separated by i negative labels. i can take the values 0 to $(p - 1)$.

Error intervals: Let $[a_1, b_1] \cup [a_2, b_2] \cup \dots [a_p, b_p]$ be the target concept C . Let $\epsilon > 0$. We can assume that $\Pr[a_i, b_i] > \epsilon/(k + 1)$. The actual error on the training set is gap between the target concept $[a_i, b_i]$ and the learned concept $[a'_i, b'_i]$. Each of the error regions r_i can occur with a probability $\epsilon/2(k + 1)$. (Factor of 2 for $[a_i, a'_i]$ and $[b'_i, b_i]$).

Proof of PAC learnability :

If $\text{error}(h_S) > \epsilon$, then either the union of the intervals misses one of the regions r_i or $\Pr[b_i, a_{i+1}] > \epsilon/(k + 1)$. Thus by union bound, we have

$$\begin{aligned} \Pr[\text{error}(h_S)] &\leq \Pr[\exists i \in [1, p] : h_S \text{ misses } r_i] + e^{-m\epsilon/(k+1)} \\ &\leq 2k(1 - \frac{\epsilon}{2(k+1)})^m + e^{-m\epsilon/(k+1)} \\ &\leq 2ke^{-m\epsilon/2(k+1)} + e^{-m\epsilon/(k+1)} \\ &\leq (2k + 1)e^{-m\epsilon/2(k+1)} \end{aligned}$$

Setting $\delta = 0$ and solving the RHS will give us a bound for m .

$$(2k + 1)e^{-m\epsilon/2(k+1)} = \delta \tag{1}$$

$$\frac{2k + 1}{\delta} = e^{m\epsilon/2(k+1)} \tag{2}$$

$$\log \frac{2k + 1}{\delta} = \frac{m\epsilon}{2(k + 1)} \tag{3}$$

$$m = \frac{2(k + 1)}{\epsilon} \log \frac{2k + 1}{\delta} \tag{4}$$

Therefore, for probability of error to be less than ϵ m should be greater than the value obtained in (4).

When p is 2 :

Substituting $k = 2$ in the equation (4) above gives us $m \geq \frac{6}{\epsilon} \log \frac{5}{\delta}$