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Q 1 Define RSA algorithm and explain setup briefly.

RSA is public key encryption algorithm (asymmetric) based on the difficulty of factoring large number steps:

1. chose two prime number  $p$  and  $q$
2. Compute  $n = p * q$
3. compute Euler's totient  $\varphi(n) = (p-1)(q-1)$
4. choose a public key  $e$  such that  $1 < e < \varphi(n)$  and  $\gcd(e, \varphi(n)) = 1$
5. Compute the ~~private~~ key  $d$ , where  $d * e = 1 \pmod{\varphi(n)}$
6. encryption  $C = m^e \pmod{n}$
7. decryption  $m = C^d \pmod{n}$

Q 2: What the homomorphic property in cryptography?

Solve:

It means operation on ciphertexts correspond to operation on plaintexts

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Q.3 Explain why RSA is multiplicatively homomorphy?

Solve:

because:  $E(n) = m^e \text{ mod } n$

$$E(m) = m^e \text{ mod } n$$

$$\Rightarrow E(m_1) \times E(m_2) = (m_1^e \text{ mod } n) (m_2^e \text{ mod } n)$$

→

$$= ((m_1^e \text{ mod } n) (m_2^e \text{ mod } n)) \text{ mod } n = (m_1 \times m_2)^e \text{ mod } n$$

$$= E(m_1 \times m_2)$$

Q.4 what the advantage and risk of homomorphy encryption?

Solve:

Advantage: compute on encrypted data

1. compute on encrypted data

2. privacy include and data analysis

Risk:

1. vulnerable to manipulate (malleability) and choose ciphertext attacks

2. slow for large data

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## 5. why use mod in RSA ?

Solve

To keep result within range prevent overflow  
and ensure reversibility,  
It the base of RSA security and math properties

6. Compute manually for small case ( $p=3, q=11$ ).  
 $e=7, m_1=2, m_2=5$ )

$$n = p \times q \Rightarrow n = 3 \times 11 = 33, \varphi(n) = (3-1)(11-1) = 20$$

$$d = 3^{-1} \text{ mod } 20$$

$$E(m_1) = 2^7 \text{ mod } 33 = 128 \text{ mod } 33 = 29$$

$$E(m_2) = 5^7 \text{ mod } 33 = 78125 \text{ mod } 33 = 19$$

Now multiply both ciphertext:

$$E(m_1) \times E(m_2) \text{ mod } 33 = 29 \times 19 \text{ mod } 33 = 10$$

$$E(m_1 \times m_2) = E(2 \times 5) = E(10) = 10^7 \text{ mod } 33 \\ = 10000000 \text{ mod } 33 = 10$$

Decryption

$$\text{dec}(29) = 29^3 \text{ mod } 33 = 24389 \text{ mod } 33 = 2$$

$$\text{dec}(19) = 19^3 \text{ mod } 33 = 2794 \text{ mod } 33 = 5$$

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Q7. Why is key size important for security?

\* solve:

Large keys make factoring n harder  $\rightarrow$  strong security

Q8. Explain what happens if  $e$  and  $\varphi(n)$  are not coprime.

If  $e$  and  $\varphi(n)$  are not coprime then there is no modular inverse for  $e$ , so  $d$  can't be calculated.

Without  $d$  decryption is impossible and the key pair is invalid.

Q9. Describe one-real-word application of fully homomorphic encryption.

Secure cloud data processing:

it allows storing encrypted data in the cloud and performing calculation on searches on it without decryption, keeping the data safe and private.



```
# Decrypt function
def decrypt(c,d,n):
    m= (c**d) % n
    return m

# Encrypt function
def encrypt(m,e,n):
    c= (m**e) % n
    return c

# Generate p, q, n, φ(n), and keys (e, d)

def genate_key():
    p=7
    q=11
    n=p*q
    phi=(p-1)*(q-1)
    e=17
    d=2753
    return p,q,n,phi ,e,d
```

```
def main():
    #Generate p, q, n, ϕ(n), and keys (e, d) by generate_key function
    p,q,n,phi ,e,d=generate_key()
    #choose the m1 and m2
    m1=6
    m2=8
    # Encrypt m1 and m2
    cm1=encrypt(m1,e,n)
    cm2=encrypt(m2,e,n)
    # Multiply ciphertexts and decrypt the result
    mult_cm= (cm1*cm2) % n
    dec_mult= decrypt(mult_cm,d,n)
    # Calculate expected result
    excepted_mult= (m1*m2) % n
```

```
    print("p:",p," q:",q)
    print("n:",n," φ:",phi)
    print("public key (e,n):(",e,",",n,")")
    print("private key (d,n):(",d,",",n,")\n\n")

    print("m1:",m1," m2:",m2)
    print("E(m1) = ",cm1," E(m2) = ",cm2)
    print ("E(m1)*E(m2) mod n = ",mult_cm)
    print("decrypeted rusute = ",dec_mult)
    print("expected (m1*m2 mod n) = ",excepted_mult)
    # Check homomorphic property
    if dec_mult == excepted_mult:
        print("\nHomomorphic property done")
    else:
        print("Homomorphic property failed")

if __name__ == "__main__":
    main()
```

```
*** p: 7 q: 11
n: 77 φ: 60
public key (e,n):( 17 , 77 )
private key (d,n):( 2753 , 77 )
```

```
m1: 6 m2: 8
E(m1) = 41 E(m2) = 57
E(m1)*E(m2) mod n = 27
decrypted result = 48
expected (m1*m2 mod n) = 48
```

Homomorphic property done

