

## Reference Sheet for Discrete Maths

### Propositional Calculus

Order of decreasing binding power:  $=, \neg, \wedge/\vee, \Rightarrow/\Leftarrow, \equiv/\neq$ .

**Equivales** is the only equivalence relation that is associative  
 $((p \equiv q) \equiv r) \equiv (p \equiv (q \equiv r))$ , and it is symmetric and has identity **true**.

**Discrepancy** (difference) ' $\neq$ ' is symmetric, associative, has identity '**false**', mutually associates with equivales  $((p \neq q) \equiv r) \equiv (p \neq (q \equiv r))$ , and mutually interchanges with it as well  $(p \neq q \equiv r) \equiv (p \equiv q \neq r)$ . Finally, negation commutes with difference:  $\neg(p \equiv q) \equiv \neg p \equiv q$ .

**Implication** has the alternative definition  $p \Rightarrow q \equiv \neg p \vee q$ , thus having **true** as both left identity and right zero; it distributes over  $\equiv$  in the second argument, and is self-distributive; and has the properties:

|   |  |
|---|--|
| <b>Shunting</b> $p \wedge q \Rightarrow r \equiv p \Rightarrow (q \Rightarrow r)$ | <b>Modus Ponens</b>                            |
|   | $p \wedge (p \Rightarrow q) \equiv p \wedge q$ |
| <b>Contrapositive</b> $p \Rightarrow q \equiv \neg q \Rightarrow \neg p$          | $p \wedge (q \Rightarrow p) \equiv p$          |
|   | $p \wedge (p \Rightarrow q) \Rightarrow q$     |
| <b>Leibniz</b> $e = f \Rightarrow E[z \asymp e] = E[z \asymp f]$                  |  |

It is a *linear* order relation generated by '**false**  $\Rightarrow$  **true**'; whence "from false, follows anything": **false**  $\Rightarrow p$ . Moreover it has the useful properties "(3.62) Contextualisation":  $p \Rightarrow (q \equiv r) \equiv p \wedge q \equiv p \wedge r$ —we *have* the context  $p$  in each side of the equivalence—and  $p \Rightarrow (q \Rightarrow r) \equiv p \wedge q \Rightarrow p \wedge r$ . Implication is "Sub-associative":  $((p \Rightarrow q) \Rightarrow r) \Rightarrow (p \Rightarrow (q \Rightarrow r))$ . Finally, we have " $\equiv$ -Elimination":  $(p \equiv q \equiv r) \Rightarrow s \equiv p \Rightarrow s \equiv q \Rightarrow s \equiv r \Rightarrow s$ .

**Conjunction and disjunction** distribute over one another, are both associative and symmetric,  $\vee$  has identity **false** and zero **true** whereas  $\wedge$  has identity **true** and zero **false**,  $\vee$  distributes over  $\vee, \equiv, \wedge, \Rightarrow, \Leftarrow$  whereas  $\wedge$  distributes over  $\equiv - \equiv$  in that  $p \wedge (q \equiv r \equiv s) \equiv p \wedge q \equiv p \wedge r \equiv p \wedge s$ , and they satisfy,

|                        |                                       |  |  |
|------------------------|---------------------------------------|--|--|
| <b>Excluded Middle</b> | <b>Contradiction</b>                  | <b>Absorption</b>                            | <b>De Morgan</b>                             |
| $p \vee \neg p$        | $p \wedge \neg p \equiv \text{false}$ | $p \wedge (q \vee \neg p) \equiv p \wedge q$ | $\neg(p \wedge q) \equiv \neg p \vee \neg q$ |
|                        |                                       | $p \vee (q \vee \neg p) \equiv p \vee q$     | $\neg(p \vee q) \equiv \neg p \wedge \neg q$ |

Most importantly, they satisfy the "**Golden Rule**":  $p \wedge q \equiv p \equiv q \equiv p \vee q$ .

The many other properties of these operations—such as weakening laws and other absorption laws and case-analysis ( $\sqcup$ -char)—can be found by looking at the list of *lattice properties*—since the Booleans are a lattice.

### Orders

An *order* is a relation  $\sqsubseteq : \tau \rightarrow \tau \rightarrow \mathbb{B}$  satisfying the following three properties:

|                    |  |   |
|--------------------|--|---|
| <b>Reflexivity</b> | <b>Transitivity</b>  | <b>Mutual Inclusion</b>                               |
| $a \sqsubseteq a$  | $a \sqsubseteq b \wedge b \sqsubseteq c \Rightarrow a \sqsubseteq c$ | $a \sqsubseteq b \wedge b \sqsubseteq a \equiv a = b$ |

Indirect Inclusion is like 'set inclusion' and Indirect Equality is like 'set extensionality'.

|   |  |
|---|--|
| <b>Indirect Equality (from above)</b>                                     | <b>Indirect Inclusion (from above)</b>   |
| $x = y \equiv (\forall z \bullet x \sqsubseteq z \equiv y \sqsubseteq z)$ | $x \sqsubseteq y \equiv (\forall z \bullet y \sqsubseteq z \Rightarrow x \sqsubseteq z)$ |

|   |  |
|---|--|
| <b>Indirect Equality (from below)</b>                                     | <b>Indirect Inclusion (from below)</b>   |
| $x = y \equiv (\forall z \bullet z \sqsubseteq x \equiv z \sqsubseteq y)$ | $x \sqsubseteq y \equiv (\forall z \bullet z \sqsubseteq x \Rightarrow z \sqsubseteq y)$ |

An order is *bounded* if there are elements  $\top, \perp : \tau$  being the lower and upper bounds of all other elements:

|                       |                                      |                          |  |
|-----------------------|--------------------------------------|--------------------------|--|
| <b>Top Element</b>    | $a \sqsubseteq \top$                 | <b>Bottom Element</b>    | $\perp \sqsubseteq a$                  |
| <b>Top is maximal</b> | $\top \sqsubseteq a \equiv a = \top$ | <b>Bottom is minimal</b> | $a \sqsubseteq \perp \equiv a = \perp$ |

### Lattices

A *lattice* is a pair of operations  $\sqcap, \sqcup : \tau \rightarrow \tau \rightarrow \tau$  specified by the properties:

|  |  |
|--|--|
| <b><math>\sqcup</math>-Characterisation</b>                              | <b><math>\sqcap</math>-Characterisation</b>                              |
| $a \sqsubseteq c \wedge b \sqsubseteq c \equiv a \sqcup b \sqsubseteq c$ | $c \sqsubseteq a \wedge c \sqsubseteq b \equiv c \sqsubseteq a \sqcap b$ |

The operations act as providing the greatest lower bound, 'glb', 'supremum', or 'meet', by  $\sqcap$ ; and the least upper bound, 'lub', 'infimum', or 'join', by  $\sqcup$ .

Let  $\square$  be one of  $\sqcap$  or  $\sqcup$ , then:

|   |   |  |
|---|---|--|
| <b>Symmetry of <math>\square</math></b> | <b>Associativity of <math>\square</math></b>        | <b>Idempotency of <math>\square</math></b> |
| $a \square b = b \square a$             | $(a \square b) \square c = a \square (b \square c)$ | $a \square a = a$                          |

|                                     |   |                             |  |
|-------------------------------------|---|-----------------------------|--|
| <b>Zero of <math>\square</math></b> | <b>Identity of <math>\square</math></b> | <b>Absorption</b>           | <b>Self-Distributivity of <math>\square</math></b>             |
| $a \sqcup \top = \top$              | $a \sqcup \perp = a$                    | $a \sqcap (b \sqcup a) = a$ | $a \sqcap (b \square c) = (a \square b) \square (a \square c)$ |
| $a \sqcap \perp = \perp$            | $a \sqcap \top = a$                     | $a \sqcup (b \sqcap a) = a$ |  |

|                                     |  |  |
|-------------------------------------|--|--|
| <b>Weakening</b>                    | <b>Induced Defs. of Inclusion</b>  | <b>Golden Rule</b>                               |
| <b>/ Strengthening</b>              | $a \sqsubseteq b \equiv a \sqcup b = b$  | $a \sqcap b = a \equiv b = a \sqcup b$           |
| $a \sqsubseteq a \sqcup b$          | $a \sqsubseteq b \equiv a \sqcap b = a$  | $a \sqcap b = a \sqcup b \equiv a = b$           |
| $a \sqcap b \sqsubseteq a$          |  | $a \sqcup b \sqsubseteq a \sqcap b \equiv a = b$ |
| $a \sqcap b \sqsubseteq a \sqcup b$ | <b>Monotonicity of <math>\square</math></b>  |  |
|                                     | $a \sqsubseteq b \wedge c \sqsubseteq d \Rightarrow a \square c \sqsubseteq b \square d$ |  |

### Duality Principle:

If a statement  $S$  is a theorem, then so is  $S[(\sqsubseteq, \sqcap, \sqcup, \top, \perp) := (\supseteq, \sqcup, \sqcap, \perp, \top)]$ .

## Conditionals

“**Axiom, Definition of if**” “if to  $\wedge$ ”  $P[z \Leftarrow \text{if } b \text{ then } x \text{ else } y \text{ fi}] \equiv (b \Rightarrow P[z \Leftarrow x]) \wedge (\neg b \Rightarrow P[z \Leftarrow y])$

“**Alternative definition of if**” “if to  $\nLeftarrow$ ”  $P[z \Leftarrow \text{if } b \text{ then } x \text{ else } y \text{ fi}] \equiv b \wedge P[z \Leftarrow x] \nLeftarrow \neg b \wedge P[z \Leftarrow y]$

“**Alternative definition of if**” “if to  $\vee$ ”  $P[z \Leftarrow \text{if } b \text{ then } x \text{ else } y \text{ fi}] \equiv (b \wedge P[z \Leftarrow x]) \vee (\neg b \wedge P[z \Leftarrow y])$

“**Alternative definition of if**” “if to  $\Rightarrow$ ”  $P[z \Leftarrow \text{if } b \text{ then } x \text{ else } y \text{ fi}] \equiv b \Rightarrow P[z \Leftarrow x] \equiv \neg b \Rightarrow P[z \Leftarrow y]$

If true    if true then  $x$  else  $y$  fi =  $x$   
 If false   if false then  $x$  else  $y$  fi =  $y$

### If then-constant

if  $R$  then true else  $P$  fi =  $R \vee P$   
 if  $R$  then false else  $P$  fi =  $\neg R \wedge P$

### If else-constant

if  $R$  then  $P$  else true fi =  $R \Rightarrow P$   
 if  $R$  then  $P$  else false fi =  $R \wedge P$

“**If swap**” if  $b$  then  $x$  else  $y$  fi = if  $\neg b$  then  $y$  else  $x$  fi

“**If idempotency**” if  $b$  then  $x$  else  $x$  fi =  $x$

“**If guard strengthening**” if  $b$  then  $x$  else  $y$  fi = if  $b \wedge x \neq y$  then  $x$  else  $y$  fi

“**If Distributivity**”  $P[z \Leftarrow \text{if } b \text{ then } x \text{ else } y \text{ fi}] = \text{if } b \text{ then } P[z \Leftarrow x] \text{ else } P[z \Leftarrow y] \text{ fi}$

“**If Context**” if  $b$  then  $E$  else  $F$  fi = if  $b$  then  $E[b \Leftarrow \text{true}]$  else  $F[b \Leftarrow \text{false}]$  fi

“**If junctivity**”  $(\text{if } b \text{ then } x \text{ else } y \text{ fi}) \oplus (\text{if } b \text{ then } x' \text{ else } y' \text{ fi}) = \text{if } b \text{ then } (x \oplus x') \text{ else } (y \oplus y') \text{ fi}$

## Set Theory

### Axiom (11.3) “Set membership”:

$F \in \{x \mid R \bullet E\} \equiv (\exists x \mid R \bullet F = E)$

### Theorem (11.7) “Simple Membership”:

$e \in \{x \mid P\} \equiv P[e]$

### Theorem (11.6) “Mathematical formulation of set comprehension”:

$\{x \mid P \bullet E\} = \{y \mid (\exists x \mid P \bullet y = E)\}$

### Theorem (11.9) “Simple set comprehension equality”:

$\{x \mid Q\} = \{x \mid R\} \equiv (\forall x \bullet Q \equiv R)$

### Axiom (11.13) “Subset” “Definition of $\subseteq$ ” “Set inclusion”:

$S \subseteq T \equiv (\forall e \mid e \in S \bullet e \in T)$

*More coming soon!*