



Lecture 10

CS 537- Big Data Analytics

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Map Reduce

Scheduling and Data Flow

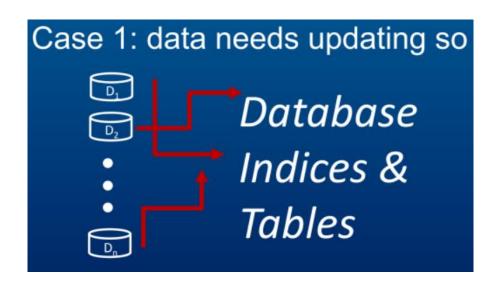
The Problem and The Solution

Problem:

Big Data -> Large amount of data stored in large amount of devices

Solution:

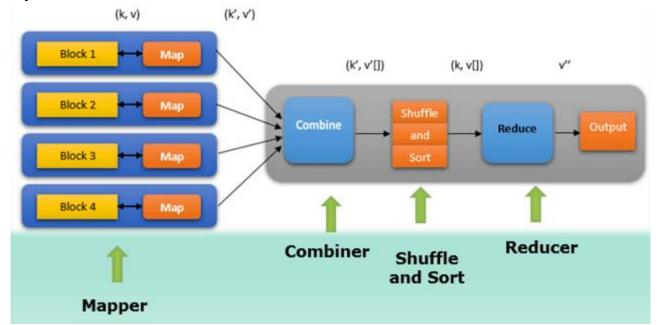
- Bring computations to data
- Possibilities:





MapReduce Framework

- User defines
 - <key, value>
 - mapper & reducer functions
- Logistics:
 - Hadoop handles the distribution and execution



MapReduce Flow

- User defines a map function
 - map() reads data and outputs <key,value>



- User defines a reduce function
 - reduce() reads <key, value > and outputs your result

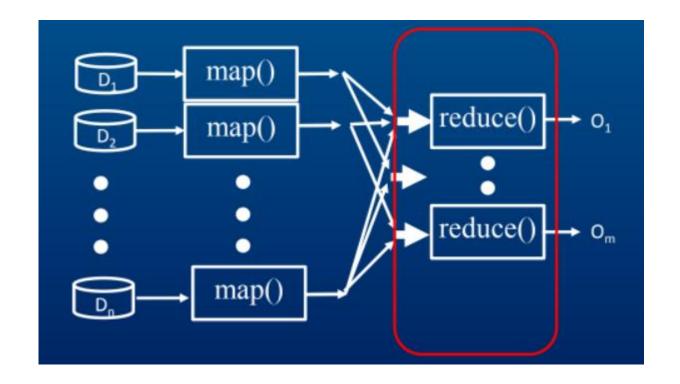


Hadoop Rule of Thumb:

- 1 mapper per data split (typically)
- 1 reducer per computer node (best parallelism)

MapReduce Flow

- Hadoop distributes map() to data
- Hadoop groups <key,value> data
- Hadoop distributes groups to reducers()

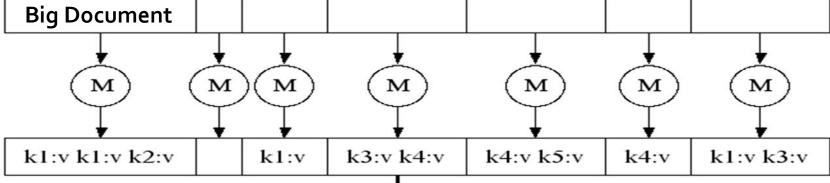


MapReduce Working Diagram

MAP:

produces a set of key-value pairs

Input



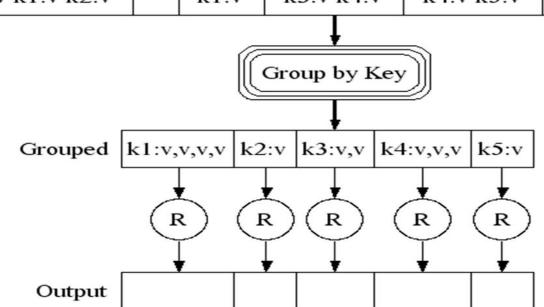
Intermediate

Group by key:

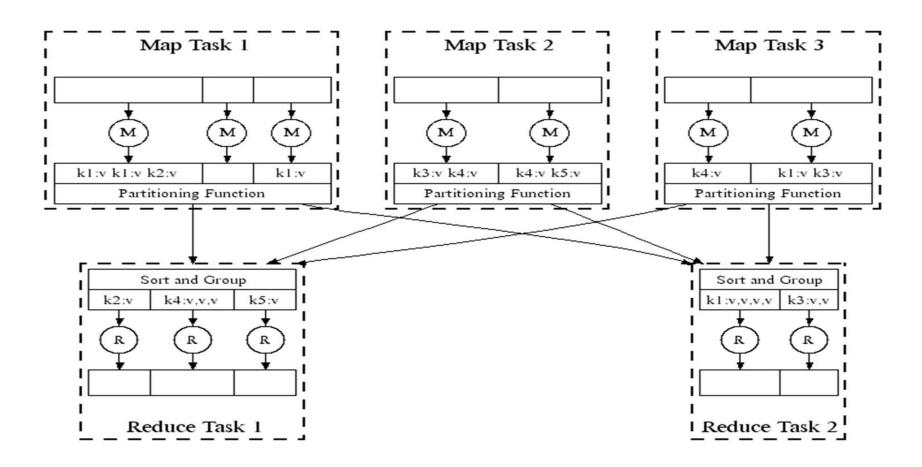
Collect all pairs with same key (Hash merge, Shuffle, Sort, Partition)

Reduce:

Collect all values belonging to the key and output



MapReduce: Parallel Processing



All phases are distributed with many tasks doing the work

Word Count Example

- Count word frequencies
- How would you count all the words in Star Wars?
- In a nutshell:
 - Get word
 - Look up word in table
 - Add 1 to count
- How would you count all the words in all the Star Wars scripts and books, blogs, and fan-fiction?



Word	Count
а	1000
far	2000
Jedi	5000
Luke	9000

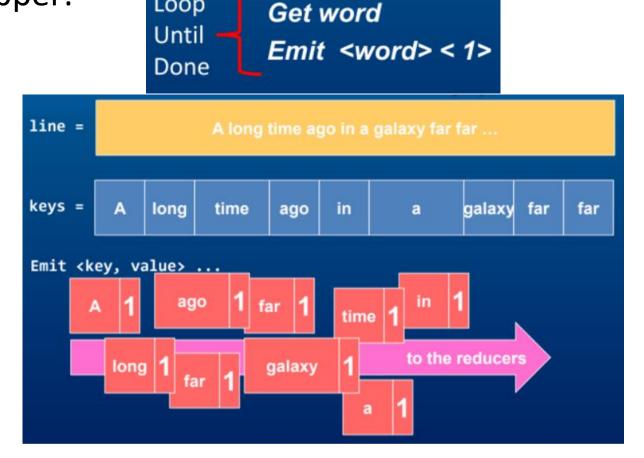
- Keep it simple (remember big data and simple aggregations)
 - Let <word, 1> be the <key, value>

Loop

The mapper:

Mappers are separate and independent

Mappers work on data parts

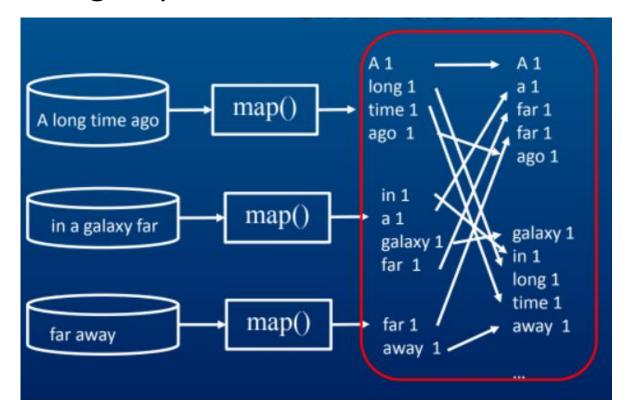


- Lets Hadoop do the hard work
 - The reducer:

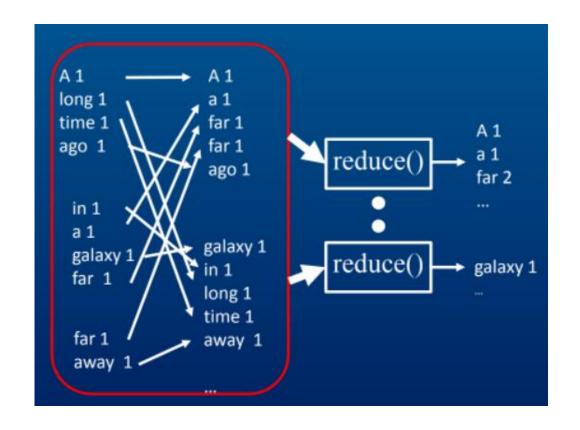
```
Cop
Over
key-
values

Get next <word><value>
If <word> is same as previous word
add <value> to count
else
emit <word> < count>
set count to 1
```

Hadoop shuffles, groups, and distributes:

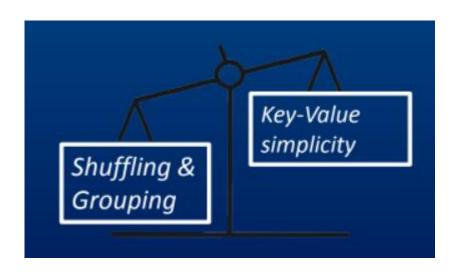


reduce() aggregates



Ideal properties

- Good key-value properties
 - Simple
 - Enables reducers to get correct output
- Good Task Decomposition:
 - Mappers: simple and separable
 - Reducers: easy consolidation



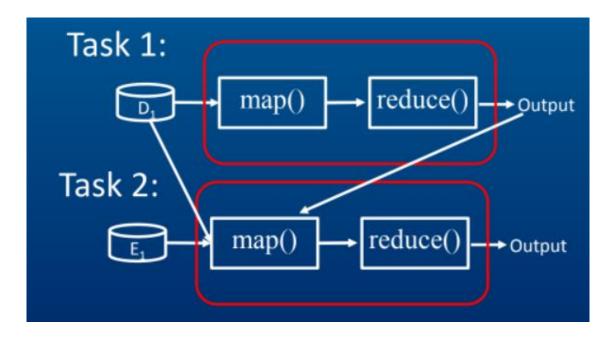
Trending Word Counts

- Let's make first example little complicated:
 - We need to calculate word count in twitter tweets by day
 - To find trending topics
 - Twitter Data:
 - Date
 - Message
 - Location
 - Other metadata
 - Tasks
 - Task 1: Get word count by day
 - Task 2: Get total word count

- 1 #PakvsNz
- 2 #SalarioRosa
- 3 #PakvsNewzealand
- 4 #Edho
 Tweet Counts N/A
- 5 Asif Ali 18K Tweets
- 6 Malik 110K Tweets
- 7 #ShamiKiFarziTrolling
- 8 Pakistan
- 9 Haris Rauf

Trending Word Counts: Task Decomposition

- For task 1 we need to use composite key:
 - Map/Reduce: <date word,count>
- For task 2 we can:
 - Reuse previous word count example
 - Use the output of task one

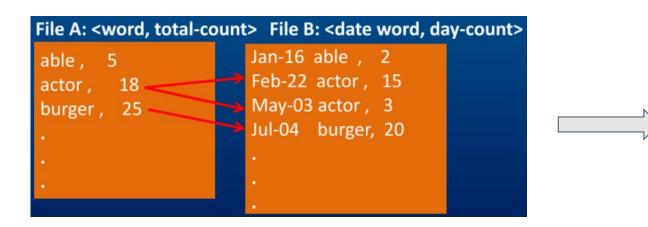


Joining Data

- Task: combine datasets by key
 - A standard data management function
 - In pseudo SQL

Select * from table A, table B, where A.key=B.key

- Joins can be inner, left or right outer
- Task: given two wordcount datasets as following:



```
File AjoinB: <word date, day-count total-count >

able Jan-16, 2 5
actor Feb-22, 15 18
actor May-03, 1 18
burger Jul-04, 20 25

.
.
```

Joining Data (Cont)

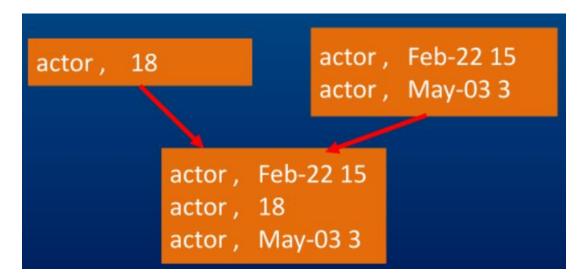
- For joining keys should be same but here:
 - File A: <word, total-count>
 - File B: <date word, day-count>
 - Word is same in both keys but date is not present in File A so we need to filter out date for key of File B
 - Now: Put Date into value field
 - File B: <word, date day-count total-count >

```
Jan-16 able , 2
Feb-22 actor , 15
May-03 actor , 3
Jul-4 burger, 20

.
.
.
```

Task Decomposition

- How will Hadoop shuffle & group these?
- Let's focus on 1 key:
- Hadoop gathers the data for a join



Task Decomposition

Reducer now has all the data for same word grouped together



Reducer can now join the data and put date back into key

```
actor, 18
actor, Feb-22 15
actor, May-03 3
```

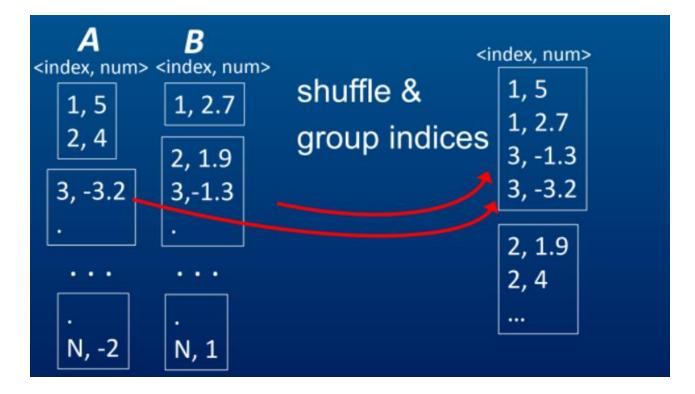
Vector Multiplication

- Task: multiply 2 arrays of N numbers
 - A basic mathematical operation
 - Let's assume N is very large
 - Data is distributed in HDFS
 - We need elements with same index together

Let <key, value> = <index, number>

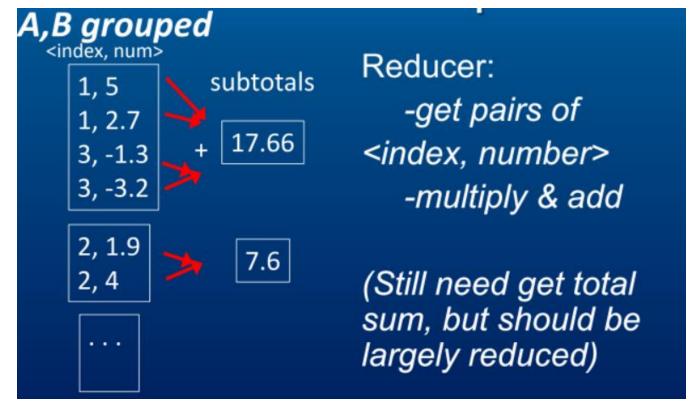
Vector Multiplication

- Lets assume we already have indexes of elements stored
- Mapper task is as following:



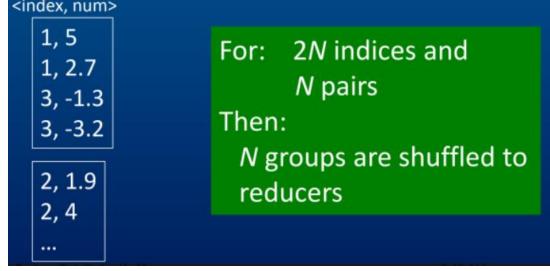
Vector Multiplication

- Lets assume we already have indexes of elements stored
- Reducer task is as following:



Computational Costs

- For Vector Multiplication
 - How many <index, number> are output from map()?
 - How many <index> groups have to be shuffled?



How many <index, number> are output from map()

How many <index> groups have to be shuffled?

Computational Costs

- We can reduce shuffling by:
 - Try: 'combine' map indices in mapper (works better for Wordcount)

<index bin, original-index number>

Or Try: use index ranges of length R

• For example, let R=10, and bin the array indices 1 2 3 4 ... 10 11 1219 20 21 (N-9) N place in bins N keys are now N/R=N/10 keys <key,value> is now

Computational Costs

Now shuffling costs depend on N/R groups

If: R=1

Then: N/R=N groups (same as before)

If: R>1

Then: N/R<N (less shuffling to do)

```
Trade-offs:

If:
    size of (N/R) ↑
    Then:
    shuffle costs ↑

But:
    reducer complexity ↓
```

Note: Matrix multiplication needs row-index and col-index in the keys

MapReduce: Environment

Map-Reduce environment takes care of:

- Partitioning the input data
 - Scheduling the program's execution across a set of machines
- Performing the group by key step
- Handling node failures
- Managing required inter-machine
- communication

MapReduce: Environment

- Input and final output are stored on the distributed file system (DFS):
 - Scheduler tries to schedule map tasks "close" to physical storage location of input data
- Intermediate results are stored on local FS of Map and Reduce workers
- Output is often input to another MapReduce task

MapReduce: Coordination Master

- Master node takes care of coordination:
 - Task status: (idle, in-progress, completed)
 - o Idle tasks get scheduled as workers become available
 - When a map task completes, it sends the master the location and sizes of its R intermediate files, one for each reducer
 - Master pushes this info to reducers
 - Master pings workers periodically to detect failures

MapReduce: Dealing with Failures

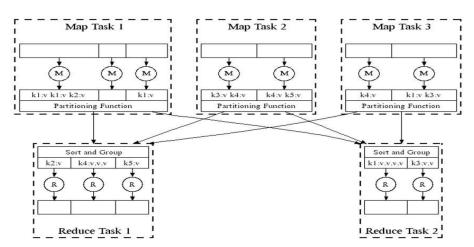
- Map worker failure
 - Map tasks completed or in-progress at worker are reset to idle.
 - Idle tasks eventually rescheduled on other worker(s)
- Reduce worker failure
 - o Only in-progress tasks are reset to idle
 - Idle Reduce tasks restarted on other worker(s)
- Master failure
 - MapReduce task is aborted and client is notified

How many Map and Reduce Jobs

- Suppose we have M map tasks, R reduce tasks
- Rule of thumb:
 - Make M much larger than the number of nodes in the cluster
 - One DFS chunk per map is common
 - Improves dynamic load balancing and speeds up recovery from worker failures
 - Usually R is smaller than M
 - O Because output is spread across R files

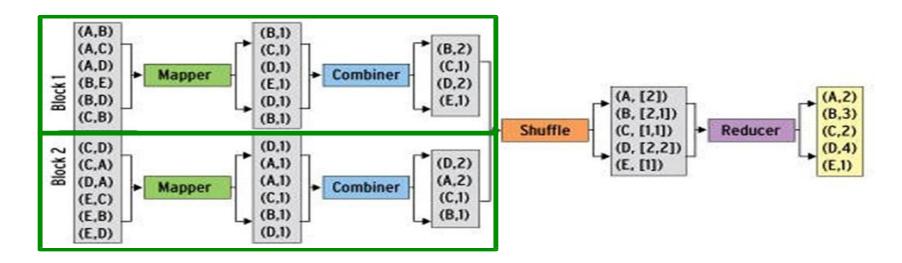
Combiners

- Often a Map task will produce many pairs of the form (k,v1), (k,v2), ... for the same key k
 - O E.g., popular words in the word count example
 - Can save network time by pre-aggregating values in the mapper:
- combine(k, list(v1))
 v2
- Combiner is usually same as the reduce function



Combiners (Cont.)

- Back to our word counting example:
 - Combiner combines the values of all keys of a single mapper (single node):



Much less data needs to be copied and shuffled!

Combiners (Cont.)

- Combiner trick works only if reduce function is commutative and associative.
- Sum:

$$2 + (5 + 7) = (2 + 5) + 7$$

- Average
- Median

Partition Function

- Want to control how keys get partitioned
 - The set of keys that go to a single reduce worker
- System uses a default partition function:
 - O Hash (key) mod R
- Sometimes useful to override the hash function:
 - E.g., hash (hostname(URL)) mod R ensures URLs from a host end up in the same output file

Limitations of MapReduce

- Must fit <key, value> paradigm
- Map/Reduce data not persistent
- Requires programming/debugging
- Not interactive

That's all for today.