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Example: greatest common divisor (gcd) Where is the recursive structure?? How to solve gcd (a, b) in terms of gcd (a', b') for smaller inputs? Claim: common divisors of (a,b)

are the same as the common divisors of (b,r) where r = a 2 b. Recall "division alsorithm"; a = 9.6 + r where $0 \le r \le b$ Suppose ala & alb. Then a=nd, b=n'd. Then $d \mid r : r = \alpha - \gamma \cdot b$ $= md - \gamma m'd = (m - \gamma m')d$ other direction: dlb + l/r => dla. Two things to notice: D gcd (a,b) = gcd (b,r) ② トきり. In C+++ int gcd (int a, int b) { if (a3b = >0) return bi