

# Recursion

very much like mathematical induction.

Inductive Proofs: trying to prove some statement for all natural #s.

① Base case: explicitly demonstrate truth for small value of  $n$ .  
Say  $n=1$ .

② Assume truth for all values  $< n$ .  
Then, show this would  $\Rightarrow$  truth for  $n$ .

③ ① & ② together  $\Rightarrow$  truth for all  $n \in \mathbb{N}$ .

Proof for  $n \geq 5$  look like?

$$\textcircled{1} \Rightarrow t(1), t(1) \xRightarrow{\textcircled{2}} t(2) \Rightarrow t(3)$$

$$\dots \xRightarrow{\textcircled{2}} t(5)$$

Recursive functions: just like induction.

① Prove your function works on small input. Say with an if statement: if  $(n == 1)$  return answer.

② Assume your function works for my input

of size  $< n$ , & build the solution for inputs of size  $n$  by calling the same function on smaller inputs ( $< n$ ).

(3) Conclude your function works for all inputs.

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Examples: write a <sup>recursive</sup> function to compute  $n!$ .

```
int fac(int n) {  
    if (n == 0) return 1;  
    return n * fac(n-1);  
}
```

$$\begin{array}{rcll} \text{fac}(4) & = & 4 * \text{fac}(3) & \leftarrow 24 \\ & & \parallel & \\ & & 3 * \text{fac}(2) & \leftarrow 6 \\ & & \parallel & \\ & & 2 * \text{fac}(1) & \leftarrow 2 \\ & & \parallel & \\ & & 1 * \text{fac}(0) & \leftarrow 1 \\ & & \parallel & \\ & & 1 & \end{array}$$

Naively, this looks worse than a loop:

takes a linear amount of stack space

(we have to save all of the other copies of  $n$ , as well as the points where you "left off" (return address).)

(Note: for simple recursive functions like this, an optimizing compiler (e.g. `g++ -O2`) will turn this into a loop for you. C.f. "tail recursion".)

Example: What is the output of  $f(4)$ , where

```
void f(int n) {  
    if (n == 0) {  
        cout << "0 ";  
    }  
    f(n-1);  
    cout << n << " ";  
}
```



