

**Dalhousie University Faculty of Computer Science**  
**Design and Analysis of Algorithms**  
**Assignment 6 CSCI 3110 Due: 15 Nov 2010**

1. (a) Ex. 2.4 p. 71  
(b) Ex. 2.5(f) and (i)  
(c) Ex. 2.12 p. 73
2. (a) Ex. 2.25  
(b) Ex. 2.29
3. (a) Ex. 2.16 p. 73  
(b) Consider an  $m \times n$  array  $A$  of integers. Each row and column of the array are in ascending order:  
 $A[i, j] \leq A[i, j + 1]$ ,  $i = 1 \dots n$ ,  $j = 1 \dots m - 1$  and  $A[i, j] \leq A[i + 1, j]$ ,  $i = 1 \dots n - 1$ ,  $j = 1 \dots m$ .  
Design an algorithm that determines if a given integer  $x$  is in  $A$ , if so, return it's position otherwise return a -1. Explain your algorithm with words and diagram(s). Show that your algorithm terminates and find its worst case complexity. *Hint: Start by comparing  $x$  with  $A[1, m]$*
4. (a) Draw the recursion tree to guess a solution and prove your guess for

$$T(n) = T(n - a) + T(a) + cn$$

for constants  $a \geq 1$  and  $c > 0$ . Assume  $T(n) = \Theta(1)$  for  $n \leq a$ .

- (b) Draw the recursion tree to guess a solution and prove your guess for

$$T(n) = T(\alpha n) + T((1 - \alpha)n) + cn$$

for constants  $0 < \alpha < 1$  and  $c > 0$ .

- (c) Suppose that the splits at each level of quicksort are in the proportion  $1 - \alpha$  to  $\alpha$  where  $0 < \alpha \leq 1/2$  is a constant. Show that the minimum depth of a leaf in the recursion tree is approx.  $-\lg n / \lg \alpha$  and the maximum depth is approx  $-\lg n / \lg(1 - \alpha)$
5. (a) Give an algorithm that sorts an array of size  $n$ , only containing the elements 1,2 and 3, in linear time. Give a plausible argument that your algorithm sorts the array in linear time  
(b) Construct an example input for which quicksort will use  $\Omega(n^2)$  comparisons when the pivot is chosen as the median of the first, last and middle elements of the sequence.