

Construction and Verification of Software

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MIEI - Integrated Master in Computer Science and Informatics
Consolidation block

Lecture 7 - Arrays in Separation Logic

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Verifast Example - Bag

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Breakpoint reached.

Bag0.java _assume.javaspec _list.javaspec _nat.javaspec _quantifiers.javaspec _bitops.javaspec _atomics.javaspec java.lang.javaspec

```
public class Bag {  
    int store[];  
    int nelems;  
  
    /*@  
        predicate BagInv(int n) =  
            store |-> ?s  
            &* & nelems |-> n  
            &* & s != null  
            &* & 0 <= n &* & n <= s.length  
            &* & s[0..n] |-> ?elems  
            &* & s[n..s.length] |-> ?others  
    @*/  
  
    public Bag(int size)  
        //@ requires size >= 0;  
        //@ ensures BagInv(0);  
    {  
        store = new int[size];  
        nelems = 0;  
    }  
  
    boolean add(int v)  
        //@ requires BagInv(_);  
}
```

Local	Value
size	size
this	this

Steps

- Implicit superclass constructor call
- Verifying call
 - Consuming assertion
 - Producing assertion
 - Executing statement
 - Executing statement

Assumptions

- $0 \leq \text{size}$
- !(this = 0)
- $\text{length}(\text{elems}) = \text{size}$
- $\text{all_eq}(\text{elems}, 0)$
- !(array = 0)
- $\text{array.length}(\text{array}) = \text{size}$

Heap chunks

- Bag_nelems(this, 0)
- Bag_store(this, array)
- java.lang.array_slice<int32>(array, 0, s

Verifast Example - Bag

- Fields must be considered in separate heap chunks, pure conditions can be added to assertions and predicates.

```
public class Bag {  
  
    int store[];  
    int nelems;  
  
    /*@  
    predicate BagInv(int n) =  
        store |-> ?s  
        &*& nelems |-> n  
        &*& s != null  
        &*& 0<=n &*& n <= s.length  
        &*& array_slice(store,0,n,?elems)  
        &*& array_slice(store,n,s.length,?others)  
    ;  
    @*/  
    ...  
}
```

Verifast Example - Bag

- Fields must be considered in separate heap chunks, pure conditions can be added to assertions and predicates.

```
public class Bag {  
  
    int store[];  
    int nelems;  
  
    /*@  
    predicate BagInv(int n) =  
        store |-> ?s  
        &*& nelems |-> n  
        &*& s != null  
        &*& 0<=n &*& n <= s.length  
        &*& s[0..n] |-> ?elems  
        &*& s[n..s.length] |-> ?others  
    ;  
    @*/  
    ...  
}
```

Verifast Example - Bag

- Fields must be considered in separate heap chunks, pure conditions can be added to assertions and predicates.

```
int get(int i)
    //@ requires BagInv(?n) &*& 0 <= i &*& i < n;
    //@ ensures BagInv(n);
{
    return store[i];
}

int size()
    //@ requires BagInv(?n);
    //@ ensures BagInv(n) &*& result >= 0 ;
{
    return nelems;
}
```

Verifast Example - Bag

```
public Bag(int size)
    //@ requires size >= 0;
    //@ ensures BagInv(0);
{
    store = new int[size];
    nelems = 0;
}

boolean add(int v)
    //@ requires BagInv(_);
    //@ ensures BagInv(_);
{
    if(nelems < store.length) {
        store[nelems] = v;
        nelems = nelems + 1;
        return true;
    } else {
        return false;
    }
}
```

Verifast Example - Bag

```
public Bag(int size)
    //@ requires size >= 0;
    //@ ensures BagInv(0);
{
    store = new int[size];
    nelems = 0;
}

boolean add(int v)
    //@ requires BagInv(?n);
    //@ ensures BagInv(n+1); // Does not hold, why?
{
    if(nelems < store.length) {
        store[nelems] = v;
        nelems = nelems + 1;
        return true;
    } else {
        return false;
    }
}
```

Verifast Example - Bag

```
public
//@ r
//@ e
{
store
nelem
}

boolean
//@ r
//@ e
{
if(ne
stor
nele
retu
} els
retu
}
```

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Cannot prove dummy == (dummy + 1)

Bag0.java _assume.javaspec _list.javaspec _nat.javaspec _quantifiers.javaspec _bitops.javaspec _atomics.javaspec java.lang.javaspec

```
predicate BagInv(int n) =
  store |-> ?s
  &* & nelems |-> n
  &* & s != null
  &* & 0 <= n &* & n <= s.length
```

Local Value

n	(dummy + 1)
s	s
this	this

Bag0.java _assume.javaspec _list.javaspec _nat.javaspec _quantifiers.javaspec _bitops.javaspec _atomics.javaspec java.lang.javaspec

```
nelems = 0;
}

boolean add(int v)
  //@ requires BagInv(_);
  //@ ensures BagInv(_);
{
  //@ open BagInv(?n);
  if(nelems < store.length) {
    store[nelems] = v;
    nelems = nelems + 1;
    //@ close BagInv(n+1);
    return true;
  } else {
    //@ close BagInv(n+1);
    return false;
  }
}

int get(int i)
```

Local Value

n	dummy
this	this
v	v

Steps

- Executing statement
- Executing second branch
- Executing statement
- Executing statement
- Consuming assertion
- Consuming assertion

Assumptions

- !(this = 0)
- !(this = 0)
- !(s = 0)
- 0 <= dummy
- dummy <= arraylength(s)
- !(s = 0)

Heap chunks

- Bag_nelems(this, dummy)
- java.lang.array_slice<int32>(s, 0, dum
- java.lang.array_slice<int32>(s, dummy

Managing arrays in SL

```
/*@  
predicate AccountInv(Account a;int b) = a.balance |-> b &*& b >= 0;  
@*/
```

```
public class Account {  
    int balance;  
  
    public Account()  
    //@ requires true;  
    //@ ensures AccountInv(this,0);  
    {  
        balance = 0;  
    }  
    ...  
}
```

Managing arrays in SL

- The bank holds an array of accounts...

```
public class Bank {  
    Account store[];  
    int nelems;  
    int capacity;  
  
    Bank(int max)  
    {  
        nelems = 0;  
        capacity = max;  
        store = new Account[max];  
    }  
    ...  
}
```

Managing arrays in SL

- And implements a couple of operations...

```
public class Bank {  
  
    Account store[];  
    int nelems;  
    int capacity;  
  
    ...  
    Account retrieveAccount()  
    {  
        Account c = store[nelems-1];  
        store[nelems-1] = null;  
        nelems = nelems-1;  
        return c;  
    }  
    ...  
}
```

Managing arrays in SL

- And implements a couple of operations...

```
public class Bank {  
  
    Account store[];  
    int nelems;  
    int capacity;  
  
    ...  
    void addnewAccount()  
    {  
        Account c = new Account();  
        store[nelems] = c;  
        nelems = nelems + 1;  
    }  
    ...  
}
```

Managing arrays in SL

```
/*@  
predicate AccountP(unit a, Account c; unit b) = AccountInv(c, ?n) &*& b == unit;  
@*/
```

```
public class Bank {
```

```
/*@  
predicate BankInv(int n, int m) =  
    this.nelems |-> n &*&  
    this.capacity |-> m &*&  
    m > 0 &*&  
    this.store |-> ?accounts &*&  
    accounts.length == m &*&  
    0 <= n &*& n <= m &*&  
    array_slice_deep(accounts, 0, n, AccountP, unit, _, _) &*&  
    array_slice(accounts, n, m, ?rest) &*& all_eq(rest, null) == true;  
@*/  
}
```

array slice assertions

```
predicate array_slice<T>(
    T[] array,
    int start,
    int end;
    list<T> elements);
```

- `array_slice(a,s,l,v)`:
 - represents the footprint of array `a[s..l-1]`
- `v` is a list of the array “values” `v_i` such that `a[i] l-> v_i`
- `v` is an immutable pure value (like a OCaml list)
- `array_slice(a,s,l,v)` is equivalent to the assertion
 - $v = \{ v_s, \dots v_{l-1} \}$
 - $a[s] \text{ l-> } v_s \ \&^* \& \ a[s+1] \text{ l-> } v_{s+1} \ \&^* \& \dots \ \&^* \& \ a[l-1] \text{ l-> } v_{l-1}$

array slice assertions

```
predicate array_slice_deep<T, A, V>(
    T[] array,
    int start,
    int end,
    predicate(A, T; V) p,
    A info;
    list<T> elements,
    list<V> values);
```

- `array_slice_deep(a, s, l, P, info, v, s):`
as in the (simple) `array_slice`
`v` is the list of the array “values” `v_i` such that `a[i] l-> v_i`,
the predicate `P(info, v_i; o_i)` holds for each `v_i`
and `s` is the list of all values `o_i`

Managing arrays in SL

```
public class Bank {  
  
    Account store[];  
    int nelems;  
    int capacity;  
  
    Bank(int max)  
    //@ requires max>0;  
    //@ ensures BankInv(0,max);  
    {  
        nelems = 0;  
        capacity = max;  
        store = new Account[max];  
    }  
    ...  
}
```


Managing arrays in SL

```
public class Bank {  
  
    Account store[];  
    int nelems;  
    int capacity;  
  
    Account retrieveLastAccount()  
    //@ requires BankInv(?n,?m) &*& n>0;  
    //@ ensures  BankInv(n-1,m) &*& AccountInv(result,_);  
    {  
        Account c = store[nelems-1];  
        store[nelems-1] = null;  
        // code does not compile without this! Why ?  
        nelems = nelems-1;  
        return c;  
    }  
}
```

Managing arrays in SL

```
public class Bank {  
  
    Account store[];  
    int nelems;  
    int capacity;  
  
    void addnewAccount()  
    //@ requires BankInv(?n,?m) &*& n < m;  
    //@ ensures  BankInv(n+1,m);  
    {  
        Account c = new Account();  
        store[nelems] = c;  
        //@ array_slice_deep_close(store, nelems, AccountP, unit);  
        nelems = nelems + 1;  
    }  
}
```

array slice “lemmas”

```
lemma void array_slice_deep_close<T, A, V>(
    T[] array, int start, predicate(A, T; V) p, A a);
requires
    array_slice<T>(array, start, start+1, ?elems) &*&
    p(a, head(elems), ?v);
ensures
    array_slice_deep<T,A,V>(array, start, start+1, p, a, elems, cons(v,nil));
```

- incorporates the spec of an array element in a (singleton) slice spec into a (singleton) slice_deep spec
- there are other lemmas, that join together slices
- verifast is usually able to apply lemmas automatically, but not always, in that case the programmer needs to “help”, by calling the needed lemmas.

array slice “lemmas”

```
lemma void array_slice_split<T>(T[] array, int start, int start1);
```

requires

```
array_slice<T>(array, start, ?end, ?elems) &*&
```

```
start <= start1 &*& start1 <= end;
```

ensures

```
array_slice<T>(array, start, start1, take(start1 - start, elems)) &*&
```

```
array_slice<T>(array, start1, end, drop(start1 - start, elems)) &*&
```

```
elems == append(take(start1 - start, elems), drop(start1 - start, elems))
```

- this “lemma” splits one array slice assertion into two (sub) array slice assertions.

array slice “lemmas”

← → ↻ 🏠 Secure | <https://people.cs.kuleuven.be/~bart.jacobs/verifast/examples/rt/Object.javaspec.html>

```
package java.lang;

import java.util.*;

/*@

inductive unit = unit;

inductive pair<a, b> = pair(a, b);

fixpoint a fst<a, b>(pair<a, b> p) {
  switch (p) {
    case pair(x, y): return x;
  }
}

fixpoint b snd<a, b>(pair<a, b> p) {
  switch (p) {
    case pair(x, y): return y;
  }
}

fixpoint t default_value<t>();

inductive boxed_int = boxed_int(int);
fixpoint int unboxed_int(boxed_int i) { switch (i) { case boxed_int(value): return value; } }

inductive boxed_bool = boxed_bool(boolean);
fixpoint boolean unboxed_bool(boxed_bool b) { switch (b) { case boxed_bool(value): return value; } }

predicate array_element<T>(T[] array, int index; T value);
predicate array_slice<T>(T[] array, int start, int end; list<T> elements);
predicate array_slice_deep<T, A, V>(T[] array, int start, int end, predicate(A, T; V) p, A info; list<T> elements

lemma_auto void array_element_inv<T>();
  requires [?f]array_element<T>(?array, ?index, ?value);
  ensures [f]array_element<T>(array, index, value) &*& array != null &*& 0 <= index &*& index < array.length;
```