

Boğaziçi University
Department of Industrial Engineering
IE 313 Supply Chain Management
Assignment 3

There are various formulations of the travelling salesman problem. When $x_{ij} = 1$ if the salesperson goes from city i directly to city j ; 0 otherwise (x_{ii} are not defined or are very large numbers), then the formulation can be developed as:

$$\begin{aligned} (1) \quad & \min \quad \sum_{i=1}^n \sum_{j=1}^n c_{ij} x_{ij} \\ (2) \quad & \text{s.t.} \quad \sum_{j:j \neq i} x_{ij} = 1 \quad j = 1, 2, \dots, n \\ (3) \quad & \sum_{i:i \neq j} x_{ij} = 1 \quad i = 1, 2, \dots, n \\ & x_{ij} \in \{0, 1\} \quad \forall i, j \end{aligned}$$

This is an assignment model. The solution may lead to *sub-tours*. To avoid sub tours let S be a non-empty strict subset of $N = \{1, 2, \dots, n\}$ and add the following constraint.

$$(4) \quad \sum_{i \in S} \sum_{j \in S} x_{ij} \leq |S| - 1 \text{ for } S \subset N, 2 \leq |S| \leq n - 1$$

Constraints (2), (3), (4) (the sub tour elimination constraint), together with binary restrictions on the variables, and objective (1) is a valid formulation of the TSP. The trouble with this formulation is there are too many sub tour elimination constraints to write. A strategy that is usually employed is to solve the assignment model, detect sub tours, add only relevant sub tour elimination constraints and continue in this manner until you get a solution without a sub tour.

An alternative formulation is the sequential formulation of Miller, Tucker, Zemlin (1960) (MTZ). They define a new decision variable u_i to indicate the sequence in which city i is visited ($i = 2, \dots, n$). They replace the sub tour elimination constraints (4) by:

$$(5) \quad u_i - u_j + nx_{ij} \leq n - 1, \quad i \neq j, i, j = 2, 3, \dots, n$$

(1) subject to (2), (3), (5) together with binary restrictions on x_{ij} variables, nonnegative u_i 's, is also a valid formulation of the TSP. A slightly stronger version of the MTZ constraints is given by Desrochers and Laporte (1991) as:

$$(6) \quad u_i - u_j + (n - 1)x_{ij} + (n - 3)x_{ji} \leq n - 2, \quad i \neq j, i, j = 2, 3, \dots, n$$

which yields (1) subject to (2), (3), (6). These formulations solve small instances of the TSP efficiently but may not be able to give a solution for large instances (since its linear programming relaxation is weak and it may take a long time for an IP solver to find an integer solution).

Given this background,

- (1) Choose a formulation and a strategy to solve TSP. Download TR81 cities data from the course web page.
- (2) Solve the TSP for randomly chosen 15 cities. Report your findings.
- (3) Solve the TSP for 81 cities. Report your findings.
- (4) There is a 5-point bonus if you find the optimal TSP tour and plot it on the map of Turkey.