

On the directed Oberwolfach problem with tables of even lengths

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A simple example

The setting: Consider a conference with 16 participants. To facilitate networking, the organizing committee decides to host 15 banquets. The banquet hall has 3 tables that seat 4, 4, and 8 participants.

The problem: The organizing committee needs a set of 15 seating arrangements (one for each banquet) such that each participant is seated to the right of every other participants exactly once.

Is this possible?

Construction of a seating arrangement

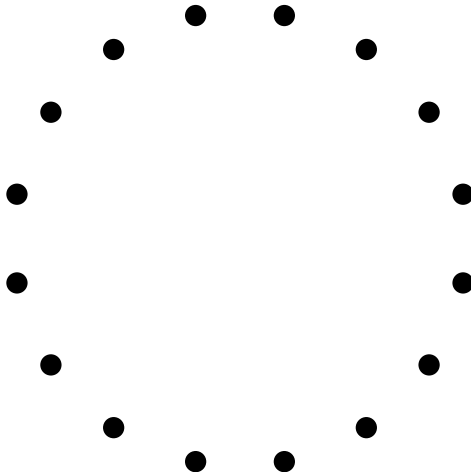


Figure: The 16 participants (one for each vertex).

Construction of a seating arrangement

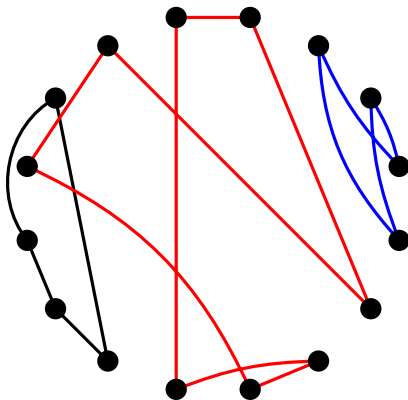


Figure: One seating arrangement with two tables of length 4 and one table of length 8.

Construction of a seating arrangement

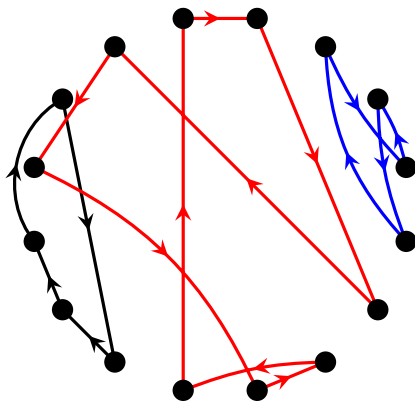


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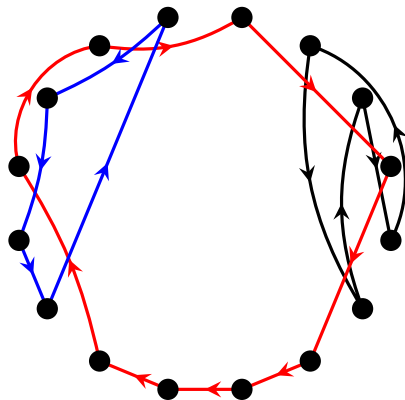


Figure: Another seating arrangement with two tables of length 4 and one table of length 8.

The directed Oberwolfach problem

The setting: Consider a conference with n participants. To facilitate networking, the organizing committee decides to host $n - 1$ banquets. The banquet hall has t round tables that sit m_1, m_2, \dots, m_t participants such that $m_1 + m_2 + \dots + m_t = n$.

The problem: The organizing committee needs a set of $n - 1$ seating arrangements (one for each banquet) such that each participant is seated **to the right** of every other participants exactly once.

Is this possible?

The complete symmetric digraph

Definition

The **complete symmetric digraph**, denoted K_n^* , is the digraph on n vertices in which for every pair of distinct vertices x and y , there are arcs (x, y) and (y, x) .

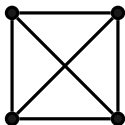


Figure: The complete graph K_4 .

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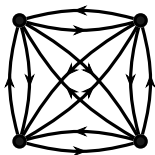


Figure: The complete symmetric digraph K_4^* .

2-factorization

Definition

A **directed** $[m_1, m_2, \dots, m_t]$ -**factor** of digraph G is a spanning subdigraph comprised of disjoint directed cycles of length m_1, m_2, \dots, m_t .

Definition

A **directed** $[m_1, m_2, \dots, m_t]$ -**factorization** of digraph G is a decomposition of G into $[m_1, m_2, \dots, m_t]$ -factors.

The graph-theoretic formulation of the directed OP

Problem $(OP^*(m_1, m_2, \dots, m_t))$

Let $m_1, m_2, \dots, m_t \geq 2$. If $m_1 + m_2 + \dots + m_t = n$, does K_n^* admit a directed $[m_1, m_2, \dots, m_t]$ -factorization?

When all tables are of length m , we write $OP^*(m^t)$.

A solution to $OP^*(m_1, m_2, \dots, m_t)$ is also a resolvable Mendelsohn design with blocks of size m_1, m_2, \dots, m_t .

Background (free solutions!)

Corollary (Kadri and Šajna (2025))

If the original Oberwolfach problem with n participants and tables of lengths $3 \leq m_1 \leq m_2 \leq \dots \leq m_t$ with n being odd, then $OP^(m_1, m_2, \dots, m_t)$ also has a solution.*

Background

Theorem (Bermond, Germa, and Sotteau (1979); Tillson (1980), Bennett and Zhang (1990); Adams and Bryant (Unpublished); Abel, Bennett, and Ge (2002); Burgess and Šajna (2014); Burgess, Francetić, and Šajna (2018); L-M (2024))

The $OP^(m^t)$ has a solution except when
 $(m, t) \notin \{(3, 2), (4, 1), (6, 1)\}$.*

The directed OP has been completely resolved when all tables are of the same length.

Background

Theorem (Bermond, Germa, and Sotteau (1979); Tillson (1980), Bennett and Zhang (1990); Adams and Bryant (Unpublished); Abel, Bennett, and Ge (2002); Burgess and Šajna (2014); Burgess, Francetić, and Šajna (2018); L-M (2024))

The $OP^(m^t)$ has a solution except when*

$$(m, t) \notin \{(3, 2), (4, 1), (6, 1)\}.$$

Theorem (Kadri and Šajna (2025), Quirion (Unpublished))

The $OP^(m_1, m_2, \dots, m_t)$ has a solution when $n \leq 17$*

Background

Theorem (Zhang and Du (2005))

The $OP^(3^t, s)$ has a solution when $s \in \{4, 5\}$ for all $t \in \mathbb{Z}^+$.*

Theorem (Kadri and Šajna (2025))

Let $m_1 < m_2$. The $OP^(m_1, m_2)$ has a solution except possibly when $m_1 \in \{4, 6\}$ and m_2 is even.*

Theorem (Horsley and L-M (2024+))

Let $m_1 < m_2$. The $OP^(m_1, m_2)$ has a solution when $m_1 \in \{4, 6\}$ and $m_2 \geq 8$ is even.*

Bipartite 2-factors

The original Oberwolfach problem has been completely resolved for bipartite 2-factorizations.

Theorem (Häggkvist (1985))

The graph $K_{2n} - I$ admits an $[m_1, m_2, \dots, m_t]$ -factorization when each m_i is even, $m_i \geq 4$, and $2n \equiv 2 \pmod{4}$.

Note that $K_{2n} - I$ is the complete graph on $2n$ vertices with a 1-factor removed.

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Theorem (Bryant and Danziger (2011))

The graph $K_{2n} - I$ admits an $[m_1, m_2, \dots, m_t]$ -factorization when each m_i is even, $m_i \geq 4$, and $2n \equiv 0 \pmod{4}$.

Note that $K_{2n} - I$ is the complete graph on $2n$ vertices with a 1-factor removed.

Our results

Theorem (Burgess, Danziger, and L-M (2025+))

Let m_1, m_2, \dots, m_t be positive even integers and $2n \equiv 2 \pmod{4}$, then $OP^(m_1, m_2, \dots, m_t)$ has a solution except for $OP^*(6)$.*

Strategy

Step 1: Decompose K_{2n}^* into $n - 2$ spanning subdigraphs that fall into one of two isomorphisms classes G_1 and G_2 .

Step 2: Show that G_1 and G_2 both admit a $[m_1, m_2, \dots, m_t]$ -factorization.

First class of digraphs

Below, we show a spanning digraph of K_{14}^* .

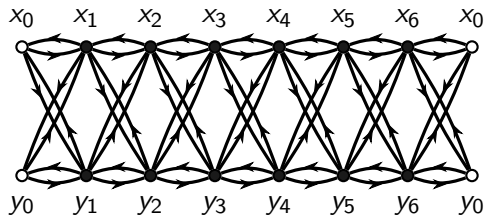


Figure: The first directed graph $G_1 = \vec{C}_7[2]$.

First class of digraphs

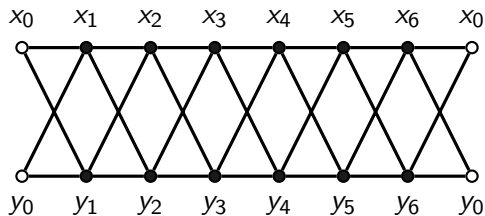


Figure: The underlying graph of $\vec{C}_7[2]$ written $C_7[2]$.

A directed version of Häggkvist Lemma

Theorem (Häggkvist Lemma (Häggkvist (1985)))

Let m_1, m_2, \dots, m_t be even integers such that $m_i \geq 4$. The graph $C_r[2]$ admits an undirected $[m_1, m_2, \dots, m_t]$ -factorization.

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Corollary

Let m_1, m_2, \dots, m_t be even integers such that $m_i \geq 4$. The graph $\vec{C}_r[2]$ admits a $[m_1, m_2, \dots, m_t]$ -factorization.

We still need to derive a solution that includes tables of lengths two!

A directed version of Häggkvist Lemma

Lemma

Let m_1, m_2, \dots, m_t be even integers such that $m_i \geq 2$. The graph $\vec{C}_r[2]$ admits a $[m_1, m_2, \dots, m_t]$ -factorization.

Proof: Case 1: We have $k \geq 2$ tables of length 2.

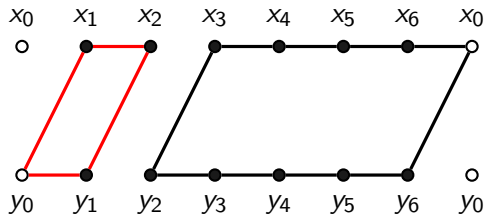


Figure: A undirected $[4, 10]$ -factor of $C_7[2]$.

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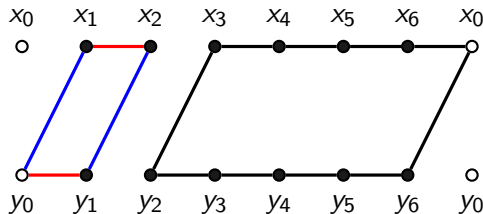


Figure: A 1-factorization of C_4 in blue and grey.

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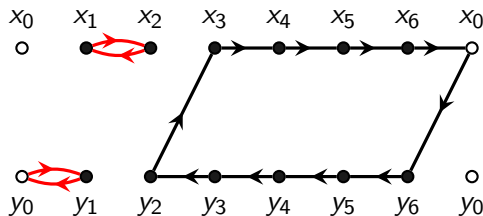


Figure: A directed $[2, 2, 10]$ factor of $\vec{C}_6[2]$.

A directed version of Häggkvist Lemma

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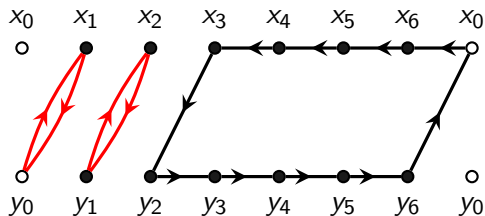


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A directed version of Häggkvist Lemma

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Proof: Case 2: We have one table of length 2.

Here an explicit construction is needed. The construction follows a similar reasoning as Häggkvist's lemma.



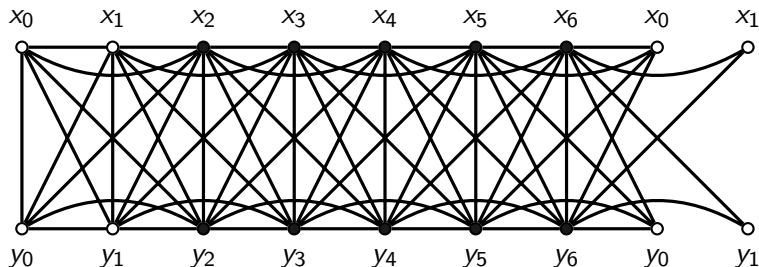
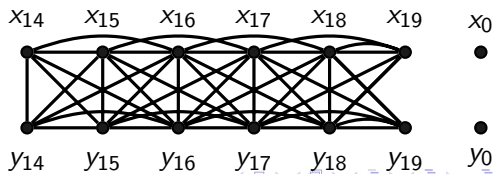
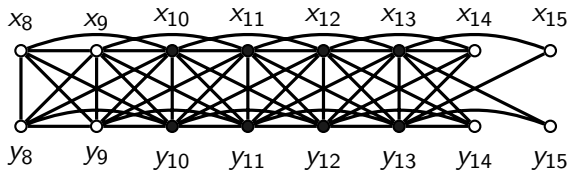
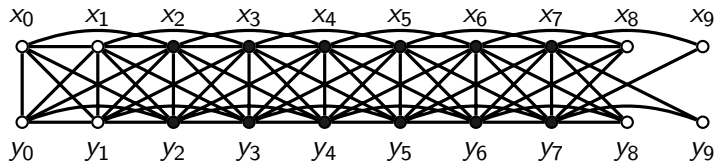
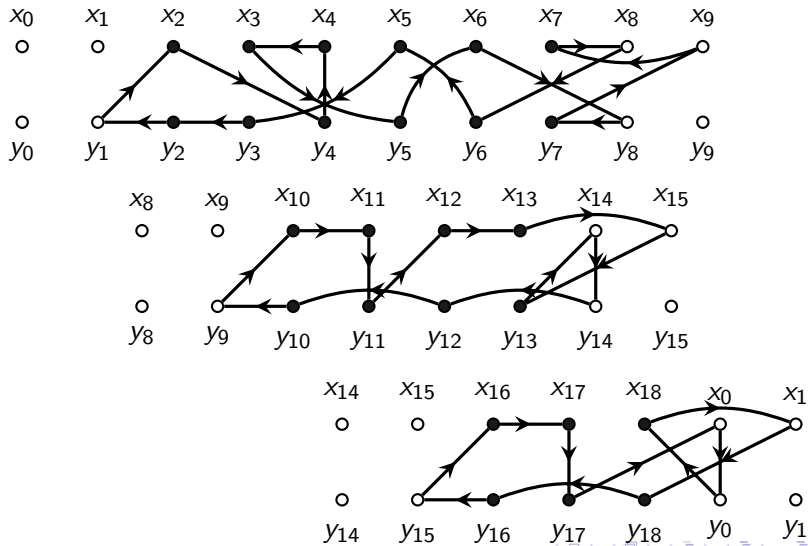
The second spanning subdigraph of K_{2n}^* 

Figure: The underlying graph of G_2 .

Each edge represents a pair of arcs, one for each direction. This underlying graph can be described as a wreath product (lexicographic product) of the circulant $X(n, \{\pm 1, \pm 2\})$ with K_2

The construction of a $[16, 12, 10]$ -factor of K_{38}^* 

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Long tables

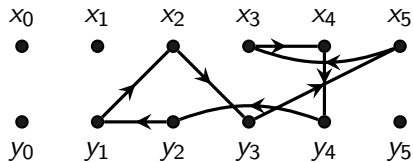
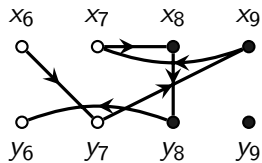
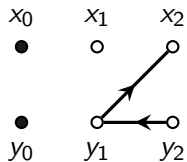
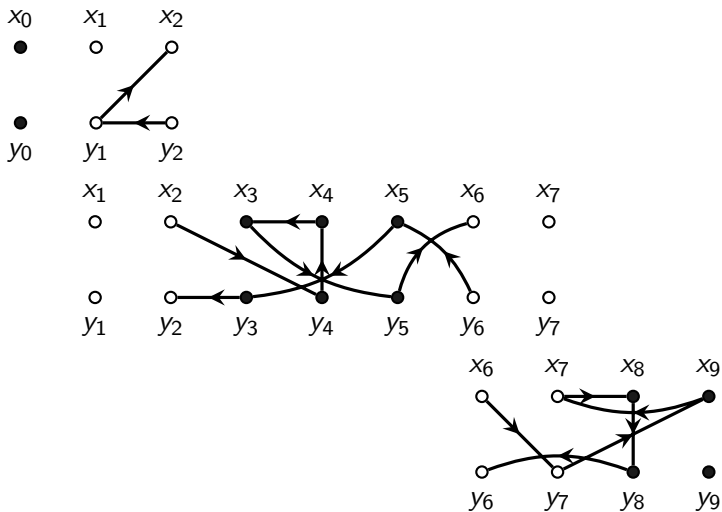


Figure: A cycle of length 8.

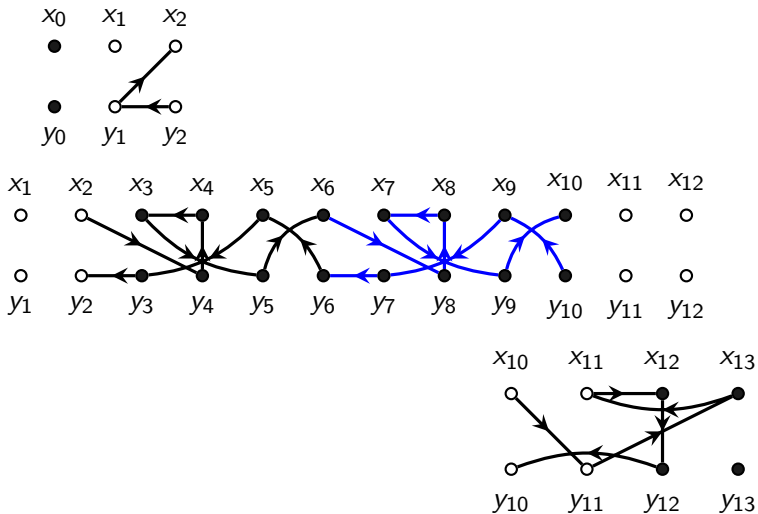
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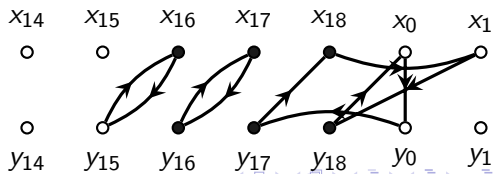
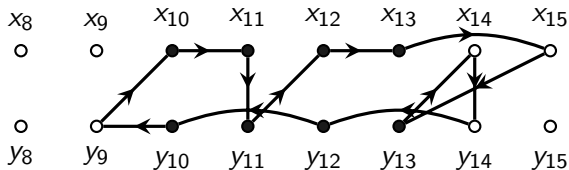
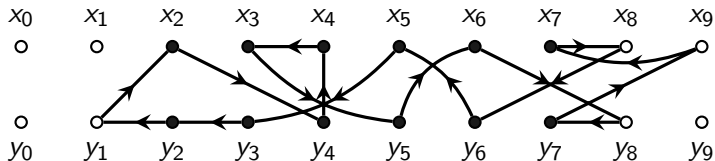
Long tables



Long tables



Special cases- A $[16,12,6,2,2]$ -factor



Future work

Problem: Can we take a similar approach for the case $2n \equiv 0 \pmod{4}$?

Answer: Yes. However, the second spanning digraph in the decomposition of K_{2n}^* is different.

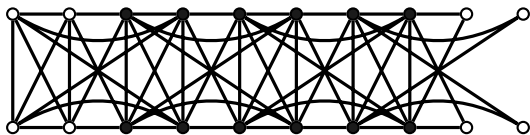


Figure: The underlying graph of G_2 .

Thank you!

The 2026 Prairie Discrete Math Workshop:

- Set to take place on May 7th and 8th in Regina;
- Students and post-docs welcome! (Some travel funding may be available)