

qliciq sang's PCS229 PORTFOLIO

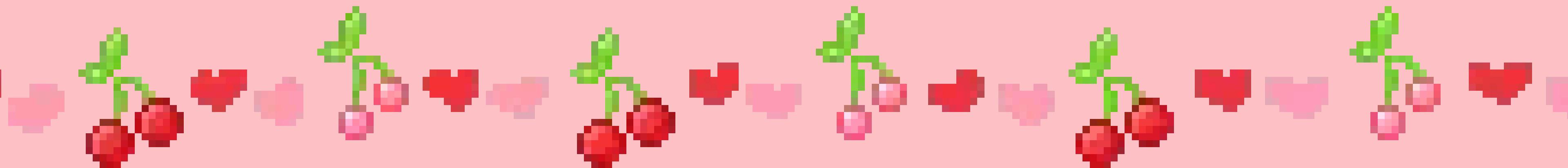


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OOP, CLASSES, & INHERITANCE

Class: Blueprint for objects

Object: Instance of a class

Inheritance: Allows subclasses to use or override methods/attributes from a parent class

Encapsulation: Hide internal state, expose functionality via public methods

Polymorphism: Same interface, different underlying forms

BIG O

Object creation: $O(1)$

Method calls: $O(1)$

```
class Student:  
    def __init__(self, name):  
        self.name = name  
    def greet(self):  
        print(f"Hello, my name is  
{self.name}.")
```

- Defined Shape, Triangle, Rectangle with `__slots__`
- Overrode methods like `foobar()` and `__str__()`
- Demonstrated polymorphism using `processShape(shape)`
- Used `__eq__` to compare objects
- i.e. class: Dog, subclasses: different breeds of dogs
(GeeksForGeeks.org)

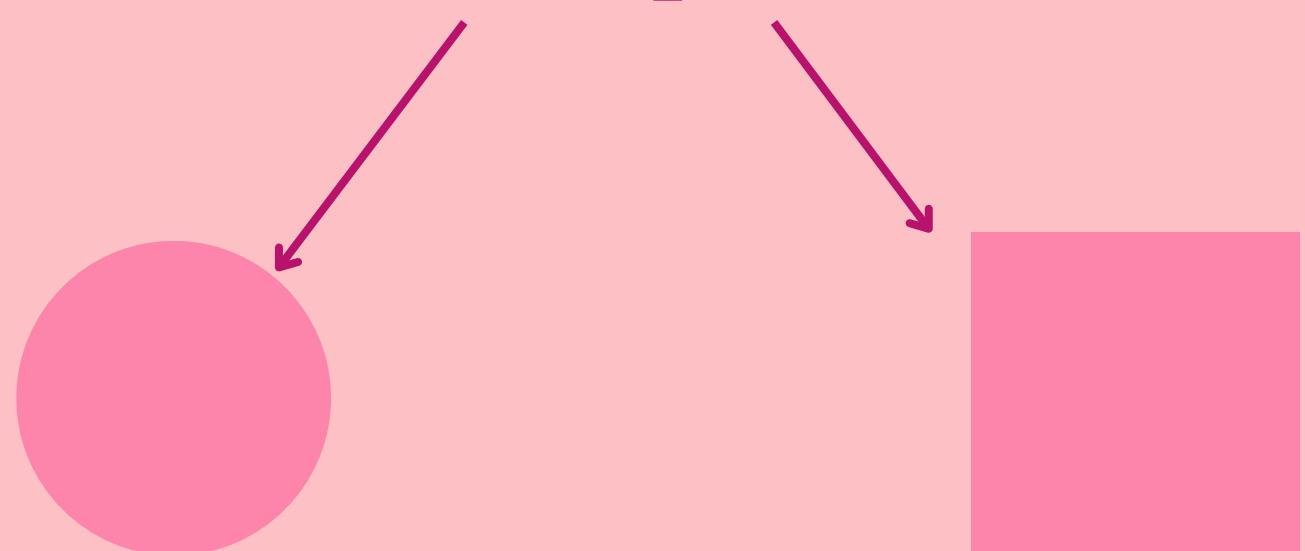
Practice question

Q: Define a Book class with attributes title, author, and isbn, and a method `__eq__` to compare two books based on ISBN. Write the class definition and `__eq__` method.

A:

```
class Book:  
    def __init__(self, title, author, isbn):  
        self.title = title  
        self.author = author  
        self.isbn = isbn  
  
    def __eq__(self, other):  
        return self.isbn == other.isbn
```

Shape



BINARY TREE SEARCH

BSTs maintain order: $\text{left} < \text{node} < \text{right}$

Operations include: `insert(e)`, `delete(e)`, `search(e)`

Can become **unbalanced** and degrade to linked-list behavior

BIG O

Average case: $O(\log n)$

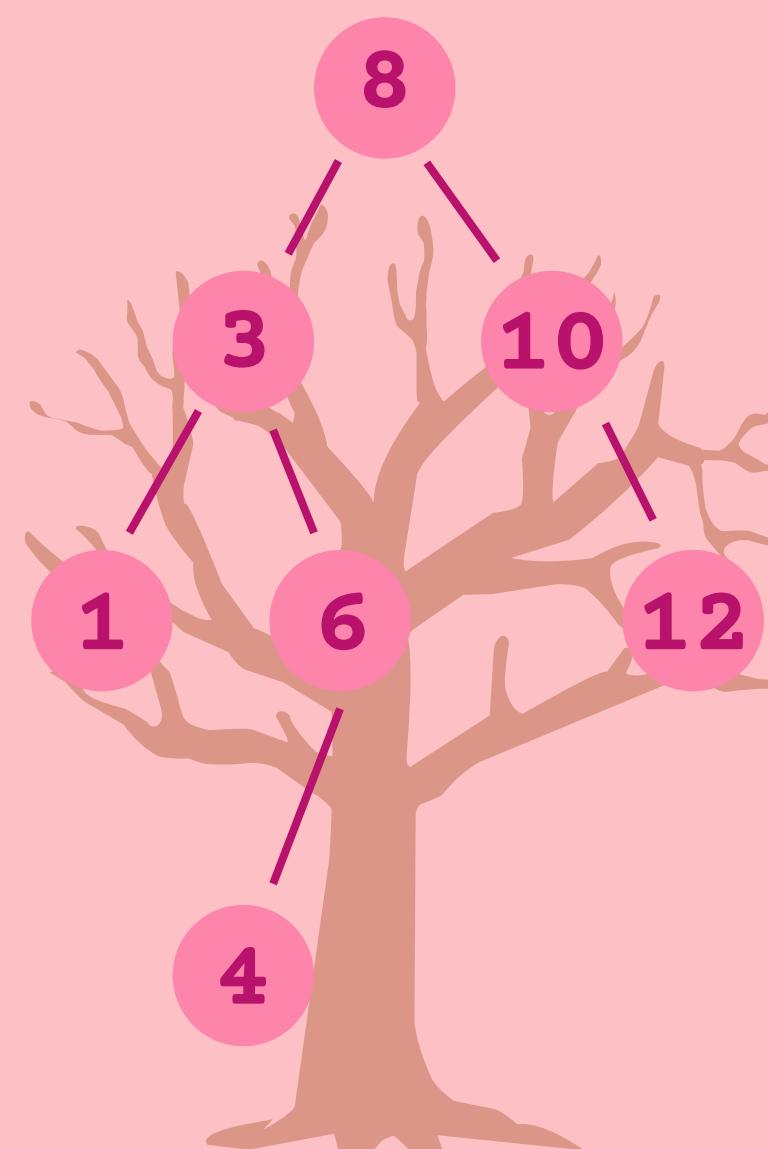
Worst case (unbalanced): $O(n)$

- Wrote insert and search functions for BSTs
- Compared balanced vs unbalanced trees
- Practiced inorder, preorder, and postorder traversal

practice question

Q: What's the worst-case time complexity for searching in a BST, and when does it occur?

A: $O(n)$, when the tree is unbalanced and degenerate (e.g., each node only has a right child).



BIG-O ALGORITHM ANALYSIS

Big O: Describes algorithm efficiency

Amortized Analysis: Average performance over time (e.g., dynamic array resizing)

BIG O

Sequential search: $O(n)$

Binary search: $O(\log n)$

Merge sort: $O(n \log n)$

Append to dynamic array: **Amortized $O(1)$**

- Compared selection sort vs merge sort vs binary search
- Counted operations to estimate runtime
- Analyzed Python list `.append()` using amortized analysis

```
def find_item(my_list, item):  
    for i in range(len(my_list)):  
        if my_list[i] == item:  
            return i  
    return None
```

practice question

Q: What is the Big-O runtime of the following function, and why?

```
def sum_pairs(nums):  
    for i in nums:  
        for j in nums:  
            print(i + j)
```



A: $O(n^2)$: Nested loop over the same list causes quadratic growth in operations as input size increases.

STACK

Operations: push(e), pop(), top(), is_empty(), len()

- Built a custom Stack class
- Used stack for DFS in maze lab
- Related stacks to Python's call stack for recursion

BIG O

Push/Pop/Top: O(1)

```
class Stack:  
    def __init__(self):  
        self.data = []  
    def push(self, e):  
        self.data.append(e)  
    def pop(self):  
        return self.data.pop()
```



**you can only eat the
pancake on the top, it
would be hard to eat the
one on the bottom first!**

Practice question

Q: What is the output of this code?

```
s = Stack()  
s.push(1)  
s.push(2)  
print(s.pop())  
s.push(3)  
print(s.top())
```

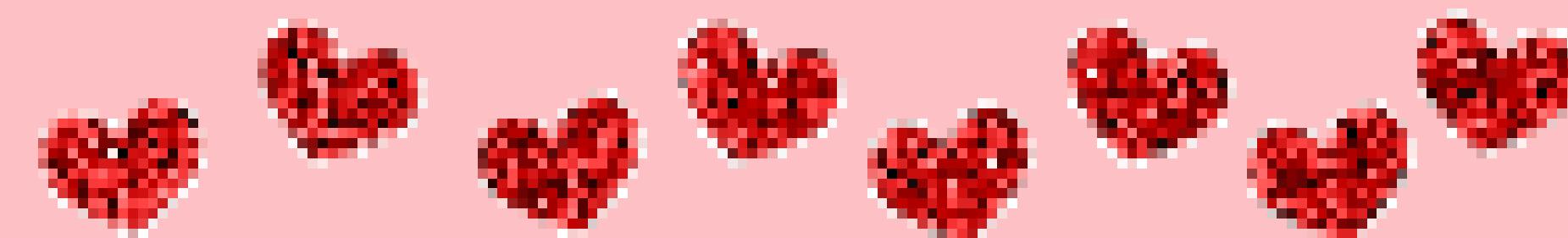
A:

2
3

Because stack is LIFO: 2 is popped, then 3 is the new top.

Uses:

- When a function runs, the computer uses a stack to store its info (like local variables).
- Once the function finishes, that info is removed (iQuanta).



QUEUE

Operations: enqueue(e), dequeue(), front(), is_empty()

BIG O

Enqueue/Dequeue: $O(1)$

- Implemented a queue with deque
- Used queue for BFS in maze lab to find shortest paths

Uses:

To run multiple programs simultaneously
(GeeksForGeeks)

it would only be fair if
the girl in the pink got
on the bus first!



Practice question

Q: You are implementing a queue using a Python list. Which operation is slower and why: enqueue() using append() or dequeue() using pop(0)?

A: dequeue() using pop(0) is slower – $O(n)$ – because it requires shifting all remaining elements left after removal.

DEQUE

Operations: add_first(e), add_last(e), delete_first(), delete_last()
All ops: $O(1)$

- Used deque to check palindromes
- Applied deque in flexible BFS traversal scenarios

Uses:

Cache: most and least recently used both accessible (Quora)



look! you can drink the juice from the top AND the bottom!

Practice question

Q: Consider a deque implemented via a doubly linked list. What is the Big-O of inserting at both ends, and why?

A: $O(1)$ for both add_first() and add_last() – because insertion only requires updating a constant number of pointers at the head or tail.

PRIORITY QUEUE

Operations: insert(e), remove_min(), min()

BIG O

Insert: $O(\log n)$

Remove Min: $O(\log n)$

- Used priority queue (min-heap) for A* in maze lab
- Compared cost functions: $g(n)$, $h(n)$, $f(n)$

NAME	FLIGHT	GATE	SEAT	DEPARTURE	NAME	FLIGHT	GATE SEAT	DEPARTURE	QR CODE
FLYING Company									

Practice question

Q: You are maintaining a priority queue using a min-heap. What is the cost of removing the highest-priority item and why?

A: $O(\log n)$ – removing the root requires bubbling down the last element to restore the heap-order property.

HEAPS

- Explored min-heap operations with heapq
- Tracked heap structure during A* pathfinding

```
# Insert: Add to end, bubble up  
# Remove min: Replace root with last, bubble down
```

Insert: $O(\log n)$,
Remove Min: $O(\log n)$,
Heapify: $O(n)$

Uses:
Scheduling
priorities
(StackOverflow)

Practice question

Q: Given a min-heap represented as an array: [2, 5, 3, 9, 6], show the array after inserting 1 and then reheapifying.

A:

Insert 1 → [2, 5, 3, 9, 6, 1]
Bubble up → [1, 5, 2, 9, 6, 3]



HASHING

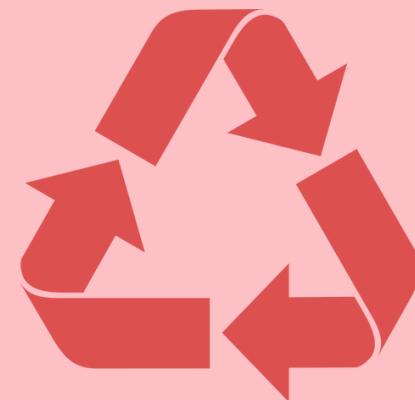
Operations: insert(k, v), get(k), delete(k)

BIG O $O(1)$ average,
 $O(n)$ worst

- Used dictionaries to track visited positions in maze
- Learned how hash collisions can degrade performance

Uses:

Password managers, compare password input to expected password (builtin)



practice question

Q: Explain why Python's `in` operator on a dictionary is $O(1)$ on average. What would cause it to degrade?

A: Average case: $O(1)$ due to direct indexing via hash functions. Worst case: $O(n)$ if many collisions occur and all items land in the same bucket (poor hash function or hash flooding).

RECURSION & MERGE SORT

BIG O

$O(n \log n)$ in all cases

- Wrote recursive merge sort
- Used recursion in DFS

Practice question

Q: What is the time complexity?

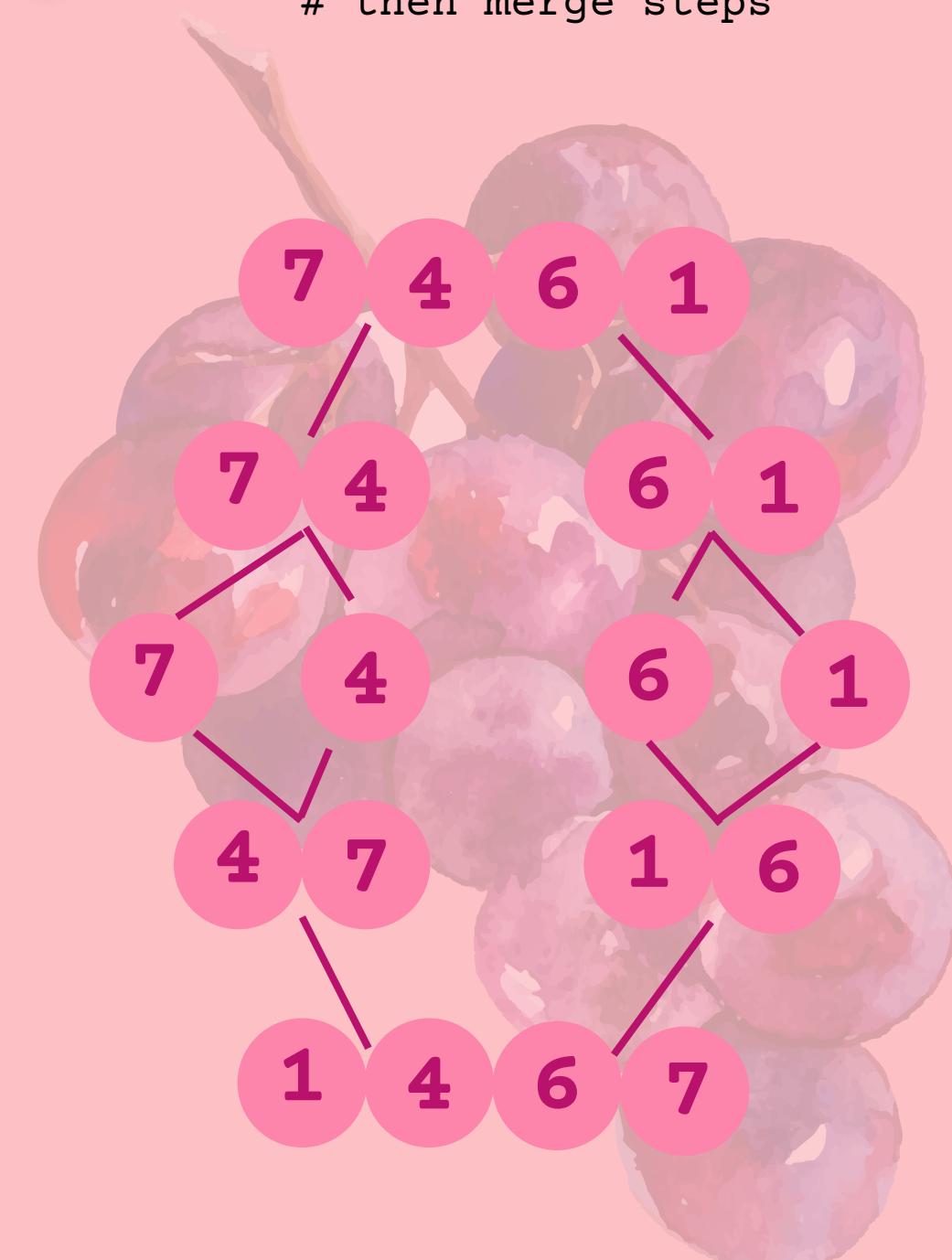
```
def func(n):
    if n <= 1:
        return 1
    return func(n-1) + func(n-1)
```

A: $O(2^n)$: Each call branches into two more calls, leading to exponential growth.

Uses:

Alphabetizing a pile of resumes (Medium).

```
def merge_sort(arr):
    if len(arr) > 1:
        mid = len(arr)//2
        L = arr[:mid]
        R = arr[mid:]
        merge_sort(L)
        merge_sort(R)
    # then merge steps
```



MAZE LAB

DFS: Uses a stack. Goes deep first. Fast, but not guaranteed to find the shortest path.

BFS: Uses a queue. Explores level by level. Guaranteed to find the shortest path.

A*: Uses a priority queue (via a heap). Combines cost-so-far ($g(n)$) + estimated cost to goal ($h(n)$) to make smart choices.

Manhattan Distance: Used as a heuristic $h(n)$ for A*. It estimates how far we are from the goal.

Explored Dictionary: Stored previously visited positions as keys (based on Position objects). This made lookup fast and avoided revisiting cells.

Pitfalls:

- Forgetting to check if a neighbor is already explored before pushing/enqueuing
- Not updating the cost properly in A*
 - DFS is simple and fast but not optimal
 - BFS is great when all steps cost the same and you need the shortest path
 - A* is ideal when steps have costs or when a heuristic helps prioritize

Practice question

Q: Why is A* faster and more effective than BFS in many maze-solving tasks?

A: Because it uses a heuristic (like Manhattan distance) to guess which paths are more promising, which reduces unnecessary exploration.

References

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