

Peer-to-Peer Systems and Security

Team 23, SS21
Midterm Report

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July 2021

1 Changes to assumptions made in the initial report

We continued to work on our understanding of Gossip and at the same time refined the flow chart, which we also used in the initial report. Since we kept major previous versions of it, it is possible to see how the project matured over the last months. The current version can be seen in figure 1.

1.1 Asynchronous I/O

Since the initial report, we replaced all uses of multi threading with Pythons AsyncIO (Asynchronous I/O). This helped to avoid race conditions without the need to implement potentially more complex synchronization. Additionally, AsyncIO is very pleasant to use.

1.2 Variables in config

We revised our understanding of some of the variables regarding peer connections in the config. Even though there was no definition given in the initial report, this had influence on the current version of our flow chart, which was given in the initial report. We removed `cache_size`, since its definition in the specification and in an answer on Moodle (27. may) were perceived as confusing and inconsistent. We decided to use `min_connections` instead, combined with `max_connections` and `degree`, which are specified in section 2.1.4.

2 Architecture

2.1 Logical structure

To provide a clear and understandable architecture, we decided to use a class based approach. A representation of these classes can be seen in figure 2. Besides these classes we also used smaller modules, which provide some functionality,

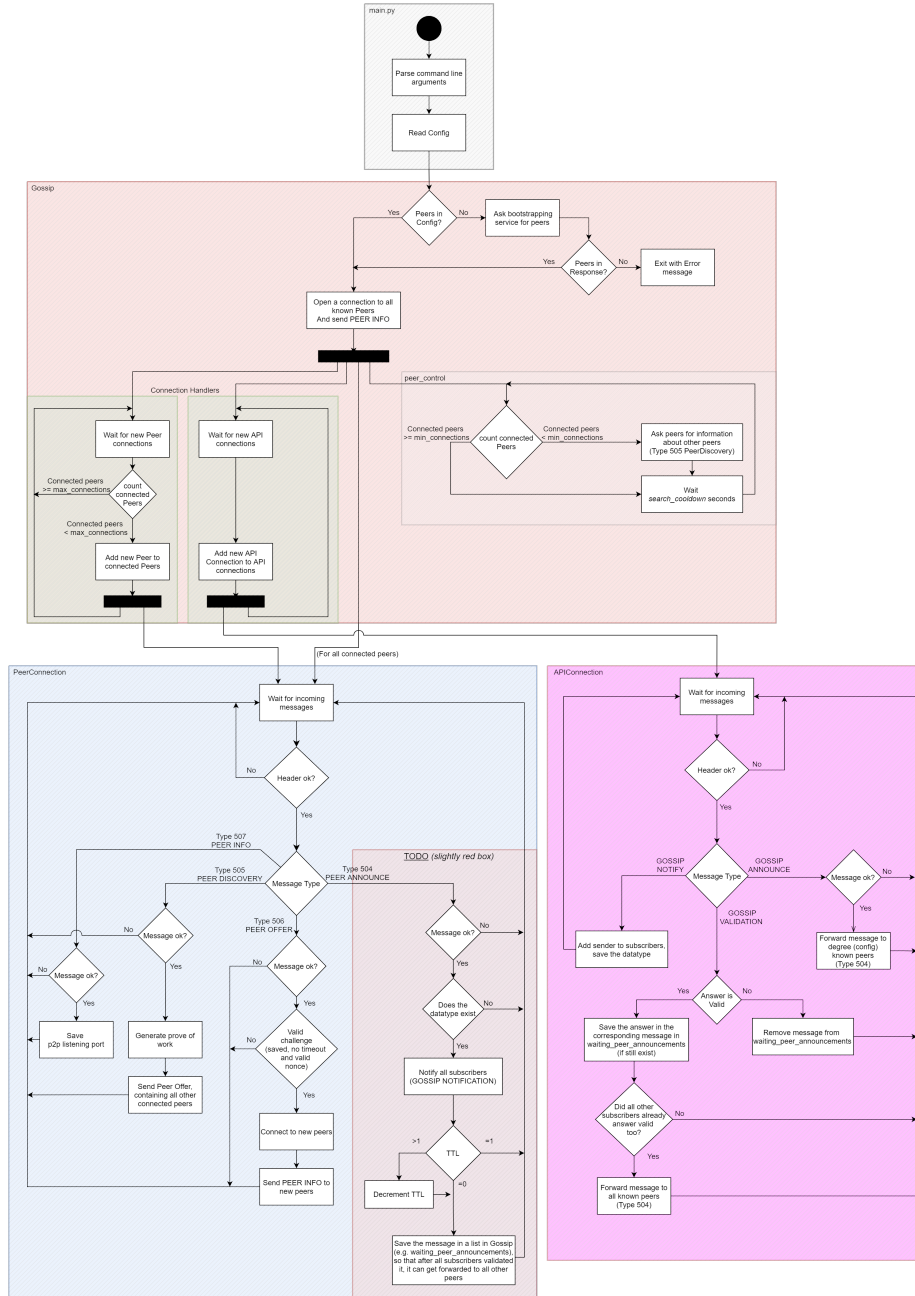


Figure 1: Flow chart (v4) for Gossip. Colored boxes represent modules/classes. Note that the boxes are a simplification and some tasks may be in box A but use functionality of class represented by box B.

which is used in multiple parts of the program. All current classes and modules will be shortly introduced in the following.

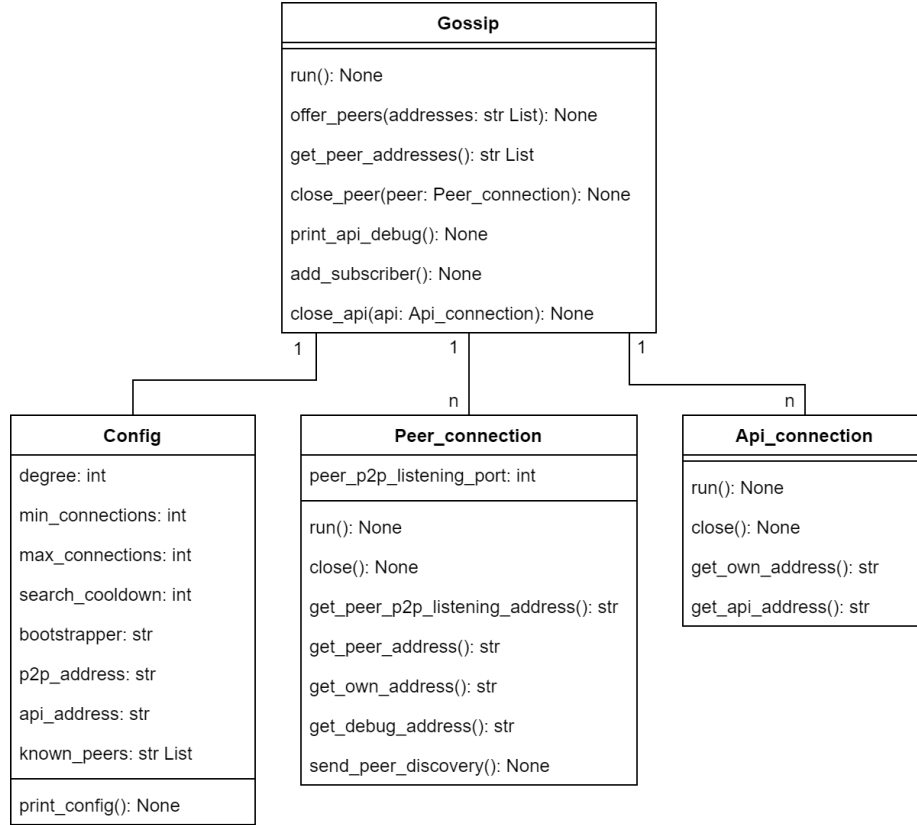


Figure 2: Gossip UML class diagram

2.1.1 Gossip class

An instance of the Gossip class represents an instance of the Gossip module. It manages the startup and searches for new peers if the amount of connected peers is below a specified number (see `peer_control`). Gossip knows all connected peers and API connections, which themselves use gossip to communicate with each other. It also knows about every datatype and subscriber. Logically it is at the very top of the hierarchy as it includes all the other classes.

2.1.2 Peer_connection class

An instance of Peer_connection represents a connection to another peer. The class provides the functions to communicate with the other peer and a `run()` function, which waits for incoming messages from the connected peer.

2.1.3 `Api_connection` class

`Api_connection` works similarly to its peer counterpart, representing an active connection to a single API.

2.1.4 `Config` class

The `Config` provides an abstraction for config files. It provides parsing and allows for quick access to config parameters. The `Config` module also contains a `'config.config'` variable, which is used as a blueprint for parsing config files. It allows quick and easy addition of new parameters, setting default values and types.

Variables:

- **degree**: Number of peers this peer exchanges information with. Relevant for PEER ANNOUNCE.
- **min_connections**: Minimum amount of alive connections this peer should try to keep.
- **max_connections**: Maximum amount of alive connections this peer should try to keep.
- **search_cooldown**: In this interval it is checked whether we have **min_connections** peers. If not start search.
- **bootstrapper**: One trustworthy bootstrapping node, used as a fallback.
- **p2p_address**: Listening ip and port number for other Gossip peers.
- **api_address**: Listening ip and port number for API users.
- **known_peers**: List of peers to bootstrap through, if it is empty consult **bootstrapper**.

2.1.5 `Connection handler` module

The connection handler module is used to wait for and accept new incoming connections from other peers or API. For this, Gossip starts two separate tasks with the peer address or API address specified in the given config respectively.

2.1.6 `Proof of work producer` module

The proof of work producer module (`pow_producer`) provides the functionality required to solve proof of work challenges.

2.1.7 Packet parser module

The packet parser module is currently divided into two different modules, one for parsing and constructing API related messages and the other for parsing and constructing peer related messages. They will likely be merged in the final version.

2.2 Process architecture & Networking

As mentioned in 1.1 we changed our current implementation to exclusively use Python's AsyncIO instead of threading. However, we may use threading or multiprocessing for generating proof of work in a future version.

2.3 Security measures

As soon as malformed messages from a peer are received we terminate the connection. We can exclude transmission errors as we use TCP.

3 Protocol/-s

Number	Name	Purpose
504	PEER ANNOUNCE	spreading of knowledge
505	PEER DISCOVERY	discovering new peers
506	PEER OFFER	spreading information about other peers
507	PEER INFO	tells peer which port is listening on my side

Table 1: Peer Message Overview

3.1 PEER ANNOUNCE

Figure 3 shows this message. Spreads the GOSSIP ANNOUNCE to peers. In case we receive a GOSSIP ANNOUNCE from an API user we transform it into this message and send it to our peers. In case we receive a PEER ANNOUNCE from a peer we keep spreading it in general. If we have a subscriber to this datatype, we transform the PEER ANNOUNCE into a GOSSIP NOTIFICATION. If we receive a negative PEER VALIDATION from the API that received the NOTIFICATION, we do not propagate it further.

The TTL field functions equally to the GOSSIP ANNOUNCE TTL field.

The id is a temporary addition from a previous sketch, we plan to rework it as soon as we further thought about how gossip is spreading its messages. It's whole purpose is to prevent message loops by using the id to remember that we already encountered this packet.

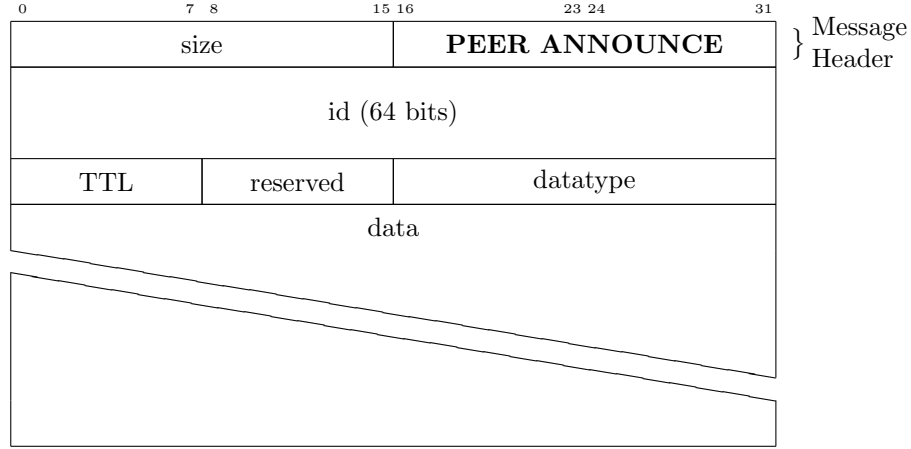


Figure 3: PEER ANNOUNCE message format

3.2 PEER DISCOVERY

Figure 4 shows this message. This and the next message (PEER OFFER) deal with the problem of keeping connection to degree amount of peers. We want to create a pool of peers, limited by max_connections in the config, from which we can choose degree many peers. Degree is also specified in the config.

To find degree many peers, do the following:

```
# n = amount of connected peers
while running
    if n < min_connections
        1. Send PEER DISCOVERY to all connected peers
        2. Peers respond (PEER OFFER) with a list
           containing all their known peers, add amount to n
        3. connect to peers in all given lists in a random order
           until we reach min_connections.
    sleep for search_cooldown seconds
```

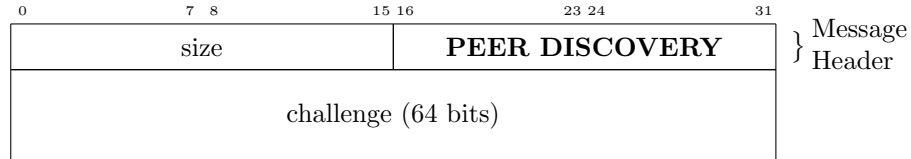


Figure 4: PEER DISCOVERY message format

3.3 PEER OFFER

Figure 4 shows this message. Answer to a PEER DISCOVERY message. Includes the senders peer list

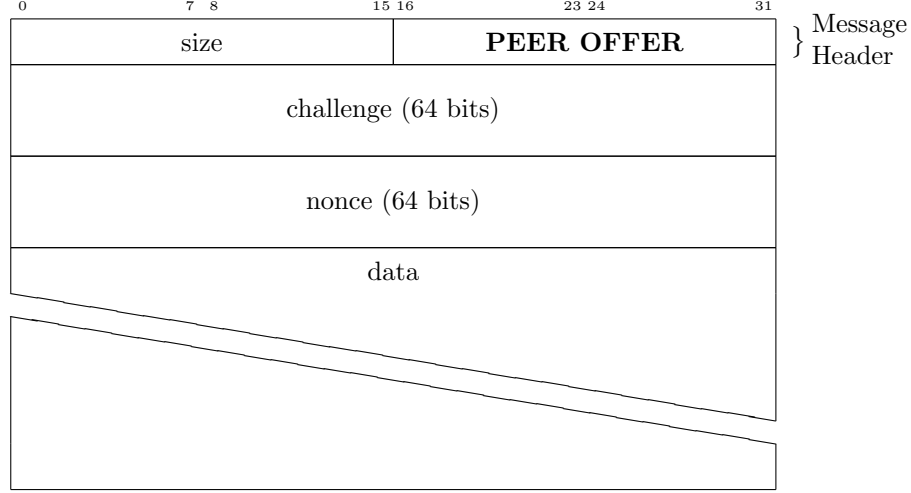


Figure 5: PEER OFFER message format

3.4 PEER INFO

Figure 4 shows this message. Informs the receiver of the senders p2p listening port for new incoming p2p connections. This is necessary since the port used to connect to the other peer is dynamic and not equal to the p2p address port. This port number is also spread further in the network through each peers list when answering with PEER OFFER to PEER DISCOVERY.

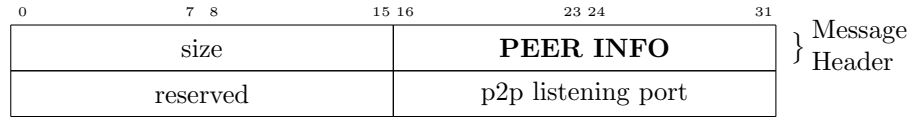


Figure 6: PEER INFO message format

4 Future Work

PoW based identities on parallel connection establishment when connecting to new peer (potentially given to us from one bad actor) to counteract Sybill and Eclipse attacks. Currently we have PoW in the request for new peers to neighbours, we plan to move it to the connection establishment.

Secure Channel between Peers: Pubkey exchange, Pubkeys as long-term keys

to this Identity, Session key establishment (DH), usage of AEAD for encryption and authentication of messages.

More tests in form of unittests for utility functions and whole system tests with mocks.

5 Workload Distribution

5.1 Liebald

- Peer_connection class implementation
- Packet parsing and generation for PEER DISCOVERY, PEER OFFER and PEER INFO
- Gossip class: basic structure and functionality related to Peer_connection
- Config class implementation
- Smaller modules: connection_handler and pow_producer

5.2 Ploch

- Api_connection class implementation
- Packet parsing and generation for GOSSIP ANNOUNCE, GOSSIP NOTIFY, GOSSIP NOTIFICATION and GOSSIP VALIDATION
- Gossip class: functionality related to Api_connection
- API Packet parser unit tests

6 Effort spent for the project

- Meetings/together: approximately 21 hours
- Liebald: approximately 30 hours
- Ploch: about 27-30 hours