

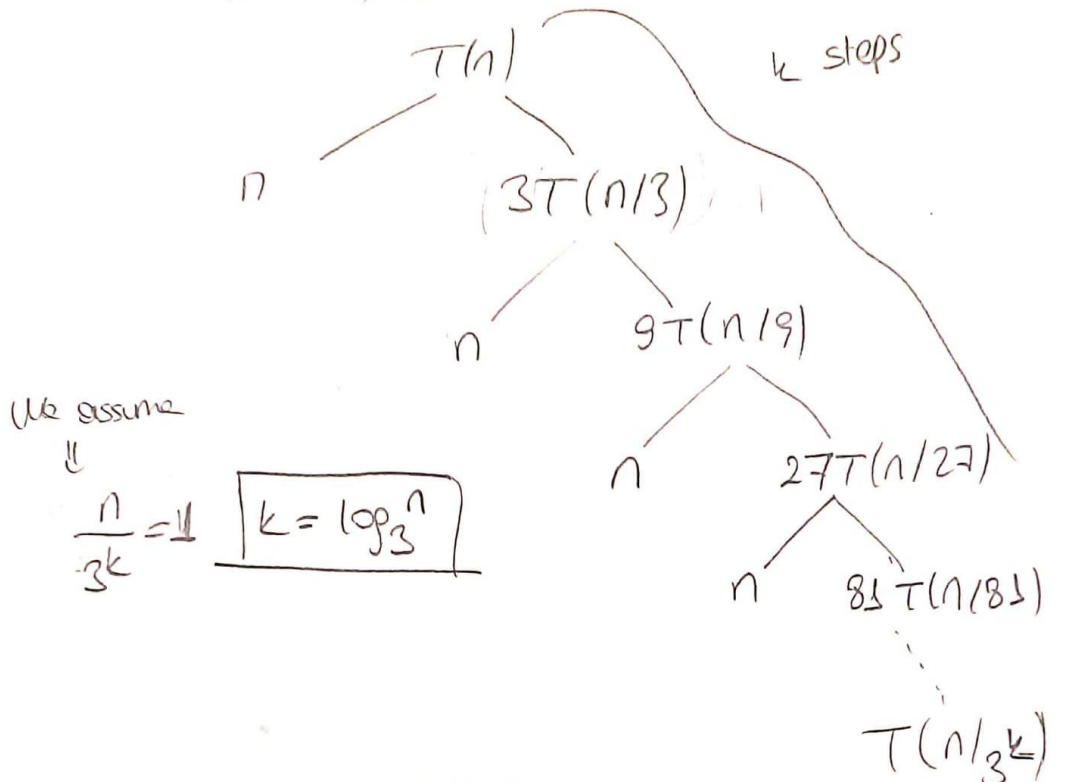
1.A-Find the asymptotic running time for.

A-) $T(n) = 3T(n/3) + n$, where $T(1) = 1$ and n is an exact power of 3

Q1-A

$$T(n) = \begin{cases} T(1) = 1 & n=1 \\ 3T(n/3) + n & n>1 \end{cases} \quad \left. \begin{array}{l} \text{Dividing} \\ \text{functions} \end{array} \right\}$$

(1)



A-

$$T(n) = 3T(n/3) + n$$

$$T(1) = 1$$

$$T(n) = 3T(n/3) + n$$

$$T(n/3) = 3T(n/9) + n/3$$

$$T(n/9) = 3T(n/27) + n/9$$

$$T(n) = 3[3T(n/9) + n/3] + n$$

$$T(n) = \underline{9T(n/9)} + 2n$$

$$T(n) = 9[3T(n/27) + n/9] + 2n$$

$T(n) = 27T(n/27) + 3n$
$T(n) = 9T(n/9) + 2n$
$T(n) = 3T(n/3) + n$

$$\Rightarrow 3^k T(n/3^k) + k \cdot n$$

$$k = \log_3 n$$

$$T(n) = 3^{\log_3 n} T(n/3^{\log_3 n}) + \log_3 n \cdot n$$

$$T(n) = n \cdot T(1) + \boxed{\log n \cdot n}$$
$$\underline{\underline{O(\log n \cdot n)}}$$

B-) $T(N) = 3T(N/2) + 1$

B-)

$$T(n) = 3T(n/2) + 1$$

$$T(1) = 1$$

$$T(n) = \begin{cases} 1 & n=1 \\ 3T(n/2)+1 & n>1 \end{cases}$$

$$T(n) = 3T(n/2) + 1 \quad T(n) = 3[3T(n/4) + 1] + 1$$

$$T(n/2) = 3T(n/4) + 1 \quad T(n) = 9T(n/4) + 4$$

$$T(n/4) = 3T(n/8) + 1 \quad T(n) = 9[3T(n/8) + 1] + 4$$

$$T(n) = 27T(n/8) + 13$$

To make $T(n/2^k) = T(1)$

$$\boxed{\begin{matrix} 2^k = n \\ \log_2 n = k \end{matrix}}$$

$$\downarrow \\ 3^k T(n/2^k) + (3 \cdot 2^{k-1} + 1)$$

$$T(n) = 3^{\log_2 n} T(n/2^k) + (3 \cdot 2^{k-1} + 1)$$

$$= n^{\log_2 3} T(n/2^{\log_2 n}) + \left(\frac{3}{2} \cdot 2^{\log_2 n} + 1 \right)$$

$$= n^{\log_2 3} T(n/n) + \frac{3}{2}n + 1$$

$$\log_2 3 > 1 \rightarrow n^{1.234} \rightarrow \boxed{O(n^2)} //$$

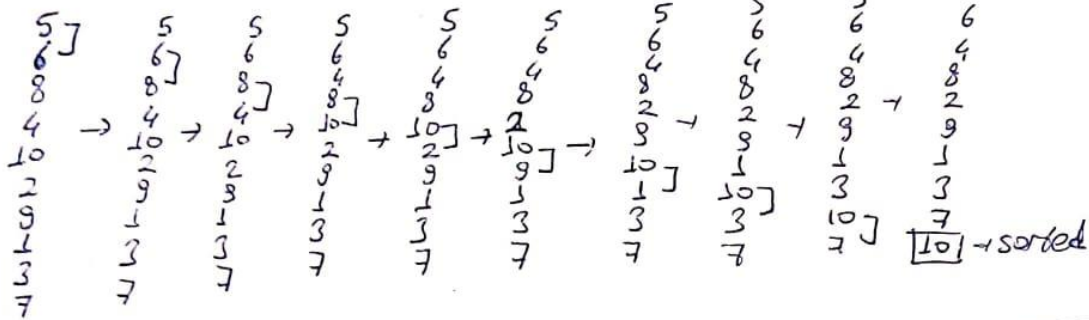
1.B-Tracing with Bubble and Selection Sort

Bubble Sort Tracing

[5, 6, 8, 4, 10, 2, 9, 1, 3, 7]

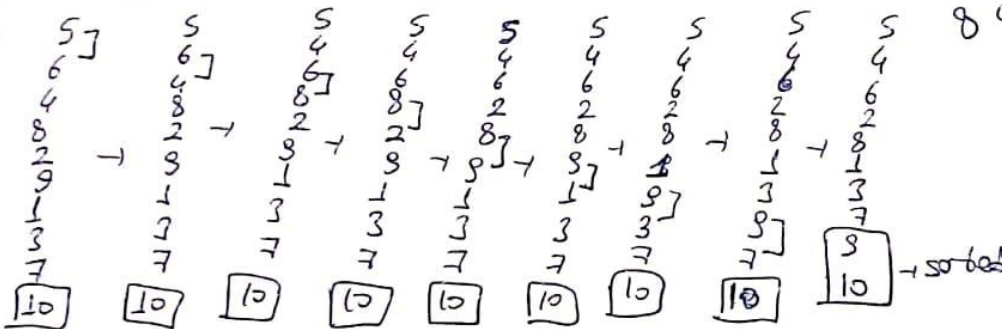
Pass 1

9 comparison

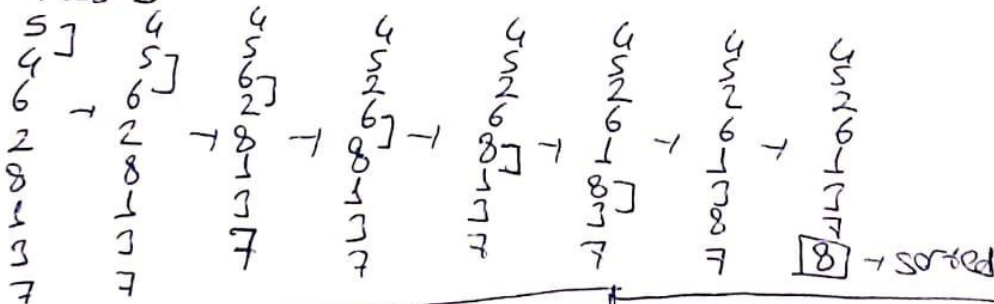


Pass 2

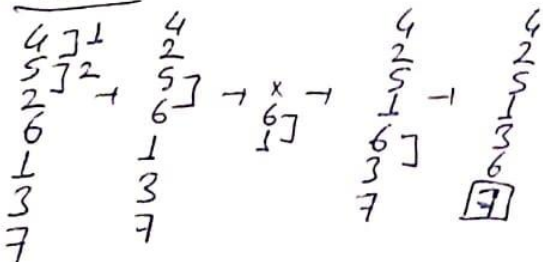
8 comparison



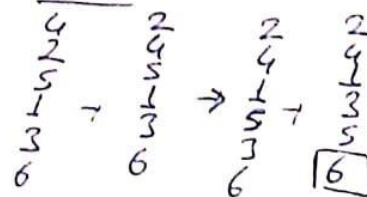
Pass 3



Pass 4



Pass 5



Pass 6

2 2 2
4 1 1
1 4 3
3 3 4
5 5 5

Pass 7

2 1
1 2
3 3
4 4

sorted

Selection Sort Tracing

5 6 8 4 (10) 2 9 1 3 7 |
← unsorted sorted

5 6 8 4 7 2 (9) 1 3 | 10
←

5 6 (8) 4 7 2 3 1 | 9 10
←

5 6 1 4 (7) 2 3 | 8 9 10
←

5 (6) 1 4 3 2 | 7 8 9 10
←

(5) 2 1 4 3 | 6 7 8 9 10
←

3 2 1 (4) | 5 6 7 8 9 10
←

(3) 2 1 | 4 5 6 7 8 9 10
←

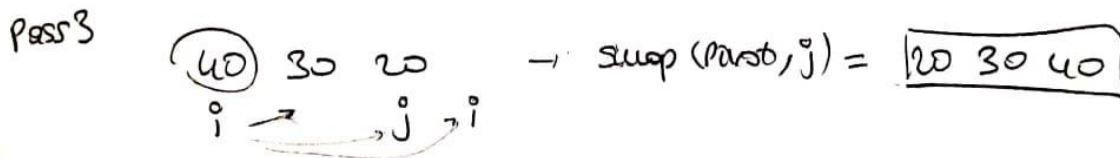
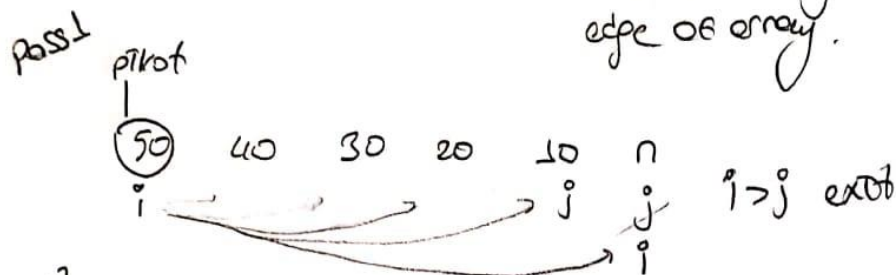
1 2 3 4 5 6 7 8 9 10 // sorted

1.C-Recurrence Relation of Quick Sort at its worst case

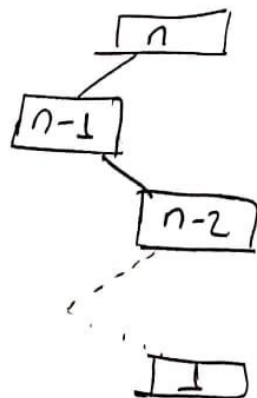
Question 1-C)

If elements are ascending or descending order, quick sort will be $O(n^2)$ which is the worst case.

$\begin{matrix} 50 & 40 & 30 & 20 & 10 \\ 10 & 20 & 30 & 40 & 50 \end{matrix}$
 Both of them will end with $O(n^2)$.
 Since partitioning happens at the edge of array.



No. of comparison to partition quick sort



$$\frac{n(n-1)}{2} \text{ comparison} = O(n^2)$$

$$T(n) = T(n-1) + 1$$

Question 1-C

$$T(n) = \begin{cases} 1 & n=0 \\ T(n-1) + O(n) & n > 0 \end{cases}$$

$$T(n) = T(n-1) + O(n)$$

$$T(n-1) = T(n-2) + O(n)$$

$$T(n-2) = T(n-3) + O(n)$$

$$T(n) = T(n-2) + 2 \cdot O(n)$$

$$T(n) = T(n-3) + 3 \cdot O(n)$$

\vdots k times

$$T(n) = T(n-k) + k \cdot O(n)$$

$$n-k=0$$

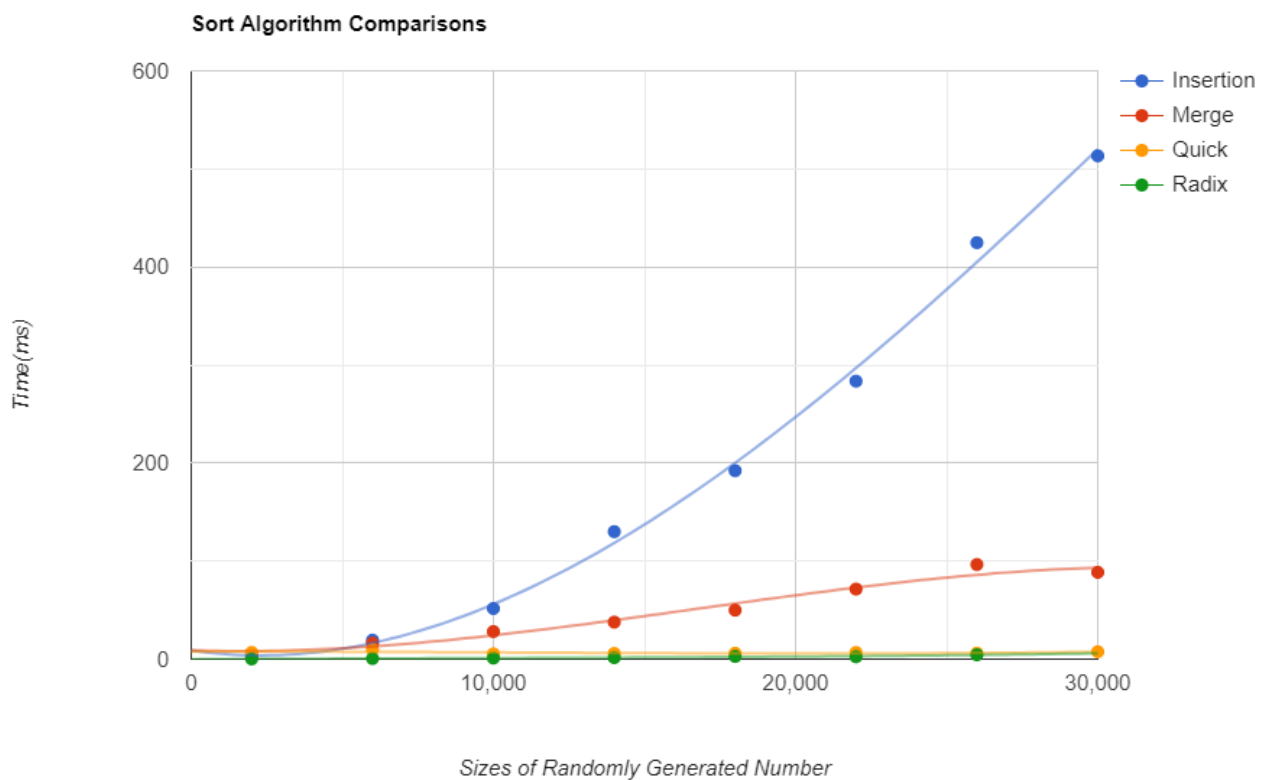
$$k=n$$

$$T(n-n) + n \cdot O(n) =$$

$$T(0) + n^2 = n^2 + 1 = \boxed{O(n^2)}$$

Performance Analysis Comparisan Between Four Sorts

Sizes	2000	6000	10000	14000	18000	22000	26000	30000
Insertion	2.41430ms	19.3366ms	51.6722 ms	130.031 ms	192.18 ms	283.524 ms	424.76 ms	513.306 ms
Merge	5.6893ms	16.188 ms	28.1506 ms	37.7412 ms	49.9368 ms	71.4521 ms	96.4917 ms	98.5735 ms
Quick	6.868ms	9.8599 ms	5.3258 ms	5.7591 ms	5.9589 ms	6.7796 ms	6.0239 ms	7.5471 ms
Radix	0.2036 ms	0.6156 ms	1.1023 ms	1.5501 ms	2.7517 ms	2.4031 ms	4.4148 ms	4.8569 ms



Insertion – Average = $O(n^2)$, Worst = $O(n^2)$

Merge – Average = $O(n \log(n))$, $O(n \log(n))$

Quick – Average = $O(n \log(n))$, worst = $O(n^2)$

Radix – Average = $O(d \cdot (n+b))$ where b is base, d is number of digits.

The worst sorting algorithm was insertion sort just like expected. Therefore, for this condition, empirical and theoretical values suit. The best algorithm was radix sort as it works with $O(n)$ complexity. All sort algorithms take more time with respect to array size. However, for a few cases, there is an exception. Sometimes even though the size of the array gets bigger, the runtime gets smaller. But that is also possible since we create randomly generated arrays. Some random series of numbers may create easier sorting for quick as it is number sensitive in a way. In quick sort order of numbers and their places directly affect the runtime. However, a sort algorithm like merge is not affected by the order of numbers. Therefore, there is no fluctuation in its runtime results. Merge algorithm divides the array till it creates sub-arrays with size 1. After it divides, it starts to merge them by comparing. Thus, it is not affected by the position of values in a given random array.

Another critical outcome is the difference between the runtime increasing with respect to the array size. As all algorithms depend on the size parameter, they are increasing but the difference is getting bigger because they have different growth rates. By considering their average cases, the difference between insertion and other algorithms is understandable. Insertion grows with n^2 . The interesting point is why the quick sort is so fast although it has also the worst case with n^2 too? Because worst case of the quick occurs when the array is sorted or sorted in descending order. As we are randomly generated our arrays, we don't see quick's worst-case scenario. Therefore, it is normal to have such a difference between quick and merge with an array size of 30000. For small sizes arrays, their differences are not that huge. The best algorithm was radix sort by considering experimental results. Radix sort controls digit by digit to sort the array. In case of the given array is sorted, merge sort won't affect it. But insertion sort is affected positively since the sorted array is its best case. After it divides the array into two as; unsorted and sorted, it doesn't have a value to put sorted side as the other unsorted side is also sorted array.

Quicksort is affected in a bad way since the sorted array is its worst case. As we choose our pivot value from the beginning, we traverse the whole array to understand our pivot is in the right position. Radix sort is also not affected since it is not value-position sensitive.

Notes: I used `auto start1 = chrono::high_resolution_clock::now();` when I collecting data and filling the table given above. However, when I try to run my program on dijkstra it gives error. Therefore, I measure the runtime with `clock_t.start()` method. But, in newer version merge sort is always slower than insertion unlike the expectataion. I have also taken the permission of Can Hoca and TA to put my answers in handwritten for Question 1.

```
run: ConsoleApplication2.o sorting.o
    g++ ConsoleApplication2.o sorting.o -o run

ConsoleApplication2.o: ConsoleApplication2.cpp
    g++ -c ConsoleApplication2.cpp

sorting.o: sorting.cpp sorting.h
    g++ -c sorting.cpp

clean:
    rm *.o run
```

```

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

```

Part A: Time Analysis of Insertion Sort

Array Size	Time Elapsed	compCount	moveCount
2000	0.3ms	1999000	1064338
6000	1.9ms	17997000	9115449
10000	5.6ms	49995000	25061625
14000	11.1ms	97993000	49687562
18000	17.8ms	161991000	81023982
22000	26ms	241989000	121835995
26000	39ms	337987000	169592679
30000	47ms	449985000	224426774

Part B: Time Analysis of Merge Sort

Array Size	Time Elapsed	compCount	moveCount
2000	5ms	115100	439040
6000	14.7ms	408670	1516160
10000	24.2ms	726520	2672320
14000	33.9ms	1057730	3872320
18000	42.6ms	1403420	5104640
22000	52.3ms	1734380	6384640
26000	61.1ms	2099940	7664640
30000	70.7ms	2462880	8944640

Part C: Time Analysis of Quick Sort

Array Size	Time Elapsed	compCount	moveCount
2000	1ms	27841	40821
6000	2ms	88494	131652
10000	3ms	166514	220776
14000	4ms	239904	320001
18000	4ms	327125	450306
22000	6ms	413113	501600
26000	7ms	541657	671949
30000	8ms	657952	763521

Part D: Time Analysis of Radix Sort

Array Size	Time Elapsed	compCount	moveCount
2000	0.2ms	0	0
6000	0.6ms	0	0
10000	1ms	0	0
14000	1.3ms	0	0
18000	1.7ms	0	0
22000	2.1ms	0	0
26000	2.5ms	0	0
30000	3ms	0	0

```

[eren.gunaltili@dijkstra CS202]$ emacs Makefile
[eren.gunaltili@dijkstra CS202]$ make
g++ -c ConsoleApplication2.cpp
g++ -c sorting.cpp
g++ ConsoleApplication2.o sorting.o -o run
[eren.gunaltili@dijkstra CS202]$ ./run
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
-----
Part A: Time Analysis of Insertion Sort
Array Size      Time Elapsed      compCount      moveCount
2000            1ms             1999000        1019807
6000            5ms             17997000       9132828
10000           13ms            49995000       24988035
14000           26ms            97993000       49024247
18000           44ms            161991000      81726530
22000           65ms            241989000      121665227
26000           90ms            337987000      169924018
30000           120ms           449985000      226948089
-----
Part B: Time Analysis of Merge Sort
Array Size      Time Elapsed      compCount      moveCount
2000            1ms             115150         439040
6000            1ms             407420         1516160
10000           2ms             726030         2672320
14000           2ms             1058490        3872320
18000           4ms             1398860        5104640
22000           4ms             1744760        6384640
26000           5ms             2099940        7664640
30000           6ms             2468610        8944640
-----
Part C: Time Analysis of Quick Sort
Array Size      Time Elapsed      compCount      moveCount
2000            0ms             24220          39621
6000            0ms             87346          122394
10000           0ms             157332         215823
14000           0ms             253502         318912
18000           10ms            325071         417498
22000           0ms             452515         559863
26000           10ms            519098         577059
30000           10ms            643395         731142
-----

```

3-Dijkstra Machine Screenshot

Part C: Time Analysis of Quick Sort				
Array Size	Time Elapsed		compCount	moveCount
2000	0ms	24220	39621	
6000	0ms	87346	122394	
10000	0ms	157332	215823	
14000	0ms	253502	318912	
18000	10ms	325071	417498	
22000	0ms	452515	559863	
26000	10ms	519098	577059	
30000	10ms	643395	731142	

Part D: Time Analysis of Radix Sort				
Array Size	Time Elapsed		compCount	moveCount
2000	0ms	0	0	
6000	1ms	0	0	
10000	2ms	0	0	
14000	3ms	0	0	
18000	3ms	0	0	
22000	4ms	0	0	
26000	4ms	0	0	
30000	6ms	0	0	

4-Remaining screenshot