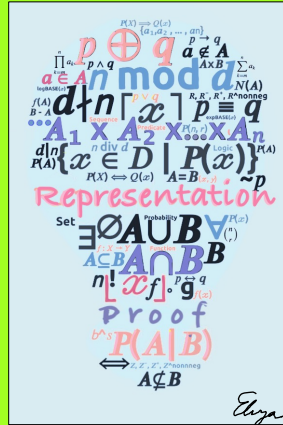


Discrete Structures



Lecture 25: Trees

Susan L. Epstein



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Last time

- ★ **Graph theory has important real-world applications**
- Connectivity
- Euler and Hamilton paths and circuits
- Some proofs rely on computers to address their myriad cases



Trees have important real-world applications



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Today's outline

- Basic ideas and examples
- Properties of trees
- Applications of trees

The material on graphs and trees in CSCI 150 supersedes your text

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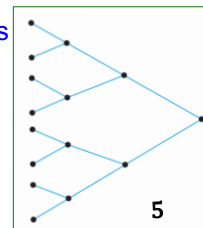
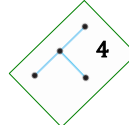
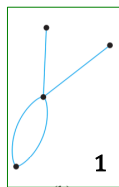
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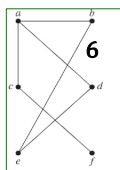
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Basic definition

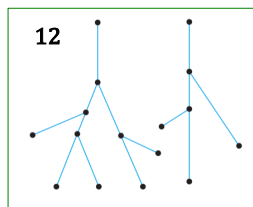
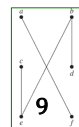
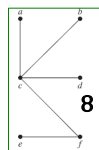
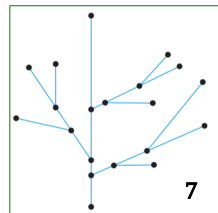
Tree = connected simple undirected graph without circuits



Which of these are trees?



[Rosen, 2019]



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More on trees

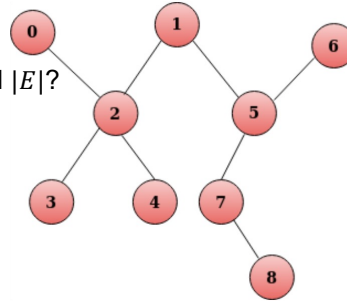
Equivalent definitions:

Any **simple undirected graph** that

- Is **connected** and **acyclic** (contains no cycles/circuits)
- Is **acyclic** and would contain a simple cycle if **any edge were added** to it
- Is **connected** and would become disconnected when **any edge is removed** from it
- Where any pair of distinct vertices is connected by **exactly one path**

For a **tree** $G = (V, E)$

- What is the relationship between $|V|$ and $|E|$?
 $|V| - |E| = 1$
- What is the total degree of a tree?
 $2(|V| - 1)$



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Vertices and edges

Theorem: $P(n)$: a tree with n vertices has $n - 1$ edges.

Proof by mathematical induction:

We must show that $P(n)$ is true for all $n \geq 1$.

Basis: $P(1)$: a tree with 1 vertex has $1 - 1$ edges. Such a tree has 0 edges, so $P(1)$ is true.

Inductive step: Assume for some k that $P(k)$ is true.

By substitution $P(k + 1)$: a tree with $k + 1$ vertex has k edges.

We must show that $P(k + 1)$ is true.

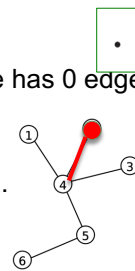
Let T be a tree with $|V| = k + 1$.

Because T is finite, there must be some $v \in V$ that is a leaf.

Let T' be the tree formed when both v and the edge from its parent are deleted from T .

By the inductive step, T' has $k - 1$ vertices, so T must have $k - 1 + 1$ vertices and $P(k + 1)$ is true.

Since we have proved the basis step and the inductive step, the theorem is true. **QED**



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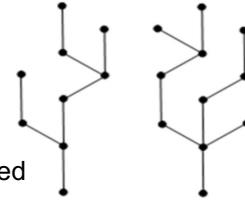
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Forests

Equivalent definitions for forest

- Undirected graph in which any two vertices are connected by **at most one path**
- Unconnected acyclic graph all of whose connected components are trees
- A disjoint union of trees $G = \bigcup_i T_i(V_i, E_i)$
where $V_i \cap V_j = \emptyset$ for $i \neq j$



Special cases

- An empty forest
- A single tree
- An edgeless graph



- How many trees in a forest $G = \bigcup_i T_i(V_i, E_i)$?
In each tree $T_i(V_i, E_i)$, $|V_i| - |E_i| = 1$ so
 G has $|V| - |E|$ trees

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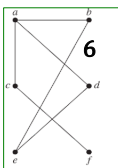
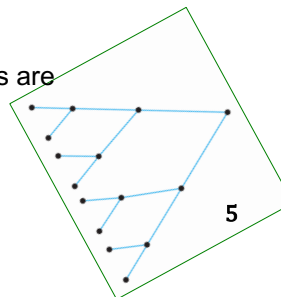
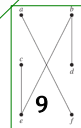
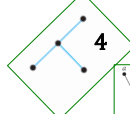
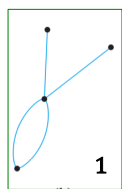
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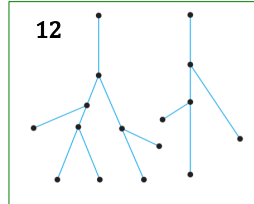
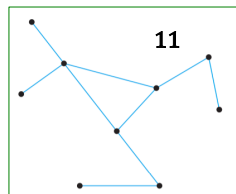
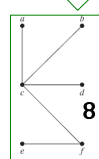
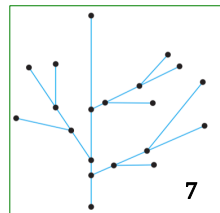
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Find the forests

Forest = undirected graph in which any two vertices are connected by at most one path



[Rosen, 2019]



Which of these are forests?

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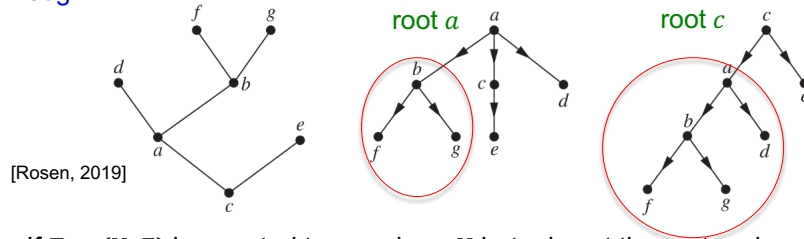
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Rooted tree

- **Rooted tree** = tree with **exactly 1** designated vertex (its **root**) and every edge is directed away from that root



- If $T = (V, E)$ is a rooted tree and $v \in V$ but v is not the root and $u \in V$ has an edge directly to v , then u is the **parent** of v and v is the **child** of u .
- Vertices with the same parents are **siblings**
- **Ancestors** of v are its parent and those on the path from the root to v
- **Descendants** of v are vertices that have v as an ancestor
- **Leaf** = node with **no children**
- **Internal node** = vertex that is neither the root nor a leaf
- **Subtree** rooted at vertex $v = v$, its descendants, and the edges incident to them

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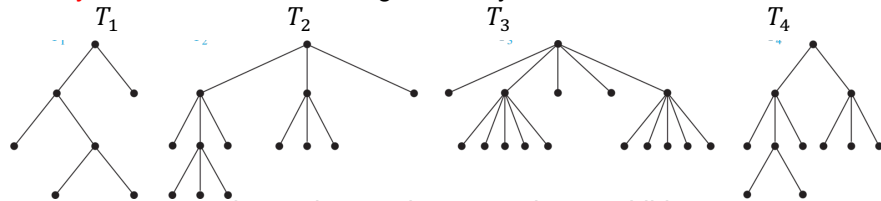
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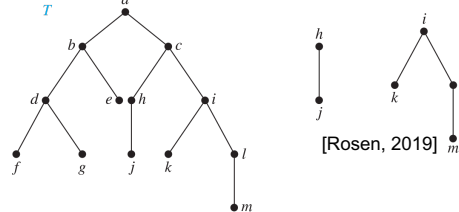
Arity of a tree

- **Arity** of tree T = maximum degree of any of its vertices



[Rosen, 2019]

- **m -ary tree** = no internal vertex has more than m children
- **Binary tree** = m -ary tree with $m = 2$
- **Full m -ary tree** = every non-leaf has exactly m children
- In a binary tree, a vertex with 2 children has
 - A **left child** and a **right child**
 - A **left subtree** rooted at its left child and a **right subtree** rooted at its right child



[Rosen, 2019]

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Today's outline

- ✓ Basic ideas and examples
- Properties of trees
- Applications of trees

The material on graphs and trees in CSCI 150 supersedes your text

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Another characterization of a tree

Theorem: An undirected graph is a tree **iff** there is a unique simple path between every pair of its vertices.

Proof:

Part 1: Let $T = (V, E)$ be a tree and $x, y \in V$ be any pair of distinct vertices in T . By definition of tree, T is connected and therefore has a path between x and y . **Assume** there are 2 distinct paths $x = pq \dots rs = y$ and $x = ab \dots cd = y$. **We will show that this assumption leads to a contradiction.** The path $x = pq \dots ryc \dots ba = x$ begins and ends at x and is a circuit. This contradicts that T is a tree. **Because the assumption led to a contradiction,** the path between x and y must be unique.

Part 2: Let $T = (V, E)$ be a graph with a unique simple path between every pair of its vertices. Then by definition, T is connected. **Assume** T has a circuit $ab \dots xc \dots yd \dots a$, that includes some pair of vertices x and y . **We will show that this assumption leads to a contradiction.** The paths $xc \dots y$ and $yd \dots ab \dots x$ are simple and distinct. **Because the assumption led to a contradiction,** T has no circuits and by definition is a tree. **QED**

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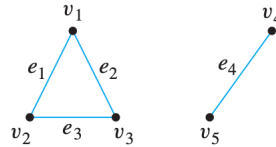
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Another characterization of a tree

- Consider a connected graph where $|V| = n$.
- How can you use $|E|$ to determine whether T is a tree?
- **Theorem:** For any $n \in \mathbb{Z}^+$, if $G = (V, E)$ is a connected graph with n vertices and $n - 1$ edges, G is a tree. (Proof is in your text.)
- What if $G = (V, E)$ is not connected; could it still be a tree?
5 vertices and 4 edges?



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Today's outline

- ✓ Basic ideas and examples
- ✓ Properties of trees
- Applications of trees

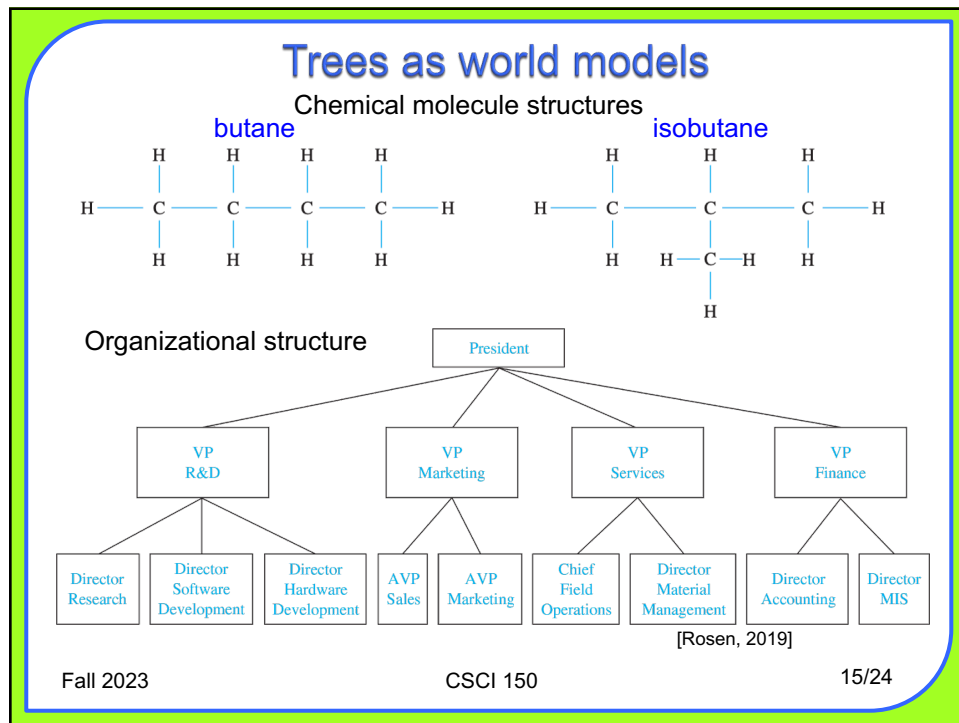
The material on graphs and trees in CSCI 150 supersedes your text

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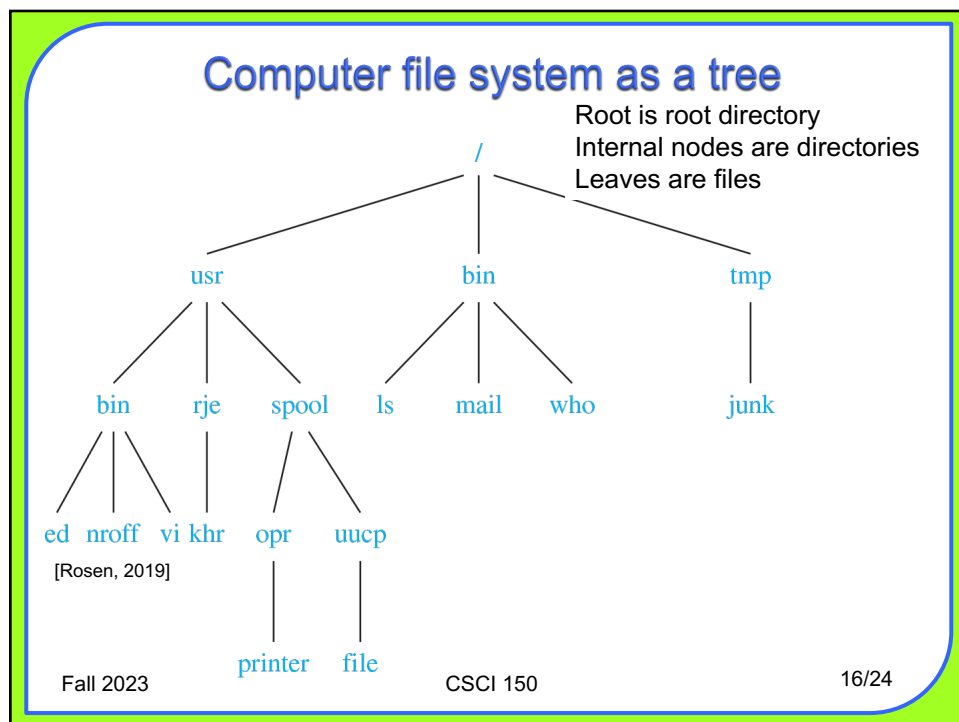
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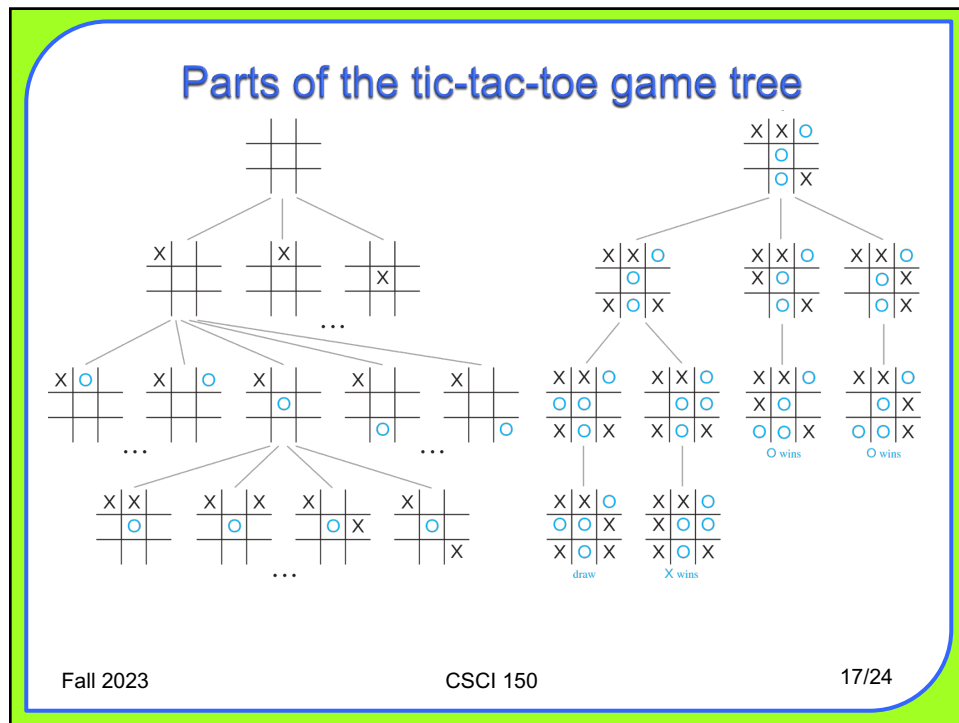
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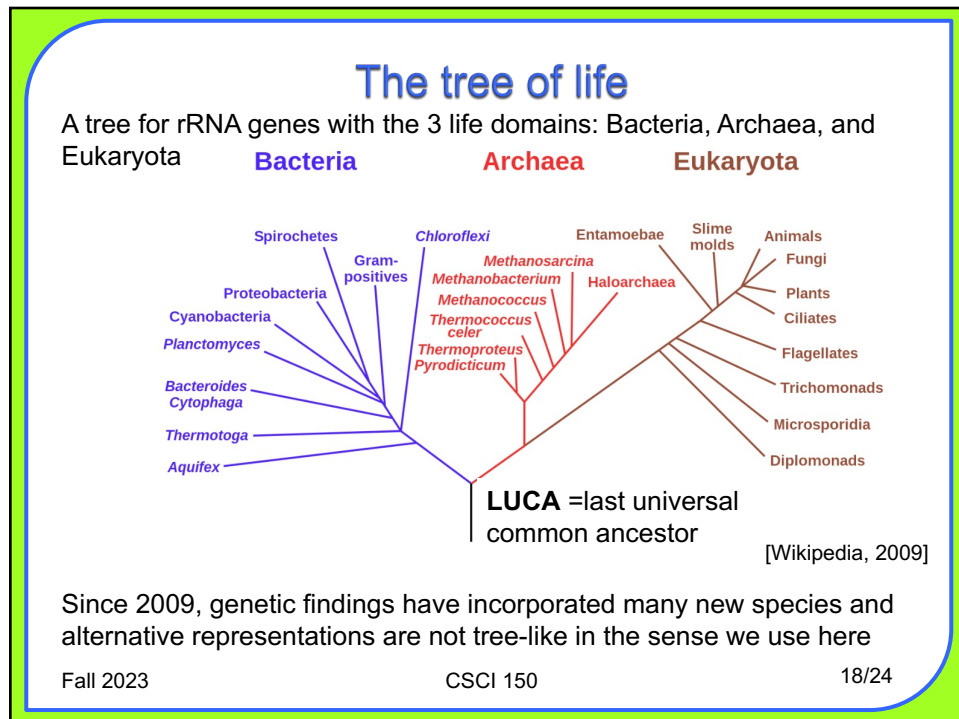
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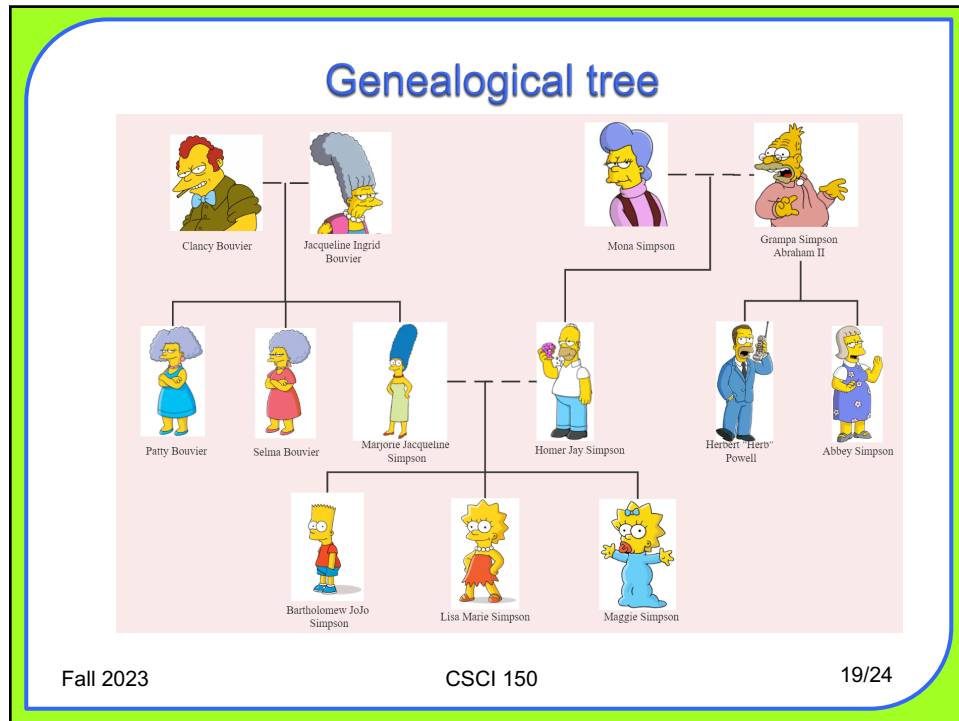
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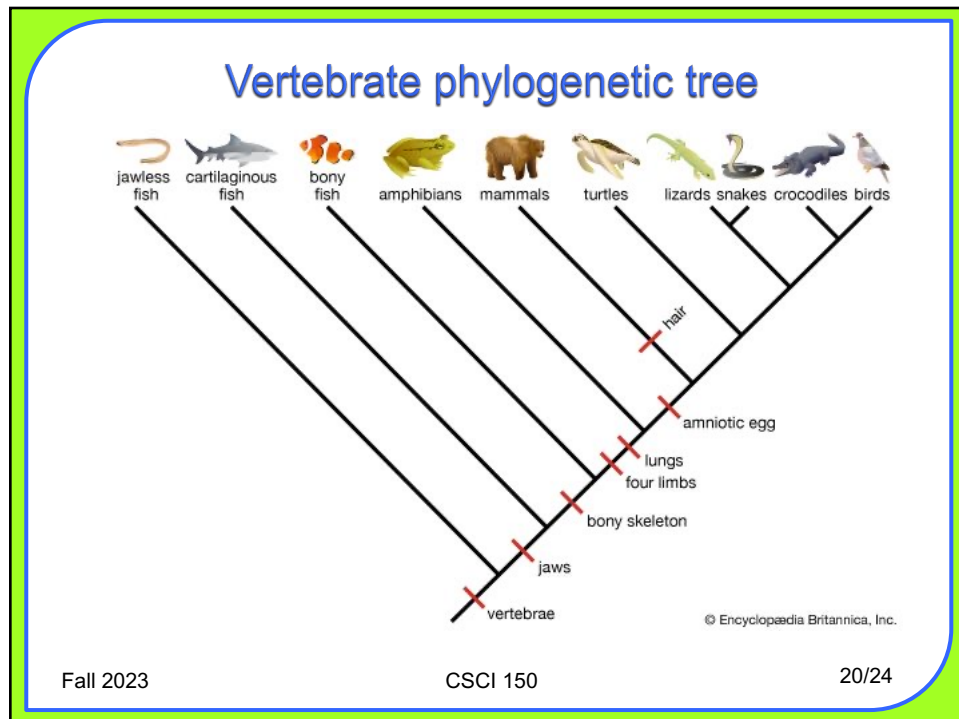
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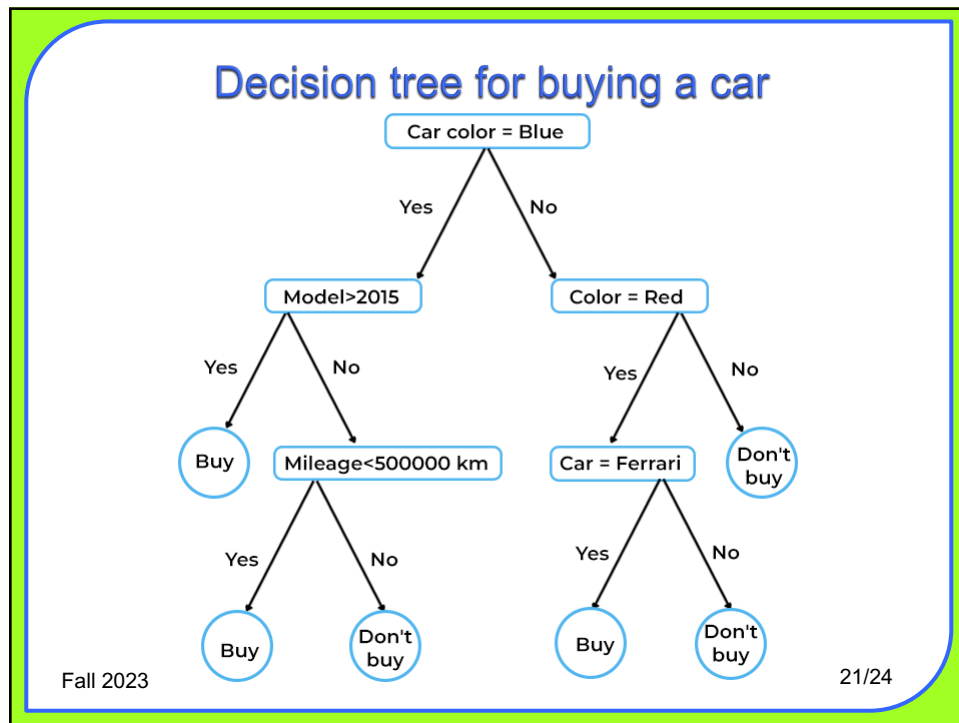
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And many more...

- Parse trees address grammatical sentence structure and support [natural language processing \(NLP\)](#)
- Balanced trees, B-trees, and B+-trees support [efficient storage, retrieval and sorting](#)
- Spanning trees support minimal [path finding](#)
- Trees support the design of [sorting algorithms](#) and [universal address systems](#)
- Trees support Huffman coding for [optimal coding of a set of symbols](#)

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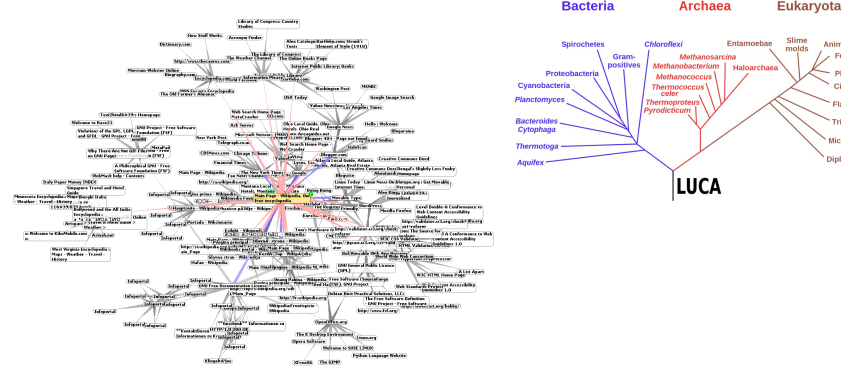
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Working with trees

- Many operations that are difficult on graphs are far easier on trees

search

traversal



- Most of those algorithms use recursion

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What you should know

- ★ Trees have important real-world applications

- Trees and forests are subclasses of graphs with special properties



Any questions?

Next up: *Ethics*

Time to finish up that Opening sheet!

Problem set 25,26 is due on Wednesday, December 13 at 11PM

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