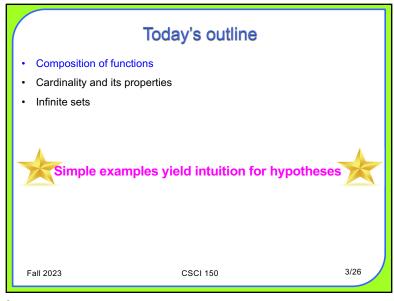
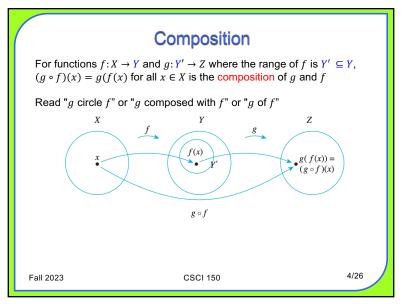


Last time ★ Functions are ubiquitous and powerful • Functions can be defined on more than numbers • Special properties of functions impact computation and storage



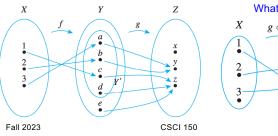




By formula:

For $f: \mathbf{Z} \to \mathbf{Z}$ with rule f(n) = n + 3 and $g: \mathbf{Z} \to \mathbf{Z}$ with rule $g(n) = n^{2}$, $(g \circ f)(n) = g(f(n)) = g(n+3) = (n+3)^2 \ \forall n \in \mathbf{Z}$ $(f \circ g)(n) = f(g(n)) = g(n^{2}) = n^{2} + 3 \ \forall n \in \mathbf{Z}$ Note that $(g \circ f)(n) \neq (f \circ g)(n)$

By arrow diagram:



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Composition with the identity function

Theorem: For function $f: X \to Y$ and identity functions $I_X: X \to X$ and $I_Y: Y \to Y$, $f \circ I_X = f$ and $I_Y \circ f = f$.

Proof:

We must show that for any $x \in X$, $(f \circ I_X)(x) = f(x)$ and $(I_Y \circ f)(y) = y$. Let x be any element of X.

By definition of composition and identity functions,

$$(f \circ I_X)(x) = f(I_X(x)) = f(x)$$

and similarly

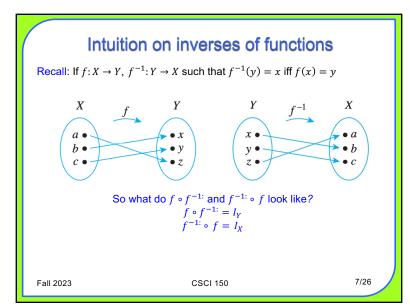
$$(I_Y \circ f)(y) = I_Y(f(y)) = f(y).$$

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QED

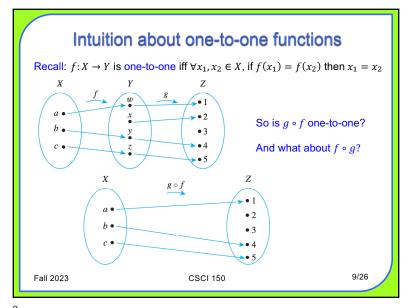
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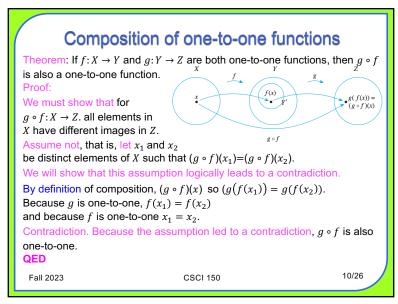
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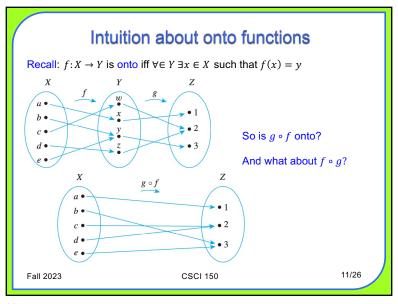


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Composition with the inverse of a function Theorem: For one-to-one and onto function $f: X \to Y$ with inverse $f^{-1}: Y \to X$ and identity function $I_v: Y \to Y$. $f^{-1}: \circ f = I_v$ and $f \circ f^{-1}: = I_x$. X = domain of fY = co-domain of ff(x) = yProof: We must show that for any $x \in X$, $(f^{-1}: \circ f)(x) = I_x(x)$ and for any $y \in Y$, $(f \circ f^{-1})(y) = I_{Y}(y).$ Let x be any element of X such that f(x) = y where $y \in Y$. By definition of inverse, $f^{-1}(y) = x$. By definition of composition, $(f^{-1} \circ f)(x) = f^{-1}(f(x)) = f^{-1}(y) = x = I_X(x)$ by definition of the identity. Let y be any element of Y such that $f^{-1}(y) = x$ where $x \in X$. By definition of inverse, $f^{-1}(y) = x$. By definition of composition, $(f \circ f^{-1})(y) = f(f^{-1}(y)) = f(x) = y = I_Y(y)$ by definition of the identity. 8/26 Fall 2023 **CSCI 150**







Composition of onto functions for V > V and $\sigma: V \to Z$ are Theorem: If $f: X \to Y$ and $g: Y \to Z$ are both onto functions then $g\circ f$ is also an we seem that under $g\circ f:X\to Z$ were element in Z is the image of some element in X. wery element in z. is the image of some element in x. if z be any element in z. By definition of not, there is some $y \in Y$ such that g(y) = z. By definition of not, there is some $x \in X$ such that f(x) = y. By definition of composition, (g * f)(x) = g(f(x)), and by substitution g * f(x) = g(f(x)) =onto function. Proof: We must show that under $g \circ f: X \to Z$ every element in Z is the image of some element in X. Let z be any element in Z. By definition of onto, there is some $y \in Y$ such that g(y) = z. By definition of onto, there is some $x \in X$ such that f(x) = y. By definition of composition, $(g \circ f)(x) = g(f(x))$, and by substitution $(g \circ f)(x) = g(f(x)) = g(y) = z$ **QED** 12/26 Fall 2023 **CSCI 150**

Today's outline

- ✓ Composition of functions
- · Cardinality and its properties
- Infinite sets

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Definitions

- A set is finite iff it is empty or can be put in one-to-one correspondence with a set of the form $\{1,2,\dots,n\}$
- An infinite set is a non-empty set that cannot be put in one-to-one correspondence with a set of the form {1,2,...,n}
- Any sets A and B have the same cardinality |A|=|B| iff there is a function from A to B that is one-to-one and onto
- An equivalence relation $\Re(x,y)$ is a binary relation that is reflexive, symmetric, and transitive on its domain
 - \mathcal{R} is reflexive on a set S iff $\forall s \in S (s, s) \in \mathcal{R}$ \leq on N
 - \mathcal{R} is symmetric on a set S iff $\forall s, t \in S$, $(s, t) \in \mathcal{R}$, $(t, s) \in \mathcal{R}$ = on \mathbf{Z}
 - \mathcal{R} is transitive on a set S iff $\forall s,t,u\in S,(s,t),(t,u)\in \mathcal{R},(s,u)\in \mathcal{R}$ > on R

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Cardinality is reflexive

Binary relation \mathcal{R} is reflexive on a set S iff $\forall s \in S (s, s) \in \mathcal{R} \leq \text{on } N$

Theorem: For any set A, cardinality is reflexive, that is, |A| = |A|.

Proof: We must show that there is a one-to-one onto function from A to A.

Let x, y be any elements of A and consider the identity function $I_A: A \rightarrow A$.

By definition of an identity function, $I_A(x) = x$ and $I_A(y) = y$.

If $I_A(x) = I_A(y)$, by substitution x = y, so I_A is one-to-one. Let y be any element of A.

Because $I_A(y) = y$, so I_A is onto.

Thus I_A is a one-to-one onto function from A to A, and |A| = |A|.

QED

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Cardinality is symmetric

Binary relation \mathcal{R} is symmetric on a set $S \forall s, t \in S, (s, t) \in \mathcal{R}$, iff $(t, s) \in \mathcal{R}$ = on Z

Lemma: Any one-to-one onto function between 2 sets has an inverse.

Proof

Let X and Y be any sets with 1-to-1 correspondence $f: X \to Y$.

Define $f^{-1}: Y \to X$ as $f^{-1}(y) = x$ iff f(x) = y. f^{-1} is a function because f is onto.

Because f is one-to-one, x is unique, and f^{-1} is f's inverse. QED

Theorem 15-22 (lecture 15, slide 22): For sets *X* and *Y* the inverse of any one-to-one correspondence is also a one-to-one correspondence.

Theorem: For any set A, cardinality is symmetric, that is, if |A| = |B| then |B| = |A|. Proof: We must show that there is a one-to-one correspondence from B to A.

By definition of equal cardinality, \exists a one-to-one onto function f from A to B. By the lemma, f^{-1} , the inverse of f, exists. By Theorem 15-22 f^{-1} is also one-to-

one and onto. Thus f^{-1} is the desired function.

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Cardinality is transitive

 \mathcal{R} is transitive on a set S iff $\forall s, t, u \in S$, (s, t), $(t, u) \in \mathcal{R}$, $(s, u) \in \mathcal{R}$

Theorem 15-22 (lecture 15, slide 22): For sets *X* and *Y* the inverse of any one-to-one correspondence is also a one-to-one correspondence.

Theorem: For any sets A, B, C, cardinality is transitive, that is, if |A| = |B| and |B| = |C| then |A| = |C|.

Proof: We must show that there is a one-to-one correspondence from ${\it A}$ to ${\it C}$.

By definition of equal cardinality, there is a one-to-one onto function f from A to B and a one-to-one onto function g from B to C.

Because (slide 10) the composition of two one-to-one functions is one-to-one, $g\circ f$ is one-to-one.

Because (slide 12) the composition of two onto functions is onto, $g\circ f$ is onto.

Thus $g \circ f$ is both onto and one-to-one and is the desired function. QED

Hence cardinality is an equivalence relation

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Today's outline

- ✓ Composition of functions
- ✓ Cardinality and its properties
- · Infinite sets

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Brace yourselves...

An infinite set and its proper subset can have the same cardinality

Consider **Z** and its proper subset $Z^{even} = \{0,2,4,...\}$.

Define the function $f: \mathbf{Z} \to \mathbf{Z}^{even}$ by the rule f(n) = 2n.

Let $a, b \in \mathbb{Z}$ such that f(a) = f(b). Then 2a = 2b, so a = b and f is one-to-one.

Let c be any element of \mathbf{Z}^{even} .

Then c is even and by definition of even, c=2m for some $m\in N$, so f(m)=c and f is onto.

Thus f establishes a one-to-one correspondence between N and its proper subset Z^{even} and $|Z| = |Z^{even}|$

But just because 2 sets are both infinite does not mean that they have the same cardinality...you have to find the function that is a one-to-one correspondence between them.

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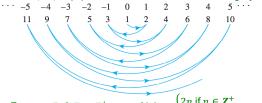
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Countability

- The paradigm for countability is the positive integers Z⁺
- A set is said to be countably infinite iff it has the same cardinality as Z⁺
- A set is said to be countable iff it is finite or countably infinite

Intuition behind $|Z^+| = |Z|$ where arrows order the counting by Z^+



For $n \in \mathbf{Z}, f: \mathbf{Z} \to \mathbf{Z}^+$ $f(n) = \begin{cases} 2n \text{ if } n \in \mathbf{Z}^+ \\ 2n + 1 \text{ otherwise} \end{cases}$

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Many sets are countably infinite

- $|N| = |N^{even}|$ even though $N^{even} \subset N$
- $|N| = |N^{odd}|$
- |N| = |Z| even though $N \subset Z$
- Let $A = \{x | x \in Z \text{ and } x \text{ mod } 3 = 0\}.$ |A| = |Z|
- Let $B = \{x | x \in Z \text{ and } x \text{ mod } 1000 = 0. |A| = |Z|$
- |A| = |B|

Cardinality equivalence proof always finds a function that is a one-to-one correspondence between two sets

Can you think of a set with cardinality larger than Z^+ ?

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Q is countable

This diagram describes a one-to-one onto function that maps $\mathbf{Z}^+ \to \mathbf{Q}^+$ so $|\mathbf{Q}^+| = |\mathbf{Z}^+|$ Follow the arrows and skip over any equivalent fractions that have already been assigned.

But $|Z^+| = |Z|$, so $|Q^+| = |Z|$

What about |Q|?

Not surprisingly, $|Q| = |Q^+|$

Are there uncountably infinite sets?

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Cardinalities and infinite sets

- Any subset of a countable set is countable
- 1874 Georg Cantor proved that *R* is uncountably infinite
- Other uncountably infinite sets
 - Number of real numbers within any interval of the real number line
 - Any set with an uncountable subset
- Cardinalities of infinite sets have assigned symbols

```
|N| = \aleph_0 read as "aleph null"
```

 $|R| = \frac{\aleph_1}{1}$ read as "aleph one"

 $\aleph_0 \leq \aleph_1$

- For any set S, $|S| < |\mathcal{P}(S)|$ even if $S = \emptyset$?
- Thus $|\mathbf{Z}| < |\mathcal{P}|\mathbf{Z}| | < |\mathcal{P}(\mathcal{P}(|\mathbf{Z}|))| < |\mathcal{P}(\mathcal{P}(\mathcal{P}(|\mathbf{Z}|)))| < \cdots$
- · This leads to infinitely many symbols for infinite cardinality

$$\aleph_0 < \aleph_1 < \aleph_2 < \dots$$

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Applications

- The set of all strings on a finite alphabet is countable (Hint: put them in alphabetical order)
- The set of all computer programs in a given programming language is countable (because they are just strings on a finite alphabet)
- The set of all functions from N⁺ to {0,1,2,3,4,5,6,7,8,9} is uncountable (proof in your text)
- · Together these indicate that

There are non-computable functions

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Proof methods (so far)

Truth table

Sequence of statements with reasons

Logic (modus ponens, modus tollens,...)

Predicate logic (quantification, vacuous truth)

Generalization from the generic particular

Proof by contradiction

Proof by contraposition

Proof by cases

Mathematical induction

Strong mathematical induction

Proof by set element

Algebraic set proof

Algebraic proof by properties of functions

Cardinality proof by one-to-one correspondence

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What you should know

- ★ Simple examples yield intuition for hypotheses
- How compositionality works
- How to build proofs with one-to-one functions, onto functions, and one-to-one correspondences
- · What cardinality is
- · How to measure the size of infinite sets

Any ques

Next up: Counting and probability

Time to finish up that Opening sheet!

Problem set 15,16 is due on Monday, November 6 at 11PM

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