

Chapter 3

Stacks and Queues

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Stacks and Queues

❑ Stack

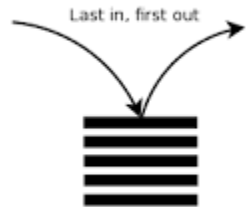
LIFO: Last-In-First-Out

FILO: First-In-Last-Out

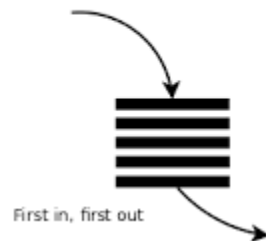
❑ Queue

FIFO: First-In-First-Out

Stack:



Queue:



stack



push

pop

queue

enqueue



dequeue

Part 1: Stacks

Stacks

- ❑ Stack

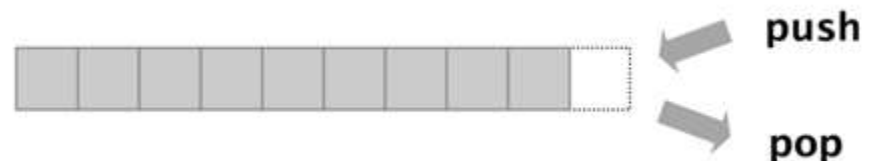
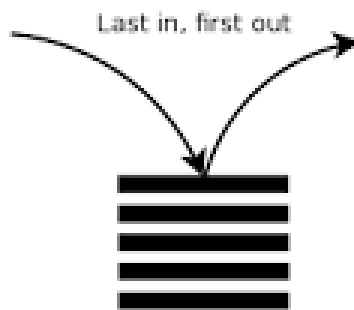
 - LIFO: Last-In-First-Out

 - FILO: First-In-Last-Out

- ❑ example: stacks of plates used in cafeterias

- ❑ $S.top$: indexes the most recently inserted element

- ❑ Underflow/Overflow



Stacks

□ Operations on stacks:

PUSH (S, x)

POP (S)

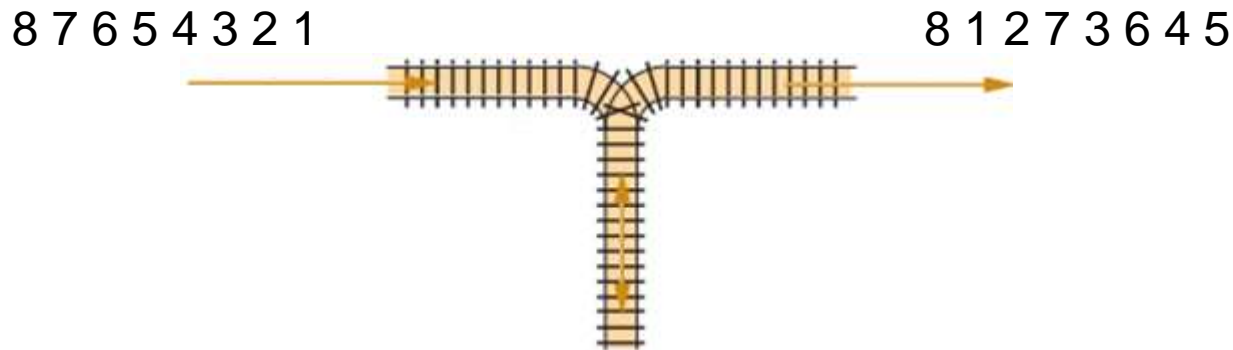
SIZE (S)

ISEMPTY (S)

TOP (S)

Stacks: Example 1

□ Stack of trains



<5,4,6,3,7,2,1,8>?

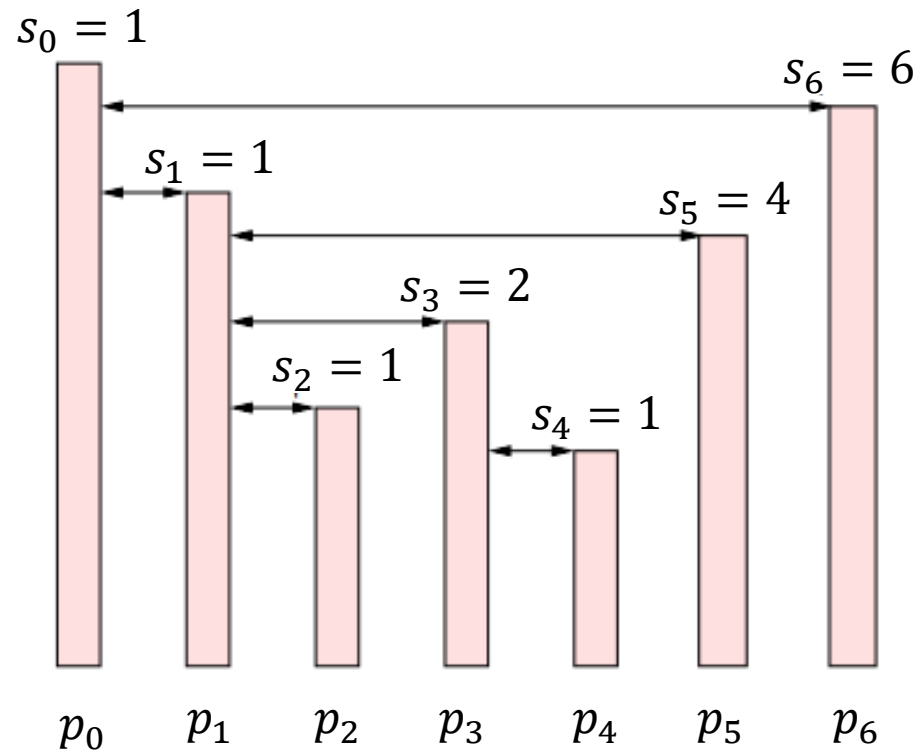
Push, Push, Push, Push, Push, Pop, Pop, Push, Pop, Pop, Push, Pop, Pop, Push, Pop

<4,3,7,8,6,2,,5,1>?

<1,8,3,6,2,7,4,5>?

Stacks: Example 2

□ Compute Spans



Stacks: Example 2 (cont.)

□ Compute Spans (cont.)

COMPUTESPANS (P)

▷ Input: n -element array P

▷ Output: n -element array S

```
1  for  $i \leftarrow 0$  to  $n - 1$ 
2      do  $k \leftarrow 0$ 
3           $done \leftarrow false$ 
4          repeat if  $P[i - k] \leq P[i]$ 
5              then  $k \leftarrow k + 1$ 
6              else  $done \leftarrow true$ 
7          until  $k = i$  or  $done$ 
8           $S[i] \leftarrow k + 1$ 
9  return  $S$ 
```


Stacks: Example 2 (cont.)

□ Compute Spans (cont.)

COMPUTESPANS2(P)

▷ we use a stack D

```
1  for  $i \leftarrow 0$  to  $n - 1$ 
2    do  $done \leftarrow \text{false}$ 
3      while ( not  $\text{ISEMPTY}(D)$  or  $done$  )
4        do if  $P[i] \geq P[\text{TOP}(D)]$ 
5          then  $\text{POP}(D)$ 
6          else  $done \leftarrow \text{true}$ 
7      if  $\text{ISEMPTY}(D)$ 
8        then  $h \leftarrow -1$ 
9        else  $h \leftarrow \text{TOP}(D)$ 
10      $S[i] \leftarrow i - h$ 
11      $\text{PUSH}(D, i)$ 
12  return  $S$ 
```

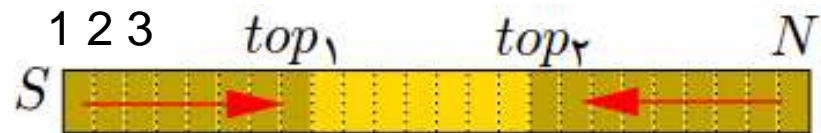
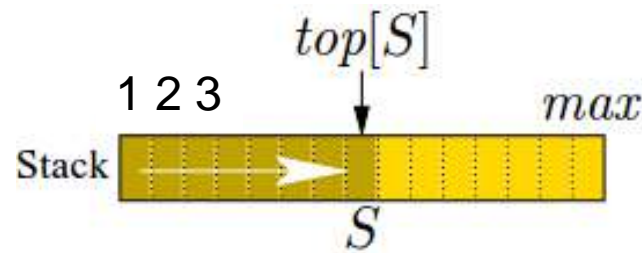
Implementation of Stack

Implementation of stack by:

- Array
- Linked list

Implementation of Stack (cont.)

❑ Implementation of Stack by Array



Implementation of Stack (cont.)

□ Implementation of Stack by Array (cont.)

SIZE(S)

1 **return** $top[S]$ ▷ assuming that initially $top[S] = 0$

ISEMPTY(S)

1 **return** $SIZE(S) = 0$

TOP(S)

1 **if** $ISEMPTY(S)$

2 **then error** ("STACK IS EMPTY")

3 **return** $S[top[S]]$

Implementation of Stack (cont.)

□ Implementation of Stack by Array (cont.)

PUSH(S, x)

```
1  if SIZE( $S$ ) =  $max$ 
2  then error ("stack is full")
3   $top[S] \leftarrow top[S] + 1$ 
4   $S[top[S]] \leftarrow x$ 
```

POP(S)

```
1  if isEmpty()
2  then error ("stack is empty")
3   $e \leftarrow S[top[S]]$ 
4   $top[S] \leftarrow top[S] - 1$ 
5  return  $e$ 
```



Implementation of Stack (cont.)

□ Implementation of Stack by Linked List

SIZE(S)

```
1 return size[ $S$ ]
```

ISEMPTY(S)

```
1 return (size[ $S$ ] = 0)
```

TOP(S)

```
1 if isEmpty( $S$ )  
2   then error ("stack is empty")  
3 return top[ $S$ ]
```

Implementation of Stack (cont.)

□ Implementation of Stack by Linked List (cont.)

PUSH(S, x)

```
1  $top[S] \leftarrow \text{ALLOCATE-NODE}(x, top[S])$   
2  $size[S] \leftarrow size[S] + 1$ 
```

POP(S)

```
1 if  $ISEMPTY(S)$   
2   then error ("STACK IS EMPTY")  
3  $n \leftarrow top[S]$   
4  $temp \leftarrow element[n]$   
5  $top[S] \leftarrow next[n]$   
6  $size[S] \leftarrow size[S] - 1$   
7  $\text{FREE-OBJECT}(n)$   
8 return  $temp$ 
```



Part 2: Queues

Elementary Data Structures

- ☐ Lists
- ☐ Stacks
- ☐ **Queues**
- ☐ Trees

Queues

- ❑ Queue:

FIFO: First-In-First-Out

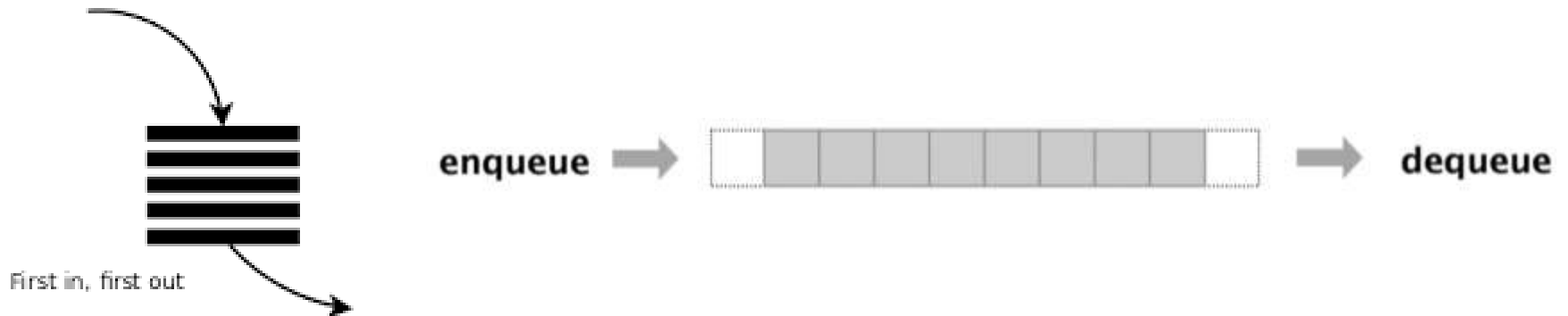


- ❑ example: a line of customers waiting to pay a cashier

- ❑ front/head

- ❑ rear/tail

- ❑ Underflows/Overflows



Queues (cont.)

□ Operations on queues:

ENQUEUE (Q, x)

DEQUEUE (Q, x)

SIZE (Q)

ISEMPTY (Q)

FRONT – ELEMENT (Q)

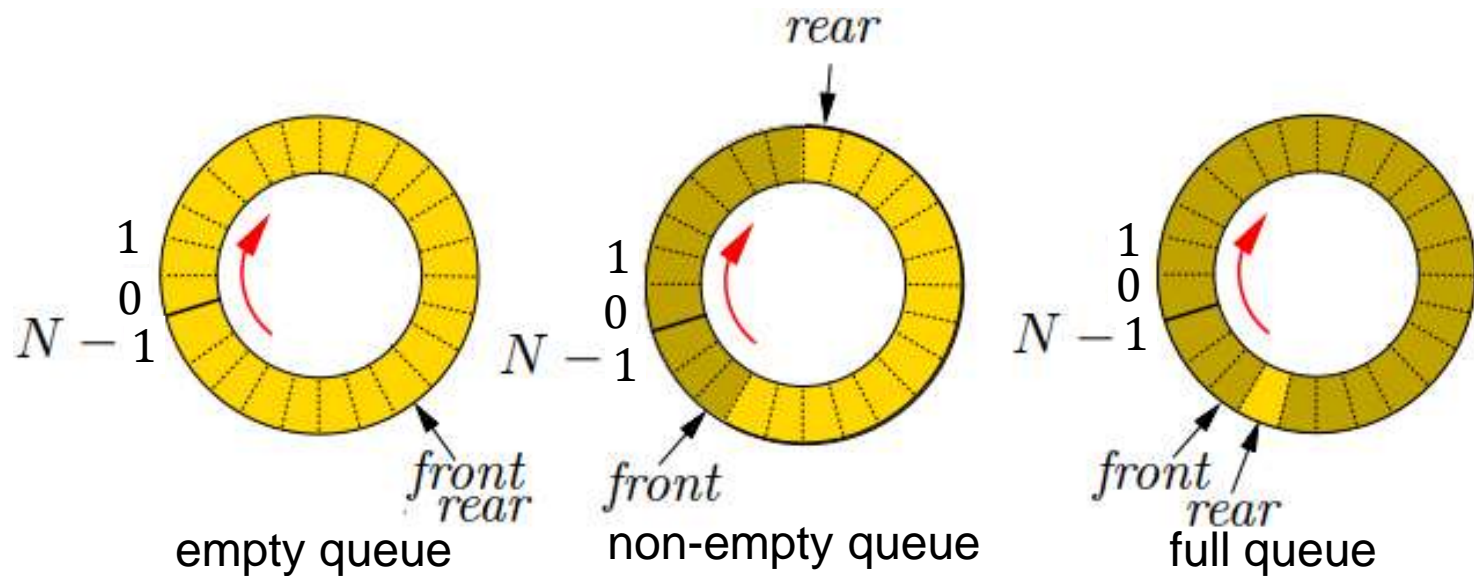
Implementation of Queue

Implementation of Queue by

- Circular Array
- Linked List

Implementation of Queue (cont.)

□ Implementation of Queue by Circular Array



At start: $front[Q] = rear[Q] = 0$

empty queue: $front[Q] = rear[Q]$

full queue: $(max - front[Q] + rear[Q]) \bmod max = max - 1$

Implementation of Queue (cont.)

□ Implementation of Queue by Circular Array (cont.)

SIZE (Q)

1 **return** $(max - front[Q] + rear[Q]) \bmod max$

ISEMPTY (Q)

1 **return** $(front[Q] = rear[Q])$

FRONT-ELEMENT (Q)

1 **if** ISEMPTY (Q)
2 **then** **error** "Queue is empty"
3 **return** $Q[front[Q]]$

Implementation of Queue (cont.)

□ Implementation of Queue by Circular Array (cont.)

ENQUEUE(Q, x)

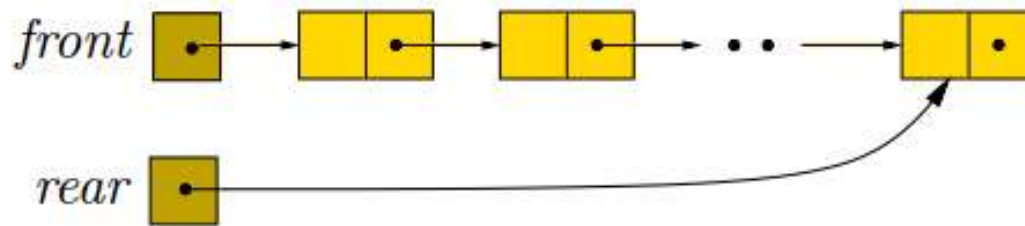
```
1  if  $\text{Size}(Q) = \text{max} - 1$ 
2    then error "Queue is full"
3   $Q[\text{rear}[Q]] \leftarrow x$ 
4   $\text{rear}[Q] \leftarrow (\text{rear}[Q] + 1) \bmod \text{max}$ 
```

DEQUEUE(Q)

```
1  if  $\text{isEmpty}()$ 
2    then error "Queue is empty"
3   $\text{temp} \leftarrow Q[\text{front}[Q]]$ 
4   $\text{front}[Q] \leftarrow (\text{front}[Q] + 1) \bmod \text{max}$ 
5  return  $\text{temp}$ 
```

Implementation of Queue (cont.)

❑ Implementation of Queue by Linked List



isEmpty(*Q*)
1 **return** *size*[*Q*] = 0

Implementation of Queue (cont.)

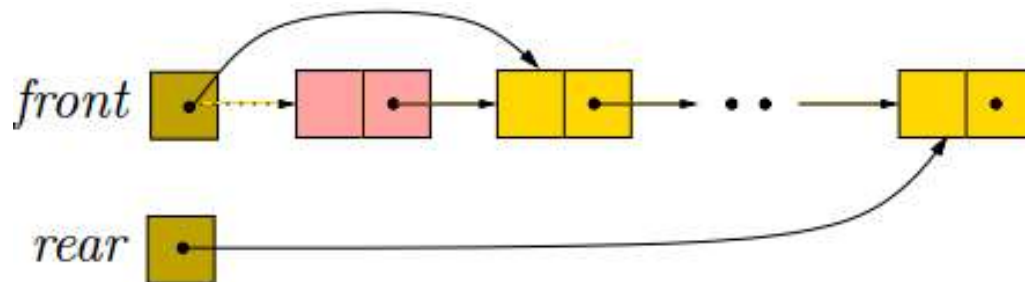
□ Implementation of Queue by Linked List (cont.)

ENQUEUE (Q, x)

1 $next[rear[Q]] \leftarrow \text{ALLOCATE-NODE}(x, \text{null})$

2 $rear[Q] \leftarrow next[rear[Q]]$

3 $size[Q] \leftarrow size[Q] + 1$



Implementation of Queue (cont.)

❑ Implementation of Queue by Linked List (cont.)

DEQUEUE(Q)

```
1  if ISEMPTY( $Q$ )  
2  then error 'QUEUE IS EMPTY'  
3   $n \leftarrow \text{front}[Q]$   
4   $x \leftarrow \text{element}[n]$   
5   $\text{front}[Q] \leftarrow \text{next}[\text{front}[Q]]$   
6  FREE-NODE( $n$ )  
7   $\text{size}[Q] \leftarrow \text{size}[Q] - 1$   
8  return  $x$ 
```

