



Struktur Data & Algoritma

(Data Structures & Algorithms)

Red Black Tree

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Version 3.0 - Internal Use Only

Motivation

- **AVL Trees take top down insertion and bottom up balancing**
 - Need recursive implementation
- **Red-Black Tree can be top-down balancing and insertion**
 - Can be implemented iteratively
- **Red-Black Tree is used as the implementation of `java.util.TreeMap` and `java.util.TreeSet`**



Objectives

- Understand the definition, properties and operations of **Red-Black Trees**.



Outline

- **Red-Black Trees**
 - Definition
 - Operation



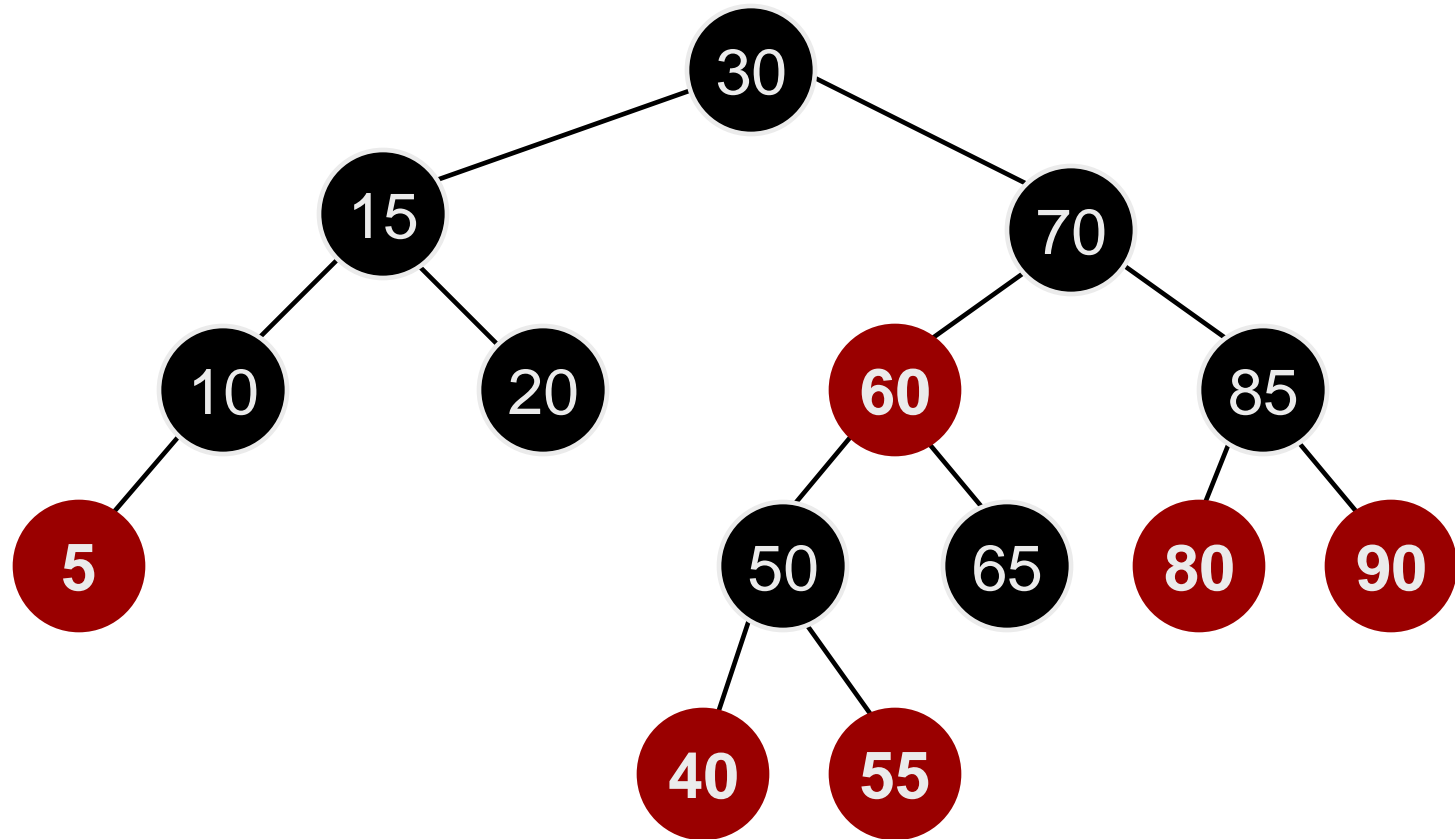
Red-Black Trees: Definition

1. Every node is colored either **red** or black
2. The root is black
3. If a node is **red**, its children must be black
 - consecutive **red** nodes are disallowed
4. Every path from a node to a null reference must contain the same number of black nodes

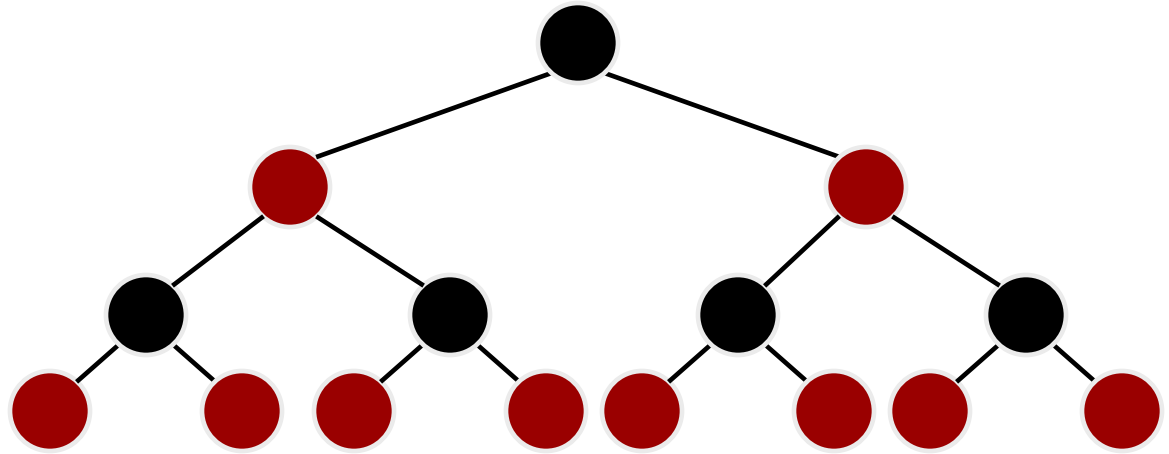
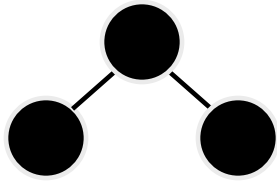


Red-Black Trees

- The insertion sequence is 10, 85, 15, 70, 20, 60, 30, 50, 65, 80, 90, 40, 5, 55



Red-Black Trees: properties



$$2^B - 1 \leq N \leq 2^{2B} - 1$$

$$2^B \leq N + 1 \leq 2^{2B}$$

$$\log 2^B = B \quad \log 2^{2B} = 2B$$

$$B \leq \log(N + 1) \quad \log(N + 1) \leq 2B$$

$$\frac{1}{2} \log(N + 1) \leq B \leq \log(N + 1)$$

$$\log(N + 1) \leq H \leq 2 \log(N + 1)$$

- **B = total black nodes from root to leaf**
- **N = total all nodes**
- **H = height**

All operation guaranteed logarithmic.

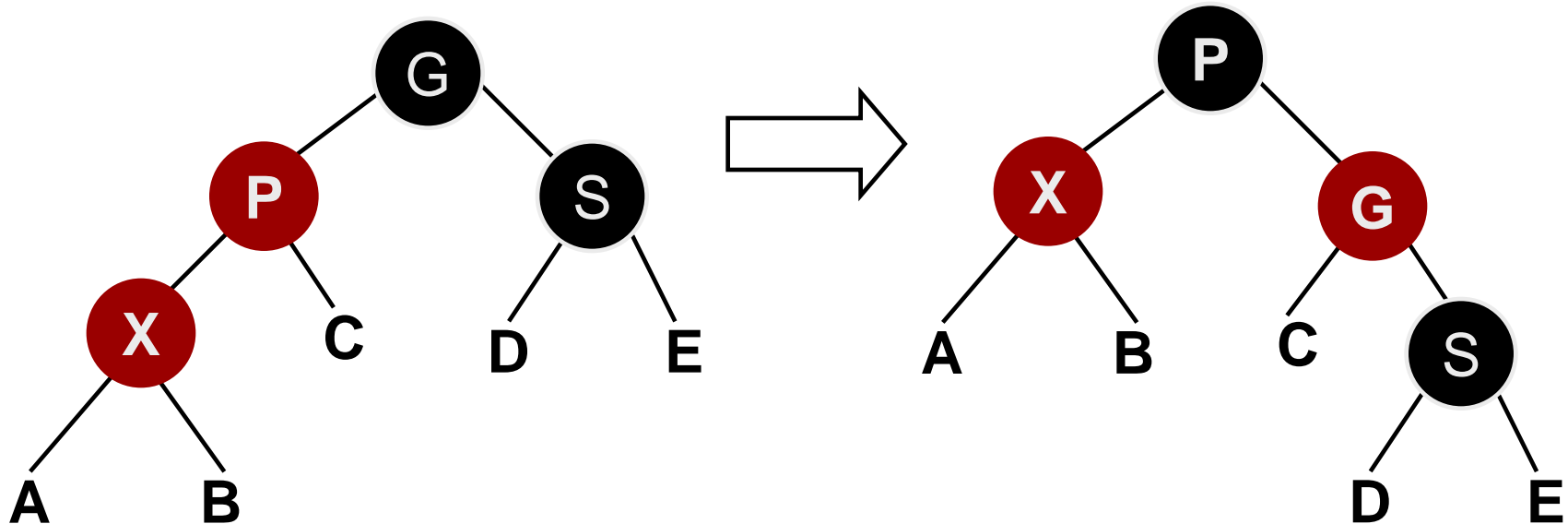


Insertion

- **A new node must be colored red**
 - why?
 - new item is always inserted as a leaf in the tree
 - if we color a new item black, then the number of black nodes from root would be different (violate property #4)
 - if the parent is black, no problem!
 - if the parent is **red**, we create two consecutive red nodes
- **Convention: null nodes are black**



Single Rotation



■ Case after insertion:

- Consecutive red
- Sibling of parent is black
- Outer node (left-left or right-right)

X: new node

P: parent

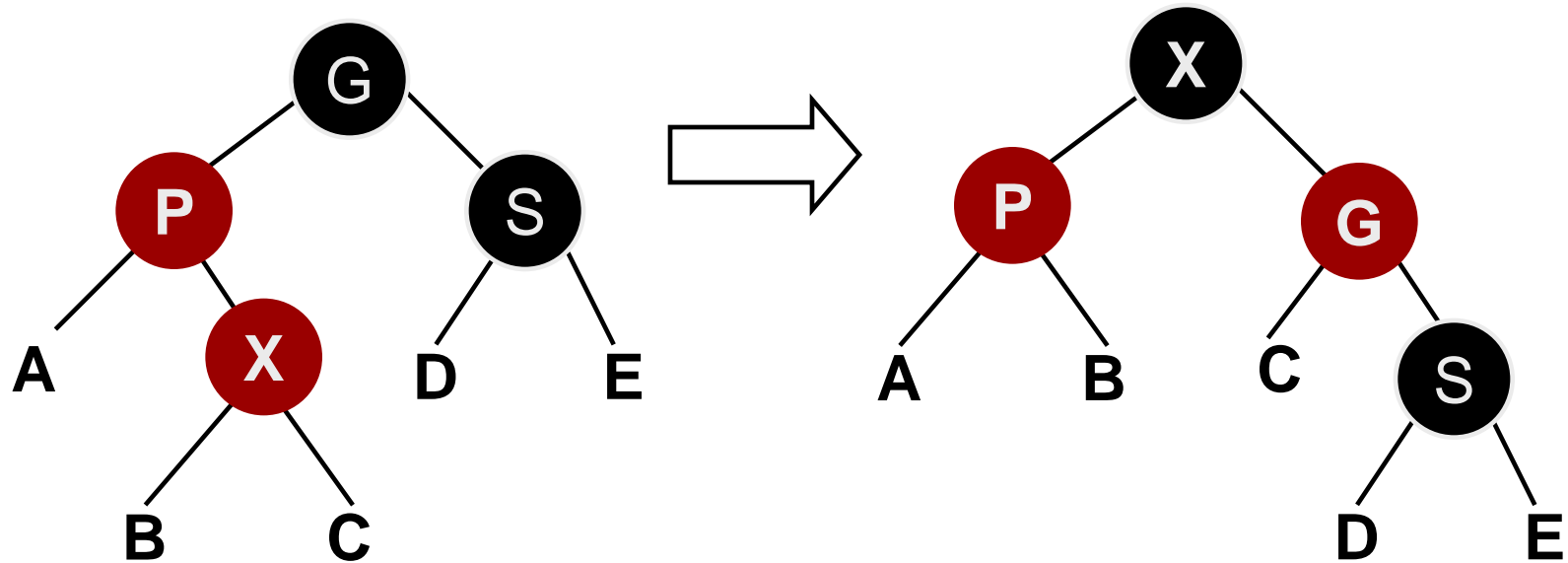
S: sibling

G: Grandparent

Maintain number of black nodes in every path



Double Rotation



■ Case after insertion:

- Consecutive red
- Sibling of parent is black
- Inner node (left-right or right-left)

X: new node

P: parent

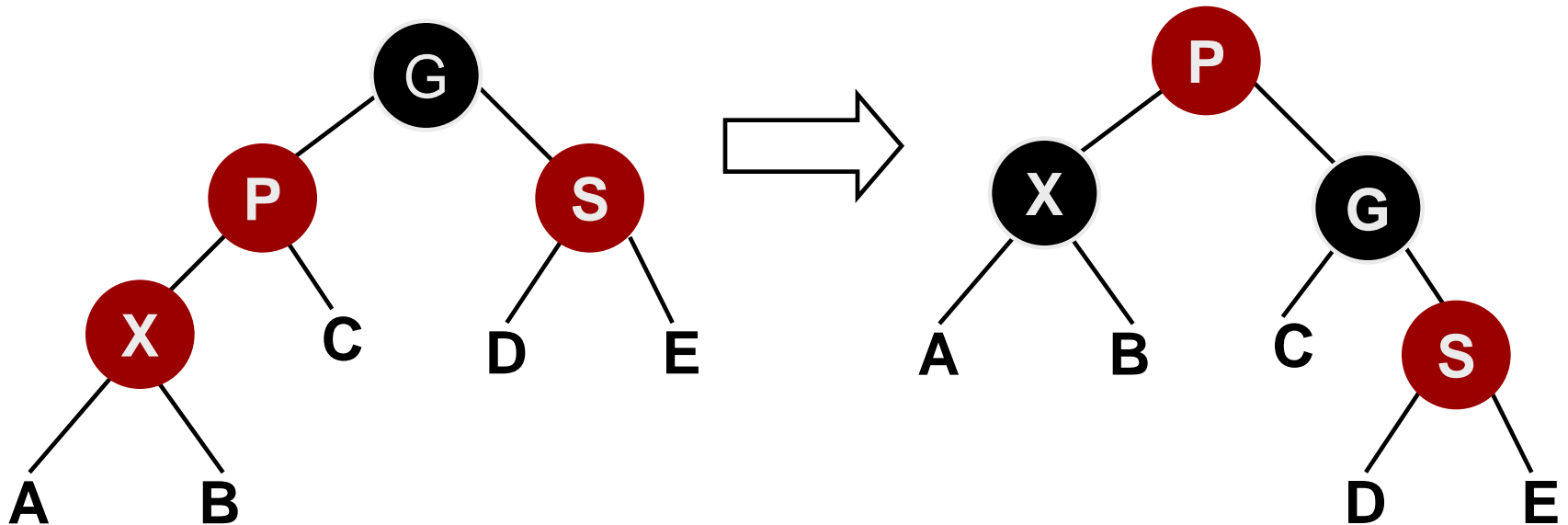
S: sibling

G: Grandparent

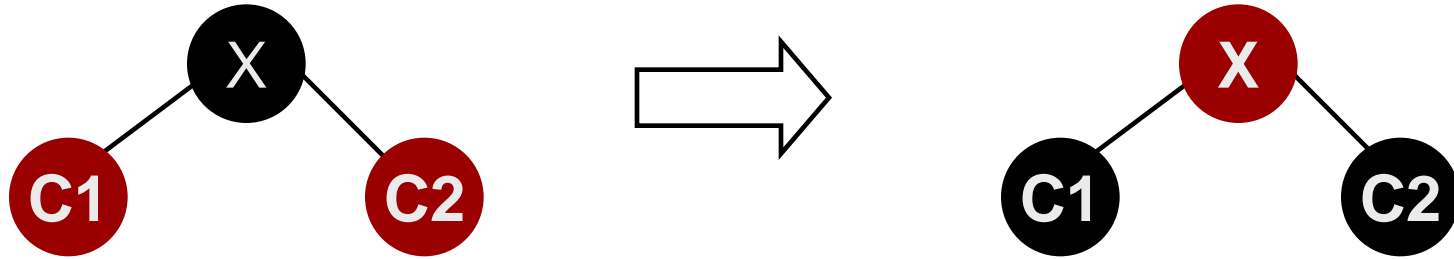


Single Rotation

- **Case after insertion:**
 - Consecutive red
 - Sibling of parent is red
 - Outer node (left-left or right-right)



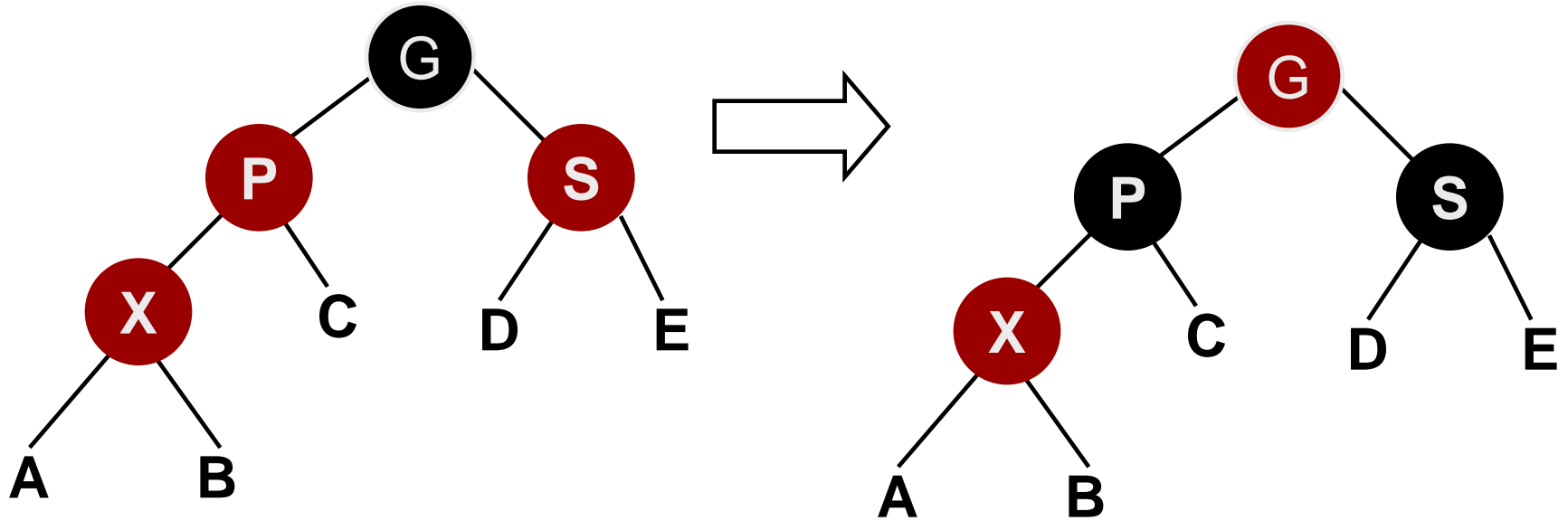
Color Flip (1)



- **During the insertion process:**
 - if we fall into black node with two red children, flip the color
 - do rotation if consecutive red occurs
 - set the root to black (If the flip process involves root, root becomes red.)



Color Flip (2)

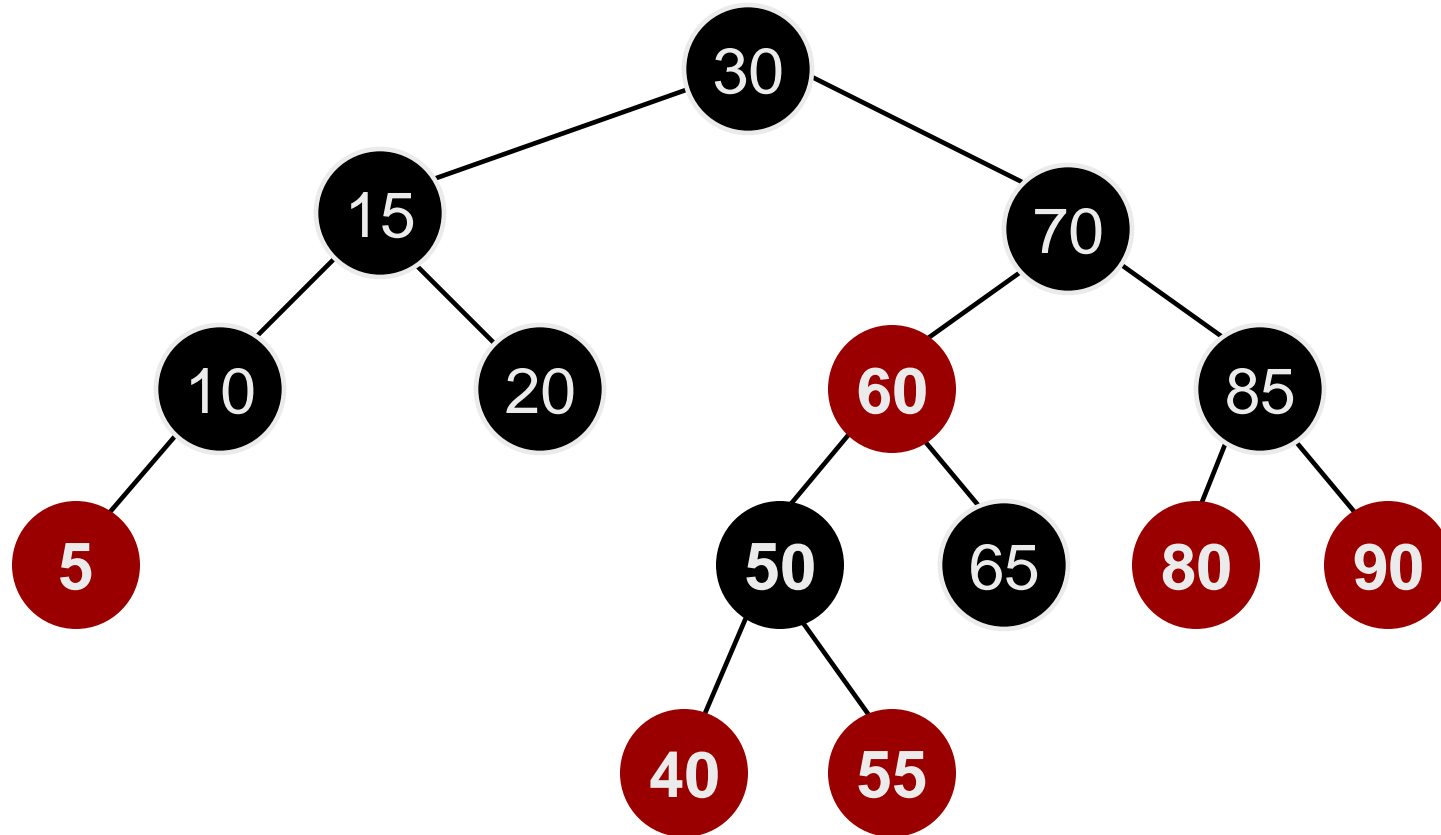


Insertion

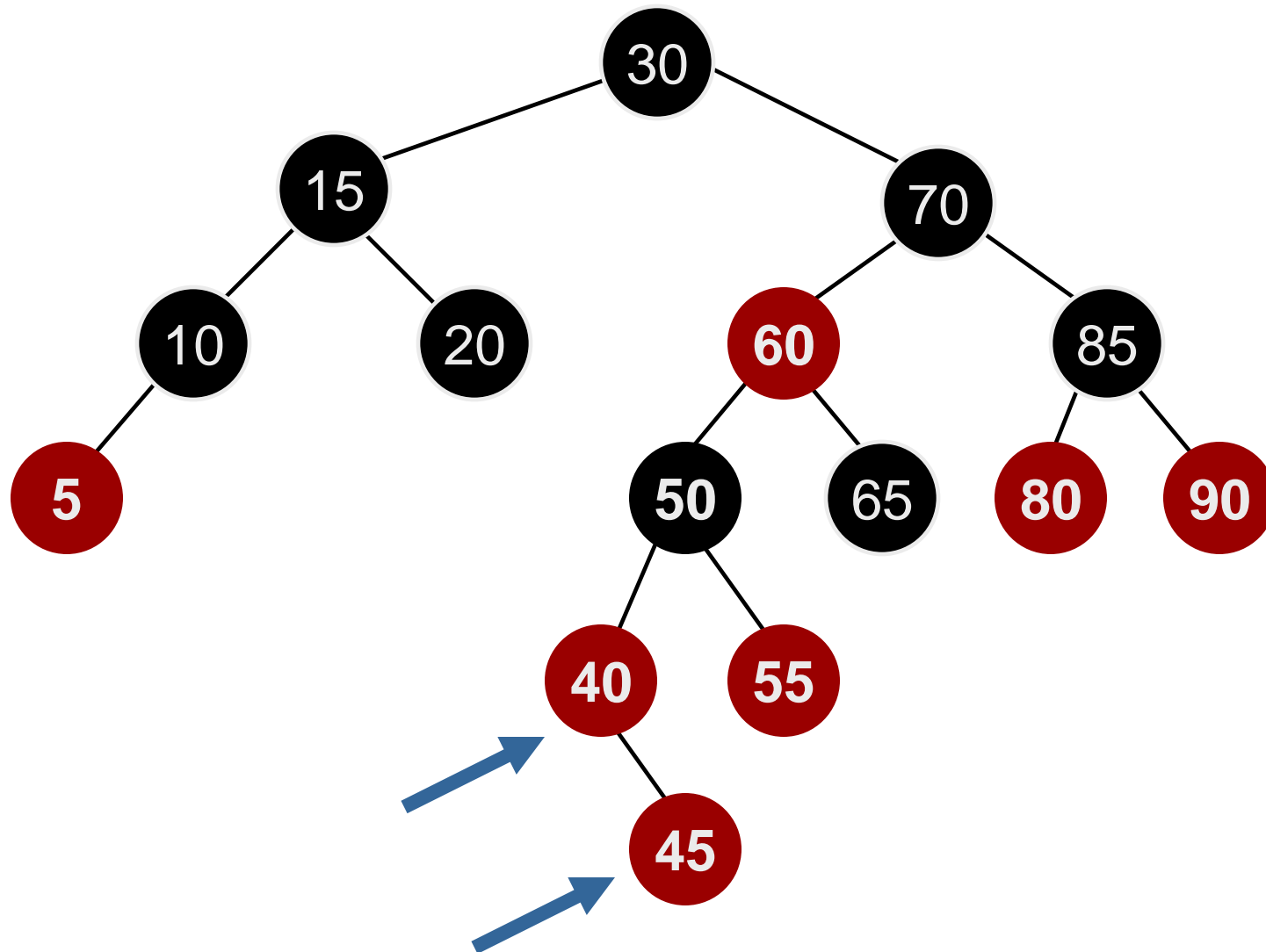
- **Top down insertion: as BST**
- **Bottom-up:**
 - check for two consecutive red nodes
 - if sibling is red → color flip
 - if sibling is black → rotation
- **Implementation: quite straight forward, similar to AVL tree**



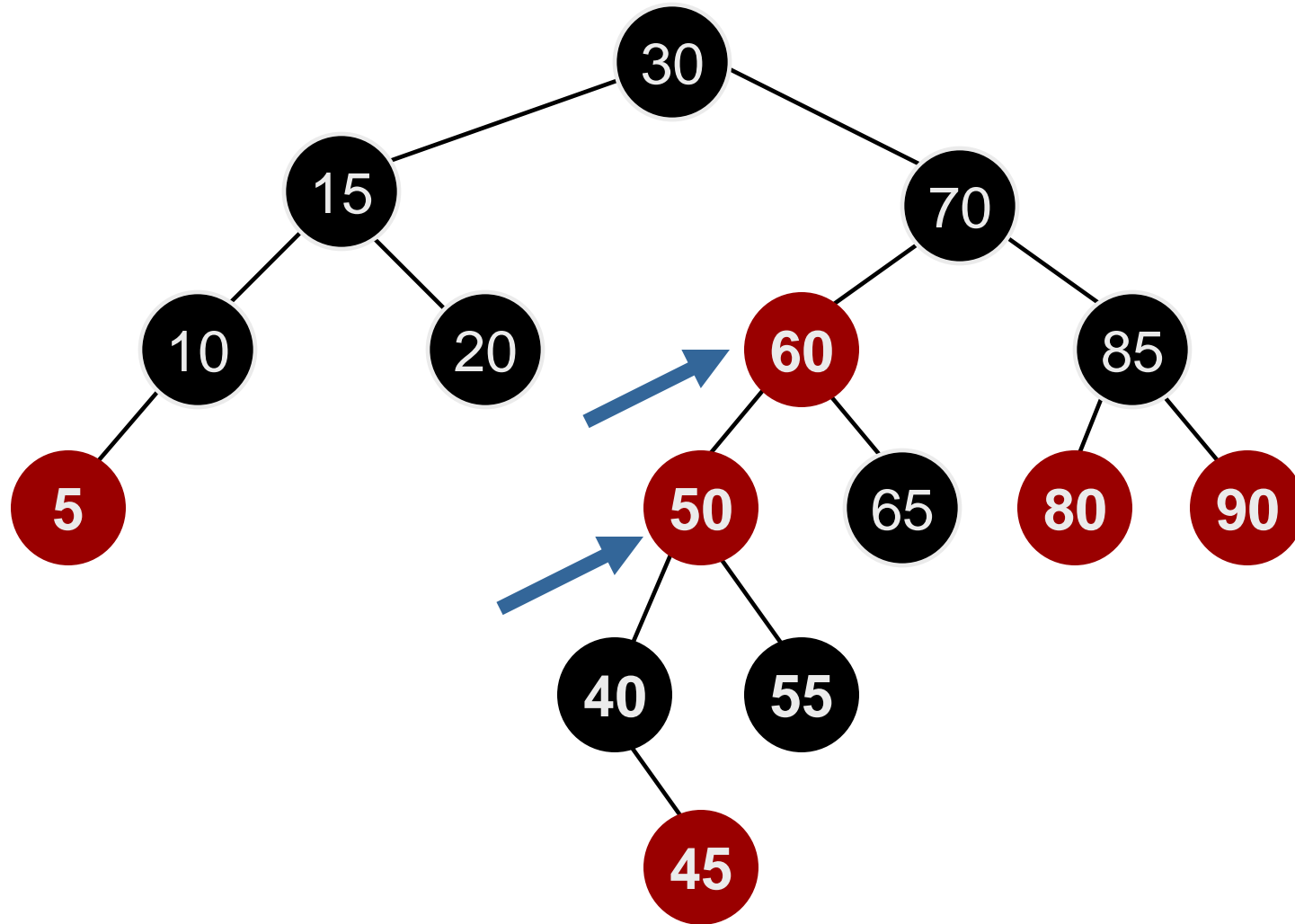
Insert 45 (original)



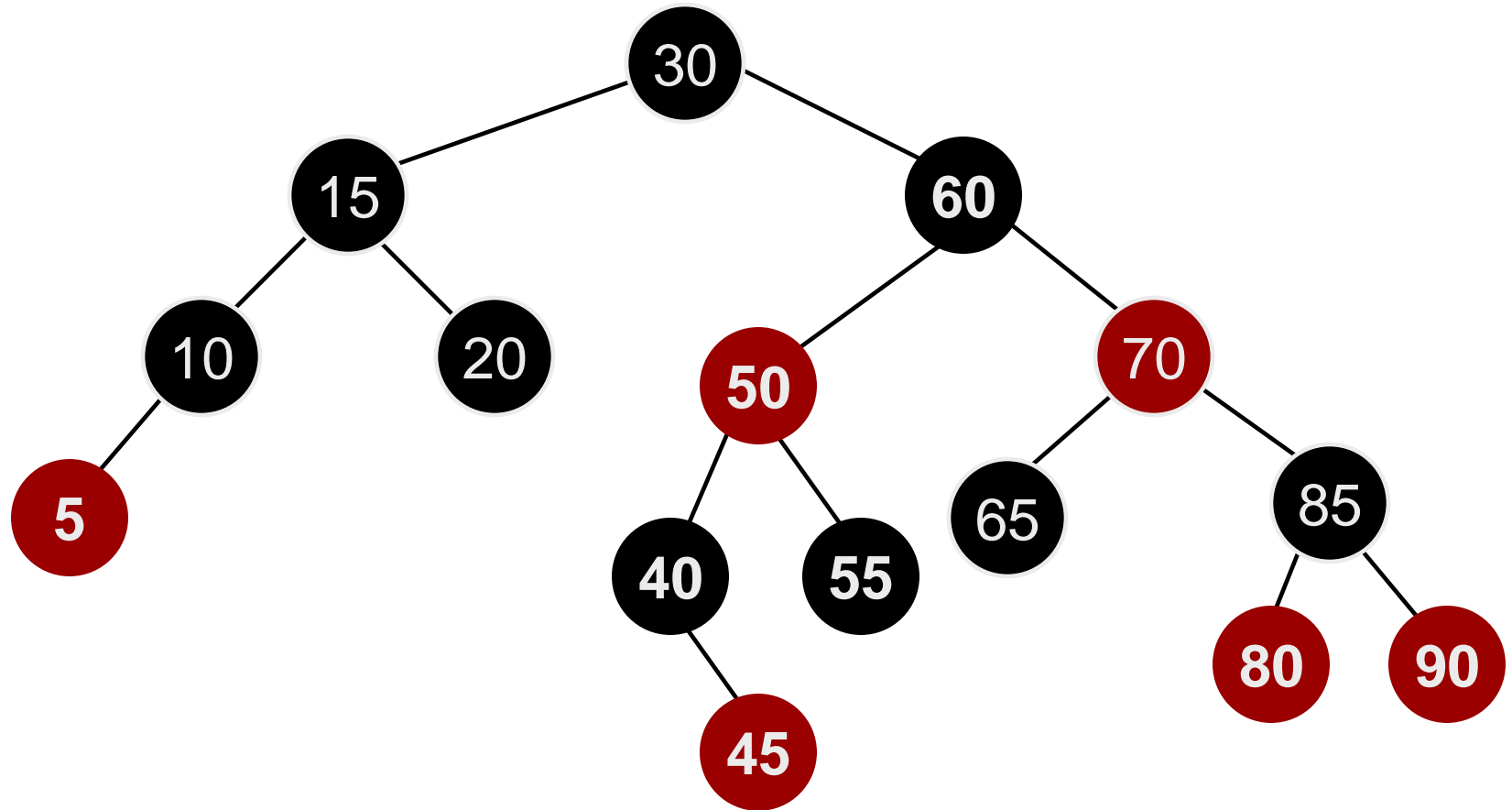
Insert 45



Insert 45 (after color flip)



Insert 45 (single rotation)



Top Down Insertion

- **Top down insertion, one pass:**

- check if a node has both children red → color flip
 - color flip might cause two consecutive red nodes
 - if after a color flip cause two consecutive red nodes → perform a single or double rotation
- insert the new red node at the leaf
 - check for two consecutive red nodes

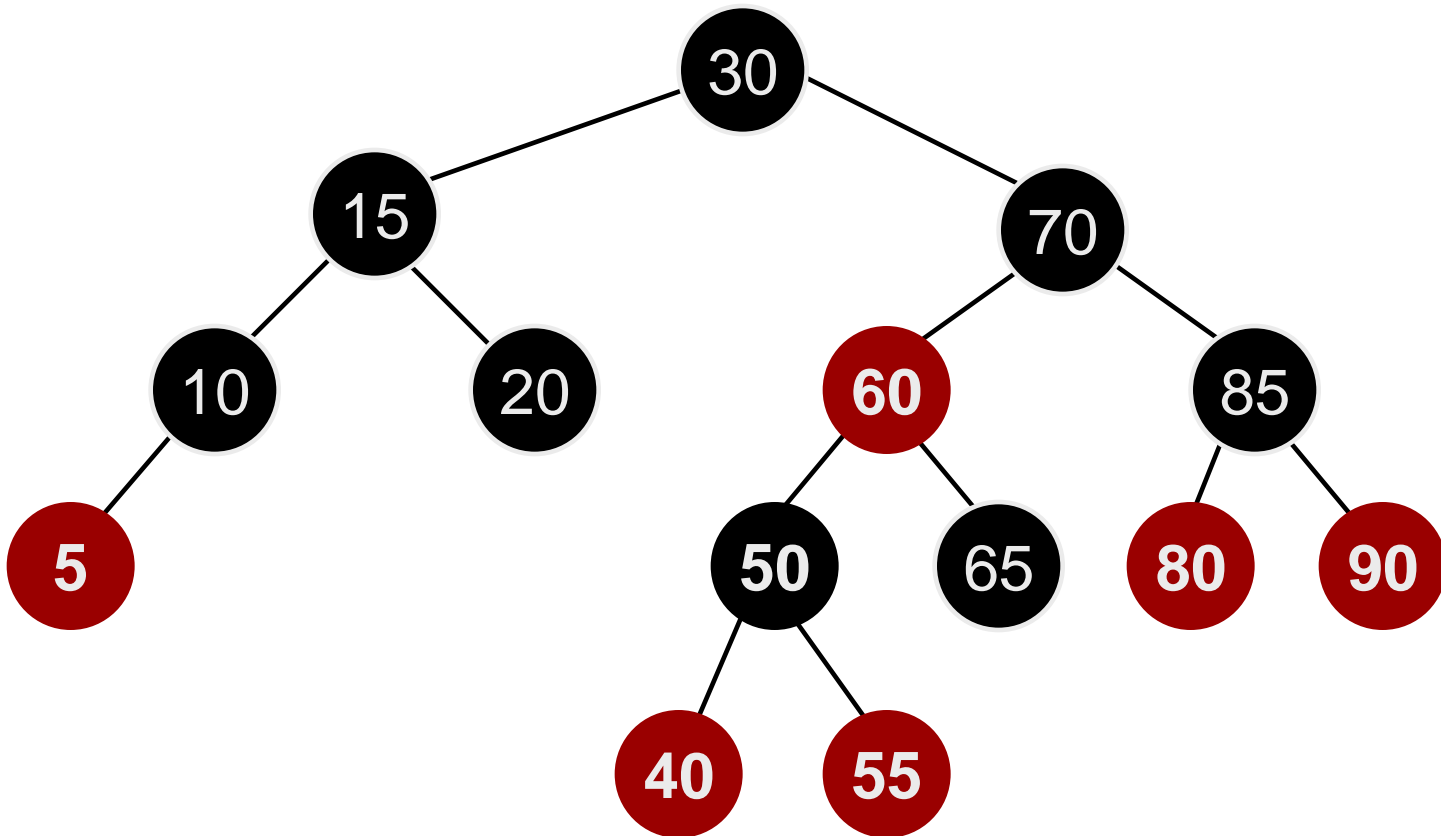
- **No need to fix on the way back to root**

- **Implementation:**

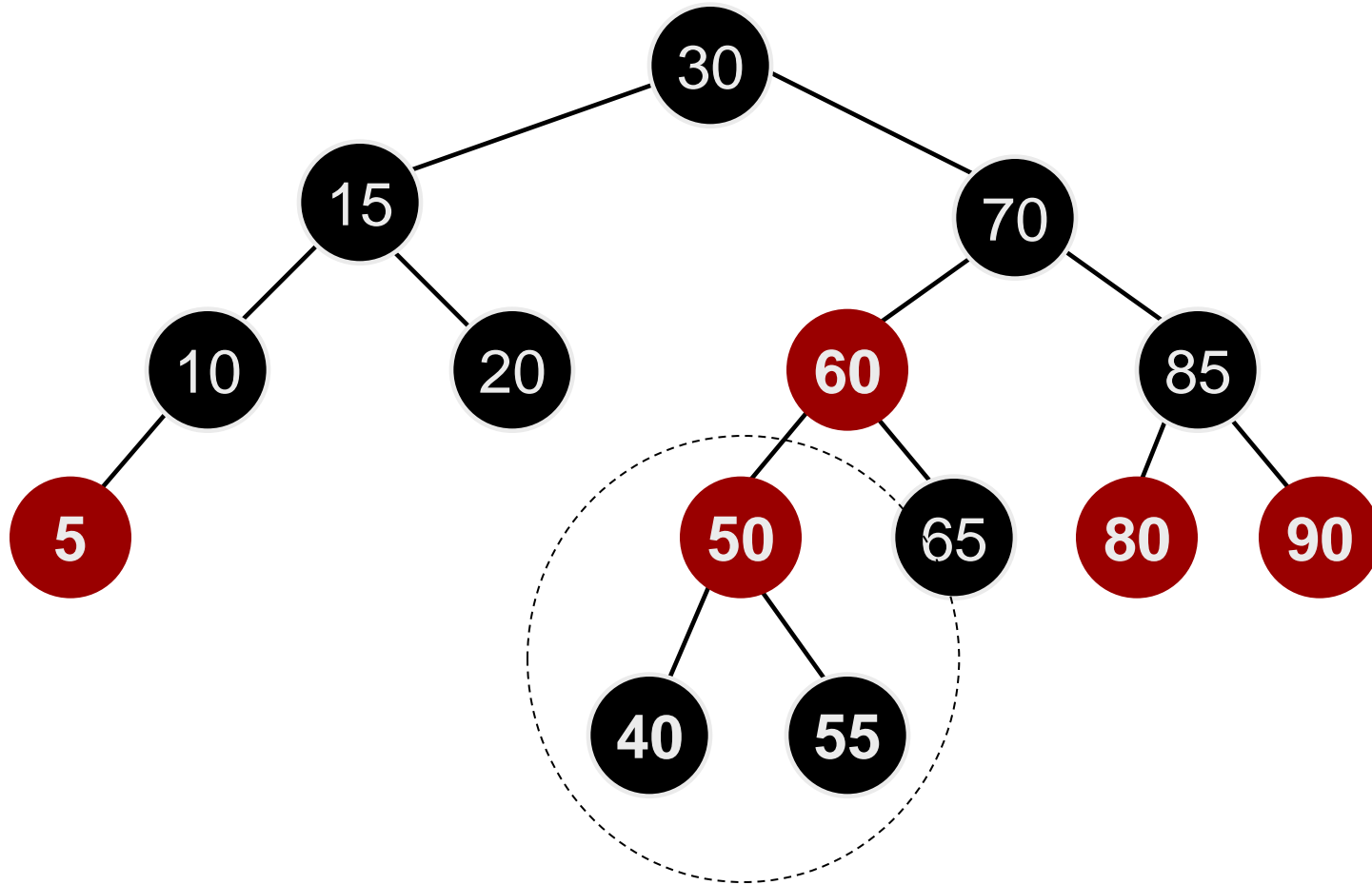
- need to store current, parent, grand parent, great grand parent.



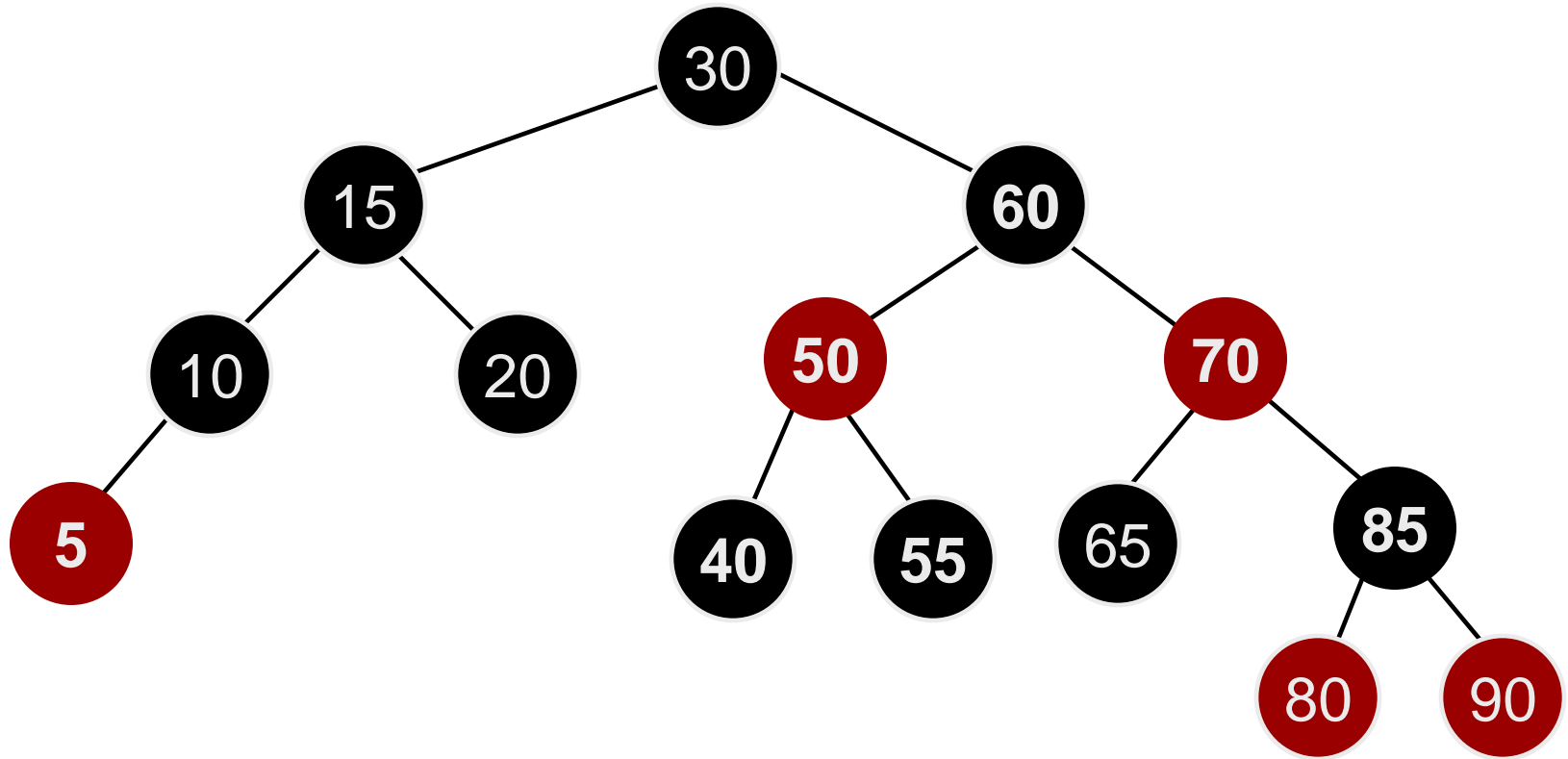
Insert 45 (original)



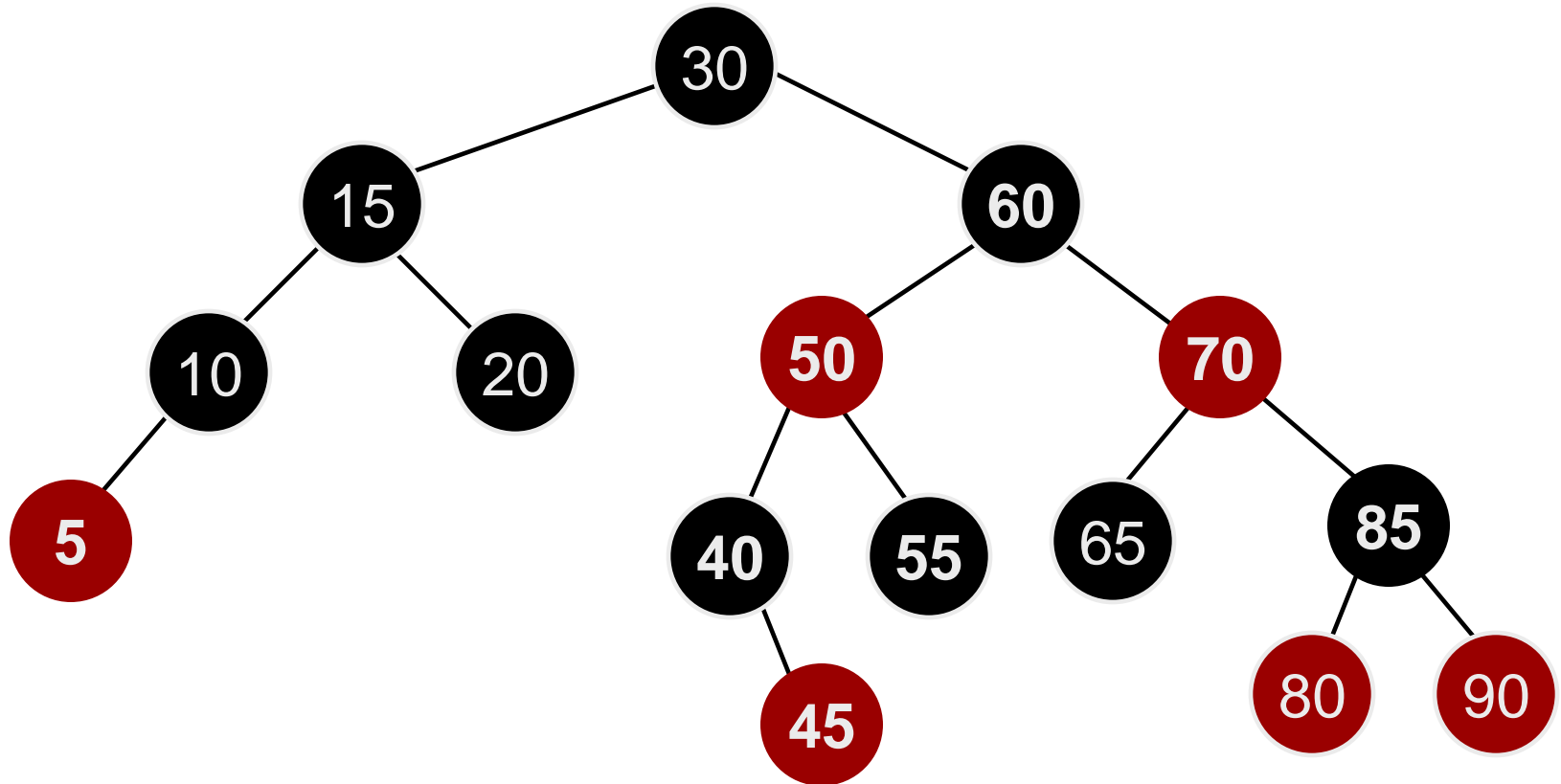
Insert 45 (after color flip)



Insert 45 (single rotation)



Insert 45



Summary

- **Red-Black trees use color as balancing information instead of height in AVL trees.**
- **An insertion may cause a local perturbation (two consecutive red nodes)**
- **The perturbation is either**
 - resolved locally (rotations), or
 - propagated to a higher level in the tree by recoloring (color flip)
- **$O(1)$ for a rotation or color flip**
- **At most one restructuring per insertion.**
- **$O(\log n)$ color flips**
- **Total time: $O(\log n)$**



Further Reading

- **Red-Black Tree Applet**

<http://webdiis.unizar.es/asignaturas/EDA/AVLTree/avl-tree.html>

