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2.1-3 - Consider the *searching problem*:

Input: A sequence of n numbers $A = \langle a_1, a_2, ..., a_n \rangle$ and a value v.

Output: An index i such that v = A[i] or the special value NIL if v does not appear in A.

Write pseudocode for Linear Search, which scans the sequence, looking for v. Using a loop invariant, prove that your algorithm is correct. Make sure that your loop invariant fulfills the necessary three properties.

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2.2-1 - Express the function
$$\frac{n^3}{1000}-100n^2-100n+3$$
 in terms of Θ notation.

 $\Theta(n^3)$

2.2-2' - Write pseudocode for Selection Sort.

```
Selection Sort(A):
    for cur = 0 to A.length:
 2
         small = A[cur]
         index = cur
 4
         for k = cur + 1 to A.length:
 5
             if A[k] < \text{small}:
                  small = A[k]
                  index = k
 7
 8
         temp = A[cur]
                                       // Swap the values
 9
         A[cur] = small
10
         A[index] = temp
```

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2.3-3 - Use mathematical induction to show that when n is an exact power of 2, the solution of the recurrence

$$T(n) = \begin{cases} 2 & \text{if } n = 2\\ 2T(\frac{n}{2}) + n & \text{if } n = 2^k \text{ for } k > 1 \end{cases}$$

is $T(n) = n \log_2 n$.

We show that T holds for n=2

$$T(2) = 2\log_2 2 = 2 * 1 = 2$$

Assuming $T(\frac{n}{2}) = \frac{n}{2} \log_2 \frac{n}{2}$

$$\begin{split} T(n) &= 2 \Big(\frac{n}{2} \log_2 \frac{n}{2}\Big) + n \\ &= 2 \Big(\frac{n}{2} (\log_2 n - \log_2 2)\Big) + n \\ &= 2 \Big(\frac{n}{2} (\log_2 n - 1)\Big) + n \\ &= (n \log_2 n - n) + n \\ &= n \log_2 n \end{split}$$

Therefore $T(n) = n \log_2 n$.

2.3-4 - We can express insertion sort as a recursive procedure as follows. In order to sort A[1..n], we recursively sort A[1..n-1] and then insert A[n] into the sorted array A[1..n-1]. Write a recurrence for the running time of this recursive version of insertion sort.

Undefined

Induction Theorem:

$$P(n_0) \land \left(\forall n_i > n + o.P(n_i) \rightarrow P(n_{i+1}) \right) \rightarrow \forall n > n_o P(n)$$

1.1-1 - Use mathematical induction to show that $T(n) = \frac{n^2 + n}{2}$ for

$$T(n) = \begin{cases} 1 & \text{if } n = 1 \\ T(n-1) + n & \text{if } n > 1 \end{cases}$$

We show that T holds for $n_0 = 1$

$$T(1) = \frac{1^2 + 1}{2} = \frac{2}{2} = 1$$

Assuming $T(n-1) = \frac{(n-1)^2 + (n-1)}{2}$ and n > 1

$$T(n-1) = \frac{(n-1)^2 + n - 1}{2}$$

$$= \frac{n^2 - 2n + 1 + n - 1}{2}$$

$$= \frac{n^2 - n}{2}$$

By the definition of T(n)

$$T(n) = \frac{T(n-1) + n}{2}$$

$$= \frac{n^2 - n}{2} + n$$

$$= \frac{n^2 - n + 2n}{2}$$

$$= \frac{n^2 + n}{2}$$

Therefore

$$T(n) = \frac{n^2 + n}{2}$$

By the induction theorem:

$$T(n) = \frac{n^2 + n}{2} \ \forall n \in \mathbb{N}, n \neq 0$$