

Discrete Math
Staff Graded Assignment-2

2023eBCS072

① The seven letters provided are a, b, c, e, f, g, h

② 5 from 7 letters need to be permuted

$$P(7, 5) = \frac{7!}{(7-5)!}$$

$$= \frac{7!}{2!}$$

$$= 2,520$$

③ There are 2 ways to interpret "these strings" in this question.

IF it refers to all strings of length 5 → (i), strings from ② → (ii)

③ (i) bahx y ←

Since 3 letters are used in "bah", x can be one of the remaining 4 letters.

After x is chosen, y can be 1 of the 3 remaining letters

So total possibilities are

$$4 \times 3 = 12 \text{ strings}$$

③ (ii) bahx y ←

Since there are no constraints regarding repetition, x and y can be any of the 7 letters, so

$$7 \times 7 = 49 \text{ possible strings}$$

④ Since d is not in the list of elements, "bad" cannot be a substring, so the possible number of strings with no repeats is

2,520 (as discussed in ②)

② Let X be a given string.

~~X_i~~ is the i^{th} letter of the string, so X_1 is the first letter

$X - X_i$ is the string X without the i^{th} letter.

XA is the string A appended to X

So, if $X = 'abcd'$, $X_2 = 'b'$, $X - X_2 = 'acd'$, $XX_2 = 'abcd b'$

R_x is the reverse of a string.

Basis step:

If length of $X = 1$,

$$R_x = X$$

Recursive step

$$R_x = R_{x-x_i} X_i$$

Using this with the given example of $abcd$

$$R_{abcd} = R_{bcd} a = dcba$$

$$R_{bcd} = R_{cd} b = dc b$$

$$R_{cd} = R_d c = dc$$

$$R_d = d$$

③ Let the 2 partitions formed by $K_{m,n}$ be M (m vertices) and N (n vertices respectively)

- ~~Every~~ Vertices in M are all ~~are~~ connected by an edge to each vertex in N , so their degree is n for each vertex in M
- Likewise all vertices in N have degree m as they are each connected to all ~~a~~ m vertices in M with an edge.

~~Case~~ i) $m > n$

In this case, all vertices in N will be listed first so the degree sequence is

$$\underbrace{m, m, \dots, m}_{n \text{ times}} \quad \underbrace{n, n, \dots, n}_{m \text{ times}}$$

ii) $n > m$

Here, the vertices in M are listed first so, the degree sequence is

$$\underbrace{n, n, \dots, n}_{m \text{ times}} \quad \underbrace{m, m, \dots, m}_{n \text{ times}}$$

iii) $n = m$

In this case, the order doesn't matter
The degree sequence is

$$\underbrace{m, m, \dots, m}_{m \text{ times}} \quad \underbrace{n, n, \dots, n}_{n \text{ times}} \quad \left| \text{ both are the same as } m = n \right.$$

④ An Euler circuit can be represented by the cycle C_i , C_i is a subgraph of G , $i = |V_G|$

For C_i to be a bipartite graph, i should be even and the two partitions are each of size $i/2$

So, if $C_i \subseteq K_{m,n}$, $i = |V_{m,n}| = m+n$, $m = n = i/2$

So $K_{m,n}$ has an Euler circuit iff $m=n$

⑥ If any one edge is removed from an Euler circuit, it becomes an Euler path. ~~to be an Euler path~~

So if $m=n$, $K_{m,n}$ has an Euler path.

Similarly, if a terminal vertex and the edge incident to it are removed, it is an Euler path for G , $V = V - v$, $E = E - e$

Say one vertex (terminal) with the corresponding edge also,

~~The~~ The vertex removed was removed from one of the partition.

So the new graph is an Euler path for either $K_{m-1,n}$ or $K_{m,n-1}$

~~to be~~ Now, both the terminal vertices lie in ~~different~~ same partitions since edges in the Euler path are odd. It starts and ends in the partition with more vertices, so if ~~another~~ terminal vertex and edge were to be removed it'd be from that partition, forming an Euler Path for $K_{m-1,n-1}$.

For every edge removed, it is an Euler path for $K_{m,n}$, $|m-n| \leq 1$

So there is an Euler path for every $K_{m,n}$ if $|m-n| \leq 1$