



Doubt Clearing Session - Part III

Foundation Course on Data Structures & Algorithm - III

Any 1D Array

→ Doubt

→ Min swaps required to group $\leq K$ element together

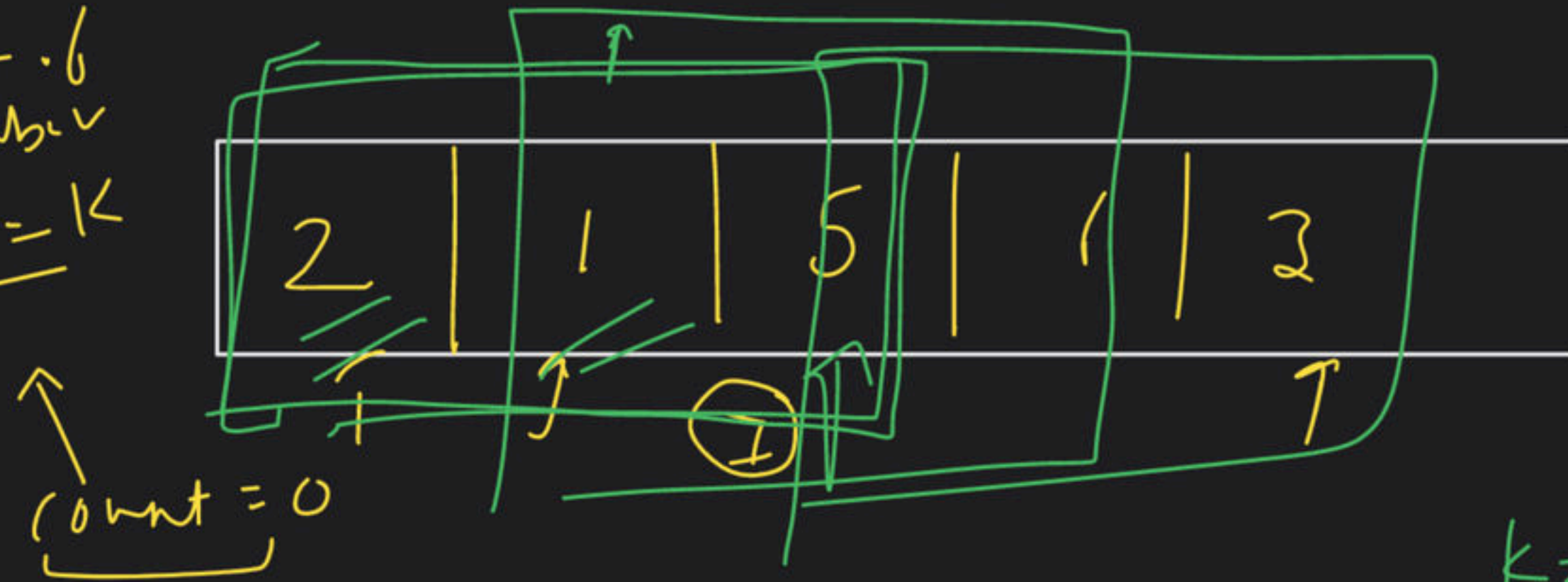
$K=3$

i/p { 2, 1, 5, 6, 3 }

→ { 2, 1, 3 }, 5, 6 }

swap = 1 ← ans

count of
number
 $\leq k$



count = 0

```
for (i → 0 < n)
{
    if (arr[i] ≤ k)
        count++;
}
```

count = 3

~~gand~~ bad = 1

$2 \leq k =$
 $1 \leq k$
 $5 \leq k$
 $\leq k$ #



Intuition → ?

approach

$\leq k$



size = row

2	1	5
---	---	---

1	5	1
---	---	---

5	1	1	3
---	---	---	---

bad = 1

1 <= 3

5 <= 3 → ✗

8 <= 3 ✗

bad = 2

5 <= 1 ✗

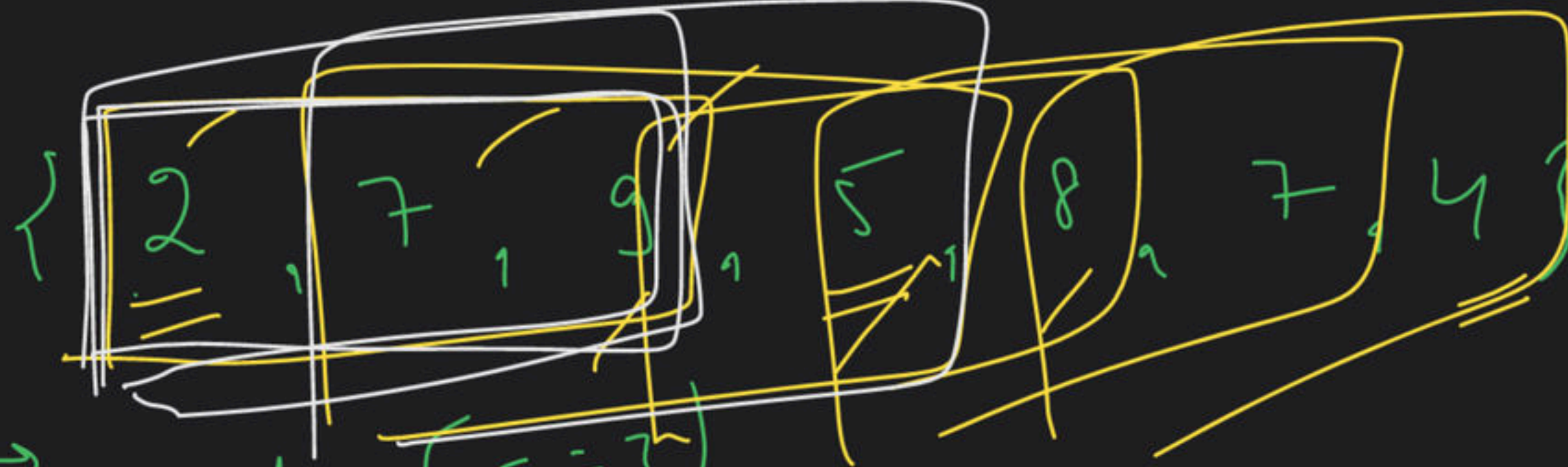
1 <= 1 ✓

3 <= 3

min

1 → ans

bad = 2



$k = 5$

① →

count (≤ 7)

count = 3

window size

~~5 | 8 | 7~~

→ bcd = 2

8 | 7 | 7

→ 2

2 | 7 | 9

→ bad → 2

$> k$ → pi

7 | 9 | 5

→ bcd → 2

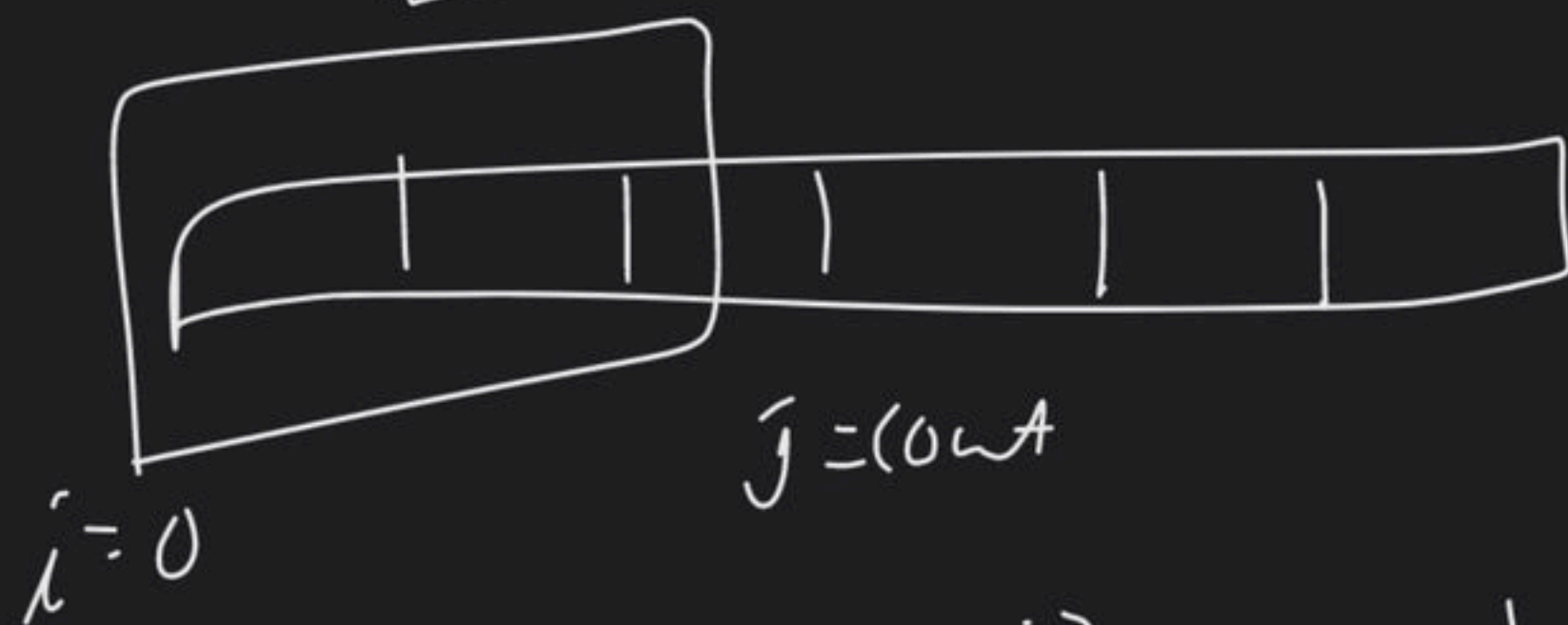
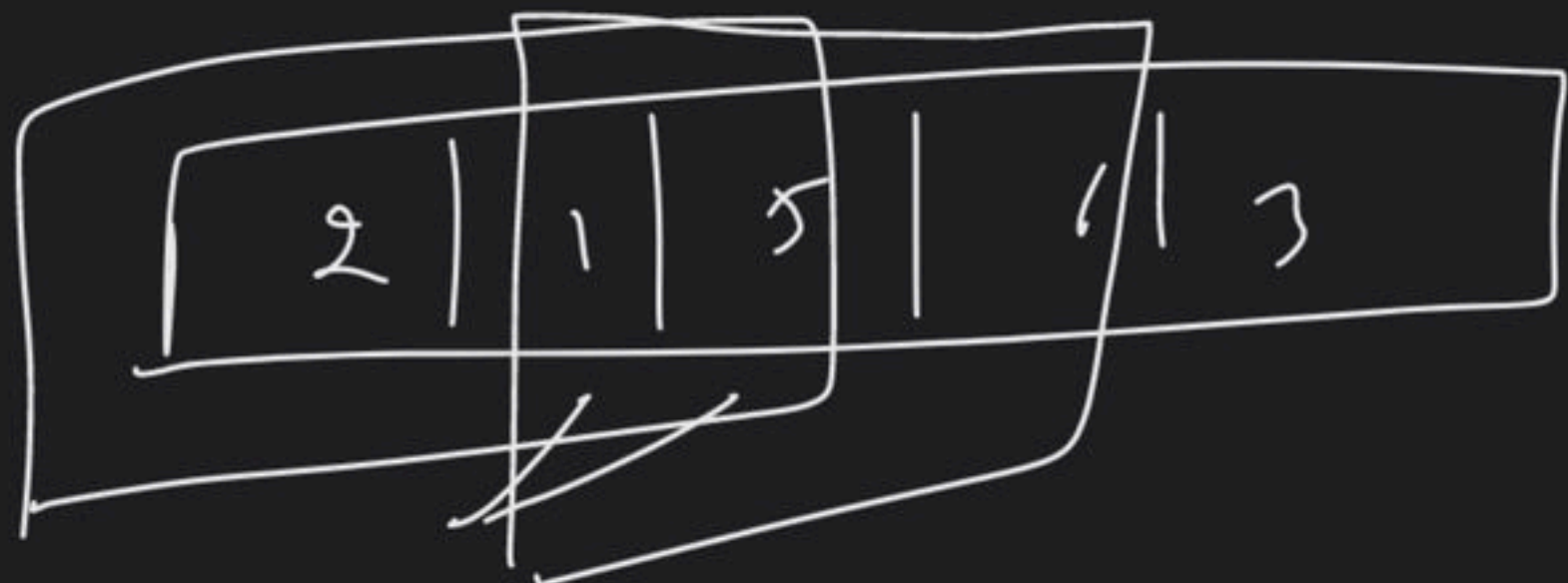
ans = 2

5 | 5 | 8

→ bad → 2

SW



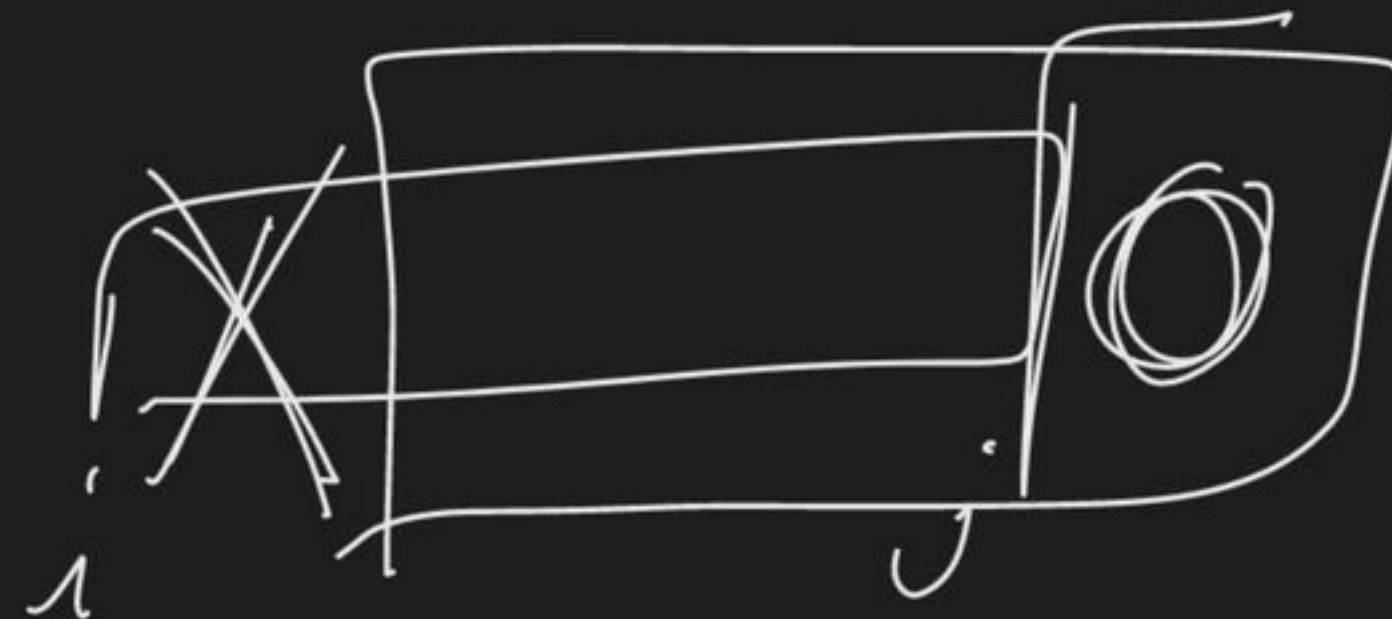


if (arr[i] > k)

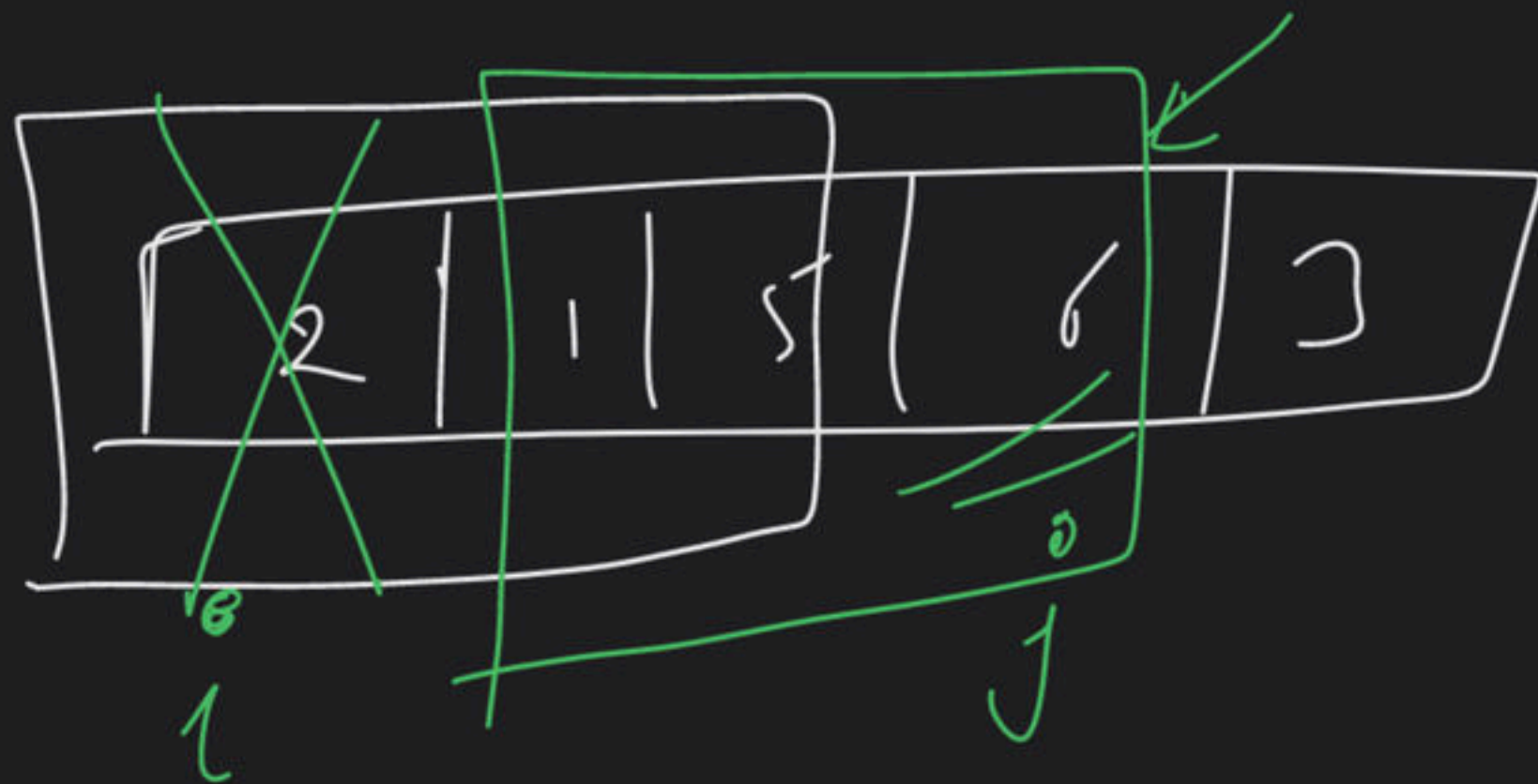
bad -

if (arr[j] > k)

bad++



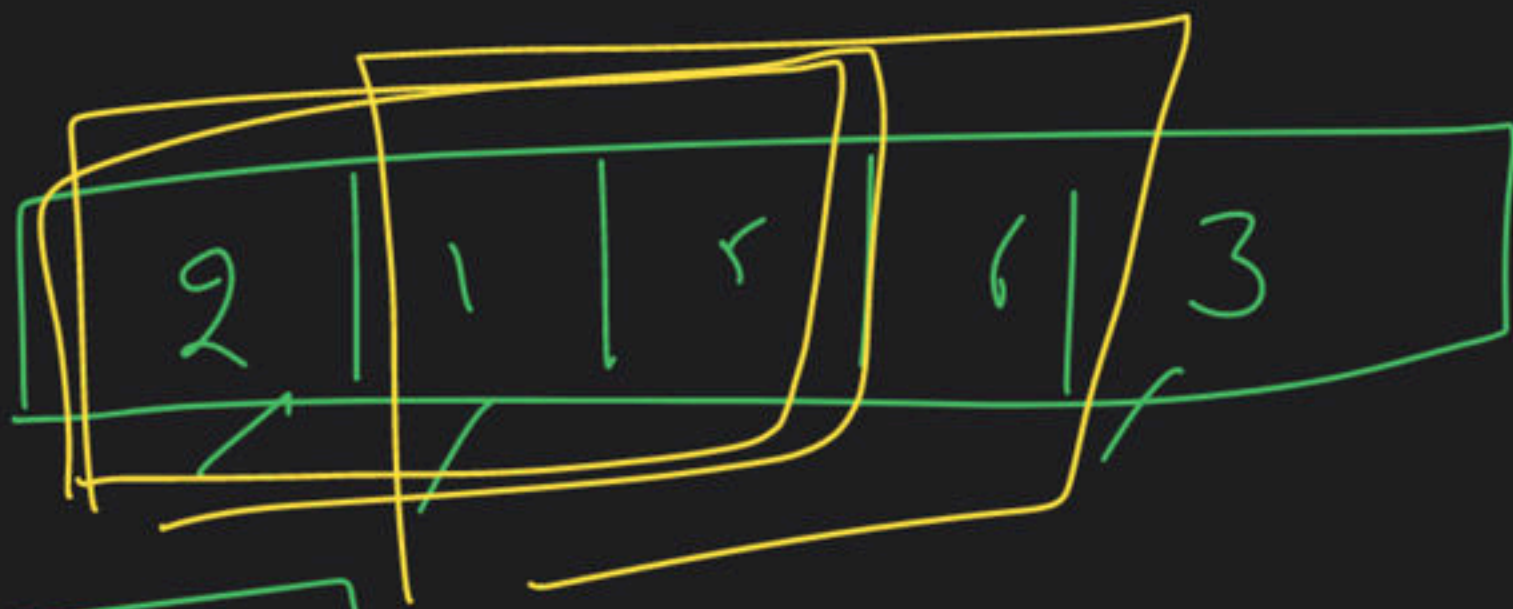
i++ j++



if ($arr[i] > k$)
 $bad--$

$ans = \min(ans, bad)$
 $\}$

if ($arr[j] > k$)
 $bad++$



$K=3$
 $\leq K$

(I)

length = 3
window size
 $\leq K$



bad \rightarrow 1

$> K$

$5 > 3$

time

ans = bad = 1

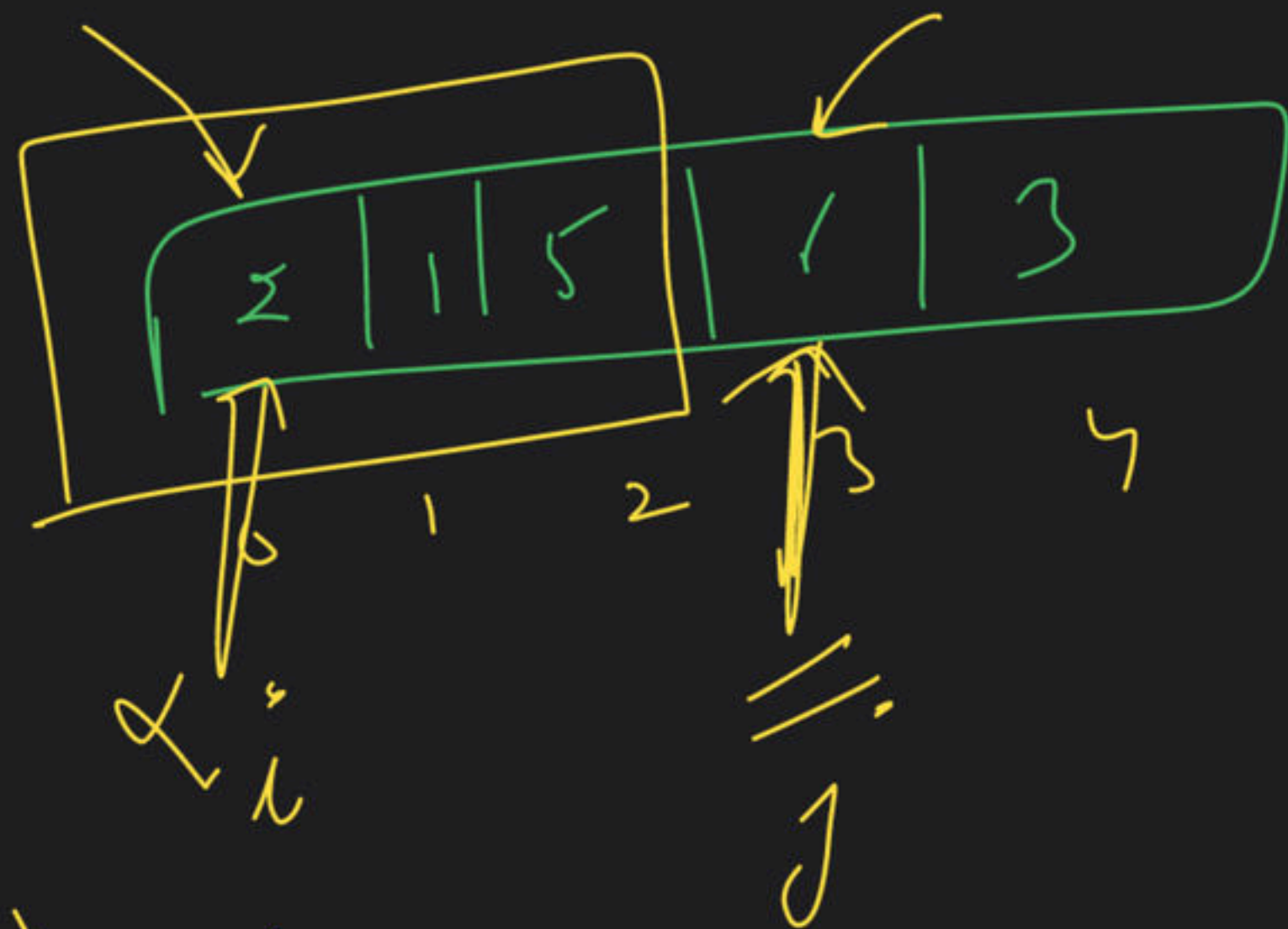
(II)



$6 > 3, 6 > 3$

bad = 2

(iii)



$$i = 0$$

$$j = \text{window} = 3$$

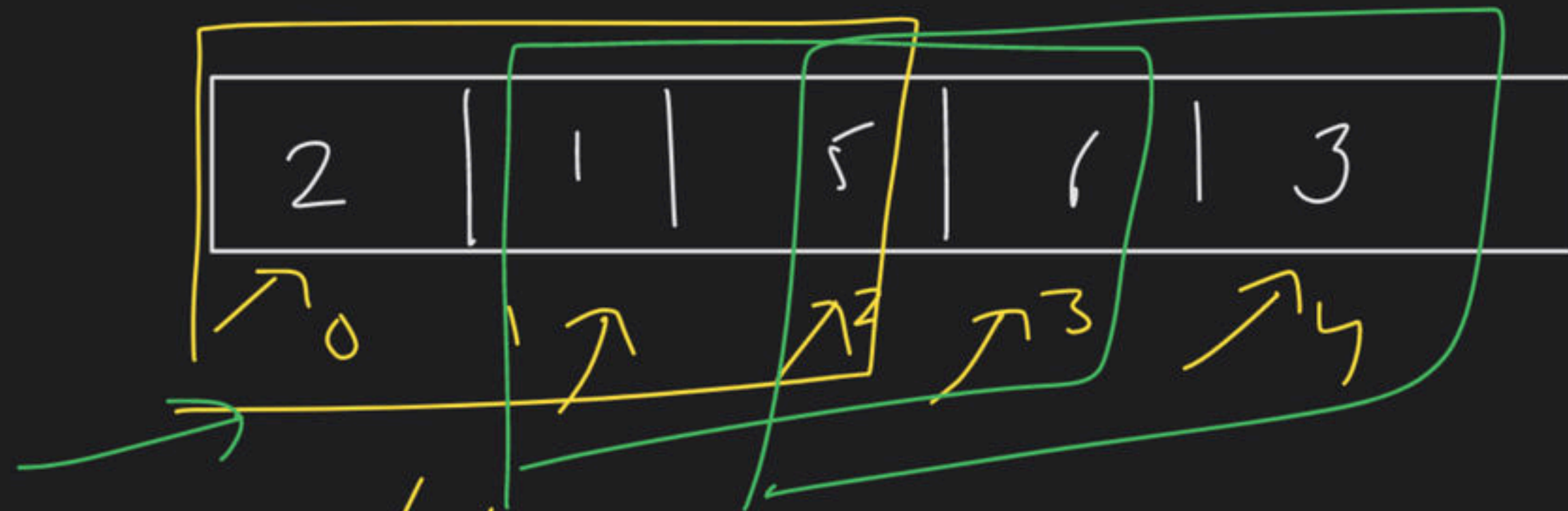
$$\frac{1}{n} \cdot n$$

$$O(n) > O(n)$$

interviewer

$$O(n) + O(n) + O(n)$$





$k = 3$

Window
size \rightarrow $hi - lo + 1 = k$

(1)

~~count = 0~~
2
3

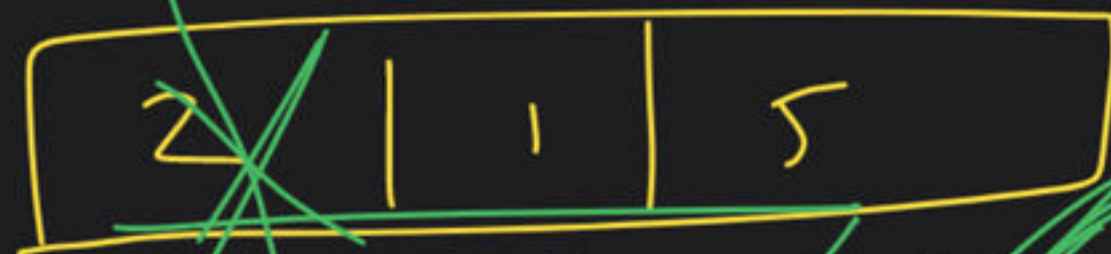
$2 \leq 3 \rightarrow T$
 $1 \leq 3 \rightarrow T$
 $5 \leq 3 \rightarrow F$

$1 \leq 3 \rightarrow F$
 $3 \leq 3 \rightarrow T$

if (arr[i] > k)
bad++

(2)

bad = 0



$k = 3$
 $2 > k \rightarrow F$
 $1 > k - 1 \rightarrow F$
 $5 > 3 \rightarrow T$

if (arr[j] > k)
bad++

~~bad = 0~~
bad = 1

$i^- \rightarrow$ demand to remove

$j^- \rightarrow$ ——— Ad

2 | 1 | 5 | 6 | 3

$\begin{matrix} 0 \\ \nearrow \\ 1 \end{matrix}$
 $\begin{matrix} 1 \\ \nearrow \\ 2 \end{matrix}$
 $\begin{matrix} -2 \\ \nearrow \\ 1 \end{matrix}$
 $\begin{matrix} 4 \\ \nearrow \\ 3 \end{matrix}$
 $\begin{matrix} 6 \\ \nearrow \\ 4 \end{matrix}$

$$2 \overset{2 > 3}{>} 1 \rightarrow F \rightarrow$$
$$6 > 3 \rightarrow T$$

bad \rightarrow 2

$$1 > 3 \rightarrow F$$

→ band = 2

1, 2, 2
 (1) → ans

→ HashMap

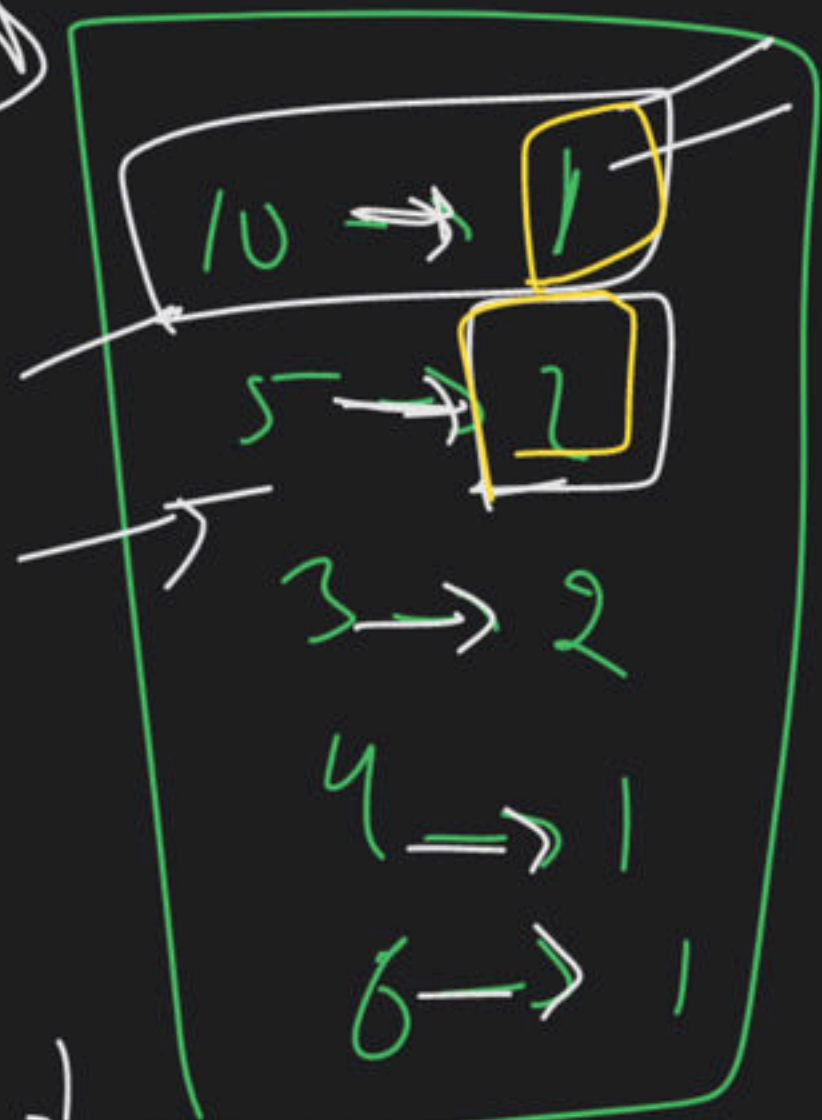
approach:-

(II)

ans = 5

$O(n)$

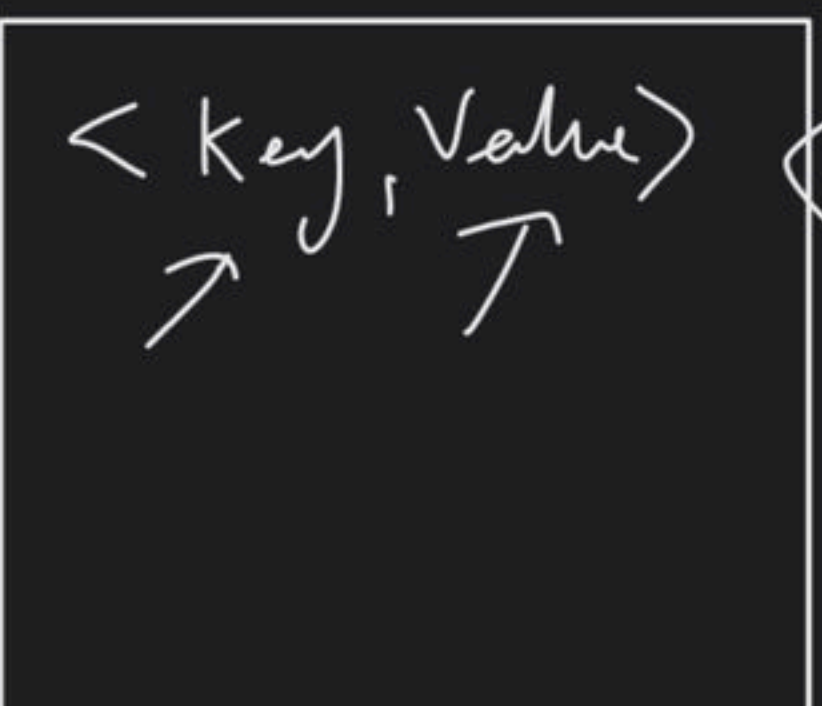
(I)



(5, 2)

HashMap

d.s



(10, 1)

10 →

5 → 2

2

2 → 5

{ 2, 2, 2, 2, 2 }

ans = 2

→ Ans:-

→ (I) → maping → < num → count >

unordered_map < int, int > map;

→ (II)

→ traverse array

→ check each element-

if (element_count > 1)

cout << ans



start row

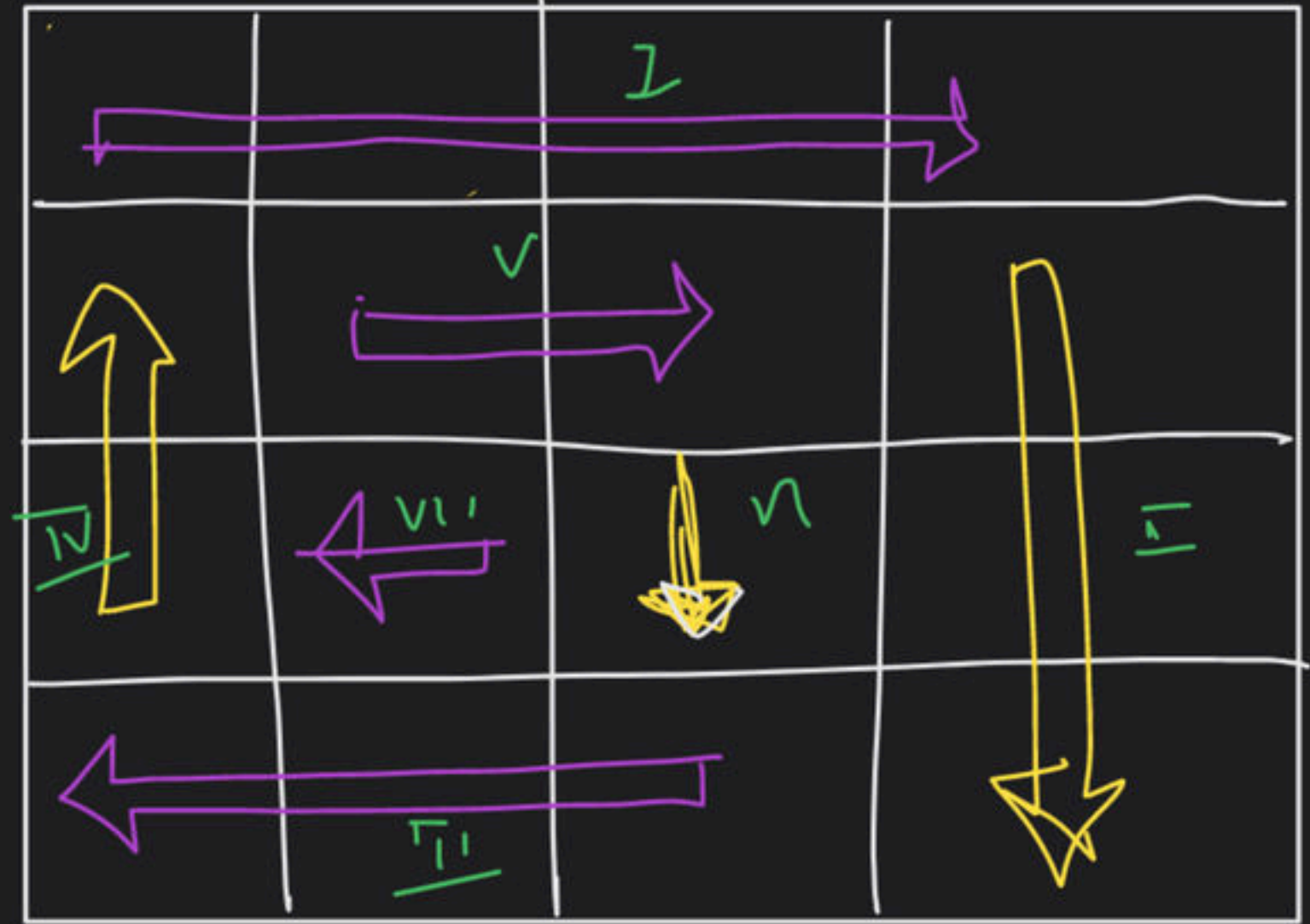


n

end row



start col end col



issue - ?

count = 0

while (count < total)

$$\text{Total} = n \times n = n^2$$

90

2 approach

10, 1, 2, 3, 4, 5

m[10]++
m[

heap <int, int> m;

m[10] → 0 / 1
<10, 1>

→ min platform problem

arrival → [a, b, c, d]

departs → [e, f, g, h]

minimum platform?

approach:

(I) sort arr / dep

(II)

arr → 9 am

11 am

dep → 1 pm

2 pm

ハ

map <int, int> m
 ↑ ↑
 key value

Key value
↓ ↓
m [10] = 3

access

m[10]

if (m[10] > 1)
 count++

Adobe / MS / Army



{ 9:00, 9:40, 9:50, 11, 15, 18 }

(BANK)

{ 9:10, 12:00, 11:20, 11:30, 15:00, 20:00 }

(I)

arr ->

{ 9:00, 9:40, 9:50, 11, 15, 18 }

dep ->

{ 9:10, 11:20, 11:30, 12:00, 15:00, 20:00 }

(II)

platform = 1 0 2 3 2 1 0 1 2 1 0

10

max -> ans = 3

can't
pair

int sum = 0;

int maxi = INT_MIN

1 condition

for (i = 0 -> n)

}

1

2

3

~~2~~
3

3

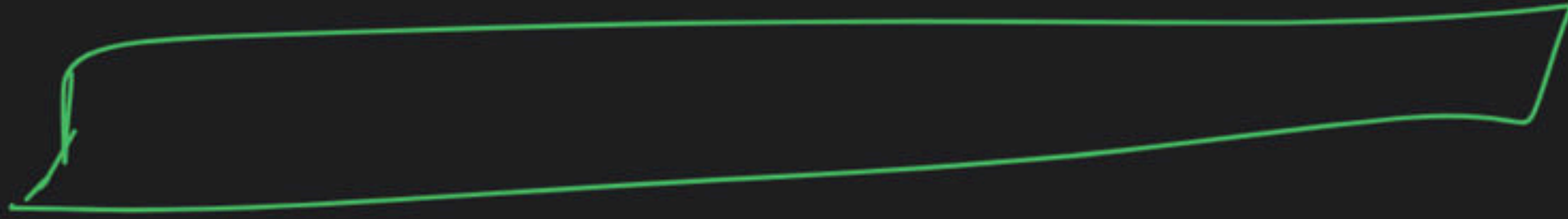
sum = sum + arr[i]

maxi = max(maxi, sum)

if (sum < 0)

sum = 0

Kade



for (i)

break for

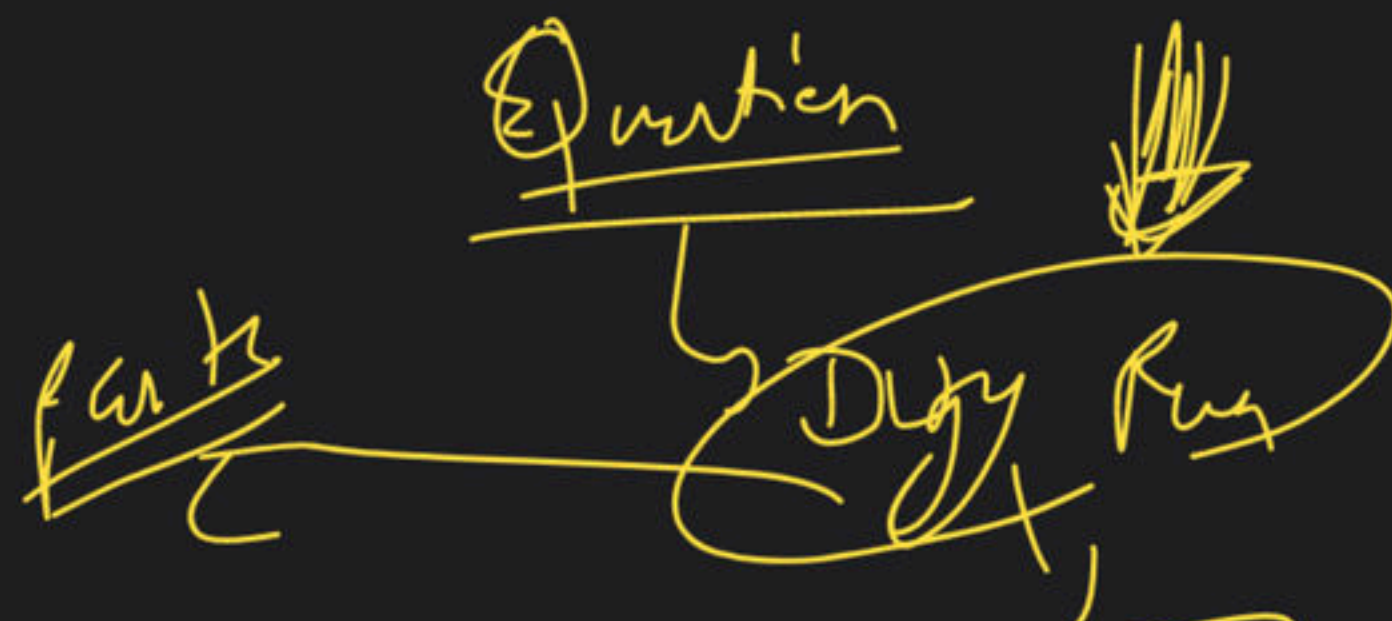
{
for (j = 0 ; j < n ; j++)

{

}

}

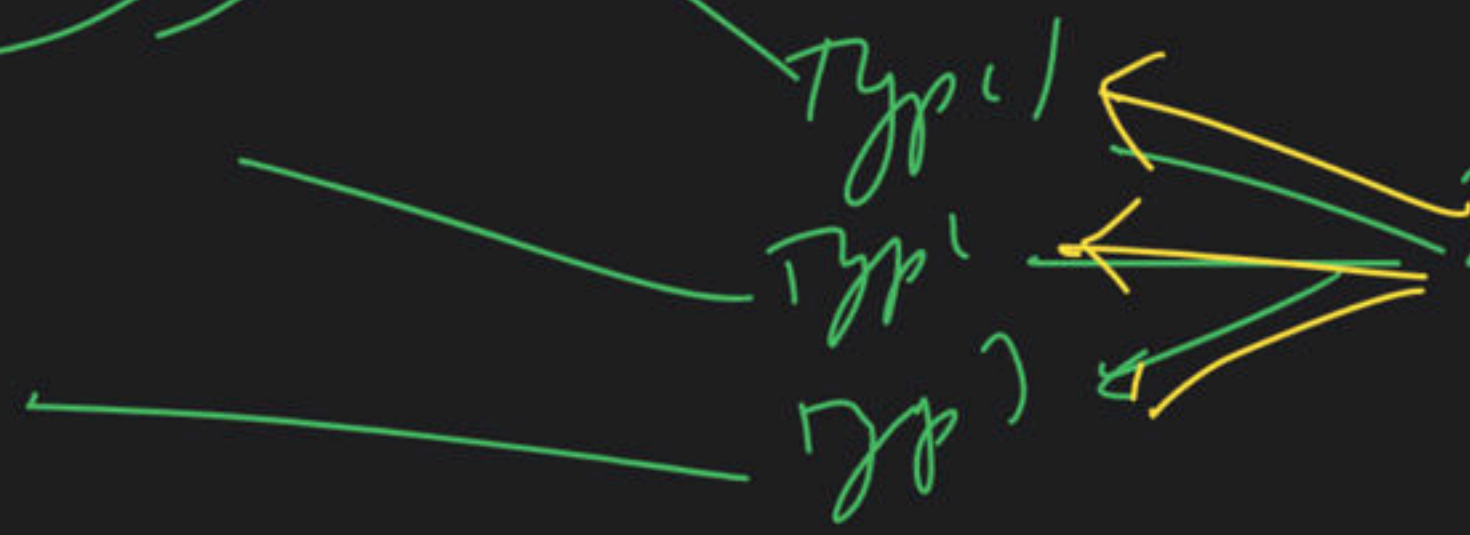
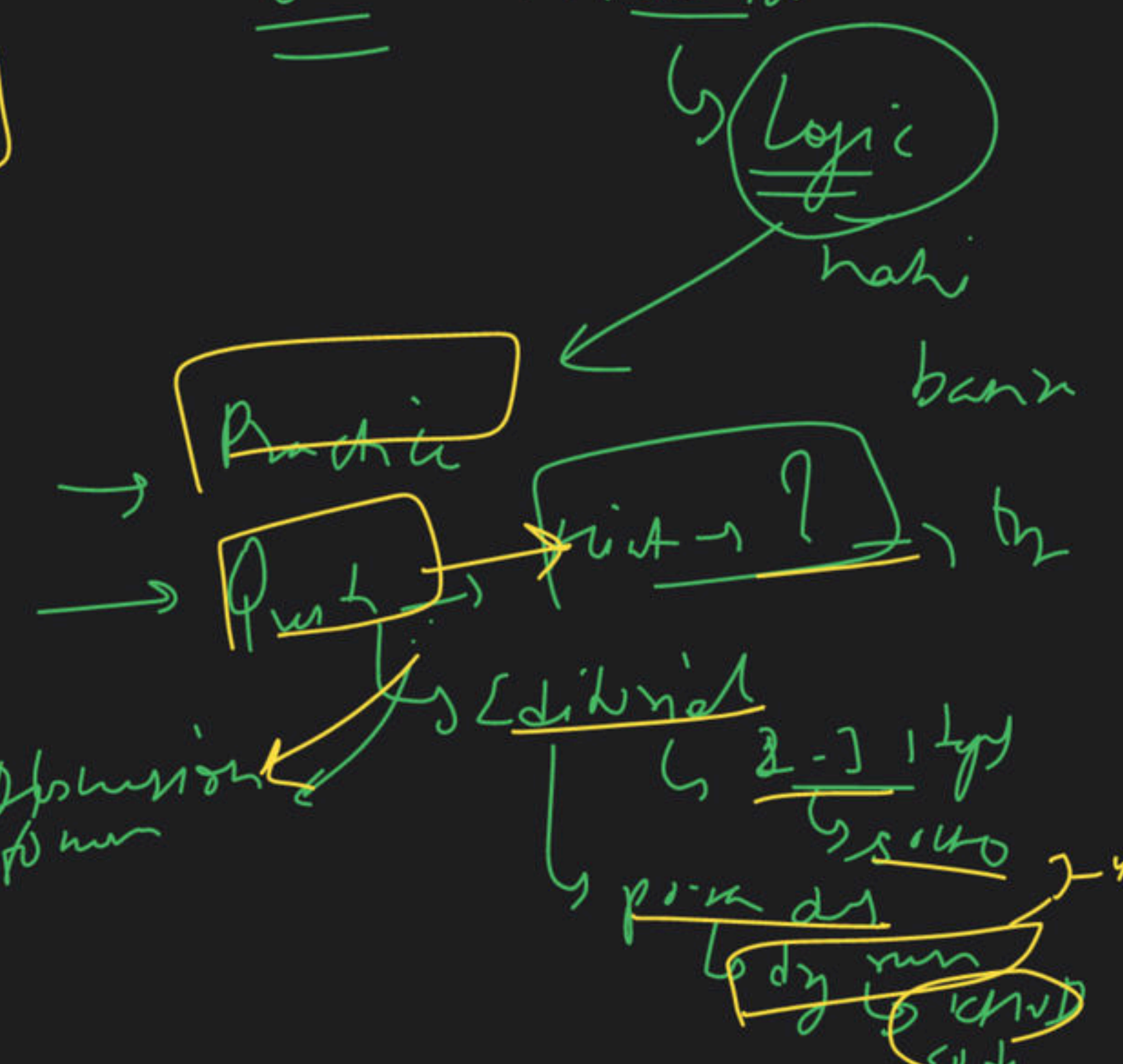
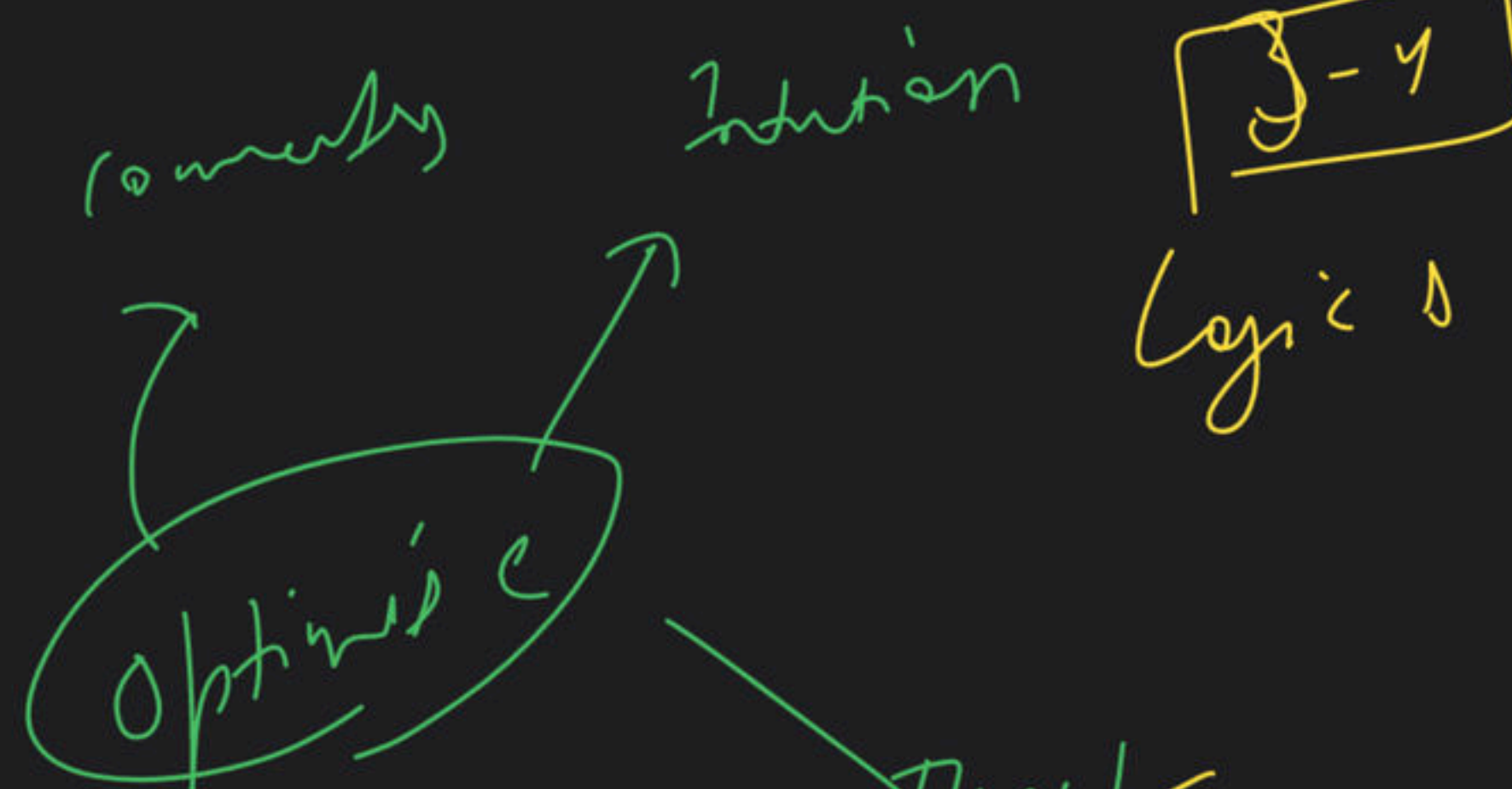
→ Downt:- →



250 m

250

50 → radio



→ Late :- →

→ minimise the height II →

$K = 2$

i/p →



$O(n^2)$

min diff to/w logit / shortest tower?

(I)

→ sort



$min = arr[0]$

$max = arr[n-1]$

$$K=2$$

→ { 1, 5, 8, 10 }

a | b | c | d

short

$$a+k$$

long
d+k

(1) Tower height → $\boxed{+k}$ $\boxed{-k}$

(1) min diff, Longest
short

$$a-k, d-k$$

$$\begin{aligned} &+k, +k \\ &+k, -k \\ &-k, +k \\ &-k, +k \end{aligned}$$

$$= \text{long} - \text{short}$$

$$= d+k - a-k$$

$$= \boxed{d-a}$$

$$= L - S$$

$$= d-k - (a-k)$$

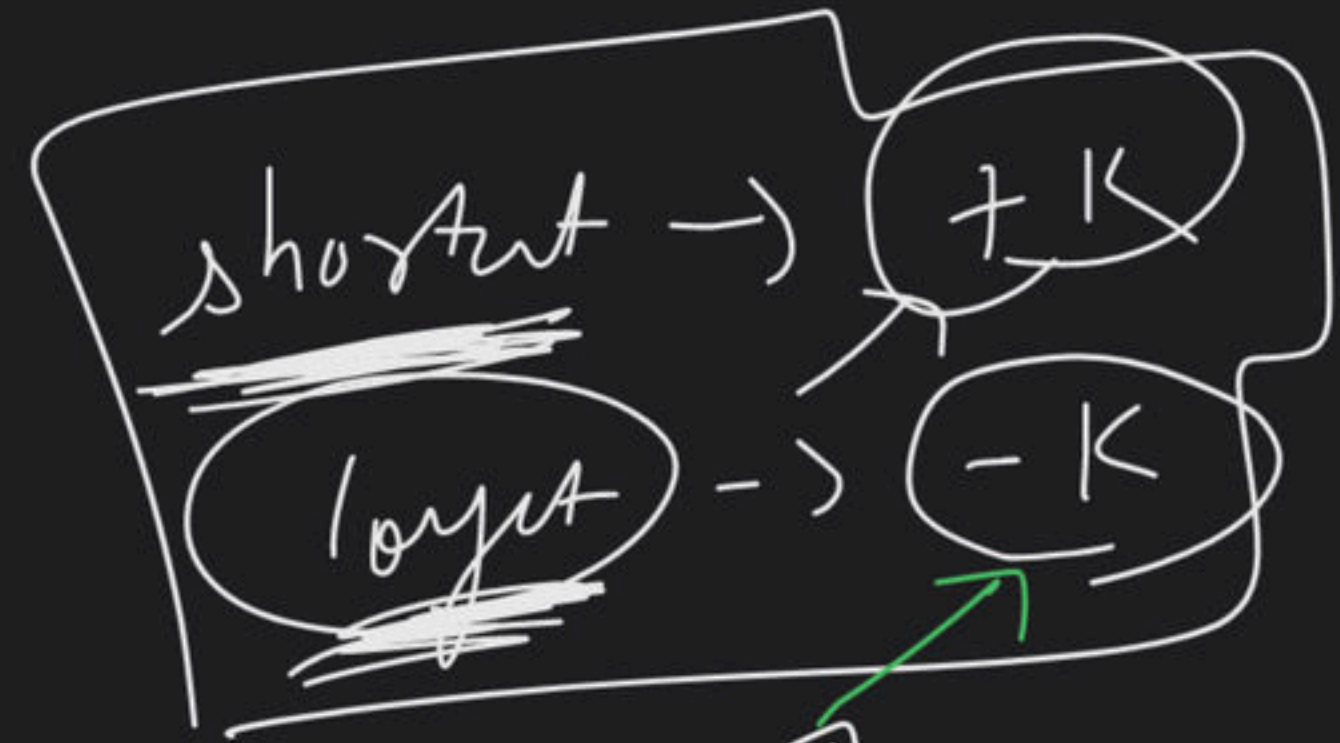
$$= d-k - a+k$$

$$= \boxed{d-a}$$

$$\begin{matrix} -K & & +K \end{matrix}$$

$$a - K \quad , \quad d + K$$

$$\begin{matrix} 1 & 1 & 1 & 1 & 1 \\ +K & +K & +K & +K & +K \\ -K & -K & -K & -K & -K \end{matrix}$$



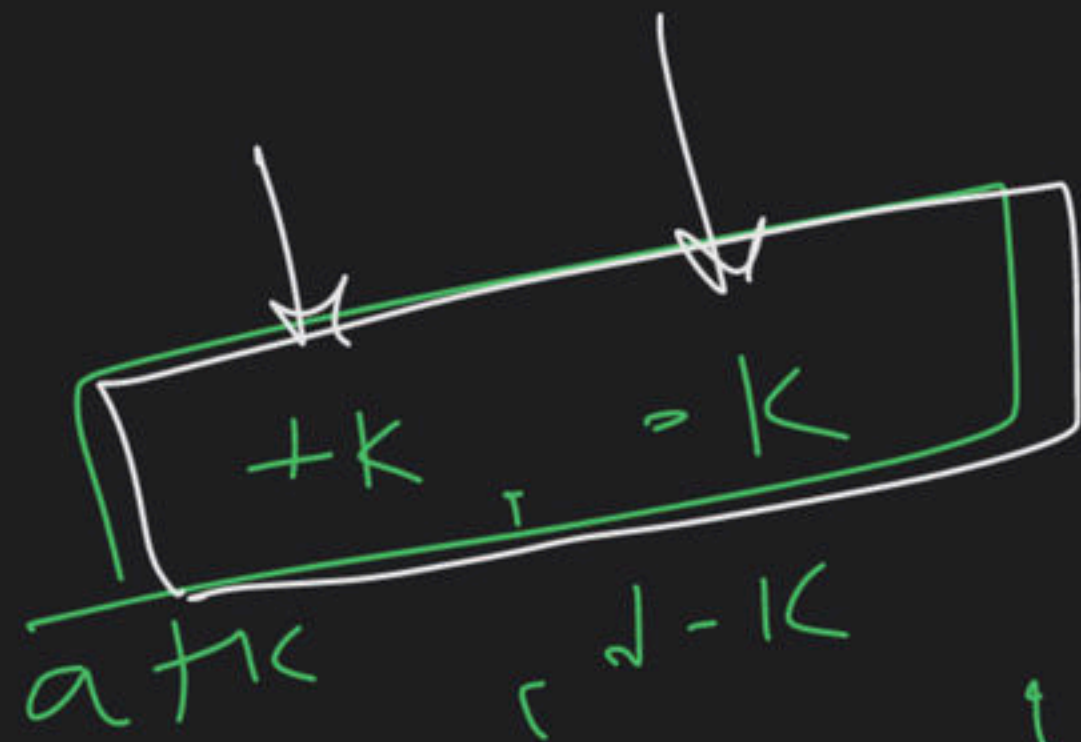
$$= L - S = d + K - (a - K)$$

$$= d + K - a + K$$

$$= d - a + 2K$$

Observation

(b) K



$$L - S \Rightarrow d - K - (a + K)$$

$$= d - K - a - K$$

$$= d - a - 2K$$

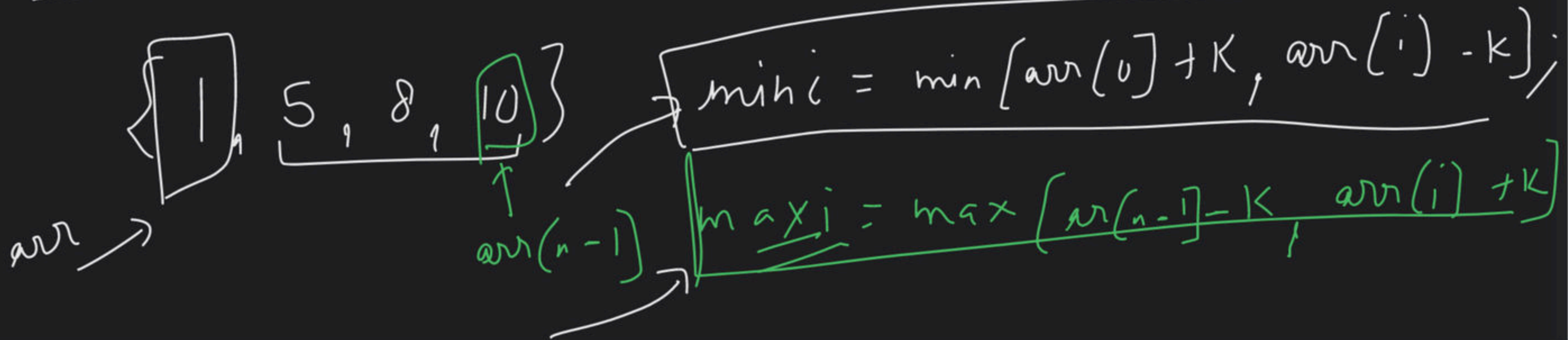
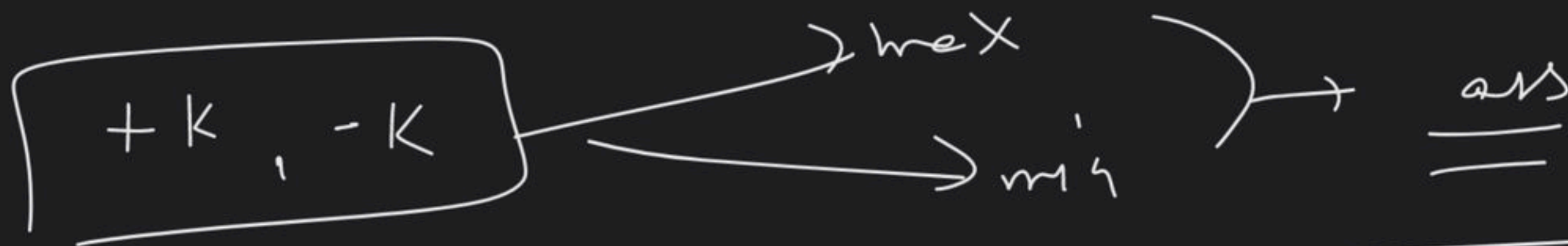
① arr

$k = 2$

shortcut $\{ 1, 5, 8, 10 \}$

$$1 + 2 = 3$$

$$10 - 2 = 8 \rightarrow 5$$



Code:-

(1)

sort (arr, arr + n); $n \log n$

30 min ↑

(2)

mini = arr[0];
maxi = arr[n-1];
ans = maxi - mini;

int n = 10;

int val = abs(n);

(3)

for (i = 1 → n) → or 0

mini = min [arr[0] + K, arr[i] - K]

maxi = max [arr[n-1] - K, arr[i-1] + K]

ans = min (maxi - mini, ans);

return ans; → $O(n \log n)$

5 examples

DRY RUN

→ report
→ dry run
→ copy & check







