

# Use of Transactions within a Reactive Microservices Environment

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#### **Assignment**

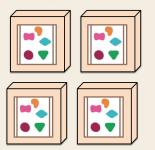
- Review the state of the art, in terms of problems of synchronous/blocking approaches for transaction management and other approaches/patterns available - taking into account the microservices context
- Propose a proof-of-concept implementation, using the Narayana transaction manager and prepare a service capable to manage transactions in the context of reactive microservices
- Prepare an example/quickstart showing the whole issues in more practical terms, proving that the transaction manager can work in an asynchronous environment

#### Microservices

A monolithic application puts all its functionality into a single process...



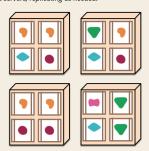
... and scales by replicating the monolith on multiple servers



A microservices architecture puts each element of functionality into a separate service...

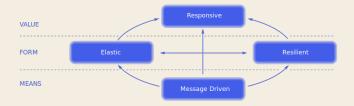


... and scales by distributing these services across servers, replicating as needed.



#### Reactive microservices

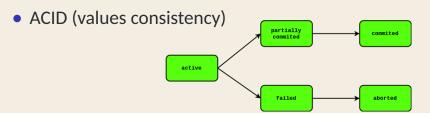
- reactive systems
- reactive programming
- reactive streams



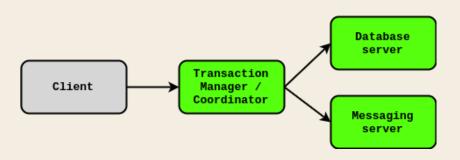
#### **Transactions**

"A transaction is a unit of processing that provides all-or-nothing property to the work that is conducted within its scope, also ensuring that shared resources are protected from multiple users" [1].

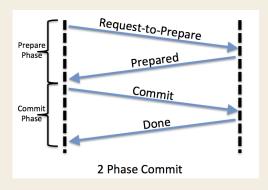
- sequence of operations
- commit or rollback



#### Distributed transactions



# Two phase commit protocol



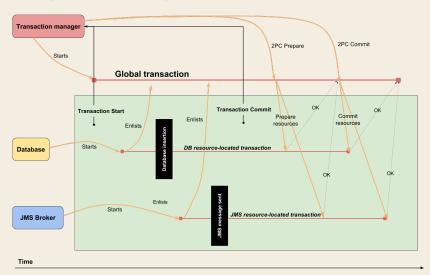
- O(n<sup>2</sup>) messages
- blocking
- coordinator single point of failure

#### Saga pattern

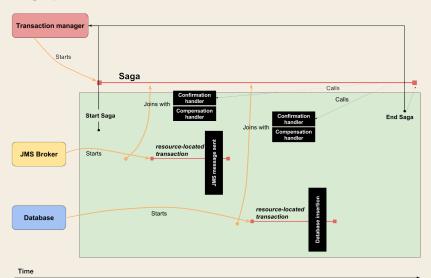
# Hector Garcia-Molina and Kenneth Salem, Princeton University, 1987

- long lived transactions
- compensations
- all-or-nothing property
  - 2PC T<sub>1</sub>, T<sub>2</sub>, ..., T<sub>n</sub> (in a single step)
     Saga
    - » success  $T_1, T_2, ..., T_n$ » failure -  $T_1, T_2, ..., T_k, C_k, C_{k-1}, ..., C_1$
- BASE (values availability)
  - Basically Available, Soft state, Eventual consistency

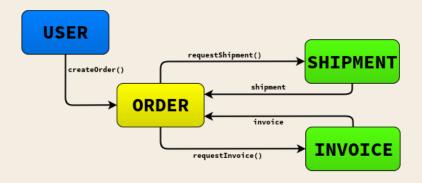
# Two phase commit protocol



### Saga pattern



# Saga implementations comparison scenario



# Saga implementations investigation

- Axon framework
- Eventuate Event Sourcing (ES)
- Eventuate Tram
- Narayana Long Running Actions (LRA)

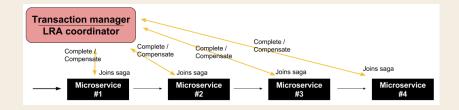
# Saga implementations comparison

| Problem                      | Axon | Eventuate<br>ES | Eventuate<br>Tram | LRA  |
|------------------------------|------|-----------------|-------------------|------|
| CQRS restriction             | Yes  | Yes             | Optional          | No   |
| Asynchronous by default      | Yes  | Yes             | No                | No   |
| Saga tracking and definition | No   | No              | Yes               | No   |
| Single point of failure      | No   | Yes             | Yes               | Yes* |
| Communication restrictions   | Yes  | Yes             | Yes               | No   |
| Distributed by default       | No   | Yes             | Yes               | Yes  |

# Saga implementations performance testing

- Axon 2 reported issues
- Eventuate ES 1 reported issue
- Eventuate Tram 1 feature request
- Narayana LRA

#### LRA executor motivation



#### LRA executor extension

- proof of concept / prototype
- asynchronous
- extensible and flexible design
- protocol / platform independent
  - further future extensions are expected
- two modules
  - LRA definitions
  - LRA executor

#### **LRA Definitions**

- LRADefinition
- Action
- fluent API

```
RESTLraBuilder.lra()
.name("testLRA")
.withAction(RESTAction
.post(new URL("http://example.com/request"))
.callbackUrl(new URL("http://example.com/callback"))
.build())
.data(42)
.callback("http://local.org")
.build();
```

```
"name": "testLRA",
  "actions": [
    "target": "http://example.com/request",
    "callbackUrl": "http://example.com/callback"
]],
  "data": 42,
  "parentLRA": null,
  "clientId": "",
  "timelimit": 0,
  "callbackUrl": "http://local.org",
  "nestedLRAs": []
}
```

#### LRA executor

- LRAExecutor
- synchronous and asynchronous executions
- AbstractLRAExecutor default implementation
  - actions are invoked in the declared order
- LRA manipulation methods
  - startLRA, completeLRA, compensateLRA
- integrated and tested (quickstart) with Narayana 5.8.1. Final

#### **Future work**

- integration in the Narayana codebase
- communication methods
- definition representations
- processing strategies

# Questions

# **Bibliography**

- [1] M. Little, J. Maron, and G. Pavlik. *Java transaction processing*. Prentice Hall, 2004.
- [2] Leslie Lamport. Lambert Preparation System. Addison-Wesley, 1986.
- [3] M. Goossens, F. Mittelbach, and A. Samarin. *The ET<sub>E</sub>X Companion*. Addison-Wesley, 1994.
- [4] Till Tantau. User's Guide to the Beamer Class Version 3.01. Available at http://latex-beamer.sourceforge.net.
- [5] https://www.martinfowler.com/articles/microservices.html
- [6] http://www.24pressrelease.com/assets/news/Propylene %20Glycol%20Solvent%2017614.jpg.
- [7] https://encrypt.co.in/2-phase-commit-protocol/

#### Opponent's review

- transaction heuristic outcomes
  - heuristic commit, rollback, mixed
  - non-atomic outcome
  - requires semantic knowledge
- LRA service performance test
  - REST requests queuing
- recovery capabilities of the executor
  - main concern failure after the marking of the participant invocation
  - idempotent requests (may be too restrictive)
  - timeouts

### Supervisor's review

- performance testing
- LRA specification relations
  - still in the draft form
  - focusing only on the coordination capabilities
  - currently only providing the REST reference implementation