

Cartesian Product X

<i>A</i>	<i>B</i>
1	2
3	4

(a) Relation *R*

<i>B</i>	<i>C</i>	<i>D</i>
2	5	6
4	7	8
9	10	11

(b) Relation *S*

<i>A</i>	<i>R.B</i>	<i>S.B</i>	<i>C</i>	<i>D</i>
1	2	2	5	6
1	2	4	7	8
1	2	9	10	11
3	4	2	5	6
3	4	4	7	8
3	4	9	10	11

(c) Result $R \times S$

Question 1: Cartesian Product

Given the following relations *R* and *S*, compute the Cartesian product $R \times S$.

Relation *R*

AB

1 2

3 4

Relation *S*

BCD

5 6 7

8 9 10

Result $R \times S$

Compute the Cartesian product and display the resulting relation.

Question 2: Cartesian Product with Different Attributes

Given the following relations *R* and *S*, compute the Cartesian product $R \times S$.

Relation *R*

XY

1 2

3 4

Relation S

YZW

2 5 6

4 7 8

Result $R \times S$

Compute the Cartesian product and display the resulting relation.

Question 3: Cartesian Product with Larger Relations

Given the following relations R and S, compute the Cartesian product $R \times S$.

Relation R

AB

1 2

3 4

5 6

Relation S

BCD

2 5 6

4 7 8

6 9 10

Result $R \times S$

Compute the Cartesian product and display the resulting relation.

Explanation of Concepts in the Questions

1. **Cartesian Product (\times):** Combines each tuple from the first relation with each tuple from the second relation, resulting in a new relation with all possible combinations.
 2. **Attributes:** The columns in relations. The resulting relation from a Cartesian product will have attributes from both input relations.
 3. **Tuples:** The rows in relations. The number of tuples in the resulting relation is the product of the number of tuples in the input relations.
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Question 1: Natural Join

Given the following relations RR and SS , compute the natural join $R \bowtie S$.

Relation RR

A B

1 2

3 4

Relation SS

B C D

2 5 6

4 7 8

9 10 11

Result $R \bowtie S$

Compute the natural join and display the resulting relation. Explain why the third tuple of SS does not appear in the result.

Question 2: Natural Join with Different Attributes

Given the following relations RR and SS , compute the natural join $R \bowtie S$.

Relation R

X Y

1 2

3 4

Relation S

Y Z W

2 5 6

4 7 8

6 9 10

Result $R \bowtie S$

Compute the natural join and display the resulting relation. Explain why the third tuple of SS does not appear in the result.

Question 3: Natural Join with Larger Relations

Given the following relations R and S , compute the natural join $R \bowtie S$.

Relation R

A B

1 2

3 4

5 6

Relation S

B C D

2 5 6

4 7 8

6 9 10

Result $R \bowtie S$

Compute the natural join and display the resulting relation. Explain why the third tuple of SS does not appear in the result.

Explanation of Concepts in the Questions

- Natural Join (\bowtie):** Combines tuples from two relations based on common attributes, and the resulting relation includes only the tuples that have matching values in these common attributes.
 - Resultant Schema:** The schema of the resulting relation is the union of the schemas of the input relations, with common attributes appearing only once.
 - Tuple Matching:** Only tuples with matching values in the common attributes are included in the result. Tuples without matches are excluded.
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Joins \bowtie

<i>A</i>	<i>B</i>	<i>C</i>
1	2	3
6	7	8
9	7	8

(a) Relation *U*

<i>B</i>	<i>C</i>	<i>D</i>
2	3	4
2	3	5
7	8	10

(b) Relation *V*

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>
1	2	3	4
1	2	3	5
6	7	8	10
9	7	8	10

(c) Result $U \bowtie V$

First tuple of *U* has two matching tuples in *V*.

Third tuple of *B* has two matching tuples in *U*.

Examples

Is the following statement correct:

a) $R \cap S = R - (R - S)$

1. Is the following statement correct?

a) $(R \cup S) - S = R - (S - R)$

2. Determine whether the following equation holds:

a) $(R - S) \cap (S - R) = \emptyset$

3. Verify if the following identity is true:

a) $(R \cap S) \cup (R - S) = R$

Let's prove each of the statements with set theory principles and provide examples.

Statement 1: $(R \cup S) - S = R - (S - R)$

Proof:

1. The left-hand side:

- $(R \cup S) - S$ represents all elements in $R \cup S$ but not in S . Since S is subtracted, only elements from R remain.
- So, $(R \cup S) - S = R - S$.

2. The right-hand side:

- $S - R$ consists of elements that are in S but not in R .
- $R - (S - R)$ means taking all elements of R and removing any that are in $S - R$.
- But $S - R$ has no elements from R , so subtracting it does not change R , meaning $R - (S - R) = R - S$.

Since both sides simplify to $R - S$, the statement is **true**.

Example:

Let $R = \{1, 2, 3\}$ and $S = \{3, 4, 5\}$:

- $R \cup S = \{1, 2, 3, 4, 5\}$
- $(R \cup S) - S = \{1, 2\}$
- $S - R = \{4, 5\}$
- $R - (S - R) = \{1, 2\}$

Since both sides are equal, the statement is **true**.

Statement 2: $(R - S) \cap (S - R) = \emptyset$

Proof:

1. $R - S$ consists of elements in R that are not in S .
2. $S - R$ consists of elements in S that are not in R .
3. The intersection of these two sets would contain elements that belong to both $R - S$ and $S - R$.
 - But by definition, $R - S$ contains elements **not in** S , and $S - R$ contains elements **not in** R .
 - There cannot be an element common to both sets.
 - Hence, their intersection must be **empty**.

Thus, the statement is **true**.

Example:

Let $R = \{1, 2, 3\}$ and $S = \{3, 4, 5\}$:

- $R - S = \{1, 2\}$
- $S - R = \{4, 5\}$
- $(R - S) \cap (S - R) = \emptyset$

Since there are no common elements, the statement is **true**.

Statement 3: $(R \cap S) \cup (R - S) = R$

Proof:

1. $R \cap S$ consists of elements common to both R and S .
2. $R - S$ consists of elements in R but not in S .
3. The union $(R \cap S) \cup (R - S)$ contains all elements that are either:
 - In both R and S (from $R \cap S$), or
 - In R but not in S (from $R - S$).
 - Since these two sets together cover all elements in R , the union must be **equal to** R .

Thus, the statement is **true**.

Example:

Let $R = \{1, 2, 3, 4\}$ and $S = \{3, 4, 5\}$:

- $R \cap S = \{3, 4\}$
- $R - S = \{1, 2\}$
- $(R \cap S) \cup (R - S) = \{1, 2, 3, 4\} = R$

Since both sides are equal, the statement is **true**.

Exercise 1

$R(A, B)$ and $S(B, C)$. Which two are equivalent. Give an example of how the different one can be different.

a. $\pi_{A,C}(R \bowtie \sigma_{B=1} S)$

b. $\pi_A(\sigma_{B=1} R) \times \pi_C(\sigma_{B=1} S)$

c. $\pi_{A,C}(\pi_A R \times \sigma_{B=1} S)$

Similar Questions:

1. Given relations $R(A, B)$ and $S(B, C)$, determine which of the following are equivalent: a.

$\pi_{A,C}(\sigma_{B=2}(R \bowtie S))$

b. $\pi_A(\sigma_{B=2}(R)) \times \pi_C(\sigma_{B=2}(S))$

c. $\pi_{A,C}(\pi_A R \bowtie \sigma_{B=2}(S))$

Provide an example where one differs from the others.

2. Consider the relations $R(A, B)$ and $S(B, C)$. Which of the following expressions are equivalent? a.

$\pi_{A,B}(R \bowtie \sigma_{B=10}(S))$

b. $\pi_A(\sigma_{B=10}(R)) \bowtie \pi_B(\sigma_{B=10}(S))$

c. $\pi_{A,B}(\pi_B(R) \bowtie \sigma_{B=10}(S))$

Justify your answer with an example.

3. For relations $R(A, B)$ and $S(B, C)$, determine if the following equivalences hold: a. $\pi_{A,C}(\sigma_{B=5}(R \bowtie S))$

b. $\pi_{A,C}(\sigma_{B=5}(R)) \times \pi_C(\sigma_{B=5}(S))$

c. $\pi_{A,C}(R \bowtie \pi_{B,C}(\sigma_{B=5}(S)))$

Prove or provide a counterexample.

Questions with Proof and Examples:

1. Prove or disprove the equivalence of the following expressions:

- a. $\pi_{A,C}(\sigma_{B=4}(R \bowtie S))$
- b. $\pi_A(\sigma_{B=4}(R)) \bowtie \pi_C(\sigma_{B=4}(S))$
- c. $\pi_{A,C}(R \bowtie \pi_{B,C}(\sigma_{B=4}(S)))$

Proof:

- The projection operation $\pi_{A,C}$ removes B, so (a) and (c) are equivalent. However, (b) decomposes the selection before joining, which may cause loss of information if B is required.
- Example:** If $R = \{(1,4), (2,4), (3,5)\}$ and $S = \{(4,10), (5,20)\}$, then (a) and (c) both return $\{(1,10), (2,10)\}$ while (b) may lose the relation context.

2. Show whether the following expressions are equivalent:

- a. $\pi_{A,C}(\sigma_{B=3}(R \bowtie S))$
- b. $\pi_A(\sigma_{B=3}(R)) \bowtie \pi_C(\sigma_{B=3}(S))$
- c. $\pi_{A,C}(R \bowtie \pi_{B,C}(\sigma_{B=3}(S)))$

Proof:

- If B is a join key, removing it before the join (as in (b)) may lead to incorrect results.
- Example:** If $R = \{(1,3), (2,3), (3,4)\}$ and $S = \{(3,7), (4,8)\}$, (a) and (c) return $\{(1,7), (2,7)\}$ but (b) fails.

3. Consider relations R(A, B) and S(B, C). Are the following transformations valid?

- a. $\pi_{A,C}(\sigma_{B=6}(R \bowtie S))$
- b. $\pi_A(\sigma_{B=6}(R)) \times \pi_C(\sigma_{B=6}(S))$
- c. $\pi_{A,C}(R \bowtie \pi_{B,C}(\sigma_{B=6}(S)))$

Proof:

- (b) represents a Cartesian product instead of a join, potentially introducing spurious tuples.
 - Example:** If $R = \{(1,6), (2,6), (3,7)\}$ and $S = \{(6,9), (7,11)\}$, (a) and (c) return $\{(1,9), (2,9)\}$, while (b) might incorrectly generate extra tuples.
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Exercise 2

Consider a relation $R(A, B)$ that contains r tuples, and a relation $S(B, C)$ that contains s tuples; and $r > 0$ and $s > 0$.

In terms of R and S the minimum and maximum number of tuples that could be in the result of each expression:

- a. $\pi_{A,C}(R \bowtie S)$
- b. $\pi_B R - (\pi_B R - \pi_B S)$
- c. $(R \bowtie R) \bowtie R$
- d. $\sigma_{A>B} R \cup \sigma_{A<B} R$

Similar Questions:

1. Consider a relation $R(A, B)$ with r tuples and a relation $S(B, C)$ with s tuples. If both relations are stored using a hash-based indexing scheme, determine the minimum and maximum number of tuples that could be produced by the following queries:
 - a. $\pi_{A,C}(R \bowtie S)$
 - b. $\pi_B R - (\pi_B R \cap \pi_B S)$
 - c. $(R \bowtie S) \bowtie T$, where $T(C, D)$ has t tuples.
 - d. $\sigma_{A>5} R \cup \sigma_{B<10} S$
 2. Given two relations $R(A, B)$ and $S(B, C)$ with r and s tuples respectively, compute the minimum and maximum possible number of tuples in the output for:
 - a. $\pi_{A,B}(R \bowtie S)$
 - b. $\pi_C (\sigma_{B>10}(R \bowtie S))$
 - c. $(R \bowtie S) - (\sigma_{B=5} R)$
 - d. $\sigma_{A=1} R \cup \sigma_{C=2} S$
 3. Let $R(A, B)$ and $S(B, C)$ be two relations with r and s tuples respectively. Assuming that attribute B is a foreign key in R referencing S , determine the result sizes for:
 - a. $R \bowtie S$
 - b. $\pi_{A,C} (\sigma_{B=1}(R \bowtie S))$
 - c. $(R \times S) - (R \bowtie S)$
 - d. $(\sigma_{A>100} R) \cap (\sigma_{B<50} S)$
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Questions with Proof and Examples:

1. **Prove or disprove:** $\pi_{A,C} (R \bowtie S) = \pi_{A,C} (\pi_A R \bowtie S)$

- **Proof:** The projection before the join may remove attributes necessary for the join condition, leading to different results. Example:
 - $R(A, B): \{(1,2), (3,4)\}$
 - $S(B, C): \{(2,5), (4,6)\}$
 - Applying π_A before join removes B, making the join impossible.

2. **Prove:** $\pi_A (\sigma_{B=10} R) \times \pi_C (\sigma_{B=10} S) = \pi_{A,C} (\sigma_{B=10}(R \bowtie S))$

- **Proof:** Since selection on B ensures that only tuples with B=10 remain, performing a join and then projecting A and C is equivalent to performing projections on each relation separately and then taking the Cartesian product.

3. **Prove or give a counterexample:** $(R \bowtie S) \bowtie T = R \bowtie (S \bowtie T)$

- **Proof:** This holds due to the associative property of joins.
- **Example:**
 - $R(A, B): \{(1,2), (3,4)\}$
 - $S(B, C): \{(2,5), (4,6)\}$
 - $T(C, D): \{(5,7), (6,8)\}$
 - $R \bowtie S$ gives $\{(1,2,5), (3,4,6)\}$, then joining with T results in $\{(1,2,5,7), (3,4,6,8)\}$.
 - $S \bowtie T$ first gives $\{(2,5,7), (4,6,8)\}$, then joining with R produces the same result.

Exercise 3

Calculate the number of Tuples for 1 and 2.

$$T(R)=100, V(R,A) =5$$

1. $\sigma_{A=10}(R)$

2. $\sigma_{A \leq 10}(R)$

Similar Questions:

Question 1:

Calculate the number of tuples for the given conditions:

$$T(R) = 200, V(R, B) = 10$$

1. $\sigma_{B=5}(R)$
 2. $\sigma_{B>5}(R)$
-

Question 2:

Given $T(R) = 500$ and $V(R, C) = 20$, determine the number of tuples for:

1. $\sigma_{C=15}(R)$
 2. $\sigma_{C \leq 15}(R)$
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Question 3:

For the relation $R(A, D)$ with $T(R) = 1000$ and $V(R, A) = 25$, calculate the tuples in the result of:

1. $\sigma_{A=7}(R)$
2. $\sigma_{A \geq 7}(R)$

Three More with Proof and Examples:**Question 4:**

Prove that for a relation $R(A, B)$ with total tuples $T(R)$ and distinct values $V(R, A)$:

$$\sigma_{A=x}(R) = \frac{T(R)}{V(R, A)}$$

Example:

If $T(R) = 300$ and $V(R, A) = 10$, then:

$$\sigma_{A=5}(R) = \frac{300}{10} = 30$$

Question 5:

Prove that for a range selection where $A \leq x$, the expected number of tuples is:

$$\sigma_{A \leq x}(R) = \frac{x}{V(R, A)} \times T(R)$$

Example:

If $T(R) = 500$, $V(R, A) = 50$, and we want $A \leq 10$:

$$\sigma_{A \leq 10}(R) = \frac{10}{50} \times 500 = 100$$

Question 6:

Show that if attribute A follows a uniform distribution in relation R(A, B), then:

$$\sigma_{A>x}(R) = \left(1 - \frac{x}{V(R, A)}\right) \times T(R)$$

Example:

Given $T(R) = 400$, $V(R, A) = 40$, find tuples for $A > 30$:

$$\sigma_{A>30}(R) = \left(1 - \frac{30}{40}\right) \times 400 = (1 - 0.75) \times 400 = 100$$

Exercise 4

Calculate the number of Tuples for 1.

$T(R)=100$. Values of A are uniformly distributed over [1,20]

1. $\sigma_{A=10}(R)$
2. $\sigma_{A \leq 10}(R)$
3. $\sigma_{A > 10}(R)$

Similar Questions:**Exercise 5**

Calculate the number of Tuples for 1 and 2.

$T(R) = 200$, $V(R, A) = 10$

1. $\sigma_{A=5}(R)$
2. $\sigma_{A \geq 8}(R)$

Exercise 6

Calculate the number of Tuples for 1, 2, and 3.

$T(R) = 500$. Values of A are uniformly distributed over [1, 50].

1. $\sigma_{A=25}(R)$
2. $\sigma_{A \leq 25}(R)$
3. $\sigma_{A > 25}(R)$

Exercise 7

Consider relation R(A, B, C), where $T(R) = 150$, and attribute A is uniformly distributed over [1, 30].

1. $\sigma_{A=15}(R)$
2. $\sigma_{A \leq 15}(R)$
3. $\sigma_{A > 15}(R)$

Examples with Solutions

Example 1

$T(R) = 100$, A uniformly distributed over $[1, 20]$.

1. $\sigma_{A=10}(R)$

- Probability of any single value of A: $1/20$
- Expected tuples: $(1/20) * 100 = 5$

2. $\sigma_{A \leq 10}(R)$

- Probability: $10/20 = 1/2$
- Expected tuples: $(1/2) * 100 = 50$

3. $\sigma_{A > 10}(R)$

- Probability: $10/20 = 1/2$
- Expected tuples: $(1/2) * 100 = 50$

Example 2

$T(R) = 200$, A uniformly distributed over $[1, 50]$.

1. $\sigma_{A=25}(R)$

- Probability: $1/50$
- Expected tuples: $(1/50) * 200 = 4$

2. $\sigma_{A \leq 25}(R)$

- Probability: $25/50 = 1/2$
- Expected tuples: $(1/2) * 200 = 100$

3. $\sigma_{A > 25}(R)$

- Probability: $25/50 = 1/2$
 - Expected tuples: $(1/2) * 200 = 100$
-

Example 3

$T(R) = 500$, A uniformly distributed over $[1, 10]$.

1. $\sigma_{A=5}(R)$

- Probability: $1/10$
- Expected tuples: $(1/10) * 500 = 50$

2. $\sigma_{A \leq 5}(R)$

- Probability: $5/10 = 1/2$
- Expected tuples: $(1/2) * 500 = 250$

3. $\sigma_{A > 5}(R)$

- Probability: $5/10 = 1/2$
 - Expected tuples: $(1/2) * 500 = 250$
-

Exercise 5

$R(A, B, C)$ is a relation. $T(R) = 10000$. $V(R, A) = 50$.
 $W = \sigma_{A=10 \text{ AND } B < 10}(R)$. What will be $T(W)$?

Similar Questions:

Question 3:

Relation $T(M, N, P, Q)$ has:

- $T(T) = 15000$
- $V(T, M) = 75$
- $V(T, P) = 50$

Define Z as:

$$Z = \sigma_{M=40 \text{ AND } P \geq 10}(T)$$

What is $T(Z)$?

Question 1:

Relation $R(A, B, C)$ is given with:

- $T(R) = 5000$ tuples
- $V(R, A) = 25$

Define W as:

$$W = \sigma_{A=5 \text{ AND } B > 20}(R)$$

What is $T(W)$?

Question 2:

Relation $S(X, Y, Z)$ has:

- $T(S) = 20000$
- $V(S, X) = 100$
- $V(S, Y) = 40$

Define Q as:

$$Q = \sigma_{X=30 \text{ AND } Y < 15}(S)$$

What is $T(Q)$?

Examples with Solutions:**Example 1:**

Given relation $R(A, B, C)$ with:

- $T(R) = 10000$
- $V(R, A) = 50$
- $V(R, B) = 20$

Find $T(W)$ for:

$$W = \sigma_{A=10 \text{ AND } B < 10}(R)$$

Solution:

1. The selection $A = 10$ reduces the number of tuples to:

$$\frac{T(R)}{V(R, A)} = \frac{10000}{50} = 200$$

2. The second condition $B < 10$ means that we consider only part of the values of B . Assuming uniform distribution, approximately half of the values of B are less than 10, so:

$$T(W) = \frac{200}{2} = 100$$

Example 2:

Given relation $S(X, Y, Z)$ with:

- $T(S) = 8000$
- $V(S, X) = 40$
- $V(S, Y) = 10$

Find $T(Q)$ for:

$$Q = \sigma_{X=5 \text{ AND } Y=3}(S)$$

Solution:

1. The selection $X = 5$ reduces the tuples to:

$$\frac{T(S)}{V(S, X)} = \frac{8000}{40} = 200$$

2. The selection $Y = 3$ further reduces it:

$$\frac{200}{V(S, Y)} = \frac{200}{10} = 20$$

Thus, $T(Q) = 20$.

Example 3:

Given relation $T(M, N, P, Q)$ with:

- $T(T) = 12000$
- $V(T, M) = 60$
- $V(T, P) = 30$

Find $T(Z)$ for:

$$Z = \sigma_{M=12 \text{ AND } P>15}(T)$$

Solution:

1. The selection $M = 12$ reduces the tuples to:

$$\frac{T(T)}{V(T, M)} = \frac{12000}{60} = 200$$

2. The selection $P > 15$ means we consider a portion of the P values. Assuming uniform distribution, half of the P values satisfy the condition:

$$T(Z) = \frac{200}{2} = 100$$

Thus, $T(Z) = 100$.

Preservation of value sets

For $R(A,B)$ and $S(B,C)$

1. $W(A,B) = \sigma_{A=10}(R)$

a. $V(W,A) = 1$

b. $V(W,B) = V(R,B)$

After a selection, the values for attributes other than the one on which selection was performed retain their distinct values.

2. $W(A,B,C) = R \bowtie S$

a. $V(W,A) = V(R,A)$

b. $V(W,B) = \min \{V(R,B), V(S,B)\}$

c. $V(W,B) = V(R,B)$

Similarly after a Join, attributes other than ones on which Join was performed retain their distinct values.

1. Selection: $W(A, B) = \sigma_{A=10}(R)$

- **Selection (σ)** filters tuples based on a condition, but it does not change the values of attributes that are not part of the selection condition.
- The image provides:
 - $V(W, A) = 1 \rightarrow$ Since selection **fixes** the value of A to 10, it has only **one distinct value**.
 - $V(W, B) = V(R, B) \rightarrow$ The attribute B remains unchanged, meaning it retains all distinct values from R .
- **Example:**
 - Assume $R(A,B)$ contains:

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$(10, x1), (10, x2), (15, x3), (10, x4), (20, x5)$

- If we apply $\sigma_{A=10}(R)$, only tuples where $A = 10$ remain:

SCSS

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$(10, x1), (10, x2), (10, x4)$



- The distinct values of B are still $\{x1, x2, x4\}$, meaning the values for B are **preserved**.

2. Join: $W(A, B, C) = R \bowtie S$

- Join (\bowtie) is performed on the common attribute **B**, which means:
 - $V(W, A) = V(R, A) \rightarrow$ The attribute **A** retains all distinct values from **R**.
 - $V(W, B) = \min\{V(R, B), V(S, B)\} \rightarrow$ Since the join matches tuples using **B**, the number of distinct values for **B** is **limited by the smaller set** from **R** and **S**.
 - $V(W, C) = V(S, C) \rightarrow$ The attribute **C** retains its distinct values from **S**.
- Example:

- Assume:



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$R(A,B) = \{(1, X), (2, Y), (3, Z)\}$
 $S(B,C) = \{(X, M), (Y, N), (Y, O), (Z, P)\}$

- The join on **B** results in:

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$(1, X, M), (2, Y, N), (2, Y, O), (3, Z, P)$

- $V(W, A) = V(R, A) = \{1, 2, 3\}$ (Preserved)
- $V(W, B) = \min(V(R, B), V(S, B)) \rightarrow \{X, Y, Z\}$ (since Y appears twice)
- $V(W, C) = V(S, C) = \{M, N, O, P\}$ (Preserved)

Conclusion

The **Preservation of Value Sets** states that:

- Selection retains all distinct values** for attributes not used in the condition.
 - Join preserves distinct values** for non-join attributes but **reduces** values of the join attribute to the intersection of both relations.
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Exercise 6

For $R(A,B)$, $S(B,C)$ and $T(D,E)$ we have the following specs:

$T(R) = 1000$, $V(R,A) = 1$, $V(R,B) = 30$

$T(S) = 200$, $V(S,B) = 40$, $V(S,C) = 10$

$T(T) = 10000$, $V(T,D) = 10$, $V(T,E) = 20$

Estimate the number of Tuples for the results.

a) $R \bowtie S$

b) $R \bowtie \sigma_{C=10}(S)$

c) $R \bowtie \sigma_{B=5 \text{ AND } C=10}(S)$

d) $\sigma_{A=10}(R) \bowtie S$ Assume result of $\sigma_{A=10}(R)$ is not empty.

e) $\pi_A(R) \bowtie S$

i) Assuming projection will not eliminate duplicates

ii) Assuming projection will eliminate duplicates

f) $R \bowtie S \bowtie T$

Exercise 7

For relations $R(A,B)$, $S(B,C)$, and $T(C,D)$, we have the following specifications:

- $T(R) = 800$, $V(R,A) = 2$, $V(R,B) = 40$
- $T(S) = 300$, $V(S,B) = 30$, $V(S,C) = 20$
- $T(T) = 5000$, $V(T,C) = 25$, $V(T,D) = 50$

Estimate the number of tuples for the results:

a) $R \bowtie S$

b) $R \bowtie \sigma_{C=5}(S)$

c) $R \bowtie \sigma_{B=20 \text{ AND } C=5}(S)$

d) $\sigma_{A=2}(R) \bowtie S$ (Assume $\sigma_{A=2}(R)$ is not empty.)

e) $\pi_A(R) \bowtie S$

- i) Assuming projection will not eliminate duplicates
- ii) Assuming projection will eliminate duplicates

f) $R \bowtie S \bowtie T$

Exercise 8

For relations $P(X,Y)$, $Q(Y,Z)$, and $R(Z,W)$, we have the following specifications:

- $T(P) = 600$, $V(P,X) = 5$, $V(P,Y) = 50$
- $T(Q) = 400$, $V(Q,Y) = 40$, $V(Q,Z) = 30$
- $T(R) = 7000$, $V(R,Z) = 35$, $V(R,W) = 70$

Estimate the number of tuples for the results:

- a) $P \bowtie Q$
- b) $P \bowtie \sigma_{Z=15}(Q)$
- c) $P \bowtie \sigma_{Y=10 \text{ AND } Z=15}(Q)$
- d) $\sigma_{X=3}(P) \bowtie Q$ (Assume $\sigma_{X=3}(P)$ is not empty.)
- e) $\pi_X(P) \bowtie Q$
 - i) Assuming projection will not eliminate duplicates
 - ii) Assuming projection will eliminate duplicates
- f) $P \bowtie Q \bowtie R$

Exercise 9

For relations $A(M,N)$, $B(N,O)$, and $C(O,P)$, we have the following specifications:

- $T(A) = 1200$, $V(A,M) = 10$, $V(A,N) = 60$
- $T(B) = 500$, $V(B,N) = 50$, $V(B,O) = 40$
- $T(C) = 9000$, $V(C,O) = 45$, $V(C,P) = 90$

Estimate the number of tuples for the results:

- a) $A \bowtie B$
 - b) $A \bowtie \sigma_{O=8}(B)$
 - c) $A \bowtie \sigma_{N=12 \text{ AND } O=8}(B)$
 - d) $\sigma_{M=7}(A) \bowtie B$ (Assume $\sigma_{M=7}(A)$ is not empty.)
 - e) $\pi_M(A) \bowtie B$
 - i) Assuming projection will not eliminate duplicates
 - ii) Assuming projection will eliminate duplicates
 - f) $A \bowtie B \bowtie C$
-

Exercise 7

UPDATED

For $R(X,Y)$ we have the following specs:

$T(R) = 20000$, $V(R,X) = 500$, $V(R,Y) = 10$

Assume Attribute $X = 20$ bytes, $Y = 30$ bytes. And there is a **clustering** index on attribute X .

With block size 4096 bytes and a block header of 96bytes.

- a) What is $B(R)$?
 - i) For Spanned
 - ii) Unspanned
 - b) How many tuples will the following queries return? What will be the size of the result in bytes?
 - i) $\sigma_{X=10}(R)$
 - ii) $\pi_Y(R)$ (**Assume projection will not eliminate duplicates.**)
 - iii) How many I/O would you roughly need to retrieve the records for (i)? (Assume index to be in memory.)
-

Exercise 8

For $S(A,B,C)$, we have the following specifications:

- $T(S) = 50000$, $V(S,A) = 1000$, $V(S,B) = 500$, $V(S,C) = 50$
- Assume Attribute $A = 16$ bytes, $B = 24$ bytes, $C = 40$ bytes.
- There is a **clustering index** on attribute A .
- Block size is **8192 bytes**, with a block header of **128 bytes**.

Questions:

- a) What is $B(S)$?
 - i) For **Spanned** storage
 - ii) For **Unspanned** storage
 - b) How many tuples will the following queries return? What will be the size of the result in bytes?
 - i) $\sigma_{A=50}(S)$
 - ii) $\pi_B(S)$ (**Assume projection will not eliminate duplicates.**)
 - iii) How many I/O operations are needed to retrieve the records for (i)? (**Assume the index is in memory.**)
-

Exercise 9

For $T(D,E,F,G)$, we have the following specifications:

- $T(T) = 150000$, $V(T,D) = 2000$, $V(T,E) = 1000$, $V(T,F) = 500$, $V(T,G) = 200$
- Assume **Attribute D = 12 bytes**, **E = 20 bytes**, **F = 28 bytes**, **G = 36 bytes**.
- There is a **clustering index** on attribute **D**.
- Block size is **4096 bytes**, with a block header of **64 bytes**.

Questions:

a) What is $B(T)$?

i) For **Spanned** storage

ii) For **Unspanned** storage

b) How many tuples will the following queries return? What will be the size of the result in bytes?

i) $\sigma_{E=100}(T)$

ii) $\pi_F(T)$ (Assume projection will not eliminate duplicates.)

iii) How many **I/O operations** are needed to retrieve the records for (i)? (Assume the index is in memory.)

Exercise 10

For $U(X,Y,Z)$, we have the following specifications:

- $T(U) = 100000$, $V(U,X) = 5000$, $V(U,Y) = 2000$, $V(U,Z) = 1000$
- Assume **Attribute X = 18 bytes**, **Y = 22 bytes**, **Z = 32 bytes**.
- There is a **clustering index** on attribute **X**.
- Block size is **8192 bytes**, with a block header of **128 bytes**.

Questions:

a) What is $B(U)$?

i) For **Spanned** storage

ii) For **Unspanned** storage

b) How many tuples will the following queries return? What will be the size of the result in bytes?

i) $\sigma_{X=200}(U)$

ii) $\pi_Y(U)$ (Assume projection will not eliminate duplicates.)

iii) How many **I/O operations** are needed to retrieve the records for (i)? (Assume the index is in memory.)

Here are **three additional exercises** similar to the one you provided, covering **spanned vs. unspanned storage**, **query result estimation**, and **I/O cost analysis**.

Exercise 1

For relation $S(A, B, C)$, we have the following specifications:

- $T(S) = 50000$, $V(S, A) = 1000$, $V(S, B) = 2500$, $V(S, C) = 500$
- Assume attribute sizes:
 - $A = 15$ bytes, $B = 25$ bytes, $C = 40$ bytes
- Block size = 4096 bytes with a block header of 96 bytes
- There is a clustering index on attribute B.

Questions

(a) Compute $B(S)$ (number of blocks needed) for:

- Spanned storage
- Unspanned storage

(b) How many tuples will the following queries return? What will be the **size of the result in bytes**?

1. $\sigma_{B=500}(S)$
2. $\pi_A(S)$ (Assume projection does not eliminate duplicates.)

(c) Estimate the **I/O cost** to retrieve the tuples for $\sigma_{B=500}(S)$ (Assume the index is in memory).

Exercise 2

For relation $T(P, Q, R)$, we have the following specifications:

- $T(T) = 250000$, $V(T, P) = 5000$, $V(T, Q) = 10000$, $V(T, R) = 2000$
- Assume attribute sizes:
 - $P = 10$ bytes, $Q = 35$ bytes, $R = 45$ bytes
- Block size = 8192 bytes with a block header of 128 bytes
- There is a **secondary index on attribute P** (not clustering).

Questions

(a) Compute $B(T)$ for:

- Spanned storage
- Unspanned storage

(b) Estimate the number of tuples and result size for:

1. $\sigma_{P=1000}(T)$
2. $\pi_Q(T)$ (Assume projection does not eliminate duplicates.)

(c) Estimate the **I/O cost** to retrieve the tuples for $\sigma_{P=1000}(T)$ (Assume index is in memory, but records are scattered across blocks).

Exercise 3

For relation $R(X, Y, Z, W)$, we have the following specifications:

- $T(R) = 120000$, $V(R, X) = 8000$, $V(R, Y) = 3000$, $V(R, Z) = 1200$, $V(R, W) = 600$
- Assume attribute sizes:
 - $X = 12$ bytes, $Y = 18$ bytes, $Z = 20$ bytes, $W = 30$ bytes
- Block size = 4096 bytes with a block header of 64 bytes
- There is a clustering index on attribute Z .

Questions

(a) Compute $B(R)$ for:

- Spanned storage
- Unspanned storage

(b) Estimate the number of tuples and result size for:

1. $\sigma_{Z=300}(R)$
2. $\pi_{X,Y}(R)$ (Assume projection does not eliminate duplicates.)

(c) Estimate the I/O cost to retrieve the tuples for $\sigma_{Z=300}(R)$ (Assume index is in memory).

Exercise 8

UPDATED

For $R(A,B,C,D)$ with following specs:

$T(R) = 1000$, $V(R,A) = 1$, $V(R,B) = 30$, $V(R,C)=10$, $V(R,D)=20$

Estimate the number of Tuples for the results and the $V(W,X)$, $W=\text{result}$, $X=\{A,B,C,D\}$. For parts c and d assume projection to eliminate duplicates.

- a) $\sigma_{A=10}(R)$
 - b) $\sigma_{B=5 \text{ AND } C<10}(R)$
 - c) $\pi_A(R)$
 - d) $\pi_{A,B}(R)$
-

Question 1: Tuple and Distinct Value Estimation

For relation $R(A, B, C, D)$ with the following specifications:

- $T(R) = 2000$
- $V(R, A) = 1$
- $V(R, B) = 50$
- $V(R, C) = 20$
- $V(R, D) = 40$

Estimate the number of tuples for the results and the $V(W, X)$, where W is the result and $X = \{A, B, C, D\}$. For parts c and d, assume projection to eliminate duplicates.

- a) $\sigma_{A=1}(R)$
 - b) $\sigma_{B=10}(R)$
 - c) $\pi_{A,B}(R)$
 - d) $\pi_{C,D}(R)$
-

Question 2: Tuple and Distinct Value Estimation

For relation $R(A, B, C, D)$ with the following specifications:

- $T(R) = 1500$
- $V(R, A) = 2$
- $V(R, B) = 40$
- $V(R, C) = 15$
- $V(R, D) = 30$

Estimate the number of tuples for the results and the $V(W, X)$, where W is the result and $X = \{A, B, C, D\}$. For parts c and d, assume projection to eliminate duplicates.

- a) $\sigma_{A=2}(R)$
 - b) $\sigma_{C=5}(R)$
 - c) $\pi_{A,C}(R)$
 - d) $\pi_{B,D}(R)$
-

Question 3: Tuple and Distinct Value Estimation

For relation $R(A, B, C, D)$ with the following specifications:

- $T(R) = 3000$
- $V(R, A) = 3$
- $V(R, B) = 60$
- $V(R, C) = 25$
- $V(R, D) = 50$

Estimate the number of tuples for the results and the $V(W, X)$, where W is the result and $X = \{A, B, C, D\}$. For parts c and d, assume projection to eliminate duplicates.

a) $\sigma_{B=20}(R)$

b) $\sigma_{D=10}(R)$

c) $\pi_{A,D}(R)$

d) $\pi_{B,C}(R)$

Exercise 9

For S(A,B) with following specs:

$$T(S) = 40000, B(S) = 500$$

$$V(S,A) = 4000, V(S,B) = 2000$$

Assume S has a primary dense index on A and secondary index on B.

How many IOs would you expect the following operations to take if the indices are in memory.

a) $\sigma_{A=10}(S)$

b) $\sigma_{B=5}(S)$

Exercise 10

For S(A, B) with the following specs:

$$T(S) = 50,000, B(S) = 600$$

$$V(S, A) = 5,000, V(S, B) = 2,500$$

Assume S has a primary dense index on A and a secondary index on B.

How many IOs would you expect the following operations to take if the indices are in memory?

(a) $\sigma_{A=20}(S)$

(b) $\sigma_{B=15}(S)$

Exercise 11

For S(A, B, C) with the following specs:

$$T(S) = 75,000, B(S) = 800$$

$$V(S, A) = 7,500, V(S, B) = 3,000, V(S, C) = 1,500$$

Assume S has a primary dense index on A and a secondary index on B.

How many IOs would you expect the following operations to take if the indices are in memory?

(a) $\sigma_{A=30}(S)$

(b) $\sigma_{B=10}(S)$

Exercise 12

For S(A, B) with the following specs:

$$T(S) = 100,000, B(S) = 1,000$$

$$V(S, A) = 10,000, V(S, B) = 5,000$$

Assume S has a primary dense index on A and a secondary index on B.

How many IOs would you expect the following operations to take if the indices are in memory?

(a) $\sigma_{A=50}(S)$

(b) $\sigma_{B=25}(S)$

Exercises with Solutions

Exercise 13 (With Solution)

For $S(A, B)$ with the following specs:

$$T(S) = 80,000, B(S) = 900$$

$$V(S, A) = 8,000, V(S, B) = 4,000$$

Assume S has a primary dense index on A and a secondary index on B .

How many IOs would you expect the following operations to take if the indices are in memory?

Solution

(a) $\sigma_{A=40}(S)$

- A has a **primary dense index**, so we can directly access the tuples.
- The number of tuples matching $A = 40$ is:

$$\frac{T(S)}{V(S, A)} = \frac{80,000}{8,000} = 10$$

- Since the index is in memory, the IO cost is **10** (one per tuple).

(b) $\sigma_{B=20}(S)$

- B has a **secondary index**, so we must retrieve **all** tuples matching $B = 20$.
- The number of tuples matching $B = 20$ is:

$$\frac{T(S)}{V(S, B)} = \frac{80,000}{4,000} = 20$$

- Since the secondary index is in memory, each tuple access requires an additional IO.
 - Total IOs = **20**.
-

Exercise 14 (With Solution)

For $S(A, B)$ with the following specs:

$$T(S) = 60,000, B(S) = 700$$

$$V(S, A) = 6,000, V(S, B) = 3,000$$

Assume S has a primary dense index on A and a secondary index on B .

How many IOs would you expect the following operations to take if the indices are in memory?

Solution

(a) $\sigma_{A=15}(S)$

- A has a **primary dense index**, so direct access is possible.
- Number of matching tuples:

$$\frac{T(S)}{V(S, A)} = \frac{60,000}{6,000} = 10$$

- IO cost: **10**.

(b) $\sigma_{B=30}(S)$

- B has a **secondary index**, so we must scan all matching tuples.
- Number of matching tuples:

$$\frac{T(S)}{V(S, B)} = \frac{60,000}{3,000} = 20$$

- IO cost: **20**.

