



## Index Structures

Records ordered by search key (may not be "key" in DB sense).

- facilitates queries on the search key

Blocks containing records therefore ordered.

- physically contiguous
- chained

On insert: put record in appropriate block if room.

- Good idea: initialize blocks to be less than full; reorganize periodically if file grows.

If no room in proper block:

- 1. Create new block; insert into proper order if possible.
- 2. If not possible, create overflow block, linked from original block.

## What is an Index?

An index is a data structure that allows us to directly locate units of data based on certain values

- Not just used for databases: also books can contain an index

Indexes for databases are used to find records that have a particular value for the indexed attribute (the "search key")

An index has to be created before it can be used

- creation often initiated by the database designer
- cost of maintenance

Different categories exist

- primary / secondary indexes
- dense / sparse indexes

## Indexes

Dense Indexes: Pointer to every record of file, ordered by search key.

- Can make sense because records may be much bigger than key-pointer pairs.
  - If index requires fewer blocks faster search through index than data file
  - Index might fit in memory, even if data file does not
- Test existence of record without going to data file.

Sparse Indexes: Keypointer pairs for only a subset of records, typically first in each block.

## Implementation of DBMS

### Example: Sequential File

Sequential File

10	
20	
30	
40	
50	
60	
70	
80	
90	
100	

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### Example: Dense Index

Dense Index

10	
20	
30	
40	
50	
60	
70	
80	
90	
100	
110	
120	

Sequential File

10	
20	
30	
40	
50	
60	
70	
80	
90	
100	
110	
120	

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### Implementation of DBMS

### Example: Sparse Index

Sparse Index

Sequential File

10	
30	
50	
70	
90	
110	
130	
150	
170	
190	
210	
230	

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### Implementation of DBMS

### Sparse vs. Dense Index

Sparse: Less index space per record can keep more of index in memory

Dense: Can tell if any record exists without accessing file

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# Multiple Levels of Index

A sparse index on a (sparse or dense) index is an option.

Good chance that 2nd or higher level indexes can be housed in main memory, so no additional disk I/O's.

Dense higher level indexes make no sense;

# DB Modifications

When we insert or delete on the data file, here are the primitive actions we might take:

1. Create or destroy an overflow block.
2. Create or destroy an empty block in the sequence of blocks belonging to the sequential file.
3. Insert a record into a block that has room.
4. Delete a record.
5. Slide a record to an adjacent block.

# Example: Second Level Index

## Sparse 2nd level

10	
90	
170	
250	

330	
410	
490	
570	

170	
190	
210	
230	

## Sequential File

10	
20	

30	
40	

50	
60	

70	
80	

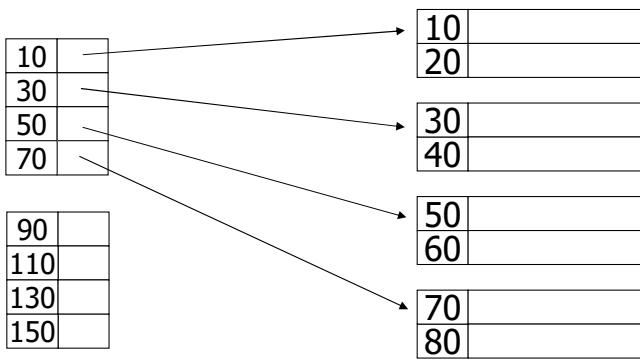
90	
100	

# Effect of Primitive Actions on Index File

Action	Dense	Sparse
Create/destroy empty overflow block	none	none
Create empty seq. block	none	insert
Destroy empty seq. block	none	delete
Insert record	insert	update(?)
Delete record	delete	update(?)
Slide record	update	update(?)

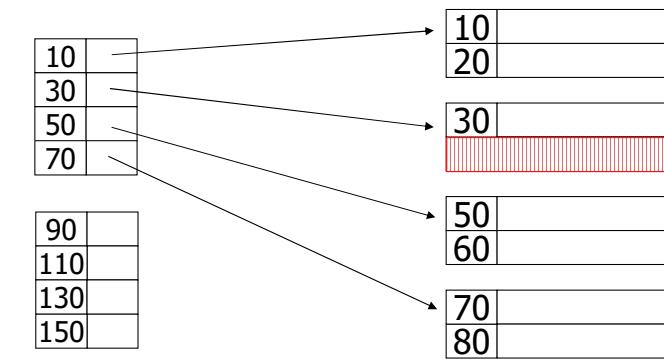
## Implementation of DBMS

### Deletion from sparse index



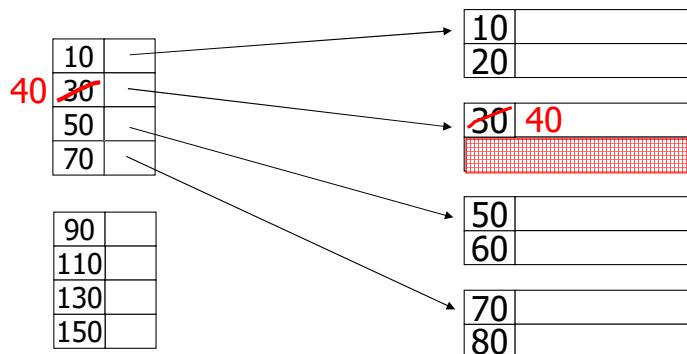
### Deletion from sparse index

– delete record 40



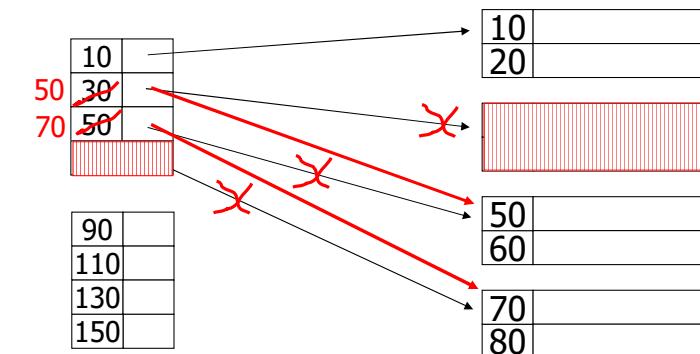
### Deletion from sparse index

– delete record 30



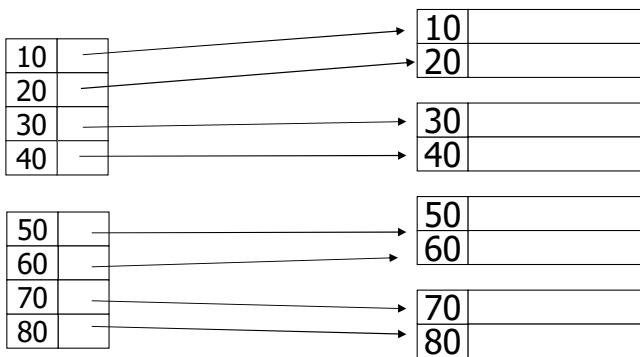
### Deletion from sparse index

– delete records 30 & 40



## Implementation of DBMS

### Deletion from dense index

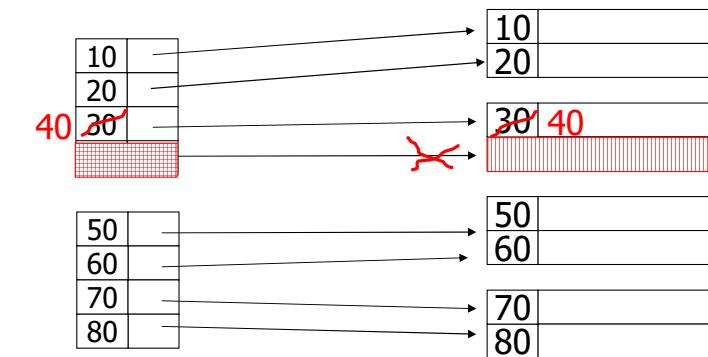


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### Deletion from dense index

- delete record 30

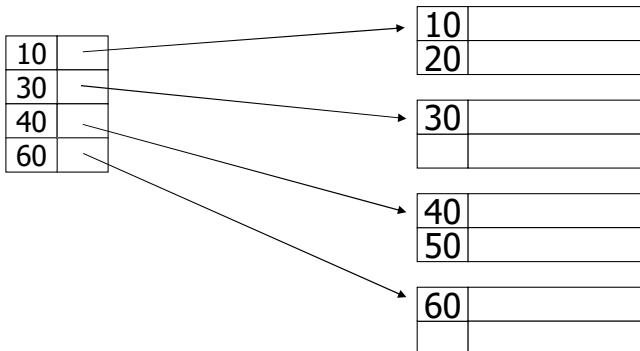


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## Implementation of DBMS

### Insertion, sparse index case



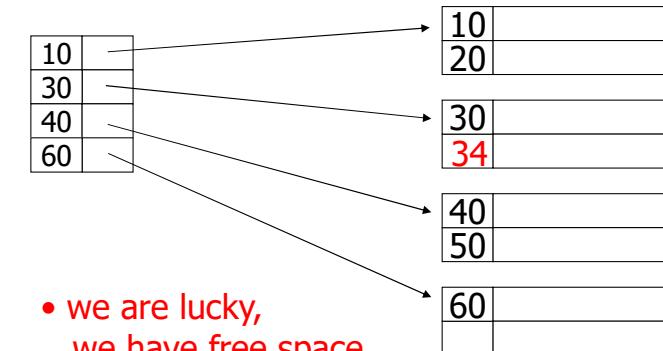
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### Implementation of DBMS

### Insertion, sparse index case

- insert record 34



• we are lucky,  
we have free space  
where we need it!

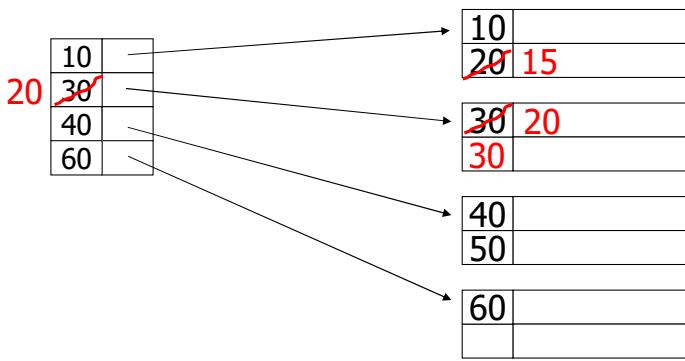
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## Implementation of DBMS

# Insertion, sparse index case

- insert record 15



- Illustrated: Immediate reorganization
- Variation:
  - insert new block (chained file)
  - update index

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## Implementation of DBMS

# Insertion, dense index case

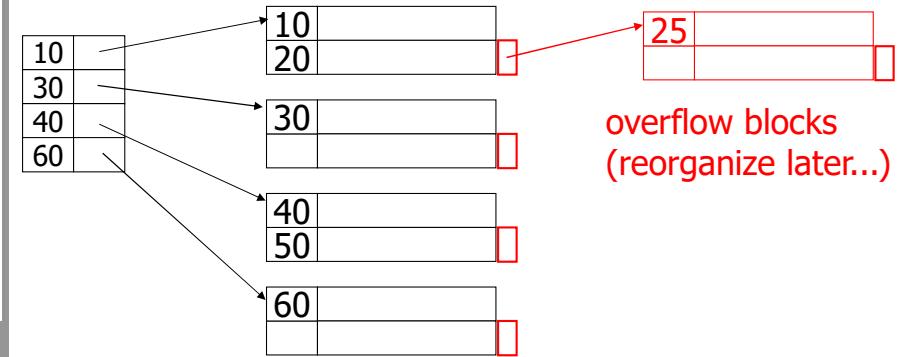
- Similar
- Often more expensive . . .

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## Implementation of DBMS

# Insertion, sparse index case

- insert record 25



overflow blocks  
(reorganize later...)

## Implementation of DBMS

# Secondary Indexes

A primary index is an index on a sorted file.

- More general: any index that "controls" the placement of records to be primary, e.g., hash table.

Secondary index = index that does not control placement, surely not on a file sorted by its search key.

- Sparse, secondary index makes no sense.
- Usually, search key is not a "key"

Multiple Levels:

- Lowest level is dense
- Other levels are sparse

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## Secondary Indexes

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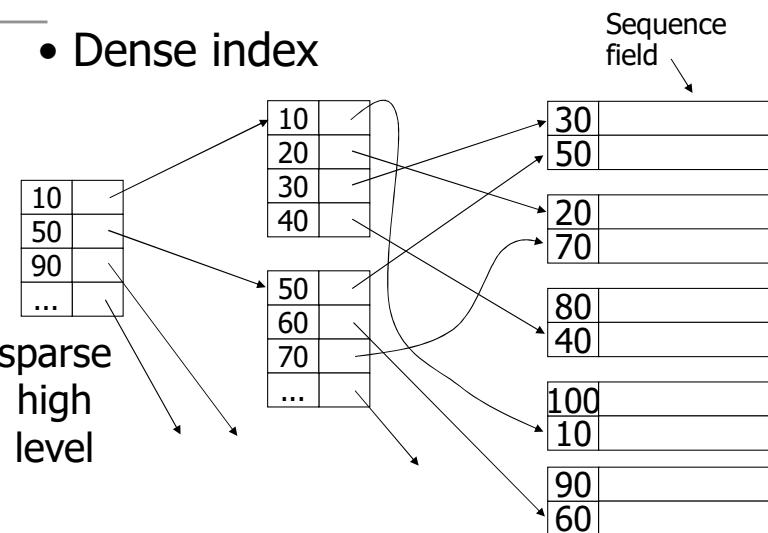
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## Secondary Indexes

- Dense index



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## Secondary Indexes

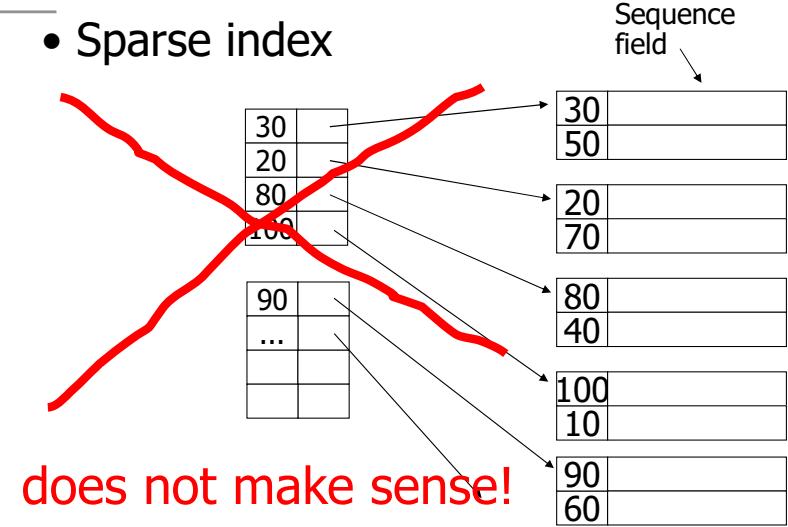
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- Sparse index



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does not make sense!

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## Duplicate values & secondary indexes

20	
10	
20	
40	
10	
40	
10	
40	
30	
40	

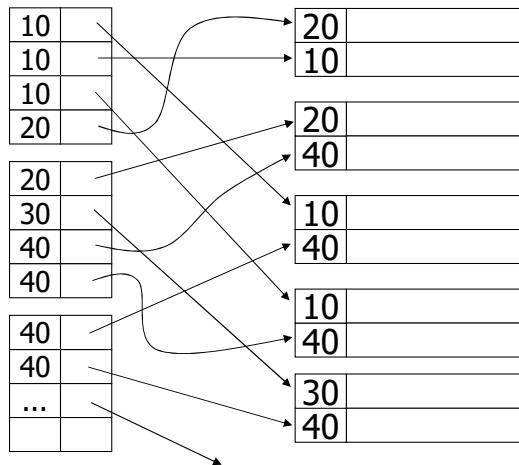
## Duplicate values & secondary indexes

one option...

### Problem:

excess overhead!

- disk space
- search time

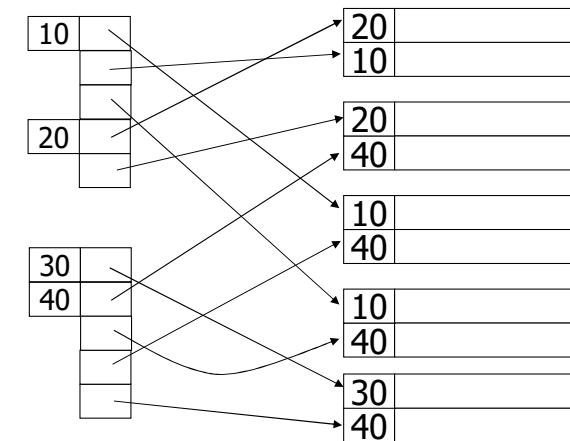


## Duplicate values & secondary indexes

another option...

### Problem:

variable size records in index!



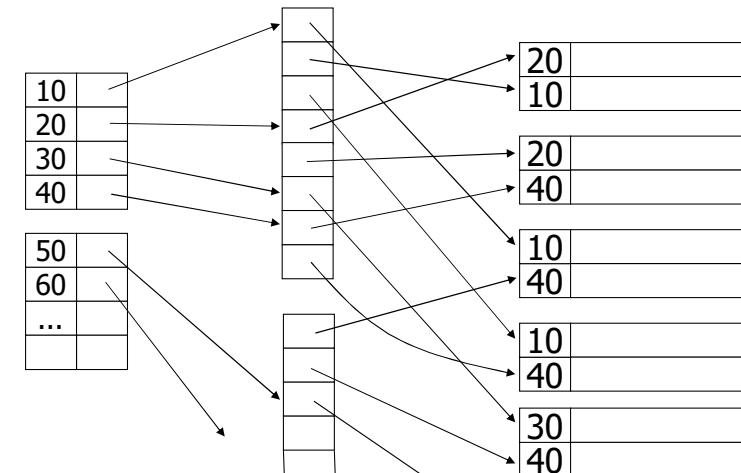
## Indirect Buckets

To avoid repeating keys in index, use a level of indirection, called buckets.

- Additional advantage: allows intersection of sets of records without looking at records themselves.

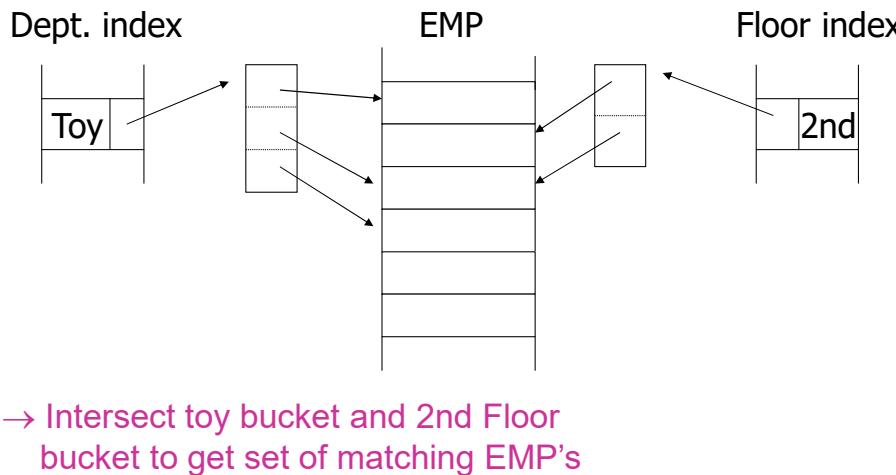
## Duplicate values & secondary indexes

buckets



## Indirect Buckets

Query: Get employees in  
(Toy Dept)  $\wedge$  (2nd floor)



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## Assessment of Conventional Indexes

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Advantage:

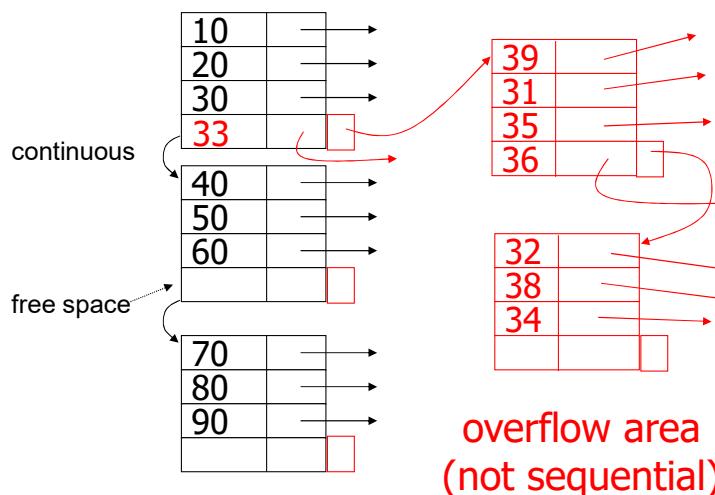
- Simple
- Index is sequential file good for scans

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Disadvantage:

- Inserts expensive, and/or
- Lose sequentiality & balance

## Example

Index (sequential)



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