Join Based on Hash

This part of the project was a lot more fun, so the quality of my code is significantly better than the one for B+ tree, mainly because it was easier to understand and plan for this part of the project. I used vectors and structures for S and R relations, and I kept them on different variables based on their location in memory or disk. This way I could make sure my memory usage does not go beyond what is available and at the same time I use the read and write function to interact with disk, so I count the number of reads and writes correctly. I also used templates so that I do not repeat code, and instead making the compiler do all the dirty work. Overall, my software architecture went better than I expected because everything is very neat and clean. I wish I could say the same about B+ tree.

In terms of generating data, I have used random generation in C++, the only issue was its boundary not reaching the values we need for the range given, and to achieve that I multiplied to random number to gather and found their remainder based on the range we need. I analyzed the distribution to make sure the result does not skew, and it had acceptable performance.

My disk I/O operations was simulated by having different vectors so I cannot accidentally read from the disk vector. It was a pretty easy strategy to make sure I end of with accurate result at the end. This part of the assignment is easier, but it is important to not miss the reads and writes. After many hours of testing I know for a fact that my approach of using two vectors helped a lot in not falling into pitfalls.

For hash function, I did two different things, but I ended up sticking to second one because it would work better for bot experiment cases. My first hash function was range based, so it would decide based on the range of B values and would divide them equally between buckets. This worked well up till 5.2 where R had its B values from 20,000 to 30,000 which the range approach would result in one overpopulated bucket not fitting in memory. In second attempt I just used remainder to number of buckets available, this way regardless of the range I would get more uniform results in each bucket allowing me to load each bucket full in memory and do a one-pass natural join on it.

The algorithm for natural join using hash bucket is simple, first you hash each relation based on their B values and store them in disk. Considering that your hash function produces uniform result, we can conclude that each hash bucket has equal number of tuples. Now the tuples in this hash in R and S may match. The good news is that now the total number of tuples to compare is divided by number of buckets so you would need less memory to load one bucket from smaller relation fully. Number of buckets is your memory size in blocks minus 1. This is so that when you hash one block remain in memory from the relation and the other n-1 blocks store the latest tuples for each hash bucket. Once each hash bucket block in memory is full, the block is moved to disk with the rest of the block belonging to the same hash bucket. Now in this

example our memory is 15 blocks and after leaving one block for reading a relation tuple, 14 blocks are left for 14 buckets. Each block can store 8 tuples so that means we can keep 8 tuples in each block. Once we have all tuples belonging to buckets, we bring the bucket from smaller relation (since our memory is small) which must fit in n-1 blocks. We use the block or blocks left to read the relation and try to find matching B values. Take note that in our case this means that each bucket should have 14*8 = 112 tuples. So this means we have only one option to bring R hash bucket in memory because 5000/14 = ~357, way above our memory capacity. However with R having 1200 tuples at most, then we will have 1200/14 = ~86 which is lower than 112 and therefore allows for a hash function that does not produce perfectly uniform result, but then it has to be good enough to make sure each bucket is below 112. The second hash function explained above that I ended up using can manage that easily.

Question	I/O Reads	I/O Writes	Total I/O
5.1	1511	761	2272
5.2	1563	788	2351

I have checked the result in table above to make sure it exactly matches the expected result based on the tuple distribution in buckets and considering that we read or write a whole block even if we mean to read from or write to portion of block.

```
Hashed S, 5000 tuples were hashed.

Bucket[0]: remainder is 0 when devided by 14 (number of buckets), and has 345 tuples
Bucket[1]: remainder is 1 when devided by 14 (number of buckets), and has 344 tuples
Bucket[2]: remainder is 2 when devided by 14 (number of buckets), and has 350 tuples
Bucket[3]: remainder is 3 when devided by 14 (number of buckets), and has 378 tuples
Bucket[4]: remainder is 4 when devided by 14 (number of buckets), and has 329 tuples
Bucket[5]: remainder is 5 when devided by 14 (number of buckets), and has 361 tuples
Bucket[6]: remainder is 6 when devided by 14 (number of buckets), and has 375 tuples
Bucket[7]: remainder is 7 when devided by 14 (number of buckets), and has 357 tuples
Bucket[8]: remainder is 8 when devided by 14 (number of buckets), and has 378 tuples
Bucket[9]: remainder is 10 when devided by 14 (number of buckets), and has 324 tuples
Bucket[10]: remainder is 11 when devided by 14 (number of buckets), and has 324 tuples
Bucket[11]: remainder is 11 when devided by 14 (number of buckets), and has 330 tuples
Bucket[12]: remainder is 12 when devided by 14 (number of buckets), and has 334 tuples
Bucket[13]: remainder is 13 when devided by 14 (number of buckets), and has 333 tuples
```

Figure 1 Hash bucket distribution for S with 5000 tuples, theoretically hash function should generate 357 tuples per bucket

```
Hashed R, 1000 tuples were hashed.

Bucket[0]: remainder is 0 when devided by 14 (number of buckets), and has 59 tuples

Bucket[1]: remainder is 1 when devided by 14 (number of buckets), and has 72 tuples

Bucket[2]: remainder is 2 when devided by 14 (number of buckets), and has 63 tuples

Bucket[3]: remainder is 3 when devided by 14 (number of buckets), and has 70 tuples

Bucket[4]: remainder is 4 when devided by 14 (number of buckets), and has 67 tuples

Bucket[5]: remainder is 5 when devided by 14 (number of buckets), and has 78 tuples

Bucket[6]: remainder is 6 when devided by 14 (number of buckets), and has 80 tuples

Bucket[7]: remainder is 7 when devided by 14 (number of buckets), and has 71 tuples

Bucket[8]: remainder is 8 when devided by 14 (number of buckets), and has 94 tuples

Bucket[9]: remainder is 10 when devided by 14 (number of buckets), and has 73 tuples

Bucket[10]: remainder is 10 when devided by 14 (number of buckets), and has 72 tuples

Bucket[11]: remainder is 11 when devided by 14 (number of buckets), and has 56 tuples

Bucket[12]: remainder is 12 when devided by 14 (number of buckets), and has 56 tuples

Bucket[13]: remainder is 13 when devided by 14 (number of buckets), and has 62 tuples
```

Figure 2 Hash bucket distribution for R with 1000 tuples with B values selected from B values of relation S, theoretically hash function should generate 71 tuples per bucket

```
Hashed R, 1200 tuples were hashed.

Bucket[0]: remainder is 0 when devided by 14 (number of buckets), and has 94 tuples

Bucket[1]: remainder is 1 when devided by 14 (number of buckets), and has 93 tuples

Bucket[2]: remainder is 2 when devided by 14 (number of buckets), and has 80 tuples

Bucket[3]: remainder is 3 when devided by 14 (number of buckets), and has 89 tuples

Bucket[4]: remainder is 4 when devided by 14 (number of buckets), and has 103 tuples

Bucket[5]: remainder is 5 when devided by 14 (number of buckets), and has 84 tuples

Bucket[6]: remainder is 6 when devided by 14 (number of buckets), and has 87 tuples

Bucket[7]: remainder is 7 when devided by 14 (number of buckets), and has 81 tuples

Bucket[8]: remainder is 8 when devided by 14 (number of buckets), and has 85 tuples

Bucket[10]: remainder is 10 when devided by 14 (number of buckets), and has 90 tuples

Bucket[11]: remainder is 11 when devided by 14 (number of buckets), and has 78 tuples

Bucket[12]: remainder is 12 when devided by 14 (number of buckets), and has 72 tuples

Bucket[13]: remainder is 13 when devided by 14 (number of buckets), and has 75 tuples
```

Figure 3 Hash bucket distribution for R with 1200 tuples with B values selected from 20,000 to 30,000, theoretically hash function should generate 86 tuples per bucket

```
Starting the natural joint operation on R (Smaller one) and S.

Hash bucket[0]: 59 R tupples were moved. 345 S tupples were processed, finding 67 matching tuple. Total Matching: 67.

Hash bucket[1]: 72 R tupples were moved. 344 S tupples were processed, finding 83 matching tuple. Total Matching: 150.

Hash bucket[2]: 63 R tupples were moved. 350 S tupples were processed, finding 68 matching tuple. Total Matching: 218.

Hash bucket[3]: 70 R tupples were moved. 378 S tupples were processed, finding 85 matching tuple. Total Matching: 303.

Hash bucket[4]: 67 R tupples were moved. 329 S tupples were processed, finding 87 matching tuple. Total Matching: 376.

Hash bucket[5]: 78 R tupples were moved. 361 S tupples were processed, finding 87 matching tuple. Total Matching: 464.

Hash bucket[6]: 80 R tupples were moved. 375 S tupples were processed, finding 89 matching tuple. Total Matching: 553.

Hash bucket[8]: 71 R tupples were moved. 375 S tupples were processed, finding 84 matching tuple. Total Matching: 637.

Hash bucket[8]: 94 R tupples were moved. 378 S tupples were processed, finding 110 matching tuple. Total Matching: 747.

Hash bucket[9]: 73 R tupples were moved. 392 S tupples were processed, finding 84 matching tuple. Total Matching: 811.

Hash bucket[10]: 83 R tupples were moved. 324 S tupples were processed, finding 89 matching tuple. Total Matching: 926.

Hash bucket[11]: 72 R tupples were moved. 380 S tupples were processed, finding 83 matching tuple. Total Matching: 1009.

Hash bucket[13]: 60 R tupples were moved. 334 S tupples were processed, finding 73 matching tuple. Total Matching: 1143.

>>>> 1143 tuples are in natural join results. Total reads: 1511, Total writes: 761, Total Operations: 2272
```

Figure 4 Natural join results for Q5.1, where R's B values are selected from S's B values. Consider that duplication was allowed. Also, the result per bucket is also shown to see how each bucket did.

```
Test 0: B value selected is 45516.But no matching B value was found in natural join of R and S.
Test 1: B value selected is 32010.But no matching B value was found in natural join of R and S.
Test 2: B value selected is 15204.
index:
          A I
         -575
                15204 -28372
  61:
               15204 -28372
         -257
Test 3: B value selected is 45519.
          A
                вΙ
index:
         -160
                45519 - 20948
 401:
Test 4: B value selected is 21875.
index:
          A
                21875 -48397
          -14
  600:
```

Figure 5 B value lookup in join result, based on Figure 16 results

```
Starting the natural joint operation on R (Smaller one) and S.

Hash bucket[0]: 94 R tupples were moved. 345 S tupples were processed, finding 6 matching tuple. Total Matching: 6.

Hash bucket[1]: 93 R tupples were moved. 344 S tupples were processed, finding 12 matching tuple. Total Matching: 18.

Hash bucket[2]: 80 R tupples were moved. 350 S tupples were processed, finding 15 matching tuple. Total Matching: 33.

Hash bucket[3]: 89 R tupples were moved. 378 S tupples were processed, finding 8 matching tuple. Total Matching: 41.

Hash bucket[4]: 103 R tupples were moved. 329 S tupples were processed, finding 8 matching tuple. Total Matching: 49.

Hash bucket[5]: 84 R tupples were moved. 361 S tupples were processed, finding 14 matching tuple. Total Matching: 63.

Hash bucket[6]: 87 R tupples were moved. 375 S tupples were processed, finding 14 matching tuple. Total Matching: 74.

Hash bucket[7]: 81 R tupples were moved. 378 S tupples were processed, finding 14 matching tuple. Total Matching: 88.

Hash bucket[8]: 89 R tupples were moved. 378 S tupples were processed, finding 11 matching tuple. Total Matching: 99.

Hash bucket[9]: 85 R tupples were moved. 392 S tupples were processed, finding 13 matching tuple. Total Matching: 117.

Hash bucket[10]: 90 R tupples were moved. 380 S tupples were processed, finding 13 matching tuple. Total Matching: 117.

Hash bucket[11]: 78 R tupples were moved. 334 S tupples were processed, finding 11 matching tuple. Total Matching: 130.

Hash bucket[13]: 75 R tupples were moved. 353 S tupples were processed, finding 8 matching tuple. Total Matching: 141.

Hash bucket[13]: 75 R tupples were moved. 353 S tupples were processed, finding 8 matching tuple. Total Matching: 149.
```

Figure 6 Natural join results for Q5.2, where R's B values are selected from 20,000 to 30,000. Consider that duplication was allowed. Also, the result per bucket is also shown to see how each bucket did.

Figure 7 printing join result as requested by Q5.2