#### Compression Schemes

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#### Flow

What? Why? How?

# What? PAC Learning $\longrightarrow$ Setting

- Domain set  $\mathcal{X}$ .
- Probability distribution over  $\mathcal{X}$ .  $\mathcal{D}$ :  $\mathcal{X} \to [0,1]$
- Label set  $\mathcal{Y}$ .  $\mathcal{Y} = \{0, 1\}$
- True labeling function  $f: f: \mathcal{X} \to \mathcal{Y}$
- Training data  $S^m$ .  $S^m = ((x_1, y_1), \dots, (x_m, y_m))$
- Concept class  $\mathcal{H}$ .  $\mathcal{H} \subseteq \mathcal{Y}^{\mathcal{X}}$
- The learner's output  $h \in \mathcal{H}$ .  $h: \mathcal{X} \to \mathcal{Y}$
- Measure of success

$$L_{\mathcal{D},f}(h) := \mathbb{P}_{x \sim \mathcal{D}}[h(x) \neq f(x)] := \mathcal{D}(\{x : h(x) \neq f(x)\})$$

# What? PAC Learnability — Finite Concept Classes

Objective:

$$\mathcal{D}^m(\{\mathcal{S}|_x: L_{D,f}(h_{\mathcal{S}}) > \epsilon\}) \le \delta$$

So:

$$m \ge \frac{\log(\frac{|\mathcal{H}|}{\delta})}{\epsilon}$$

- $\delta$ : Confidence Parameter
- $\bullet$   $\epsilon$ : Accuracy Parameter

# What? PAC Learnability — Finite VC Dimension

Objective:

$$\mathcal{D}^{m}(\{\mathcal{S}|_{x}: L_{D,f}(h_{\mathcal{S}}) > \epsilon\}) \leq \delta$$

So:

$$m \ge \max(\frac{32d}{\epsilon}\log(\frac{16e}{\epsilon}), \frac{16}{\epsilon}\log(\frac{2}{\delta}))$$
 (1)

$$m \ge \max(\frac{8d}{\epsilon}\log(\frac{8d}{\epsilon}), \frac{4}{\epsilon}\log(\frac{2}{\delta}))$$
 (2)

- $d: VCdim(\mathcal{H})$
- e: Euler's number

# What? PAC Learnability — Compression Schemes

#### Claim

Compression Schemes give weaker conditions for PAC learnability  $\iff$  A <u>lower</u> lower bound on the Sample Complexity

#### 

"Can we do better?"

#### How? PAC Learnability — Compression Schemes

- Kernel  $\kappa$ .  $\kappa: \bigcup_{m=k}^{\infty} \mathcal{S}^m \to \mathcal{S}^k$  (Compressor)
- Reconstructor  $\rho$ .  $\rho: \mathcal{S}^k \times \mathcal{X} \to \mathcal{Y} = \{0,1\}$  (Decompressor)
- $\blacksquare$   $\forall m \geq k, \mathcal{S}^k$  is a subsequence of length k of  $\mathcal{S}^m$

#### How? Compression Schemes

■ (Previously) Objective:

$$\mathcal{D}^m(\{\mathcal{S}|_x: L_{D,f}(h_{\mathcal{S}}) > \epsilon\}) \le \delta$$

■ (Now) Objective:

$$\mathcal{D}^{m}(\{\mathcal{S}|_{x}: L_{D,f}(\rho(\kappa(\mathcal{S}^{m}), x)) > \epsilon\}) \leq \delta$$

#### How? Compression Schemes $\longrightarrow$ Punchline

Let T be the collection of all k-element subsequences of the sequence (1, 2, ..., m). For any  $\bar{t} = (t_1, ..., t_k) \in T$ :

$$A_{\bar{t}} = \{ \mathcal{S}^m : \kappa(\mathcal{S}^m) = \mathcal{S}^k \}$$

$$E_{\bar{t}} = \{ \mathcal{S} \in A_{\bar{t}} : P(\{x : \rho(\kappa(\mathcal{S}), x) = f(x)\}) < 1 - \epsilon \}$$

$$U_{\bar{t}} = \{ \mathcal{S}^m : P(\{x : \rho(\mathcal{S}^k, x) = f(x)\}) < 1 - \epsilon \}$$

$$B_{\bar{t}} = \{ \mathcal{S}^m : \text{mark } \rho(\mathcal{S}^k, x_i) = f(x), \forall x_i \text{ s.t. } i \notin t \}$$

$$E_{\bar{t}} = U_{\bar{t}} \cap A_{\bar{t}} \xrightarrow{A_{\bar{t}} \subseteq B_{\bar{t}}} P(E_{\bar{t}}) \leq P(U_{\bar{t}} \cap A_{\bar{t}}) \leq \binom{m}{k} (1 - \epsilon)^{m-k}$$

#### Compression Schemes $\longrightarrow$ Results

Previously:

$$m \geq \max(\frac{8d}{\epsilon}\log(\frac{8d}{\epsilon}), \frac{4}{\epsilon}\log(\frac{2}{\delta}))$$

Now:

$$m \geq \max(\frac{4k}{\epsilon}\log(\frac{4k}{\epsilon}) + 2k, \frac{2}{\epsilon}\log(\frac{1}{\delta}))$$

#### PAC Learnability ← Compression Schemes

#### Theorem

If a Concept Class is PAC-Learnable then it has a Compression Scheme of size  $k^1$ .

#### Questions

- $\bullet$  Why are there different bounds based on VCdim?
- **2** Why is the Accuracy Parameter  $\epsilon$  turned into Confidence-like Parameter ?
- **3** Why/When is  $k \leq d$ ?
- Is there an *Information Theoric* approach for the same purpose?

Thank You! :)