

Concurrent Error Handling

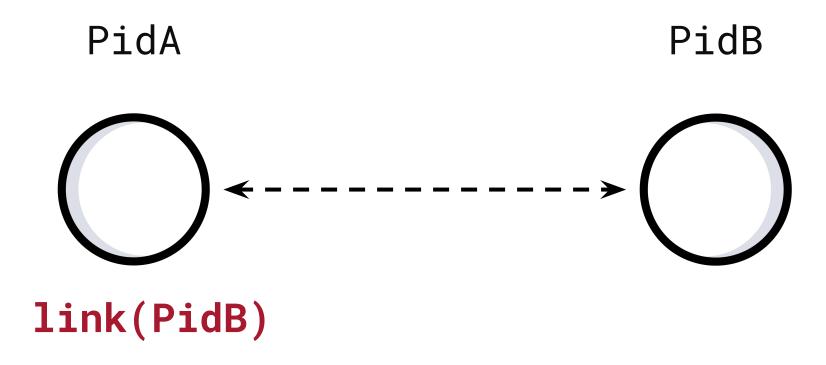


Process Error Handling

- Links
- **Exit Signals**
- Definitions
- Propagation Semantics
- Monitors
- Robust Systems
- Error Handling Example

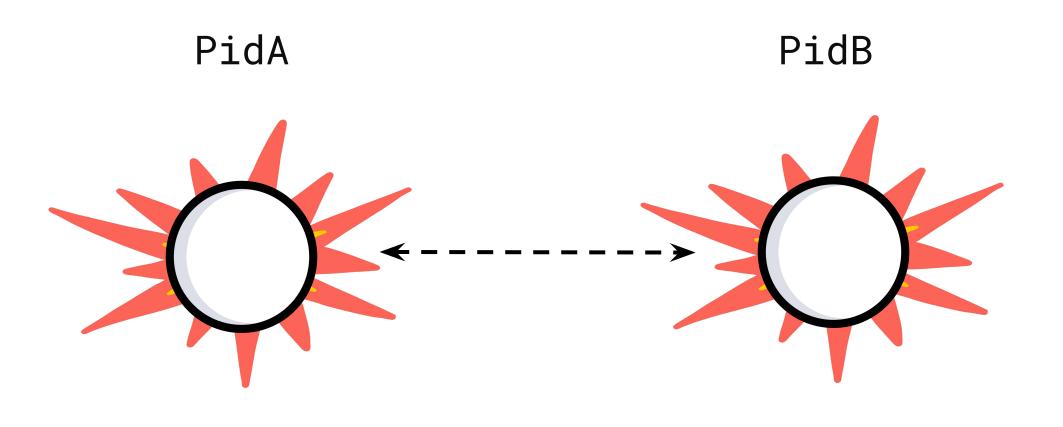


Links



- ► link/1 will create a bi-directional link between the process calling the BIF and the process PidB
- spawn_link/3 will yield the same result as calling spawn/3 followed by link/1, only that it will do it atomically





{'EXIT', PidA, Reason}

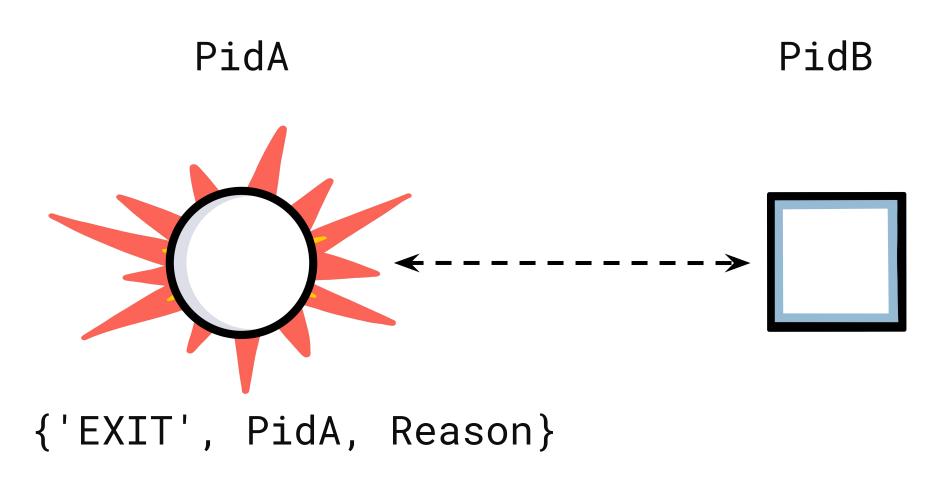
- ► Exit Signals are sent when processes terminate abnormally
- ► They are sent to all processes to which the failing process is currently linked to
- ► The process receiving the signal will exit, then propagate a new signal to the processes it is linked to



{'EXIT', PidA, Reason} PidA PidB {'EXIT', PidB, Reason} PidC

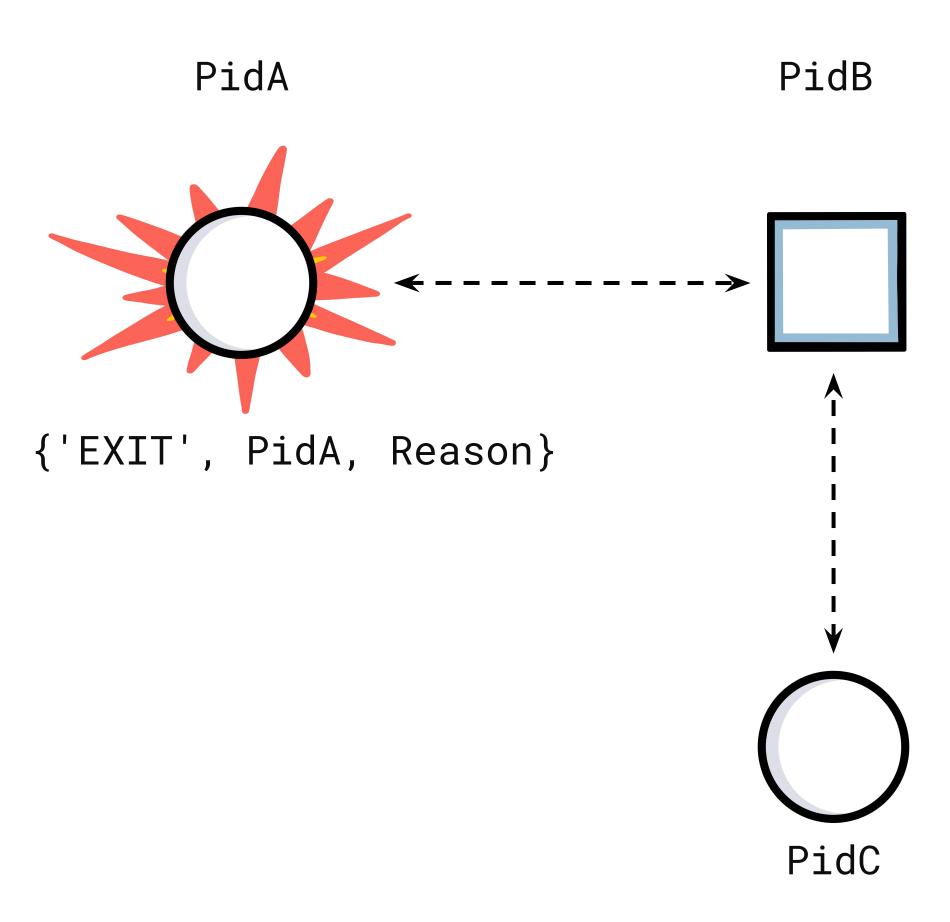
- ► When process PidA fails, the exit signals propagate to PidB
- ► From PidB, it propagates to PidC.





- Processes can trap exit signals by calling the BIF process_flag(trap_exit, true)
- Exit signals will be converted to messages of the format ('EXIT', Pid, Reason)
- ► They are saved in the process mailbox
- ► If an exit signal is trapped, it does not propagate further





- Process B marked with a double ring is trapping EXITs
- ► If an error occurs in A or C, then they will terminate.
- Process B will receive the {'EXIT', Pid, Reason} message
- ► The process that did not terminate will not be affected.



Definitions: terminology

Link

A bi-directional propagation path for exit signals set up between processes

Exit Signal

An asynchronous signal transmitted by a process upon exiting. It contains termination information

Error Trapping

The ability of a process to handle exit signals as if they were messages



Definitions: built-in functions

link(Pid)

Set a link between the calling process and Pid

unlink(Pid)

Removes a link to Pid

spawn_link(M, F, Args)

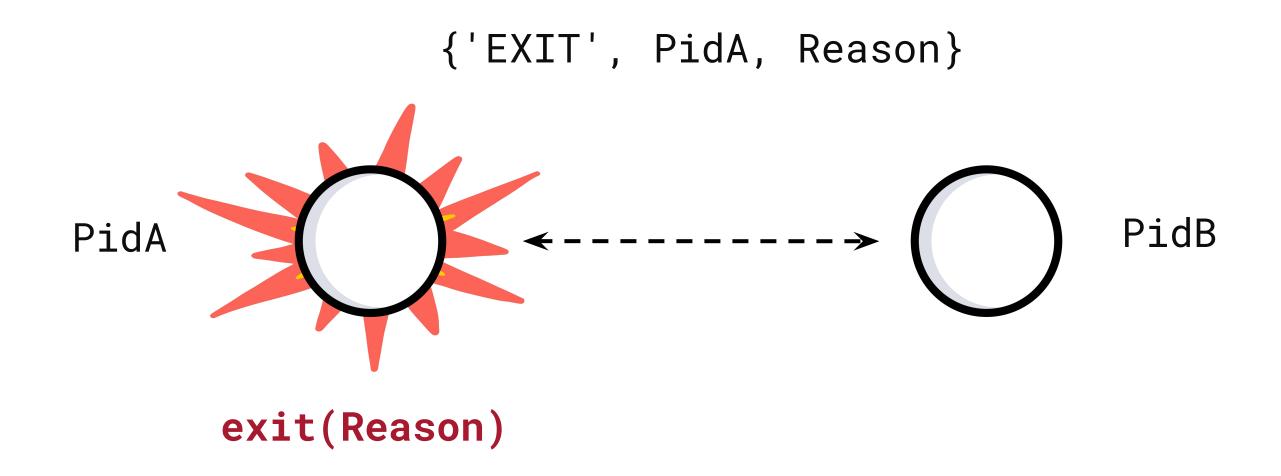
Atomically spawns and sets a link between the calling and the spawned processes.

process_flag(trap_exit, Bool)

Sets the current process to convert exit signals into exit messages



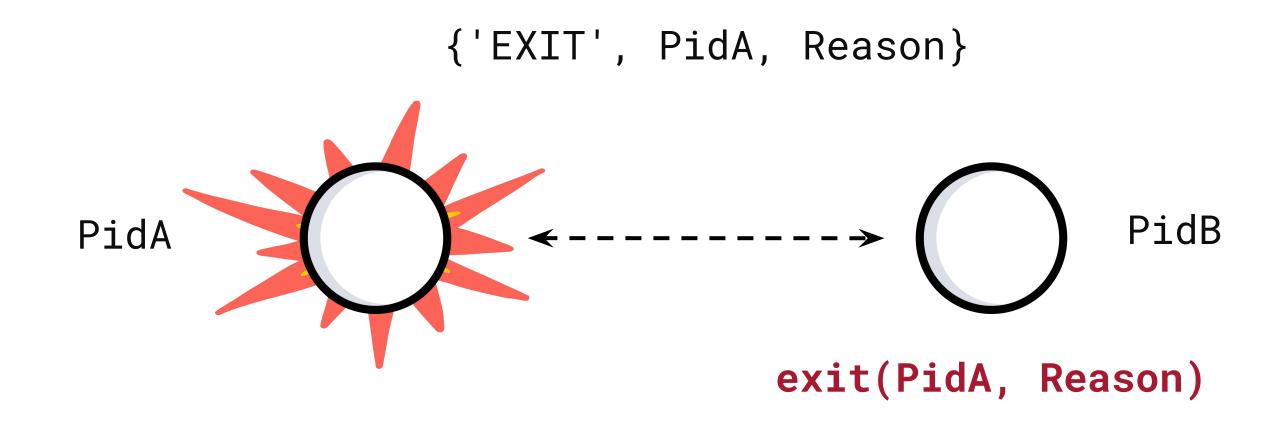
Definitions: built-in functions



- ▶ The BIF exit(Reason) terminates the process which calls it
- ▶ It generates an exit signal sent to linked processes
- ▶ The BIF exit/1 can be caught in a try ... catch.



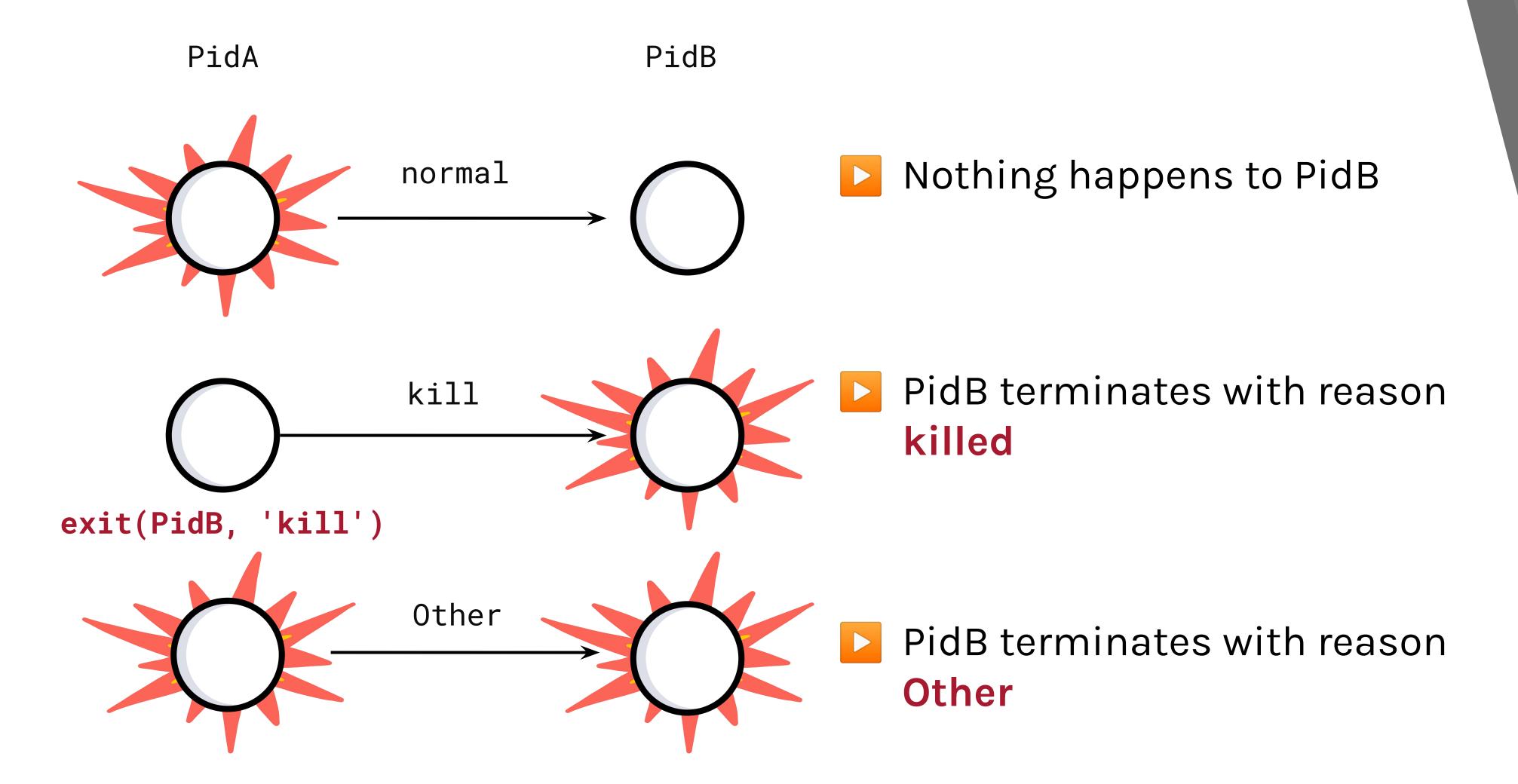
Definitions: built-in functions



- ▶ exit(Pid, Reason) sends an exit signal containing Reason to the process Pid
- ▶ If PidA is trapping exits, the signal is converted to an exit message

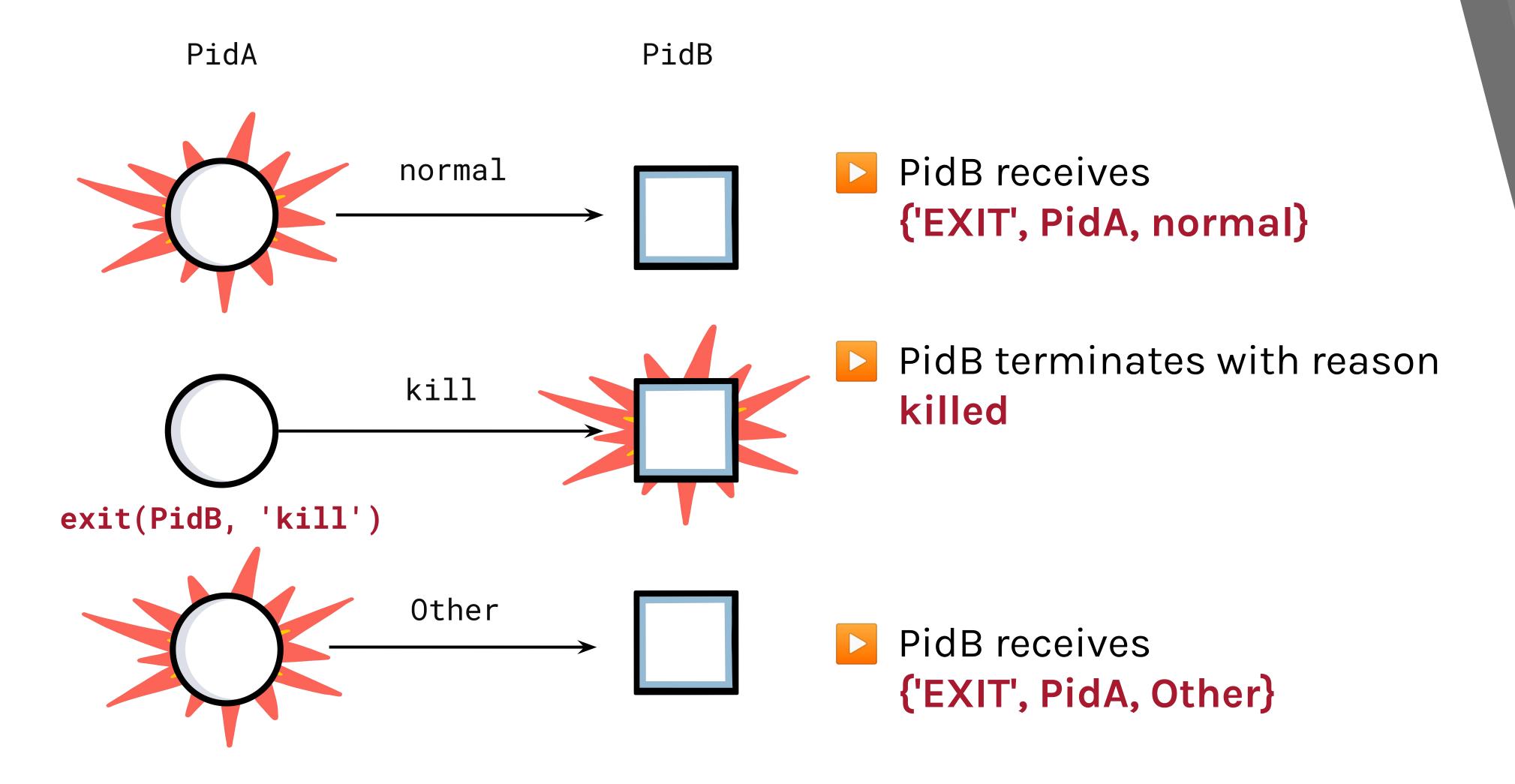


Propagation Semantics: no trapping





Propagation Semantics: trapping exits



Propagation Semantics

- When a process terminates, it sends an exit signal to the processes in its link set
- Exit signals can be normal or non-normal
- A process not trapping exits dies if it receives a non-normal one. Normal signals are ignored.
- A process which is trapping exit signals converts all incoming exit signals to conventional messages handled in a receive statement
- If the reason is kill, the process is terminated unconditionally with reason killed



Monitors



Ref = erlang:monitor(process, PidB)

- erlang:monitor/2 will the process calling the BIF monitor the process PidB. The monitoring is uni-directional and stackable.
- When process PidB dies then PidA will be sent the message {'DOWN', Ref, process, PidB, Reason}
- A monitor can be turned by calling erlang:demonitor(Ref) and erlang:demonitor(Ref, [flush])

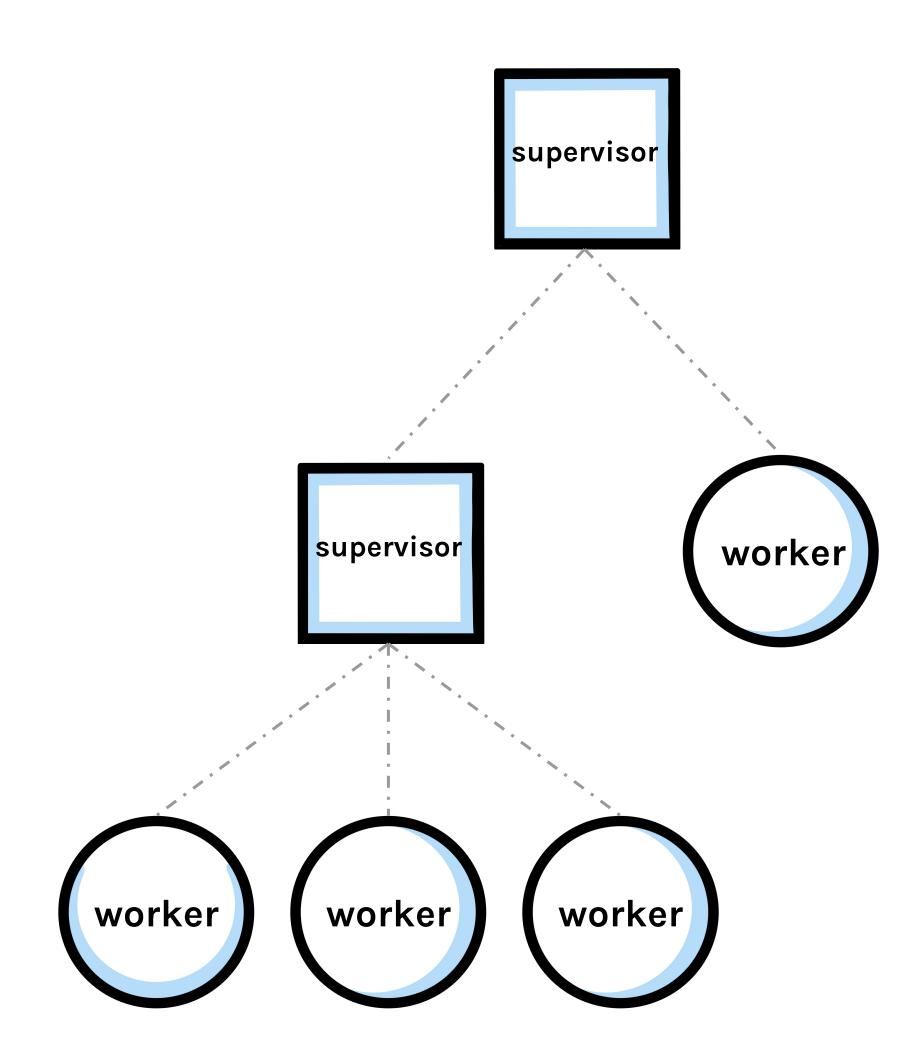


Robust Systems

- Building a system in layers can make it robust
 - Level N-1 traps and fixes errors occurring in level N
 - The leaves of the tree are workers
- In well designed systems, application programmers will not have to worry about error handling code
 - Error handling will be isolated by higher levels of the system, managed uniformly across processes
- Processes whose only task is to supervise children are called supervisors



Robust Systems



Robust systems can be designed by layering

A Robust Server

Remember the server example in the process design patterns section? The Server is unreliable!

- Let's rewrite the server making it reliable by monitoring the clients
- Let's rewrite the clients making them reliable by monitoring the server



A Robust Server

```
-module(frequency).
-export([start/0, stop/0, allocate/0, deallocate/1]).
-export([init/0]).
start() ->
    register(frequency, spawn(frequency, init, [])).
init() ->
    process_flag(trap_exit, true),
    Frequencies = {get_frequencies(), []},
    loop(Frequencies).
get_frequencies() -> [10,11,12,13,14, 15].
```

A Server Example

```
stop()
                 -> call(stop).
allocate() -> call(allocate).
deallocate(Freq) -> call({deallocate, Freq}).
%% We hide all message passing and the message protocol in
%% functional interfaces.
call(Message) ->
    Ref = erlang:monitor(process, frequency),
    frequency ! {request, self(), Message},
                                                   Keep race conditions
    receive
                                                              in mind!
       {reply, Reply} ->
          erlang:demonitor(Ref, [flush]),
          Reply;
       {'DOWN', Ref, process, Pid, Reason} ->
          exit(Reason)
    end.
reply(Pid, Message) -> Pid ! {reply, Message}.
```

A Server Example

```
%% The main server loop.
loop(Frequencies) ->
  receive
    {request, Pid, allocate} ->
       {NewFrequencies, Reply} = allocate(Frequencies, Pid),
       reply(Pid, Reply),
       loop(NewFrequencies);
    {request, Pid , {deallocate, Freq}} ->
       NewFrequencies = deallocate(Frequencies, Freq),
       reply(Pid, ok),
       loop(NewFrequencies);
    {request, Pid, stop} ->
       reply(Pid, ok);
    {'EXIT', Pid, Reason} ->
       NewFrequencies = exited(Frequencies, Pid),
       loop(NewFrequencies)
  end.
```

A Robust Server

```
%% The Internal Functions
%% Functions used to allocate and deallocate frequencies.
allocate({[], Allocated}, Pid) ->
    {{[], Allocated}, {error, no_frequencies}};
allocate({[Freq|Frequencies], Allocated}, Pid) ->
    link(Pid),
    {{Frequencies, [{Freq, Pid}|Allocated]}, {ok, Freq}}.
deallocate({Free, Allocated}, Freq) ->
    {Freq, Pid} = lists:keyfind(Freq, 1, Allocated),
    unlink(Pid),
    NewAllocated = lists:keydelete(Freq, 1, Allocated),
    {[Freq|Free], NewAllocated}.
```



A Robust Server

```
%% Help function used when a client crashes.
                                                         Keep race conditions
exited({Free, Allocated}, Pid) ->
                                                                     in mind!
  case lists:keyfind(Pid, 2, Allocated) of
    {Freq, Pid} ->
      NewAllocated = lists:keydelete(Freq, 1, Allocated),
      {[Freq|Free], NewAllocated};
    false ->
      {Free, Allocated}
  end.
```

A Server Example

```
Client
                                                                      Server
             {request, Pid, allocate}
              {reply, {error, no_frequencies} or {ok, Frequency}}
            {request, Pid, {deallocate, Frequency}}
               {'EXIT', Pid, Reason}
                                                         {reply, ok}
```

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