

Process Design Patterns

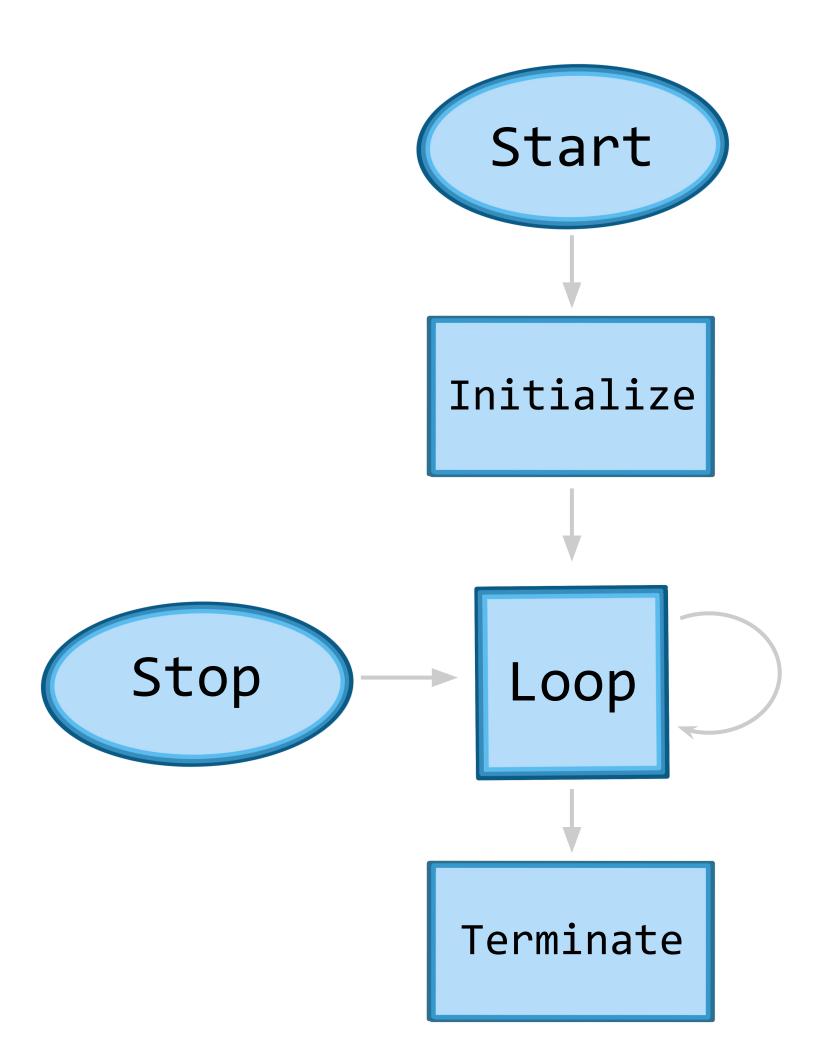
Process Design Patterns

- Client Server Models
- A Server Example
- Finite State Machines
- Event Managers
- Supervisors

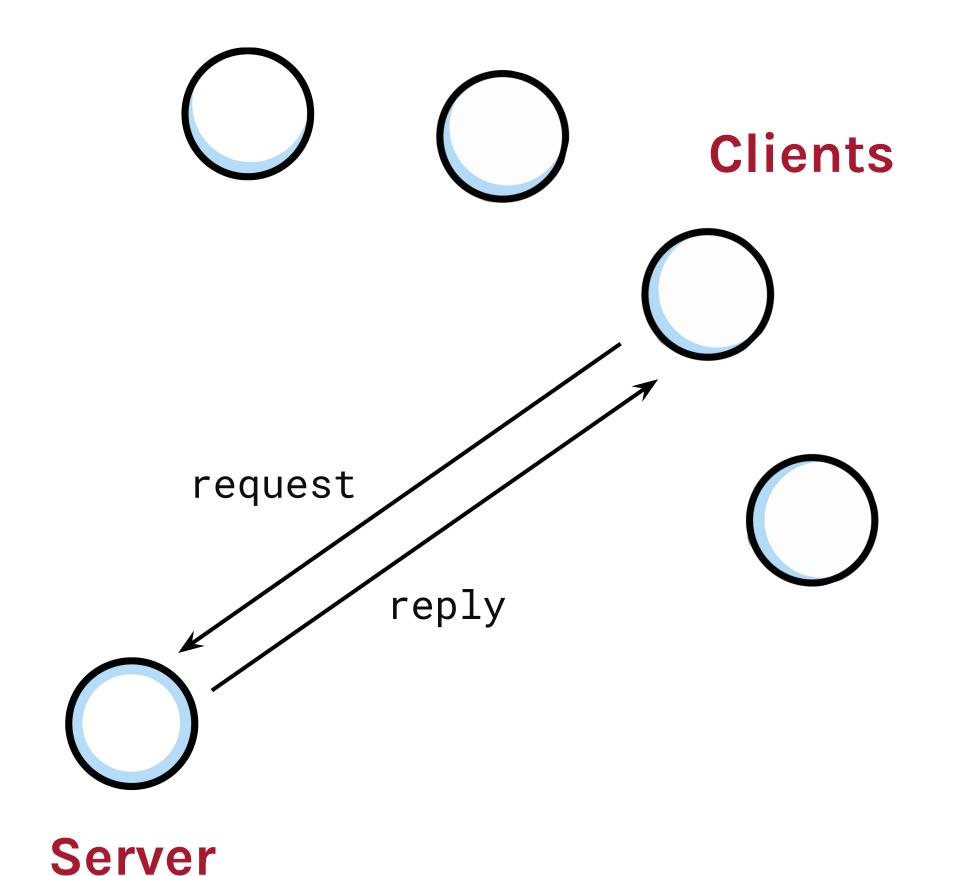


Client Server Models: process skeleton

```
start(Args) ->
   spawn(server, init, [Args])
init(Args) ->
  State = initialize_state(Args),
  loop(State).
loop(State) ->
   receive
      {handle, Msg} ->
         NewState = handle(Msg, State),
         loop(NewState);
      stop -> terminate(State)
   end.
terminate(State) -> clean_up(State).
```

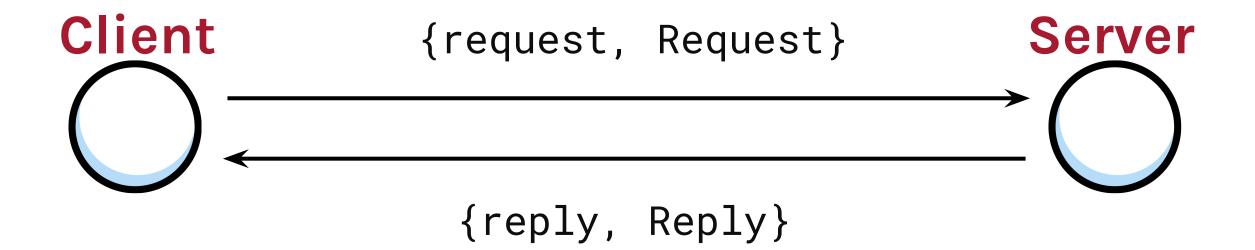


Client Server Models



- Processes can be used to implement client server solutions
- ➤ A server is usually responsible for providing a service or handling a resource
- ► Clients are the processes which use these resources

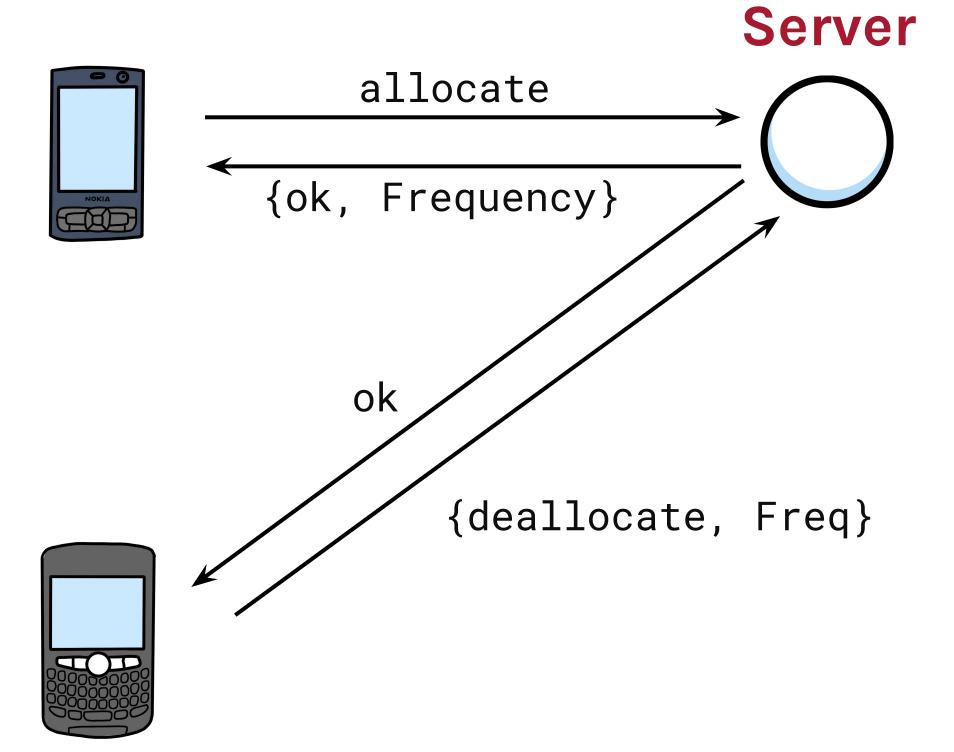
Client Server Models



- Clients make requests to the server through message passing
- ▶ Message passing is often hidden in functional interfaces
- ▶ If the client using the service needs a reply to the request, the call to the server has to be **synchronous**
- ▶ If the client does not need a reply, the call to the server can be asynchronous



Clients



► The following server is responsible for allocating and deallocating frequencies on behalf of mobile phones



```
Client
                                                                        Server
      {request, Pid, allocate}
                      {reply, {error, no_frequencies} or {ok, Frequency}}
      {request, Pid, {deallocate, Frequency}}
                                                             {reply, ok}
```



```
-module(frequency).
-export([start/0, stop/0, allocate/0, deallocate/1]).
-export([init/0]).
start() ->
    register(frequency, spawn(frequency, init, [])).
init() ->
    Frequencies = {get_frequencies(), []},
    loop(Frequencies).
get_frequencies() -> [10,11,12,13,14,15].
```

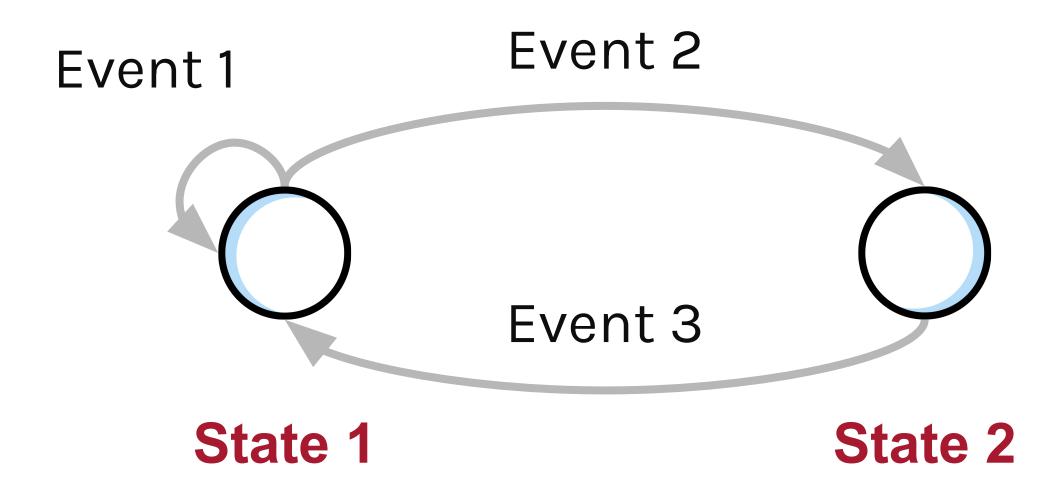
```
stop() -> call(stop).
allocate() -> call(allocate).
deallocate(Freq) -> call({deallocate, Freq}).
%% We hide all message passing and the message protocol in
%% functional interfaces.
call(Message) ->
    frequency ! {request, self(), Message},
    receive
      {reply, Reply} -> Reply
    end.
reply(Pid, Message) ->
   Pid ! {reply, Message}.
```

```
%% The main server loop.
loop(Frequencies) ->
  receive
    {request, Pid, allocate} ->
       {NewFrequencies, Reply} = allocate(Frequencies, Pid),
       reply(Pid, Reply),
       loop(NewFrequencies);
    {request, Pid , {deallocate, Freq}} ->
       NewFrequencies = deallocate(Frequencies, Freq),
       reply(Pid, ok),
       loop(NewFrequencies);
    {request, Pid, stop} ->
       reply(Pid, ok)
    end.
```

```
%% The Internal Functions
%% Functions used to allocate and deallocate frequencies.
allocate({[], Allocated}, Pid) ->
    {{[], Allocated}, {error, no_frequency}};
allocate({[Freq|Free], Allocated}, Pid) ->
    {{Free, [{Freq, Pid}|Allocated]}, {ok, Freq}}.
deallocate({Free, Allocated}, Freq) ->
    NewAllocated = lists:keydelete(Freq, 1, Allocated),
    {[Freq|Free], NewAllocated}.
```

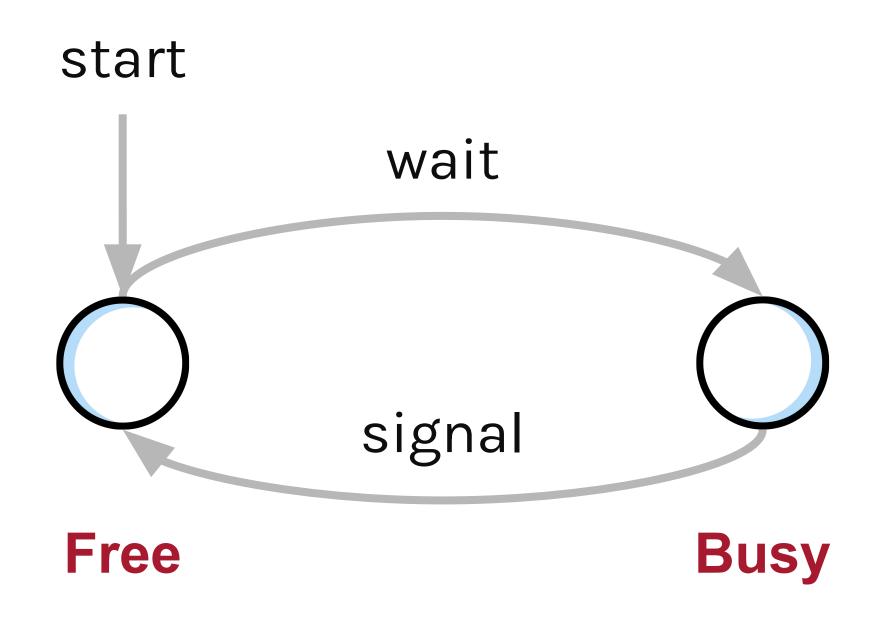


Finite State Machines



- Processes can be used to implement finite state machines
- ► Each state is represented as a tail recursive function
- ► Each event is represented as an incoming message
- Each state transition is achieved by calling the function denoting the new state

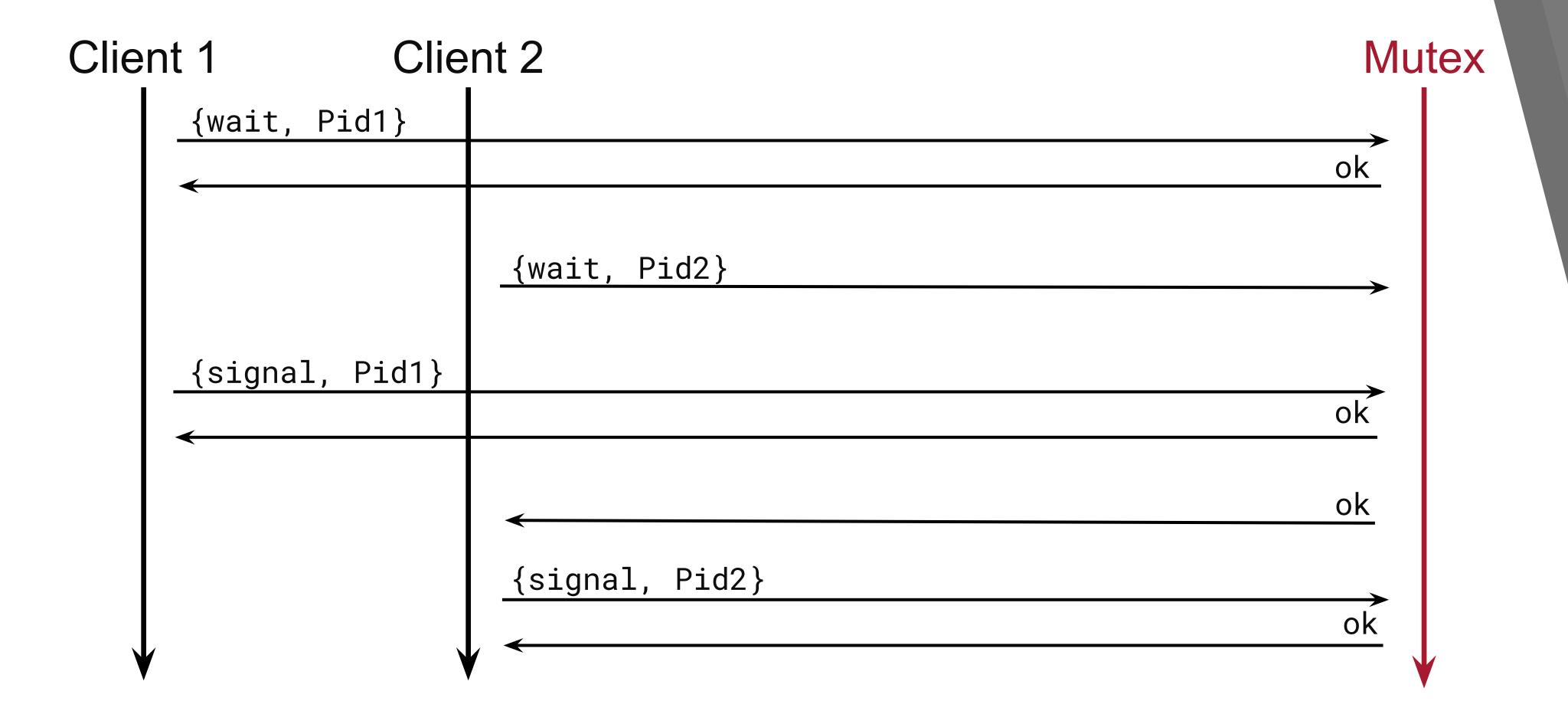




- A mutex is a program that allows multiple processes to share the same resource
- It has two states, **Free** and **Busy**
- It has two events, wait and signal
- When started it transitions to state **Free**



A Mutex Example





```
-module(mutex).
-export([start/0, stop/0]).
-export([wait/0, signal/0]).
-export([init/0]).
start() ->
    register(mutex, spawn(?MODULE, init, [])).
stop() ->
   mutex ! stop.
init() ->
    free().
```



```
wait() ->
    call(wait).
signal() ->
    call(signal).
%% We hide all message passing and the message protocol in
%% functional interfaces.
call(Message) ->
    mutex ! {Message, self()},
    receive
       {reply, Reply} -> Reply
    end.
reply(Pid, Message) ->
    Pid ! {reply, Message}.
```



```
%% The state functions.
free() ->
    receive
        {wait, Pid} ->
            reply(Pid, ok),
            busy(Pid)
    end.
busy(Pid) ->
    receive
        {signal, Pid} ->
            reply(Pid, ok),
            free()
    end.
```

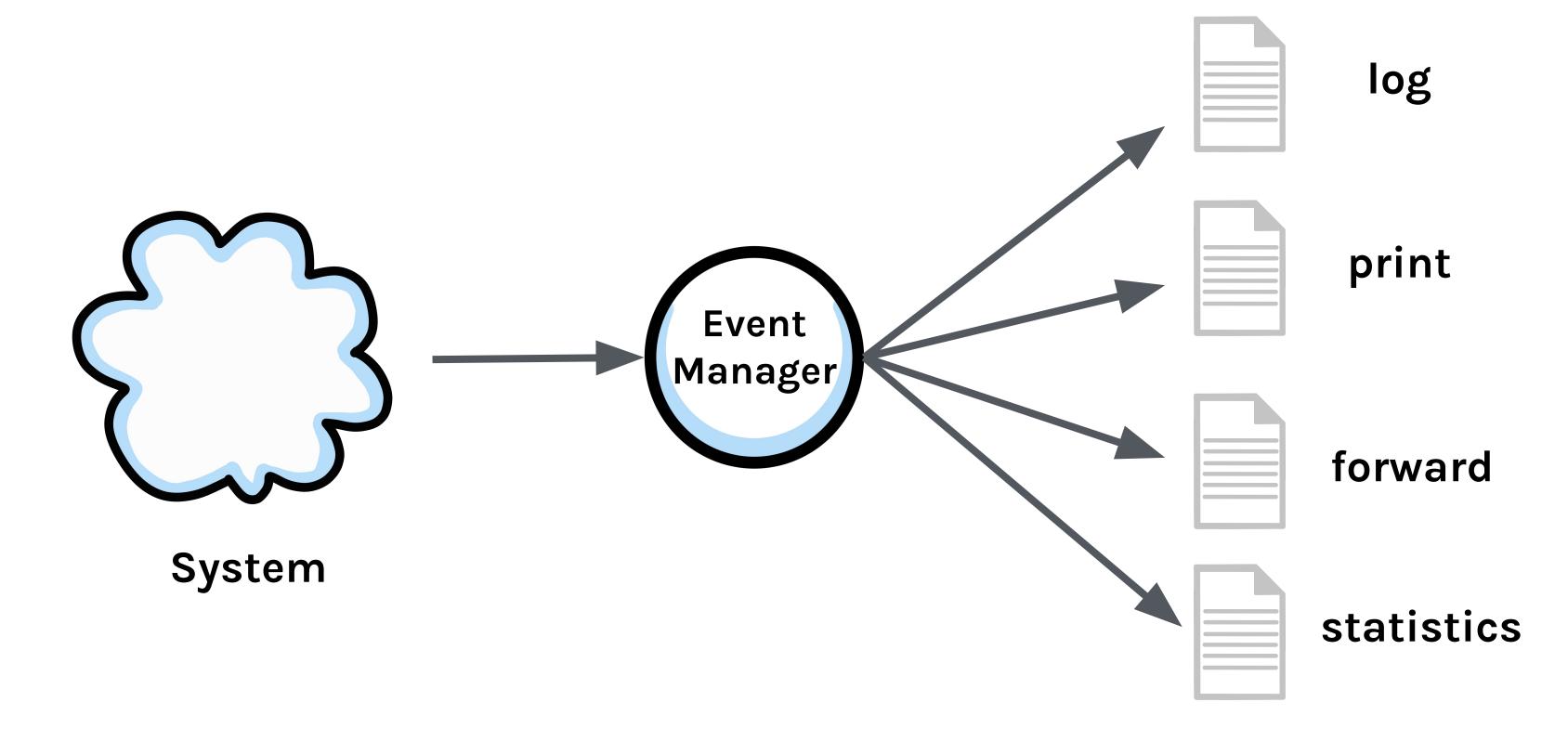


Event Managers and Handlers

- Processes can be used to implement event managers
- A manager will receive a specific type of event, e.g.
 - Alarms
 - State Changes
 - Commands
 - Errors
- When an event is received, one or more operations are applied on the event
- Some or all of the operations can be enabled and disabled during run time



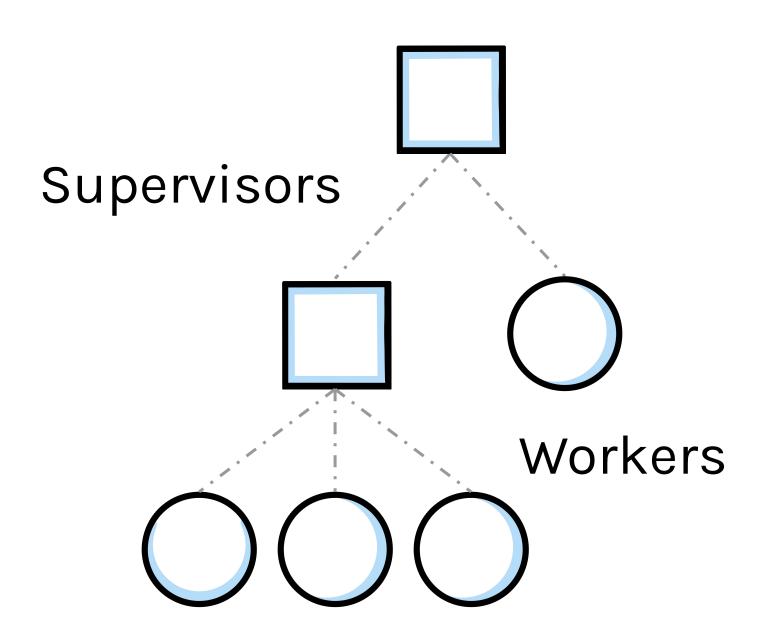
Event Managers: example



> Alarm managers are implemented as event managers with handlers

Supervisors

- Supervisors are processes whose only task is to start, monitor, and manage children.
- Child processes are either
 - Workers
 - Supervisors
- Supervisors will monitor their children
- Supervisors can restart the children when they terminate



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