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LSSC

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Abstract

- 1 Overview
- 2 Theoritical Description

2.1 Notations

In order to keep coherence with [3], the notations used are the same. A picture of n pixels is seen as a column vector in \mathbb{R}^n . The noisy picture is noted \mathbf{y} and the denoised one \mathbf{x} . The i-th pixel of \mathbf{x} is noted $\mathbf{x}[i]$ and the patch centered in $\mathbf{x}[i]$ and of size m is noted \mathbf{x}_i .

2.2 Learned Sparse Coding

The idea behind this method is to assume that the denoised picture is a signal that can be approximated by a sparse linear combinations of elements from a basis set. The basis set is called a dictionary $\mathbf{D} \in \mathbb{R}^{m \times k}$ and is composed of k elements. The denoised patches are then computed from \mathbf{D} with:

$$min_{\boldsymbol{\alpha} \in \mathbb{R}^k} ||\boldsymbol{\alpha}||_p \quad s.t. \quad ||\mathbf{y}_i - \mathbf{D}\boldsymbol{\alpha}||_2^2 \leqslant \epsilon$$
 (1)

 $\mathbf{D}\alpha$ is the estimate of the denoised patch and $||\alpha||_p$ is a regularization term that impose sparsity for α .

p is usually 0 or 1. Eq. 1 becomes NP-hard to solve when p=0 and a greedy algorithm such as Orthogonal Matching Pursuit [5] can give an approximation. With p=1, the problem is convex and solved efficiently with the LARS algorithm [6]. Experimental observations [7] have shown that the learning part is better with p=1 and the recomposition part with p=0.

 ϵ can be chosen according to the value of the estimated standard deviation of the noise.

2.3 Simultaneous Sparse Coding

3 Algorithm Description

3.1 Algorithm Overview

The first part consists in initializing a dictionary that will denoise roughly the picture.

Once the picture is denoised a first time, a clustering is made in order to regroup similar patches for further treatment.

Then, iteratively for each clusters, the dictionary is updated using simultaneous sparse coding and the cluster is denoised.

3.2 Dictionnary Initialization

The initial dictionnary is first learned offline on the 10 000 images of the PASCAL VOC'07 database using the online dictionary learning procedure of [2]. This procedure is then used on the noisy picture in order to improve the dictionary efficiency.

In fact, only a fixed number T of patches in the picture are used to update the dictionary. However they are chosen so that they are independently and identically distributed in the picture.

The algorithm corresponds then to the minimization of eq.1 on the T patches with the l_1 norm using the LARS [6] algorithm.

Input: number of iterations T, i.i.d. sampling of T patches of the noisy picture Y_T , initial dictionary $\mathbf{D}^0 \in \mathbb{R}^{m \times k}$, regularization parameter λ

Output: learned dictionary D

Initialization: $\mathbf{A}^0 \in \mathbb{R}^{k \times k} \leftarrow 0, \, \mathbf{B}^0 \in \mathbb{R}^{m \times k} \leftarrow 0$

for t = 1...T do

$$\mathbf{y}_t = \mathbf{Y}_T[\mathbf{t}]$$

Sparse coding: compute with LARS algorithm:

$$\alpha^t = \operatorname{argmin}_{\alpha \in \mathbb{R}^k} ||\alpha||_1 \text{ s.t. } ||\mathbf{y}_t - \mathbf{D}^{t-1}\alpha||_2^2 \le \lambda$$

$$\mathbf{A}^t \leftarrow \mathbf{A}^{t-1} + \boldsymbol{\alpha}^t \boldsymbol{\alpha}^{tT} \\ \mathbf{B}^t \leftarrow \mathbf{B}^{t-1} + \mathbf{y}_t^t \boldsymbol{\alpha}^{tT}$$

$$\mathbf{B}^t \leftarrow \mathbf{B}^{t-1} + \mathbf{v}_{\perp}^t \boldsymbol{\alpha}^{tT}$$

Update dictionary from \mathbf{D}^{t-1} to \mathbf{D}^t so that:

$$\mathbf{D}^{t} = \operatorname{argmin}_{\mathbf{D} \in \mathbb{R}^{m \times k}} \frac{1}{t} \sum_{i=1}^{t} (\frac{1}{2} |\mathbf{y}_{t}^{i} - \mathbf{D} \boldsymbol{\alpha}^{i}||_{2}^{2} + \lambda ||\boldsymbol{\alpha}^{i}||_{1})$$

end return \mathbf{D}^T

Algorithm 1: Online Dictionary Learning

Input :input dictionary $\mathbf{D} = [\mathbf{d}^1, \dots, \mathbf{d}^k] \in \mathbb{R}^{m \times k}$ $\mathbf{A} = [\mathbf{a}^1, \dots, \mathbf{a}^k] \in \mathbb{R}^{k \times k}, \, \mathbf{B} = [\mathbf{b}^1, \dots, \mathbf{b}^k] \in \mathbb{R}^{m \times k}$

Output: updated dictionary D

repeat

for
$$j = 1..k$$
 do
$$\begin{array}{c|c} & \text{update the } j^{th} \text{ column:} \\ & \text{if } \quad \boldsymbol{A}(j,j) = 0 \text{ then} \\ & & \\ & \text{end} \\ & \text{else} \\ & & \\ &$$

until convergence Marc: apparemment, dans le code on a une boucle sur params.updateIteration = 1. Est-ce qu'il ne vaudrait mieux pas calculer l'argmin à chaque boucle et s'arrêter lorsque c'est plus petit qu'une certaine valeur ?Yohann: je n'ai pas trop compris ce que tu voulais faire. Ca c'est le pseudo code prsent par Mairal, qui est donc une sorte de descente de gradient. Ce genre d'algo converge en thorie l'infini, en pratique on prend un grand nombre d'itration. Cependant, dans ce cas on fait une minimisation alterne d'un problem plus global (selon D puis selon alpha et on itre). Du coup, et surement cause de contraintes de temps, Mairal a fix le nombre d'itration 1 mais permet de le changer en paramtre dans son algo. Du coup j'ai voulu faire pareil que Mairal.;

return D

Algorithm 2: Dictionary Update updateDictionary Marc: Algo validé

```
Marc: TODO: choisir entre les indices ou les crochets pour les indices
Input: Input dictionary \mathbf{D} \in \mathbb{R}^{m \times k}, noisy patch \mathbf{y} \in \mathbb{R}^m, constraint \lambda \in \mathbb{R}
Output : code \alpha \in \mathbb{R}^k
—INITIALIZATION—
\mathbf{G} \in \mathbb{R}^{k \times k} \leftarrow \mathbf{D}^T \mathbf{D}
normPatch \in \mathbb{R}^+ \leftarrow ||\mathbf{y}||_2^2
oldsymbol{lpha} \in \mathbb{R}^k \leftarrow \mathbf{0}
\mathcal{A} \in [0; k]^k \leftarrow \mathbf{0}
Most\ correlated\ element
\hat{\mathbf{c}} \in \mathbb{R}^k \leftarrow \mathbf{D}^T \mathbf{y}
C \in \mathbb{R}^+ \leftarrow \max_{j=1..k}(|\hat{\mathbf{c}}_j|)
currentInd \leftarrow j s.t. \hat{\mathbf{c}}_j = C
newAtom \leftarrow True
TODO: explain the condition
if normPatch > \lambda then
 return 0
\mathbf{end}
for i = 1..k do
| LOOP, see below
end
return \alpha
                      Algorithm 3: LARS algorithm - Mairal Version computeLars
```

```
-NEW ATOM-
if newAtom then
              \mathcal{A}[i] \leftarrow \text{currentInd}
             G_A \in \mathbb{R}^{k \times i}, G_S \in \mathbb{R}^{i \times i}
             G_A[i^{th} \text{ column}] \leftarrow G[\text{currentInd}^{th} \text{ line}]
             G_S[i^{th} \text{ line}] \leftarrow G_A[\text{currentInd}^{th} \text{ line}]
              symmetrize G_S
             UPDATE G_S^{-1}
end
 —VARIABLES UPDATES—
\mathbf{u} \in \mathbb{R}^{i} \leftarrow G_{S}^{-1}(\operatorname{sgn}(\hat{\mathbf{c}}_{\mathcal{A}[j]}))_{j \in [1;i]}
ratio \in \mathbb{R}^{i} \leftarrow \left(-\frac{\alpha[\mathcal{A}[j]]}{u_{j}}\right)_{j \in [1;i]}
stepMAX \in \mathbb{R}^+_* \leftarrow min^+(ratio) (min<sup>+</sup> is the minimum between strictly positive values only)
criticalInd \in [1; k] \leftarrow \mathcal{A}[j] s.t. ratio[j] = \text{stepMAX}
C \leftarrow \hat{\mathbf{c}}[\mathcal{A}[1]] = \max_{i=1..k}(|\hat{\mathbf{c}}_i|) Marc: ça ne devrait pas être C = \max \hat{\mathbf{c}}[j]? Yohann: \mathbf{c} est
pareil (apres correction)
\gamma \in \mathbb{R}_*^+ \leftarrow \min^+ \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \text{ça fait 2 possibilités ou bien 4 avec les } \pm \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}_j = \frac{1}{2} \left(\frac{C \pm \hat{\mathbf{c}}_j}{1 \pm (G_A u)[j]}\right)_{j \text{ s.t. } \mathcal{A}[j]=0} \text{Marc} : \hat{\mathbf{c}}
?Yohann: 4 possibilites
—POLYNOMIAL RESOLUTION—
a \in \mathbb{R} \leftarrow \sum_{j \in [1;i]} \operatorname{sgn}(\hat{\mathbf{c}}[\mathcal{A}[j]]) u[j]b \in \mathbb{R} \leftarrow \sum_{j \in [1;i]} \hat{\mathbf{c}}[\mathcal{A}[j]] u[j]
c \in \mathbb{R} \leftarrow \text{normPatch} - \lambda
\Delta \in \mathbb{R} \leftarrow b^2 - ac
stepMAX2 \in \mathbb{R}^+ \leftarrow min(\frac{b-\sqrt{\Delta}}{2}, C)
 —FINAL STEP & BREAK—
\gamma \leftarrow \min(\gamma, \text{stepMAX}, \text{stepMAX2})
for j = 1..i \, do
  | \boldsymbol{\alpha}[\mathcal{A}[j]] \leftarrow \boldsymbol{\alpha}[\mathcal{A}[j]] + \gamma \mathbf{u}[j]
end
\hat{\mathbf{c}} \leftarrow \hat{\mathbf{c}} - \gamma G_A u
normPatch \leftarrow normPatch + a\gamma^2 - 2b\gamma
if |\gamma| < 1e^{-6} || \gamma = stepMAX2 || normPatch < 1e^{-6} || normPatch - \lambda < 1e^{-6} then
end
if \gamma = stepMAX then
             DOWNDATE G_S^{-1} w.r.t criticalInd
              \mathcal{A}[\text{criticalInd}] \leftarrow 0
              \alpha[criticalInd] \leftarrow 0
             newAtom \leftarrow False
             i \leftarrow i-2 Marc c'est possible que i devienne négatif ? Ça risque de poser problème non
              ?Yohann non car quand i=1 stepMAX n'est pas dfini (car il vaut min^+(\{0\}) = \emptyset) et pour
            i=2 ou plus on fait i=i-2 puis on rincrmente la fin de la boucle soit i=i-1
end
else
  \perp newAtom \leftarrow True
end
```

Algorithm 4: LARS algorithm - Mairal Version computeLars - LOOP

```
Input: Gram matrix G_S \in \mathbb{R}^{i \times i}, and its former inverse to update G_S^{-1} \in \mathbb{R}^{i-1 \times i-1}
    Output: updated G_S^{-1} \in \mathbb{R}^{i \times i}
    if i = 1 then
     return \frac{1}{G_S}
    end
    u \leftarrow G_S^{-1}G_S[i^{th} \text{ line}]

\sigma \leftarrow \frac{1}{G_s(i,i) - u.G_S[i^{th} \text{ line}]} 

G_s^{-1}(i,i) \leftarrow \sigma 

G_s^{-1}[i^{th} \text{ line}] \leftarrow -\sigma u

return G_s^{-1} \leftarrow G_s^{-1} + \sigma u u^T

Algorithm 5: Update invert algorithm updateGram Marc: Algo correspondent, mais à valider.
    Input: pseudo-Gram matrix G_A \in \mathbb{R}^{k \times i} Gram matrix G_S \in \mathbb{R}^{i \times i}, and its inverse G_S^{-1} \in \mathbb{R}^{i \times i},
    criticalInd \in [1; k], current iteration i
    Output: downdated matrices G_A \in \mathbb{R}^{k \times i-1}, G_S, G_S^{-1} \in \mathbb{R}^{i-1 \times i-1}
    \sigma \leftarrow \frac{1}{G_S^{-1}(\text{criticalInd},\text{criticalInd})}
    u \leftarrow G_S^{-1}[\text{criticalInd}^{th} \text{ line}] \text{ without its criticalInd}^{th} \text{ coefficient}
    for j = criticalInd:i-1 do
          G_A[j^{th} \text{ column}] \leftarrow G_A[(j+1)^{th} \text{ column}]
          for k=1:criticalInd-1 do
               G_S(j,k) \leftarrow G_S(j+1,k)

G_S^{-1}(j,k) \leftarrow G_S^{-1}(j+1,k)
          for k = criticalInd:i do
              G_S(j,k) \leftarrow G_S(j+1,k+1)

G_S^{-1}(j,k) \leftarrow G_S^{-1}(j+1,k+1)
    end
    G_s^{-1} \leftarrow G_s^{-1} - \sigma u u^T
  Algorithm 6: Downdate invert algorithm downdateGram Marc: Algo ne correspondant pas, j'ai
```

Algorithm 6: Downdate invert algorithm downdateGram Marc: Algo ne correspondant pas, j'ai créé une nouvelle fonction downdateGramBis dont je pense est une meilleure version. Il faudra la tester et la valider.

Glossary

Image Credits



POL (c) IPOL (there's no need to credit this image, here is used as an example.)

4 References

References

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