### **PRQL**

#### **Pipelined Relational Query Language**

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```
from animals
filter height_cm > 20
take 10
join keepers (==keeper_id)
filter keepers.first_name == "John"
sort animals.age
select {
    species = animals.species,
    keeper = keepers.name,
    f"I'm {species}, {animals.age} years old",
```

## Is it any good?

But there are transition costs!

#### **Overview**

Flaws of SQL

Language for relations

**Compiling queries** 

PRQL, the project

A deep dive into

# Flaws of SQL

### Origins of the relational model

1970, Edgar F. Codd: abstraction over data storage

→ Tuple relational calculus

1974, Donald D. Chamberlin & Raymond F. Boyce: SEQUEL

→ Not a "proper" programming language



### Not really composable

SELECT album\_id, COUNT(\*) AS track\_count
FROM tracks GROUP BY album\_id

## Not really composable

```
SELECT i.album_id, i.track_count, a.artist_id
FROM (
    SELECT album_id, COUNT(*) AS track_count
    FROM tracks GROUP BY album_id
) AS i
JOIN albums a USING (album_id)
```

```
SELECT
SUM(total)
FROM
invoices
```

```
SELECT
  total / SUM(total) OVER () AS normalized_total
FROM
  invoices
```

SELECT DISTINCT name FROM invoices

```
SELECT EVALUATE TYPE-EMPLOYEE
 WHEN "F"
    MOVE "FULL TIME" TO EMP-TYPE-PR
 WHEN "P"
    MOVE "PART TIME" TO EMP-TYPE-PR
 WHEN "C"
    MOVE "CONSULTANT" TO EMP-TYPE-PR
 WHEN OTHER
    MOVE "INVALID" TO EMP-TYPE-PR
```

Too much syntax

... but also ...

Not enough syntax

#### Name resolution

SELECT title AS title\_alias FROM albums

#### Name resolution

```
SELECT title AS title_alias
FROM albums
WHERE title_alias LIKE 'Do I Wanna %'
GROUP BY title_alias
ORDER BY title_alias
```

#### Name resolution

More rules:

- ORDER BY positionals
- Correlated subqueries
- LATERAL

#### **Relations vs scalars**

SELECT \* FROM table

SELECT count(\*) FROM table

#### **Relations vs scalars**

```
SELECT emp_id FROM emp WHERE role = 'manager'
```

#### **Relations vs scalars**

```
SELECT *
FROM emp
WHERE emp_id = (
    SELECT emp_id FROM emp WHERE role = 'manager'
)
```

#### Differences in:

- syntax (TOP vs LIMIT)
- available functions
- available data types

A class of languages

There is a standard

Slight deviations

#### Different:

- priorities
- backward compatibility guarantees
- implementation limitations

No clear & robust specification

Compilers could:

- adapt query to target database
- produce error early

Design of a new

# Language for relations

from animals

```
animal_id = 1,
 species = "Mus musculus",
 height_cm = 5,
 age = 1,
 keeper_id = 2
... and 100 others ...
```



from animals
filter height\_cm > 10

```
animal_id = 1,
 species = "Suricata suricatta",
 height_cm = 15,
 age = 3,
 keeper_id = 4
... and 14 others ...
```



```
from animals
filter height_cm > 10
take 10
```

```
from animals
filter height_cm > 10
take 10
join keepers (
    animals.keeper_id == keepers.keeper_id
```

```
from animals
filter height_cm > 10
take 10
join k = keepers (
    animals.keeper_id == k.keeper_id
```

```
from animals
filter height_cm > 10
take 10
join k = keepers (==keeper_id)
```

```
from animals
filter height_cm > 10
take 10
join k = keepers (==keeper_id)
filter k.first_name == "John"
```

```
from animals
filter height_cm > 10
take 10
join k = keepers (==keeper_id)
filter k.first_name == "John"
sort animals.age
```

```
take 10
join k = keepers (==keeper_id)
filter k.first name == "John"
sort animals.age
select {
    animals.age,
    species = animals.species,
    keeper = f"{k.first_name} {k.last_name}",
```

```
select {
    animals.age,
    species = animals.species,
    keeper = f"{k.first_name} {k.last_name}",
select {
    f"I'm {species}, {animals.age} years old",
    keeper
```

```
"I'm Suricata suricatta, 3 years old",
keeper = "John Doe"
"I'm Capra ibex, 12 years old",
keeper = "John Doe"
```





```
- Top to bottom
```

- Easy exploration
- Lazy evaluation
- Extract a variable
- Extract a function

```
from animals
filter height_cm > 10
take 10
join k = keepers (==keeper_id)
filter k.first name == "John"
sort animals.age
select {
    animals.age,
    species = animals.species,
    keeper = f"{k.first_name} {k.last_name}",
select {
    f"I'm {species}, {animals.age} years old"
    keeper
```

#### **Extract a variable**

```
from animals
filter height_cm > 20
take 10
```

#### **Extract a variable**

```
let big_animals = (
    from animals
    filter height_cm > 20
)

from big_animals
take 10
```

#### **Extract a function**

```
from animals
sort {-height_cm}
take 5
```

#### **Extract a function**

```
let take_biggest = n rel -> (
    rel
    sort {-height_cm}
    take n
from animals
take_biggest 5
```

### **Tuple relational calculus**

Relation  $\sim$  a set of tuples

$$\pi_{\text{track\_id,name,title}}(R)$$
  $\sigma_{\text{track\_id=5}}(R)$   $R * S$ 

### Basic data types

bool, int, float, str

### **Tuples**

```
{my\_int = 5, 4.2, my\_bool = true}
```

- named fields
- ▶ different types
- ► static number of fields

# Arrays

```
[1, 2, 10, -3]
```

- unnamed items
- ▶ items have the same type
- dynamic number of items

Relation := an array of tuples

```
{my_int = 5, 4.2, my_bool = true},
{my_int = -2, 6.1, my_bool = false},
{my_int = 12, 3.0, my_bool = false},
```

#### **Declarations**

let 
$$a = 5$$
  
let  $b = a + 1$ 

#### **Functions**

let add\_one = 
$$x \rightarrow x + 1$$
  
let add =  $x y \rightarrow x + y$ 

#### **Functions**

```
let five = (add_one 4)
let six = (add 4 2)
```

#### **Functions**

```
let seven = (5 | add_one | add_one)
let seven = (
    add one
    add_one
```

Transform := a function on relations

```
let animals = ...
```

```
let main = (filter (height_cm > 20) animals)
```

```
let animals = ...
let main = (animals | filter (height_cm > 20))
```

```
let animals = ...

let main = (
    animals
    filter (height_cm > 10)
)
```

```
let animals = ...
animals
filter (height_cm > 10)
```

```
from animals
filter (height_cm > 10)
```

```
from animals
filter height_cm > 10
```

std.from	std.take
std.select	std.join
std.derive	std.group
std.filter	std.window
std.aggregate	std.append
std.sort	std.loop

## **Orthogonal**

```
from expenses
filter dept == "Sales"
aggregate {total = sum cost}
filter total > 100.00
```

WHERE  $\mapsto$  filter

 $\texttt{HAVING} \mapsto \texttt{filter}$ 

# **Orthogonal**

#### Transform invariants:

- filter will not change columns
- derive & select will not change number of rows
- aggregate will produce exactly one row

```
from expenses
aggregate {total = sum cost}

[
     {total = 431.22},
]
```

```
from expenses
group dept (
   aggregate {total = sum cost}
 \{dept = "Sales", total = 331.00\},\
 {dept = "Accounting", total = 100.22},
```

```
from expenses
group dept (
   take 1
 \{dept = "Sales", id = 33, cost = 5.30\},\
 {dept = "Accounting", id = 45, cost = 12.22},
```

```
from expenses
group dept (
    sort {-cost}
   take 1
 \{dept = "Sales", id = 3, cost = 33.30\},\
 {dept = "Accounting", id = 16, cost = 12.22},
```

```
from expenses
group expenses.* (
    take 1
)
```

```
from expenses
group expenses.* (
    take 1
SELECT DISTINCT *
FROM expenses
```

#### **Nulls**

```
# PRQL
null == null # true
my_col == null
-- SQL
my_col IS NULL
```

### **Ergonomics**

```
from employees
derive {
 age = 02023-01-31 - birth_date,
 full name = f"{first name} {last name}",
 manager = reports_to ?? "No one",
 salary = 1_{000_{00}}
# is_fired = "No",
```

# Challenges of

# **Compiling queries**

# SQL as a compilation target

How is this language executed?

- X database interface
- ✓ a query language

#### The task of a database interface

Imagine a database without a query language.

SELECT \* FROM albums

... and then transform in client code.

 $\rightarrow$  super slow

### The task of a database interface

#### Extreme example:

```
SELECT COUNT(*)
FROM albums
WHERE title LIKE 'The %'
```

# The task of a database interface

#### Processing should be close to data

- minimal data transfer
- parallelism
- vectorization

# The task of a database interface

Databases are:

- execution platforms
- compilation targets

Analogous to amd64, JVM

# **Leaky abstractions**

Database interface should be transparent

Currently, this is not the case:

- invalid SQL
- sub-optimal SQL
- runtime errors

# PRQL, the project

- an opensource effort

# The compiler and its IRs

prqlc: compiler from PRQL to SQL

targets: sql.postgres, sql.sqlite, sql.duckdb, sql.mysql, sql.clickhouse

bindings for C, Python, JS, Java, .NET, PHP

# The compiler and its IRs

Don't connect, infer

#### Fail early

# **Architecture**

 $PRQL \rightarrow PL \rightarrow RQ \rightarrow SQL$ 

### Licence

Apache License 2.0

Open community

Will never monetize

# **Check it out: playground**

```
introduction.prgl
                                                                                                       output.pl.vaml
                                                                    Save
                                                                           output.sal
                                                                                        output.arrow
                                                           Rename
PRQL Playground
                                 from invoices
                                                                             WITH table 1 AS (
                                 filter invoice date >= @1970-01-16
EXTERNAL LINKS
                                 derive [
PROL Website >
                                                                                 customer id.
                                   transaction fees = 0.8,
                                                                                 total - 0.8 AS expr 0,
Book ≥
                                   income = total - transaction fees
                                                                                 total
                                                                               FROM
EXAMPLES
                                                                                 invoices
introduction.pral
                                 group customer id (
                                                                               WHERE
                                   aggregate [
let-table-0.prgl
                                                                                 invoice date >= DATE '1970-01-16'
                                     average total,
artists-0.prgl
                                     sum income = sum income.
                                    ct = count.
CHINOOK
                                                                               customer id.
albums.prgl
                                                                               AVG(total).
artists.prol
                                                                               SUM( expr 0) AS sum income,
customers.pral
                                 filter sum income > 1
                                                                               COUNT(*) AS ct
                                 sort [-sum income]
                                                                             FROM
employees.pral
                                                                               table 1
                                 take 10
genres.prql
                                                                             GROUP BY
invoice items.pral
                                                                               customer id
invoices.pral
                                                                             HAVTNG
media_types.prql
                                                                               SUM( expr \theta) > 1
                                                                             ORDER BY
playlists.prgl
                                                                               sum income DESC
playlist track.prgl
tracks.prgl
```

### **Check it out: VSCode extension**

```
_a.prgl - prgl-compiler - Visual Studio Code
          select [album id, name, unit price]
          sort I-unit price, namel
                                                                    COUNT(*) AS expr 1.
          group album id (
                                                                    album id
              track count - count.
              album price = sum unit price
              track count = sum track count.
                                                                    albums artist id
                                                                    table 1
                                                                    JOIN albums ON table 1.album id = albums.album id
      18 select [artists.name, artist price, track count]
                                                                    albums.artist id
      28 derive ava track price = artist price / track count
                                                                  artists name
In 21, Col 1 Spaces: 4 UTF-8 LF PROL @ Go Live @ Spell @ Prettier C
```

# **Check it out: prql-query - pq**

```
chinook$ pg --from tracks.csv 'select [track_id, name, bytes] | take 10'
 track id I name
                                                     bytes
            For Those About To Rock (We Salute You)
            Balls to the Wall
            Fast As a Shark
            Restless and Wild
            Princess of the Dawn
            Put The Finger On You
            Lets Get It Up
            Inject The Venom
             Snowballed
             Evil Walks
chinook$
```

### **Check it out**

pip install pyprql
install.packages("prqlr")
npm install prql
cargo add prql-compiler

https://prql-lang.org

https://github.com/PRQL/prql

https://discord.gg/TfyM755m

### Please vote and leave feedback!



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Any questions?



