Saman Rastgar TA-1 Due: Monday 09/24

Exercises 1.8

(Q1.1) Language Used=Java

- a) A lexical error will occur in java, if you to assign a char to an integer.
- b) A syntax error will occur in java, if you forget to include a semicolon at an end of a statement in java.
- c) A static semantic error will occur in java, if you access a method that has not been declared in the class.
- d) A dynamic semantic error will occur in java, if you access an element of an array that does not exist. Example: Index out of bound .
- e) An error that the compiler can neither catch nor easily generate code to catch would be if we were to use a method name as a variable.

(Q1.8) Make Analysis

The make utility searches for the target's dependency and create a chain of dependency and it compiles the chain of file in the order they have been added. But this tool does not always recompile the previously compiled code. Dependency management uses timestamp to keep track of the compiled files and starts from the part that needed to be compiled.

- ~>As mentioned make looks for dependent targets but you can always include comment on the dependent file which does not affect the compile procedure(The dependency management described is not accurate).
- ~>compiling procedure for make file fails when it can't find the file description for dependent file.

Exercises 2.6

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(Q2.1) Regular Expressions
```

- a) Strings in C are captured by: string ? ""(!{\, ", \n} | \!{\n})* "
- b) Comments in Pascal are captured by: Pascal comment ? (* (!(*) | (*!())))* *+)
- c) Numeric constants in C:

```
 C\_constant \rightarrow int\_const \mid fp\_const \\ int\_const \rightarrow (oct \mid dec \mid hex) int\_suffix \\ oct\_int \rightarrow 0 \ oct\_digit^* \\ dec\_int \rightarrow nonzero\_digitdec\_digit^* \\ hex\_int \rightarrow (0x \mid 0X) \ hex\_digithex\_digit^* \\ oct\_digit \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \\ nonzero\_digit \rightarrow 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9 \\ dec\_digit \neg \rightarrow 0 \mid nonzero\_digit \ hex\_digit \rightarrow dec\_digit \ A \mid B \mid C \mid D \mid E \mid F \mid a \mid b \mid c \mid d \mid e \mid f \\ dec\_float \rightarrow dec\_digit^* \mid .dec\_digit^* \mid E \mid e \ hex\_float \rightarrow e \mid \epsilon \ exponent \mid \epsilon \ type \rightarrow long \mid unsigned \\ long \mid longlong \mid unsigned \ longlong \ exponent \rightarrow + \mid - \mid \epsilon \ unsigned \rightarrow U \mid u \ float \rightarrow F \mid f \ long \rightarrow L \mid I \ longlong \rightarrow LL \mid II
```

d) Floating point constant in Ada:

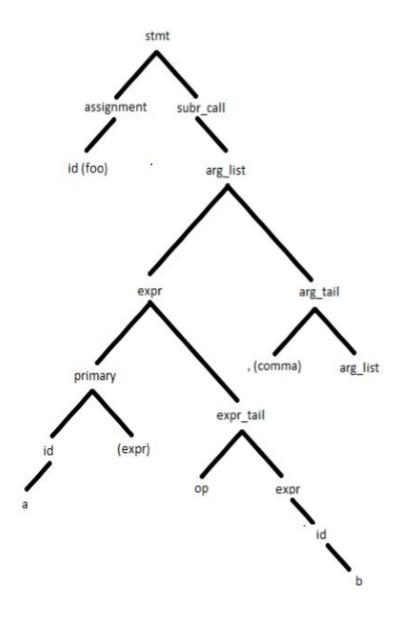
e) Showing inexact constant in scheme:

```
Digit + # * (.# * | \epsilon ) | digit* .digit + # *
```

f) nonzero digit → 1|2|3|4|5|6|7|8|9
 digit → 0 | nonzero digit
 group_type → , digit
 Number → \$ * * (0 | nonzero digit (ϵ | digit | digit digit) group_type *) (ϵ | .digit digit

(Q2.13) Grammer interpretation

(a) Construct parse tree



(b) Give a canonical (rightmost) derivation of this same string.

```
stmt \rightarrow assignment, subr\_call
subr_call → arg_list
arg_list → expr,args_tail
args_tail →,,arg_list
expr → primary,expr_tail
expr_tail \rightarrow op, expr
\mathsf{expr} \to \mathsf{id}
expr →primary, expr_tail
primary →id,expr
assignment →id
(Q2.17)
program → stmt_list $$
stmt_list → stmt_list stmt
stmt\_list \rightarrow stmt
stmt \rightarrow id:= expr
stmt \rightarrow read\ it
stmt \rightarrow write expr
expr \rightarrow term
expr \rightarrow expr \ add_op \ term
term \rightarrow factor
term → term mult_op factor
factor \rightarrow ( expr )
factor \rightarrow id
factor \rightarrow number
add_op \rightarrow + | -
mult_op \rightarrow * | /
stmt \rightarrow if condition then stmt\_list fi
→ while condition do stmt_list od
condition \rightarrow expr relation expr
relation \rightarrow <
\rightarrow >
→ <=
→ >=→ =
→ !=
```