| MPS 2200 Assignment 1 | |
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| me: | |
| his assignment, you will learn more about asymptotic notation, parallelism, ctional languages, and algorithmic cost models. As in the recitation, some our answer will go here and some will go in main.py. You are welcome to this assignment-01.md file directly, or print and fill in by hand. If you do latter, please scan to a file assignment-01.pdf and push to your github esitory. | |
| . (2 pts ea) Asymptotic notation | |
| 1a. Is $2^{n+1} \in O(2^n)$? Why or why not? . Yes $2^{n+1} = 2 \cdot 2^n, 2 \mid 3 \text{ a constant}$ | |
| 1b. Is $2^{2^n} \in O(2^n)$? Why or why not? No , Z^n grows exponentially faster $Than Z^n$ | |
| | |
| 1c. Is $n^{1.01} \in O(\log^2 n)$? $\mathcal{N}o$, as \mathcal{N} grows larger $\mathcal{N}o$ grows Much faster than $\log^2 \mathcal{N}o$ | |
| Id. Is $n^{1.01} \in \Omega(\log^2 n)$? I Yes, no matter what constant we put in front of $N^{1.01}$, it is still much larger | |
| The Is $\sqrt{n} \in O((\log n)^3)$? If No, logarithmic functions still grow S is slower than S art functions as N get large, even when cubed If Is $\sqrt{n} \in \Omega((\log n)^3)$? Yes, \sqrt{n} is still greater than $(\log n)^3$ if $\log n$ and $\log n$ is still greater than $\log n$ if $\log n$ is $\log n$ in $\log n$ and $\log n$ is $\log n$ in | for t |
| Id. Is $n^{1.01} \in \Omega(\log^2 n)$? Yes, no matter what constant we put in front of $N^{1.01}$, it is still much larger 1e. Is $\sqrt{n} \in O((\log n)^3)$? No, logarithmie functions still grow slower than sqrt functions as N ge (arge, even when cubed) 1f. Is $\sqrt{n} \in \Omega((\log n)^3)$? Yes, $\sqrt{n} \in \Omega((\log n)^3)$? | fo † |

2. SPARC to Python

Consider the following SPARC code of the Fibonacci sequence, which is the series of numbers where each number is the sum of the two preceding numbers. For example, 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610 ...

```
\begin{array}{l} foo\ x=\\ \text{ if }\ x\leq 1 \text{ then }\\ x\\ \text{ else }\\ \text{ let }(ra,rb)=(foo\ (x-1))\ ,\ (foo\ (x-2))\ \text{in }\\ ra+rb\\ \text{ end } \end{array}
```

• 2a. (6 pts) Translate this to Python code – fill in the def foo method in main.py

• 2b. (6 pts) What does this function do, in your own words? I bonacc! Number this function feturns the xth fibonacc! Number by recursively calling on Hself until it reaches the base case of XEI

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3. Parallelism and recursion

Consider the following function:

```
def longest_run(myarray, key)
    """
    Input:
        `myarray`: a list of ints
        `key`: an int
    Return:
        the longest continuous sequence of `key` in `myarray`
"""
```

E.g., $longest_run([2,12,12,8,12,12,12,0,12,1], 12) == 3$

- 3a. (7 pts) First, implement an iterative, sequential version of longest_run in main.py.
- 3b. (4 pts) What is the Work and Span of this implementation?

Work $\in O(n)$ Stan $\in O(n)$

- 3c. (7 pts) Next, implement a longest_run_recursive, a recursive, divide and conquer implementation. This is analogous to our implementation of sum_list_recursive. To do so, you will need to think about how to combine partial solutions from each recursive call. Make use of the provided class Result.
- 3d. (4 pts) What is the Work and Span of this sequential algorithm?

Work =
$$W(n)$$

Span = $S(n)$

• 3e. (4 pts) Assume that we parallelize in a similar way we did with sum_list_recursive. That is, each recursive call spawns a new thread. What is the Work and Span of this algorithm?