

Tanner Martz

Assignment #1

1A) Is $2^{n+1} \in O(2^n)$? $\Rightarrow f(n) = 2^{n+1}$, $g(n) = 2^n$

$$\lim_{n \rightarrow \infty} \frac{2^{n+1}}{2^n} = \frac{2^{n+1}}{2^n} = \frac{\cancel{2} + 2}{\cancel{2}} = \frac{\infty + 2}{\infty} = \boxed{2}$$

the limit is not equal to 0.

thus 2^{n+1} not bounded by 2^n .

1B) Is $2^{2^n} \in O(2^n)$? $\Rightarrow f(n) = 2^{2^n}$, $g(n) = 2^n$

$$\lim_{n \rightarrow \infty} \frac{2^{2^n}}{2^n} = (\ln(2)) \lim_{n \rightarrow \infty} (2^{2^n})$$

$$(\ln(2)) \infty = \infty$$

the limit is not equal to 0.

thus not upper bounded by 2^n .

1C) Is $n^{1.01} \in O(\log^2 n)$? $\Rightarrow f(n) = n^{1.01}$, $g(n) = \log_2(n)$

$$\lim_{n \rightarrow \infty} \frac{n^{1.01}}{\log_2(n)} = \lim_{n \rightarrow \infty} \frac{1.01 n^{1.01}}{\frac{1}{n \cdot \ln(2)}} = .7 n^{1.01}$$

The limit is not equal to 0.

$n^{1.01}$ is not upper bounded by $\log_2(n)$.

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Assignment 1 pt. 2

1D) Is $n^{1.01} \in \Omega(\log^2 n)$? $\Rightarrow \lim_{n \rightarrow \infty} \frac{n^{1.01}}{\log^2(n)} > 0$
 $n^{1.01}$ is lower bounded by $\log^2(n)$.

1E) Is $\sqrt{n} \in O(\log n)^2$? $\Rightarrow \lim_{n \rightarrow \infty} \frac{\sqrt{n}}{(\log n)^2}$

$\lim_{n \rightarrow \infty} \frac{\ln(2) \sqrt{n}}{\log_2(n)^2} = \frac{\ln^{3/2}}{48} \parallel \lim_{n \rightarrow \infty} \sqrt{n} = \infty$

$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$ not equal 0.

\sqrt{n} is not upper bounded by $(\log n)^2$

1F) Is $\sqrt{n} \in \Omega(\log n)^2$?

$\lim_{n \rightarrow \infty} \frac{\sqrt{n}}{(\log n)^2}$ is > 0 \uparrow

\sqrt{n} is lower bounded by $(\log n)^2$

2(A): CODE

2(B): What does the function do?

The function for the fibonacci sequence iterates through the list, summing values as it progresses. This is a recursive function.

First check x to a valid input, then call the function with values $(x-1)$ and $(x-2)$.

Sum these values and return the input.

Run the function until complete.

Tanner Maltz
Assignment 1 pt. 3

3A) CODE

3B) What is the work and span?

$$W(n) = 2C_1 + n \times C_2 \in O(n) \quad \text{where} \quad \begin{cases} C_1 = \text{initial ass.} \\ C_2 = \text{cost per iteration of the loop} \end{cases}$$

$$S(n) = 2C_1 + n \times C_2 \in O(n)$$

3C) CODE

3D) What is the work and span?

$$\begin{cases} C_1 = \text{Thread operation} \\ C_2 = \text{Conditional check} \\ C_3 = \text{List operation} \end{cases} \quad V(n) = \log n \sum_{i=0}^{\log n} (C_3 n + 2^i C_1) + C_2$$

$$W(n) \in O(n \log n) \quad \star \quad \underline{\text{work}}$$

$$S(n) = \log n \sum_{i=0}^{\log n} (\log n - i) + C_2$$

$$S(n) \in O(\log^2 n) \quad \star \quad \underline{\text{span}}$$

3E) What is the work and span of parallelized algorithm?

$$\begin{cases} C_1 = \text{Thread operation} \\ C_2 = \text{Conditional check} \\ C_3 = \text{List operation} \end{cases} \quad V(n) = \sum_{i=0}^n (i C_3) + C_2$$

$$W(n) \in O(n^2)$$

$$S(n) = \sum_{i=0}^n (i C_3) + C_2$$

$$S(n) \in O(n^2)$$