

CMPS 2200 Assignment 2  
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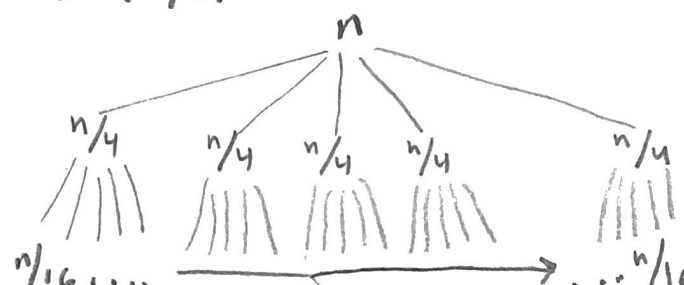
1.) a)  $W(n) = 2W(n/3) + 1$

$$W(n) = \sum_{i=0}^{\log_3 n} (2^i C_1) < n C_1$$

$\log_3 n$  {   $\begin{matrix} 1(C_1) \\ 2(C_1) \\ 4(C_1) \end{matrix}$

$W(n) \in O(n^{\log_3 2})$

b)  $W(n) = 5W(n/4) + n$

$\log_4 n$  {   $\begin{matrix} C_1 n + C_2 \\ 5(n/4 C_1 + C_2) \\ 25(n/16 C_1 + C_2) \end{matrix}$

$$W(n) = \sum_{i=0}^{\log_4 n} (5^i (\frac{n}{4^i} C_1 + C_2)) < n C_2 \log_5 n + \frac{5}{4} n C_2$$

$$W(n) \in O(n^{\log_4 5})$$

c)  $W(n) = 7W(n/7) + n$

Depth:  $\log_7 n$

lvl 1:  $C_1 n + C_2$   
lvl 2:  $7(\frac{n}{7} C_1 + C_2) = n C_1 + 7 C_2$   
lvl 3:  $49(\frac{n}{49} C_1 + C_2) = n C_1 + 49 C_2$

$$W(n) = \sum_{i=0}^{\log_7 n} (n C_1 + 7^i C_2) < n C_2 \log_7 n + 7 C_2 \in O(n \log_7 n)$$

$$W(n) \in O(n \log_7 n)$$

d)  $W(n) = 9W(n/3) + n^2$

Depth:  $\log_3 n$

lvl 1:  $C_1 n^2 + C_2$   
lvl 2:  $9(C_1 (n/3)^2 + C_2) = n^2 C_1 + 9 C_2$   
lvl 3:  $81(C_1 (n/9)^2 + C_2) = n^2 C_1 + 81 C_2$

$$W(n) = \sum_{i=0}^{\log_3 n} \frac{n^2}{3^i} + 9^i C_2 < 3 n^2 C_1 + 9 n C_2 \in O(n^2)$$

$$W(n) \in O(n^2 \log_3 n)$$

$$e.) W(n) = 8W(n/2) + n^3$$

$$\text{Depth: } \log_2 n$$

$$|v| 1: c_1 n^3 + c_2$$

$$|v| 2: 8(c_1 (n/2)^3 + c_2) = n^3/2 c_1 + 8c_2$$

$$|v| 3: 64(c_1 (n/4)^3 + c_2) = n^3/4 c_1 + 64c_2$$

$$W(n) = \sum_{i=0}^{\log_2 n} n^3/2^i c_1 + 8^i c_2 < 2n^3 c_1 + 8n c_2 \in O(n^3)$$

$$W(n) \in O(n^3 \log_2 n)$$

$$f.) w(n) = 49(n/25) + n^{3/2} \log n$$

$$\text{Depth: } \log_{25} n$$

$$|v| 1: n^{3/2} \log n + c_2$$

$$|v| 2: 49(n^{3/2}/25 \log(n/25) c_1 + c_2)$$

$$= \frac{49^{3/2}}{25} \log(n/25) c_1 + 49 c_2$$

$$w(n) = \sum_{i=0}^{\log_{25} n} \left( \frac{n^{3/2}}{25^i} \log \left( \frac{n}{25^i} \right) c_1 + 49^i c_2 \right) = \sum_{i=0}^{\log_{25} n} \left( \frac{n^{3/2}}{25^i} \log \left( \frac{n}{25^i} \right) c_1 \right) + \sum_{i=0}^{\log_{25} n} 49^i c_2$$

$$w(n) < 25 n^{3/2} (\log^2 n + \log n) + 49 n c_2$$

$$< 25 n^{3/2} (\log^2 n + \log n)$$

$$w(n) \in O(n^{3/2} \log n)$$

$$g.) w(n) = w(n-1) + 2$$

$$w(n) = \sum_{i=1}^n 2 = n(n-1) \in O(n)$$

$$w(n) \in O(n)$$

$$h.) w(n) = w(n-1) + n^c$$

$$w(n) = \sum_{i=1}^n i^c = \frac{n^{c+1}(n^c - 1)}{2} \in O(n^{c+1})$$

$$w(n) \in O(n^{c+1})$$

$$i.) w(n) = w(\sqrt{n}) + 1$$

$$w(n) = w(n^{1/2}) + 1$$

$$w(n) \in O(\log(\log(n)))$$

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2.) Algorithm A:  $W(n) = 5W(n/2) + n$

Algorithm B:  $W(n) = 2W(n-1) + 1$

Algorithm C:  $W(n) = 9W(n/3) + O(n^2)$

Algorithm A would have  $W(n) \in O(n^{\log_2 5})$

Algorithm B would have  $W(n) \in O(2^n)$

Algorithm C would have  $W(n) \in O(n^2 \log_3 n)$

work proven in #1:  
1b, 1g, 1e

Comparing  $W(n)$  A to  $W(n)$  C:  $n^{\log_2 5} = ? n^2 \log_3 n$  meaning  
 $\log n < n^{\log_2 5/2}$   
 $W(n)A < W(n)C$   
 $W(n)C$  is less work and  
therefore the best choice.