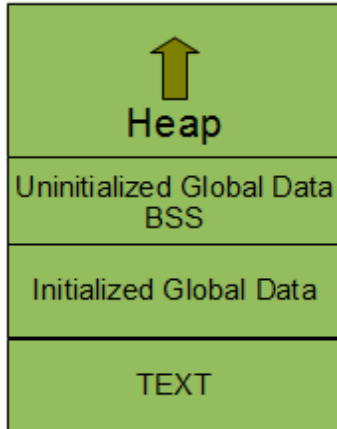
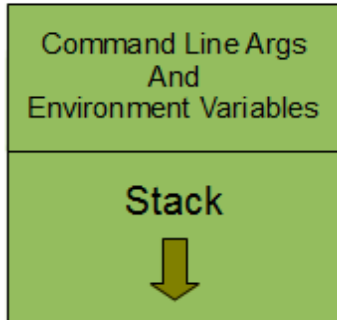


Buffer Overruns

Week 8

Process Layout

(Higher Address)



(Lower Address)

- TEXT segment
 - Contains machine instructions to be executed
- Global Variables
 - Initialized
 - Uninitialized
- Heap segment
 - Dynamic memory allocation
 - malloc, free
- Stack segment
 - Push frame: Function invoked
 - Pop frame: Function returned
 - Stores
 - Local variables
 - Return address, registers, etc
- Command Line arguments and Environment Variables

The Stack

- Contiguous block of memory containing data
- The stack pointer (SP) points to top of stack
- Bottom of stack is at a fixed address (FP- frame pointer)

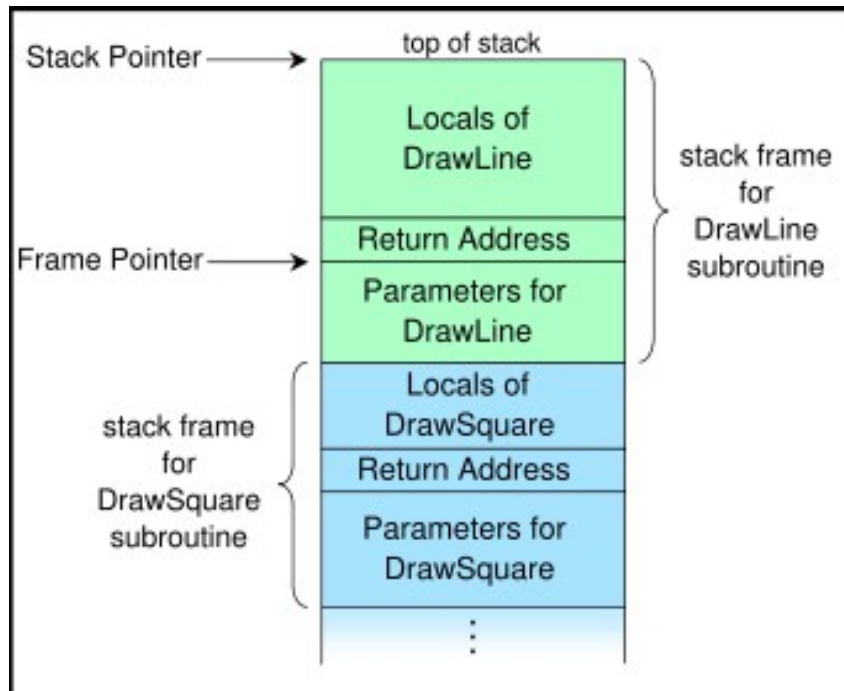


Image source : wikipedia

```
• DrawSquare(topleft, bottomright)
{
    ...
    DrawLine(p1, p2);
    FillRect(topleft, bottomright, color);
    ...
}

• DrawLine(pt1, pt2)
{
    int local;
    float slope;
    ...
}
```

The Frame

- Stack consists of logical stack frames that are pushed when calling a function and popped when returning
- Stack frame contains parameters to function, local variables, and data necessary to recover previous stack frame including the instruction pointer at the time of the function call
- Frame pointer (FP) points to fixed location within a frame and variables are referenced by offsets to the FP

Calling a Procedure

- Save previous FP
- Copies SP into FP to create the FP
- Advances SP to reserve space for the local variables

Buffer overflows

- Result of stuffing more data into a buffer than it can handle

Example

```
void function(char *str)
{
    char buffer[16];
    strcpy(buffer, str);
}
```

```
void main(int argc, char** argv)
{
    char large_string[256];
    int i;
    for( i = 0; i < 255; i++)
    {
        large_string[i] = 'A';
    }
    large_string[255] = '\0'
    function(large_string);
    printf("Hello world!\n");
}
```

bottom of
memory

top of
memory

	buffer	sfp	ret	*str
<-----	[][][][

top of
stack

bottom of
stack

What's Going to Happen?

- What happened to buffer and the other regions in the stack?
- If the character 'A' hex value is 0x41, what is the return address?
- What happens when the function returns?

Shell Code

- Now that we can modify the return address and the flow of execution, what program to execute?
- We want to spawn a shell so we can execute anything else.
- We need to place instructions into the program's address space

The Answer

- Place the code we are trying to execute in the buffer we are overflowing
- Overwrite the return address so it points back into the buffer

Prevention?

- Hardware?
- Software?

Laboratory

- Build thttpd 2.25b with the patch
 - Gunzip and untar it
 - Patch it
 - Configure it
 - Make it
- Run it on port (12400 to 12499)
 - `./thttpd -p <port>`
- Do a simple request like
 - `wget http://localhost:portnumber`

Laboratory

- Crash the web server by sending it a suitably-formatted request
 - `wget http://localhost:portnumber/AAAAAAAAAA.....AAAA` (about 7k As)
 - Where does the buffer overrun occur? Why?

Laboratory

- Run the web server under GDB and get traceback (bt) after the crash
 - `./thttpd -p <portnumber>`
 - Find the pid for thttpd: `ps -e | grep thttpd`
 - `gdb thttpd pid`
 - Send your crashing request
 - Continue and when it crashes do `bt`
 - Include this in lab8.txt
- Describe how you would build a remote exploit in the modified thttpd
 - Smashing the stack for Fun and Profit will be helpful

Stack Protection

Here is a list of commands to generate the assembly files. You can do a compile-only (-c flag) on the thttpd.c file

StackProtection

Run these inside the "src" folder

```
gcc -m32 -c -S -o thttpd-nostackprot.s -fno-stack-protector -I . -I ../ thttpd.c
```

```
gcc -m32 -c -S -o thttpd-stackprot.s -fstack-protector -I . -I ../ thttpd.c
```

Homework

- Things to think about...
 - Significance of Damage
 - Ease of Exploitation
 - Widespread
 - Ease of repair/prevention