I3304 : NETWORK ADMINISTRATION AND SECURITY
CHAPTER 1 – INTRODUCTION

DR. BASSEM HAIDAR

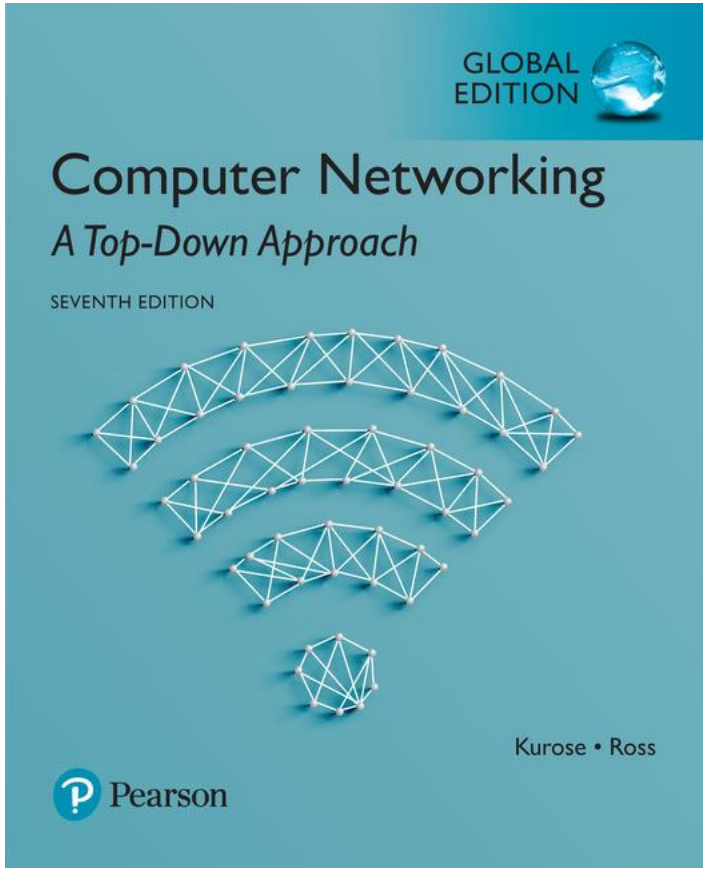
FALL 2021-2022

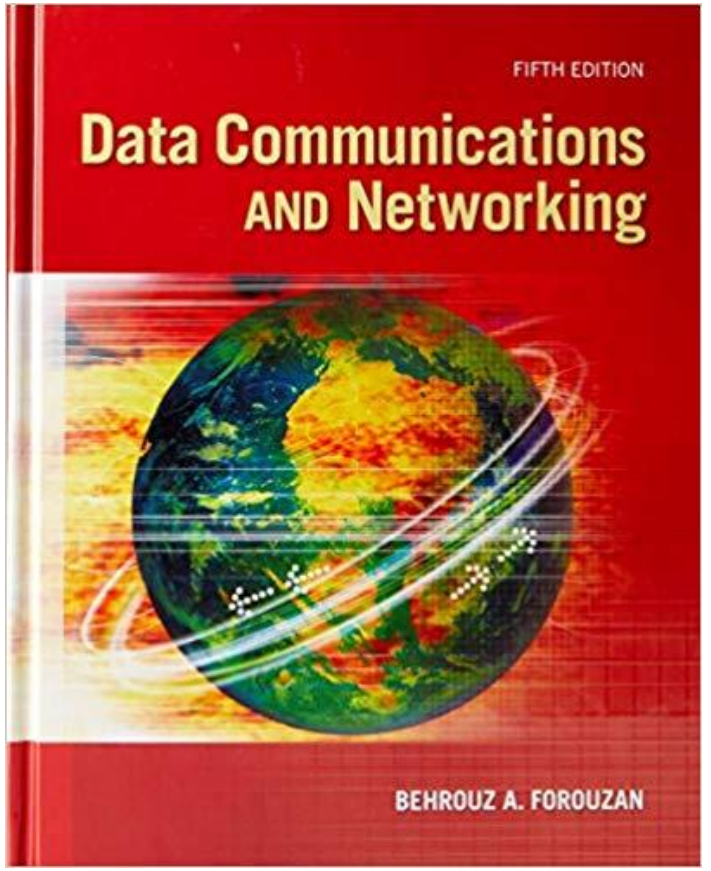
SOME INFORMATION ABOUT THE COURSE

- Instructor Information
  - Instructor: Bassem Haidar
  - Office: 236
  - Office Hours: department schedule or by appointment
  - email: bassem.haidar@ul.edu.lb

- Course Information
  - Lectures:
    - Tuesday 08:00 –9:40 for English section
    - Friday 08:00 –9:40 for French section
  - Exercises: integrated in the course
  - Lab : Monday 8:00 – 9:40

TEXT BOOKS


GLOBAL
EDITION
Computer Networking
A Top-Down Approach
SEVENTH EDITION
Kurose • Ross
Pearson

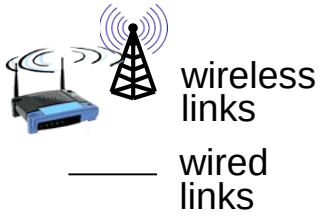

FIFTH EDITION
Data Communications
AND Networking
BEHROUZ A. FOROUZAN

Outline

- Introduction
  - Introduction to the course
  - Recall Network Basics (I2208)
- Network Layer
  - IP packet structure (Recall)
  - Static Routing
  - Dynamic Routing Algorithm
  - Dynamic Routing Protocols
  - NAT (Network Address Translation)
- Transport Layer
  - Function of the transport layer
  - UDP Protocol
  - TCP Protocol
    - Connection management
    - Flow control
    - Congestion control
- Application Layer
  - HTTP protocol
  - FTP protocol
  - Mail protocols
  - DNS
- Introduction to Security
  - Security services
  - Cryptography
  - Digital Signature
  - Principle of network security protocols

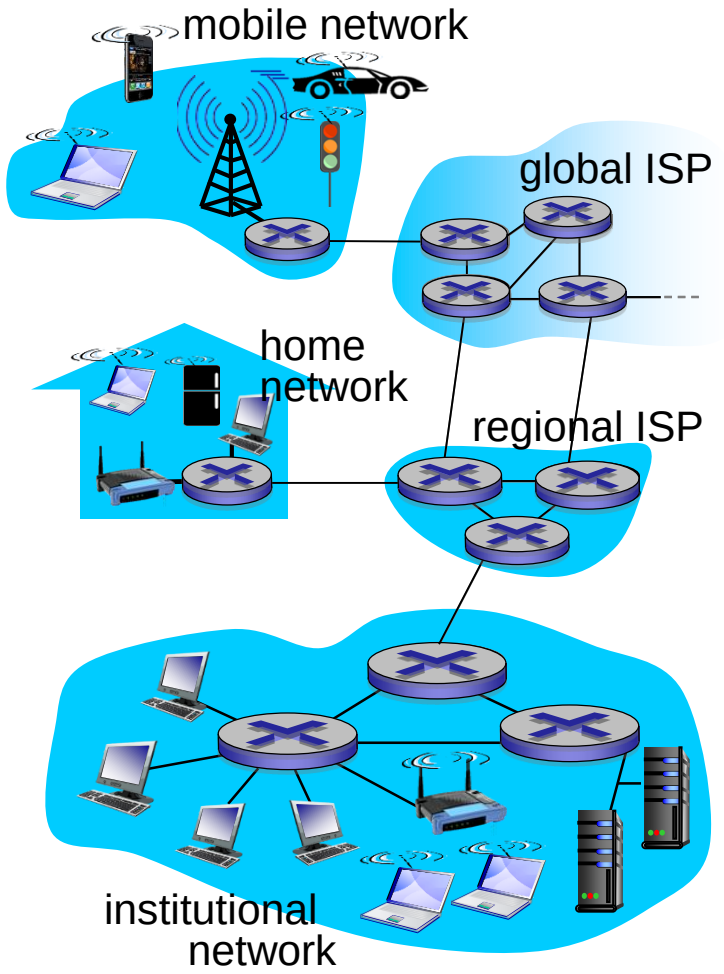
What's the Internet: "nuts and bolts" view

- billions of connected computing devices:
  - *hosts* = *end systems*
  - running *network apps*



- *communication links*
  - fiber, copper, radio, satellite
  - transmission rate: *bandwidth*

- *packet switches:* forward packets (chunks of data)
  - *routers* and *switches*



"Fun" Internet-connected devices

IP picture frame
http://www.ceiva.com/

Web-enabled toaster +
weather forecaster

Tweet-a-watt:
monitor energy use

Internet
refrigerator

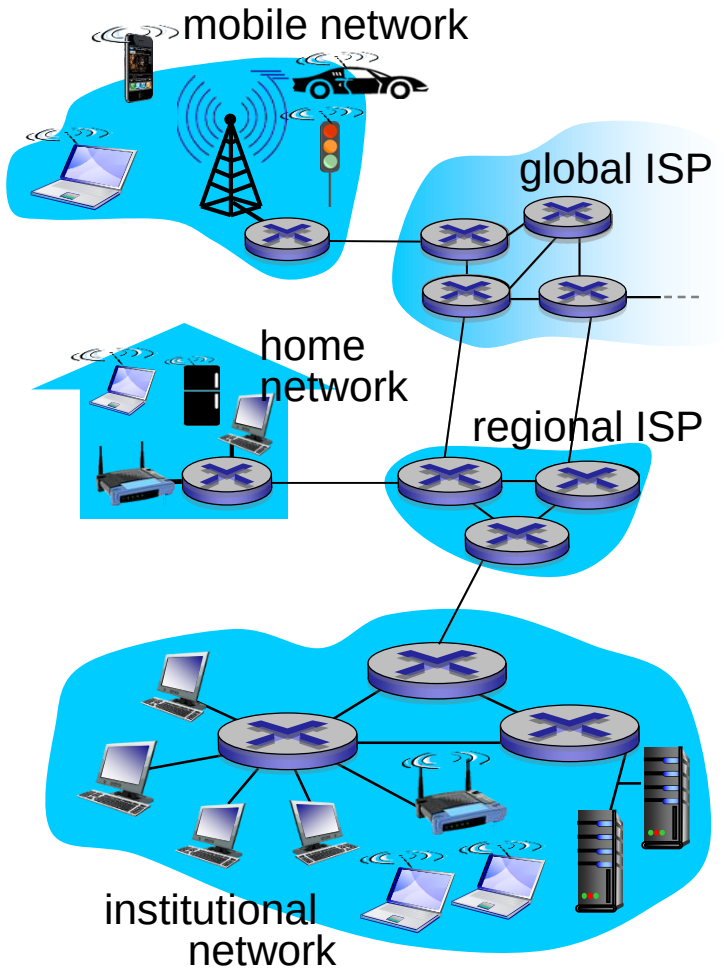
Slingbox: watch,
control cable TV remotely

sensorized,
bed
mattress

Internet phones

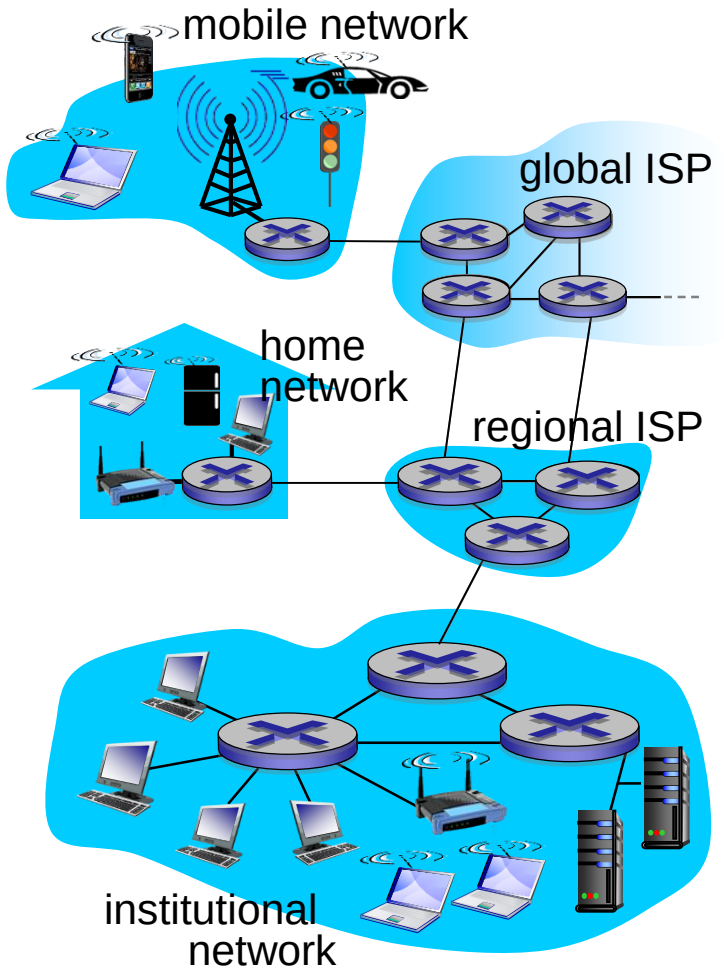
What's the Internet: "nuts and bolts" view

- *Internet:* "network of networks"
  - Interconnected ISPs
- *protocols* control sending, receiving of messages
  - e.g., TCP, IP, HTTP, Skype, 802.11
- *Internet standards*
  - RFC: Request for comments
  - IETF: Internet Engineering Task Force



What's the Internet: a service view

- *infrastructure that provides services to applications:*
  - Web, VoIP, email, games, e-commerce, social nets, …
- *provides programming interface to apps*
  - hooks that allow sending and receiving app programs to "connect" to Internet
  - provides service options, analogous to postal service



What's a protocol?

human protocols:

- "what's the time?"
- "I have a question"
- introductions

… specific messages sent

… specific actions taken when messages received, or other events

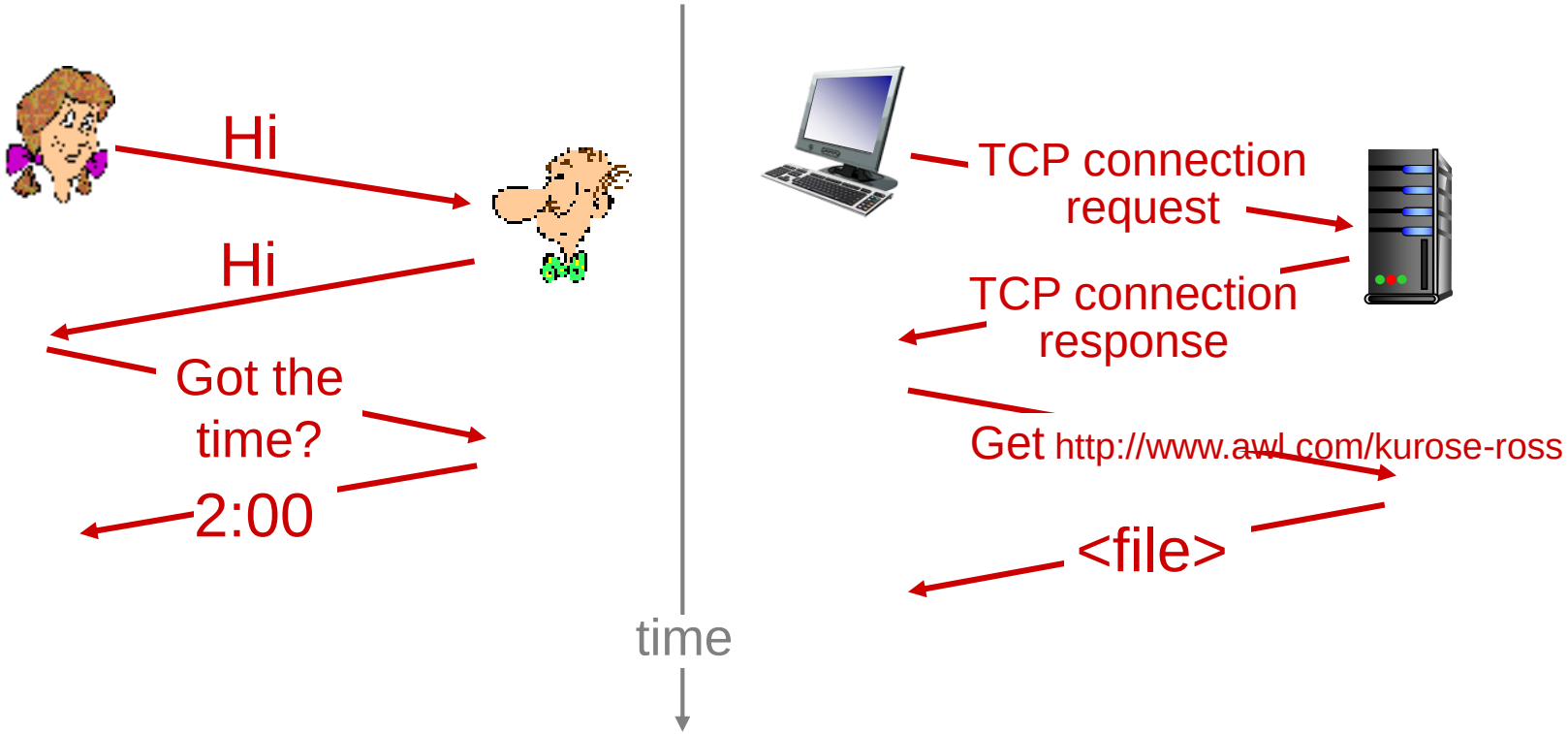
network protocols:

- machines rather than humans
- all communication activity in Internet governed by protocols

protocols define format, order of messages sent and received among network entities, and actions taken on message transmission, receipt

What's a protocol?

a human protocol and a computer network protocol:



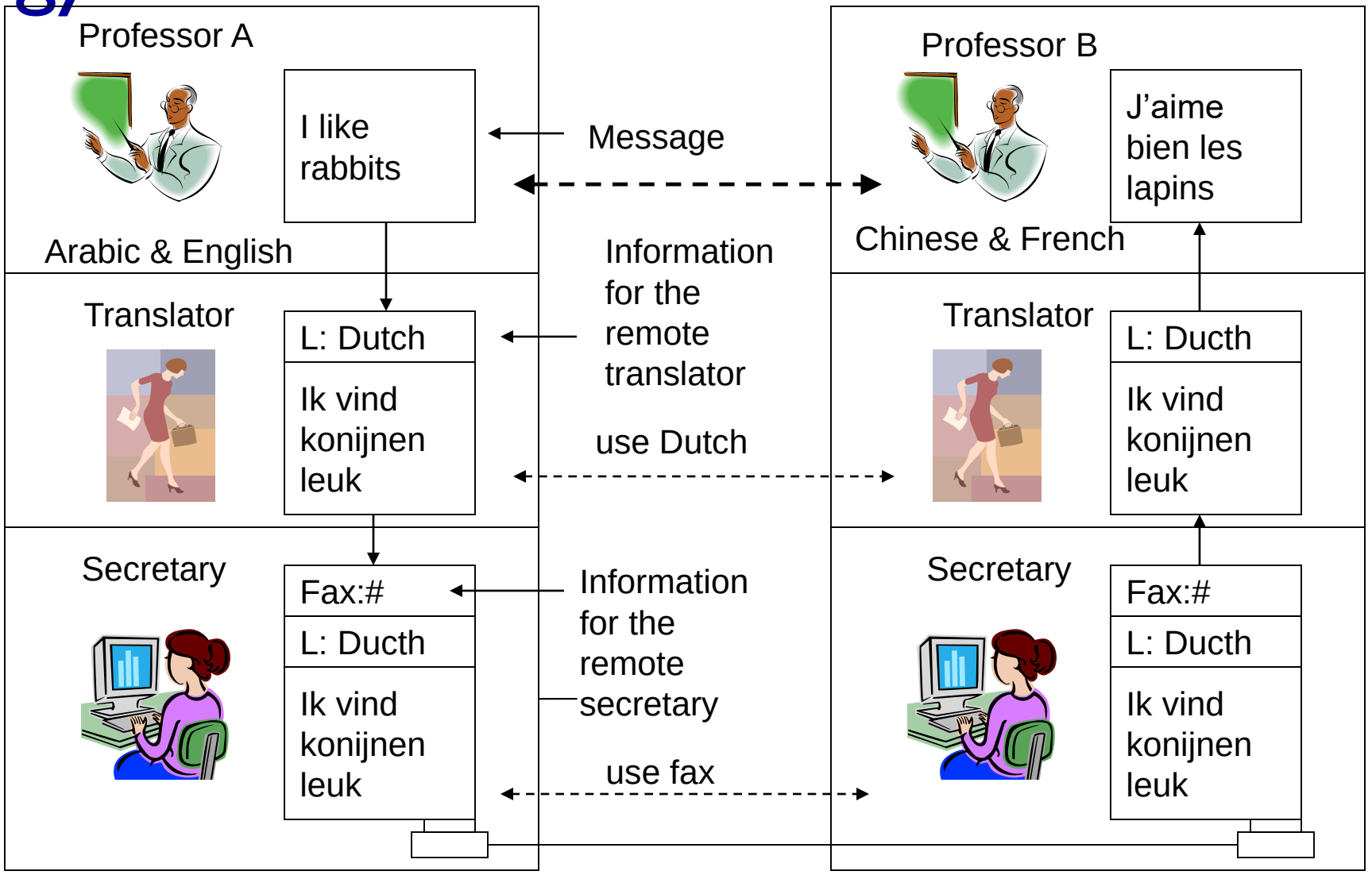
Q: other human protocols?

Why layering?

dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
  - layered *reference model* for discussion
- modularization eases maintenance, updating of system
  - change of implementation of layer's service transparent to rest of system

Analogy


Professor A
I like rabbits
Message
Professor B
J'aime bien les lapins
Arabic & English
Chinese & French
Translator
L: Dutch
Ik vind konijnen leuk
Information for the remote translator
use Dutch
Translator
L: Ducth
Ik vind konijnen leuk
Secretary
Fax:#
L: Ducth
Ik vind konijnen leuk
Information for the remote secretary
use fax
Secretary
Fax:#
L: Ducth
Ik vind konijnen leuk

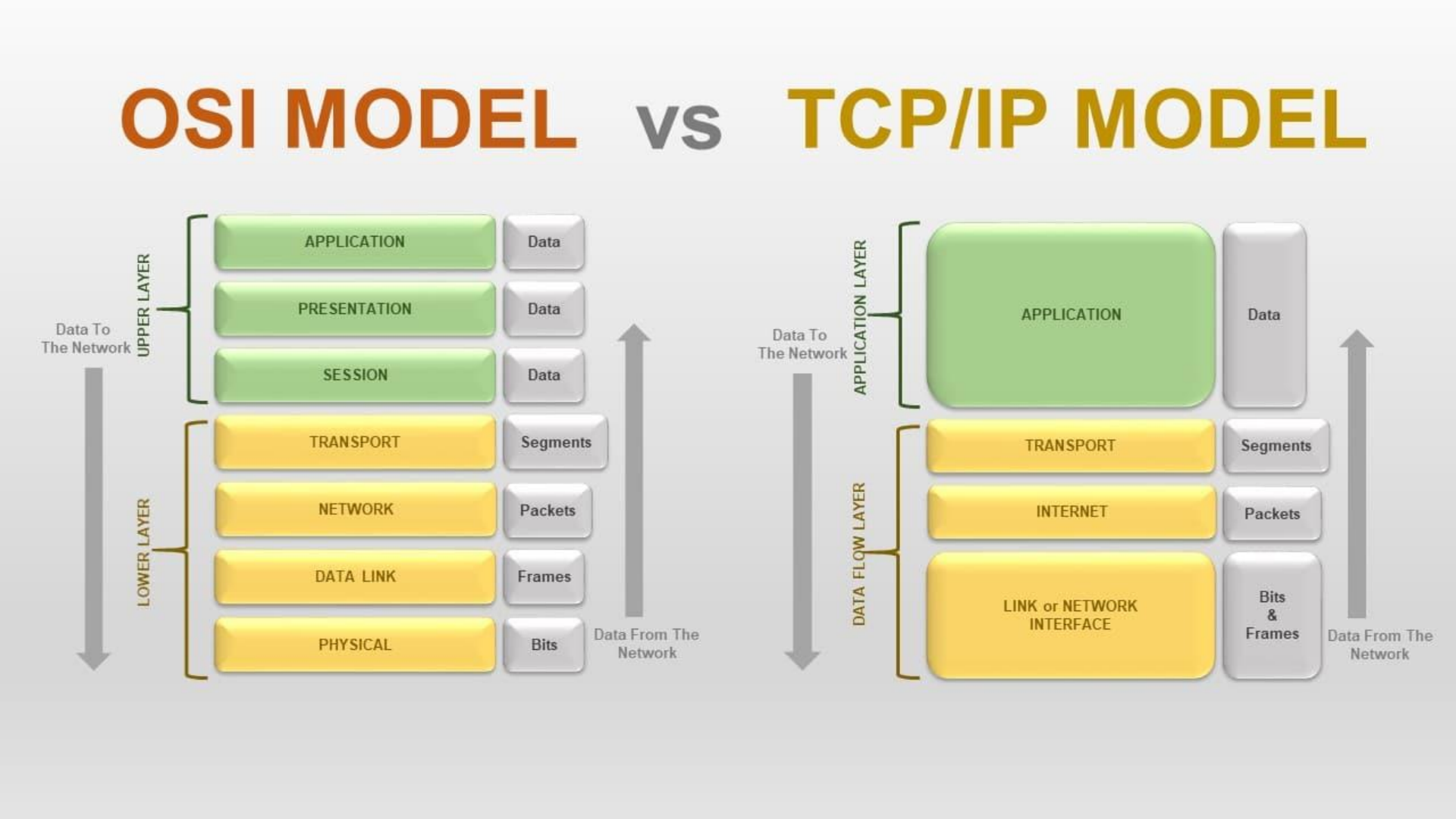
Reference Models

- There are two competing models for how the software is layered. These are the OSI and the TCP models.

- **OSI (Open Systems Interconnection)**
  - Developed by ISO (International Standards Organization)
  - 7 layers

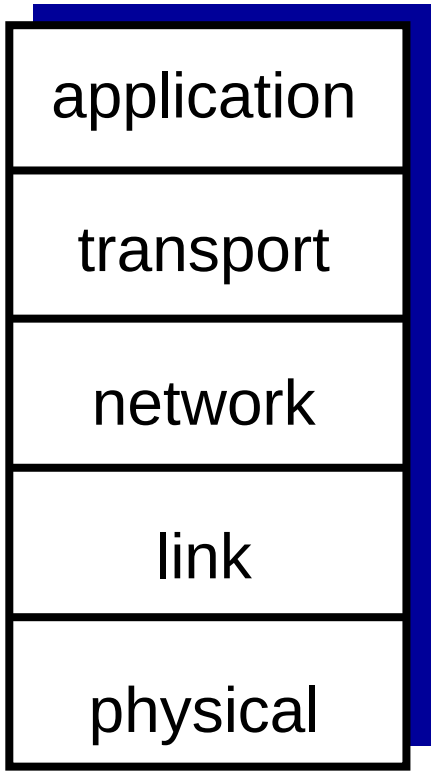
- **TCP (Transfer Control Protocol)**
  - Used in the Arpanet and in the Internet. Common mechanism that is surpassing the OSI Model.
  - 5 layers

Modèles de références



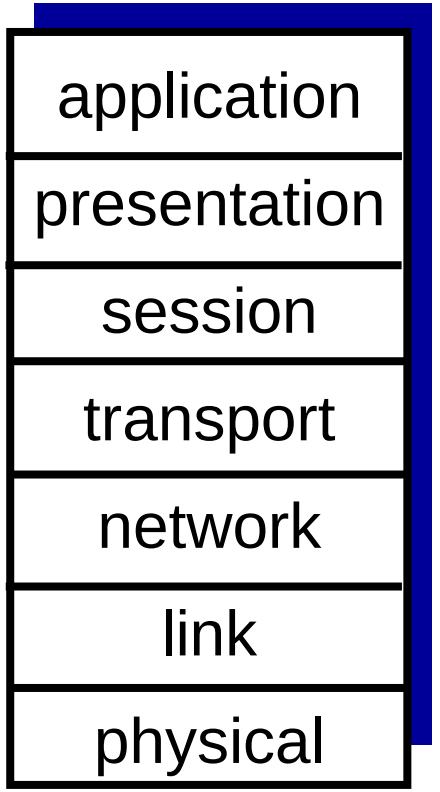
Internet protocol stack

- *application:* supporting network applications
  - FTP, SMTP, HTTP
- *transport:* process-process data transfer
  - TCP, UDP
- *network:* routing of datagrams from source to destination
  - IP, routing protocols
- *link:* data transfer between neighboring network elements
  - Ethernet, 802.111 (WiFi), PPP
- *physical:* bits "on the wire"

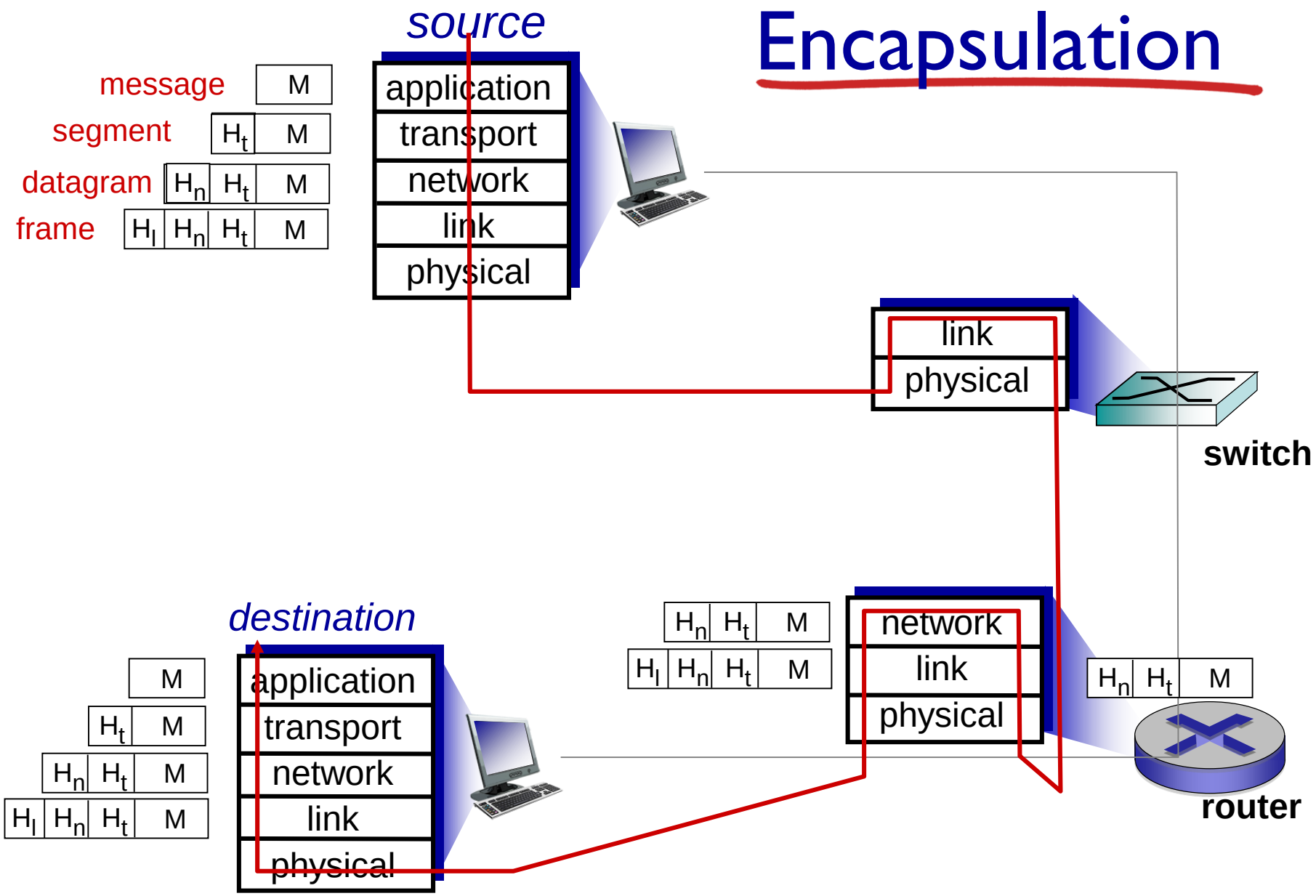


ISO/OSI reference model

- *presentation:* allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- *session:* synchronization, checkpointing, recovery of data exchange
- Internet stack "missing" these layers!
  - these services, *if needed*, must be implemented in application
  - needed?



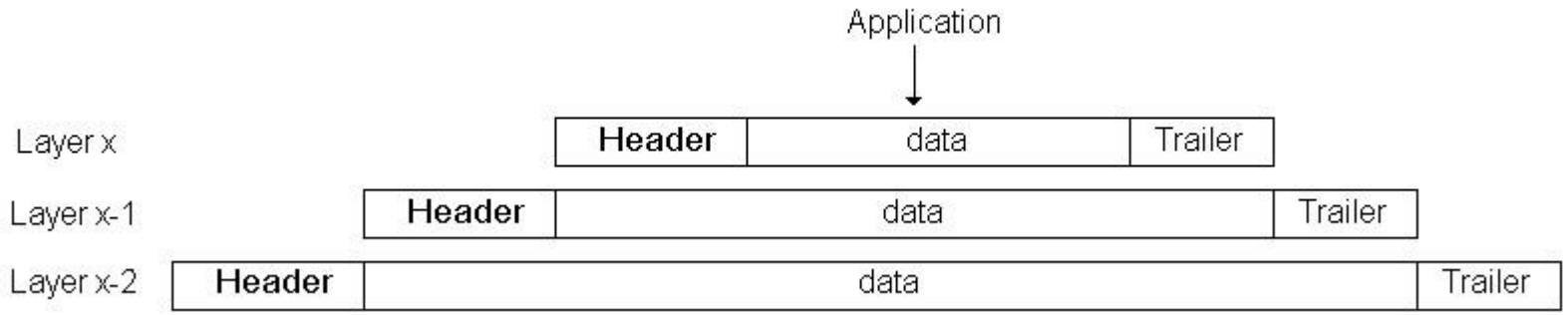
Encapsulation



Reference Models/Modèles de références

- Headers, Data, and Trailers (Entête, Données, et postambule)

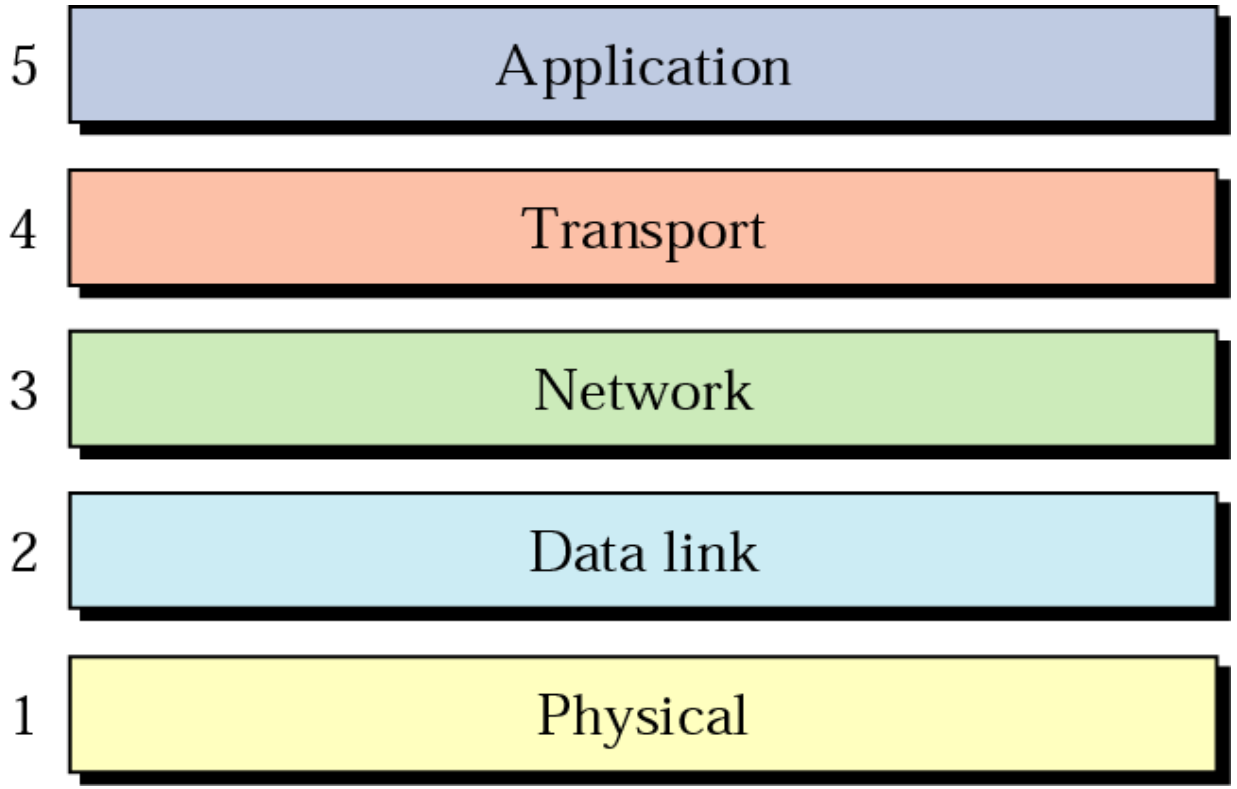
- Encapsulation

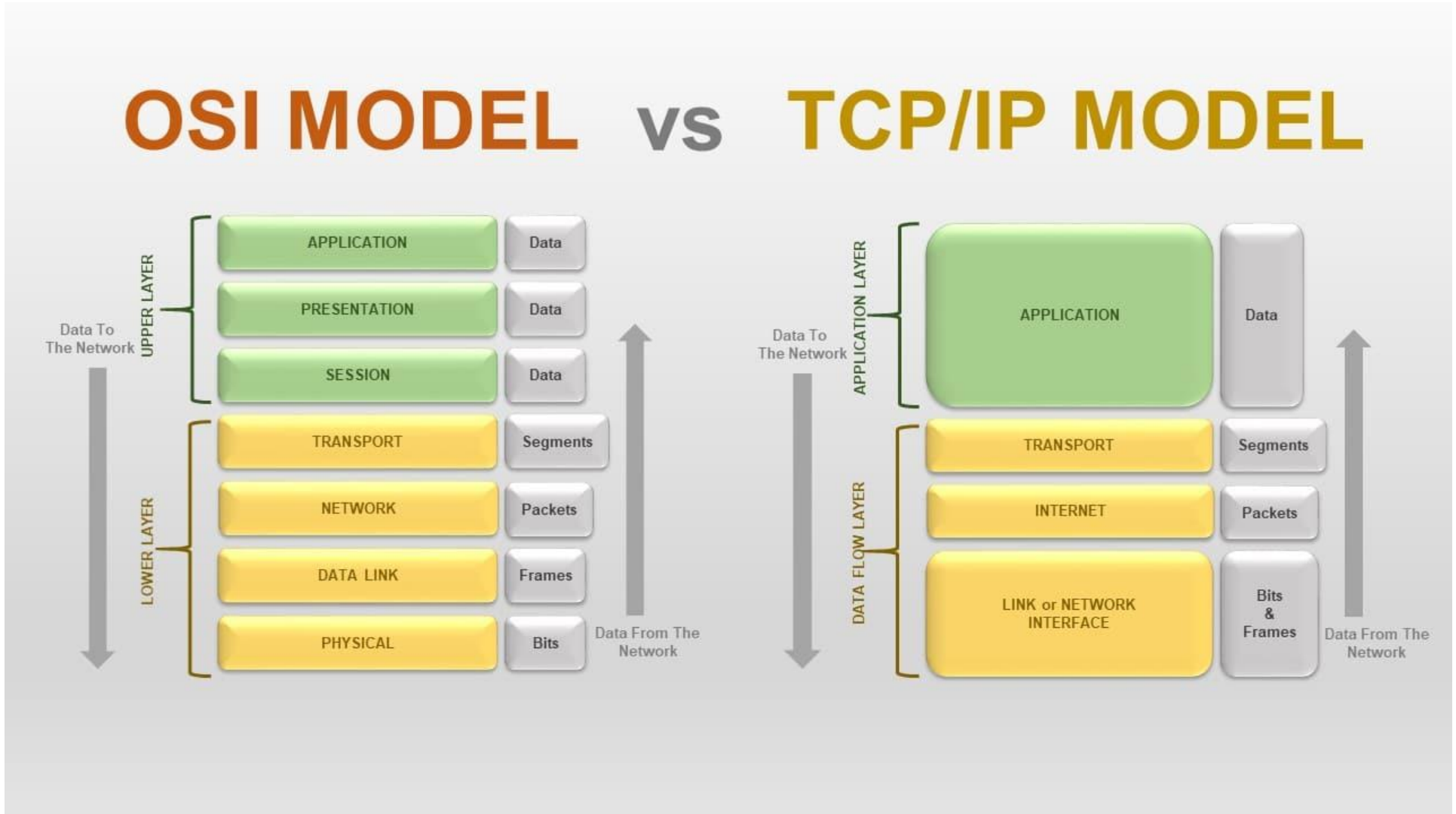


TCP/IP MODEL

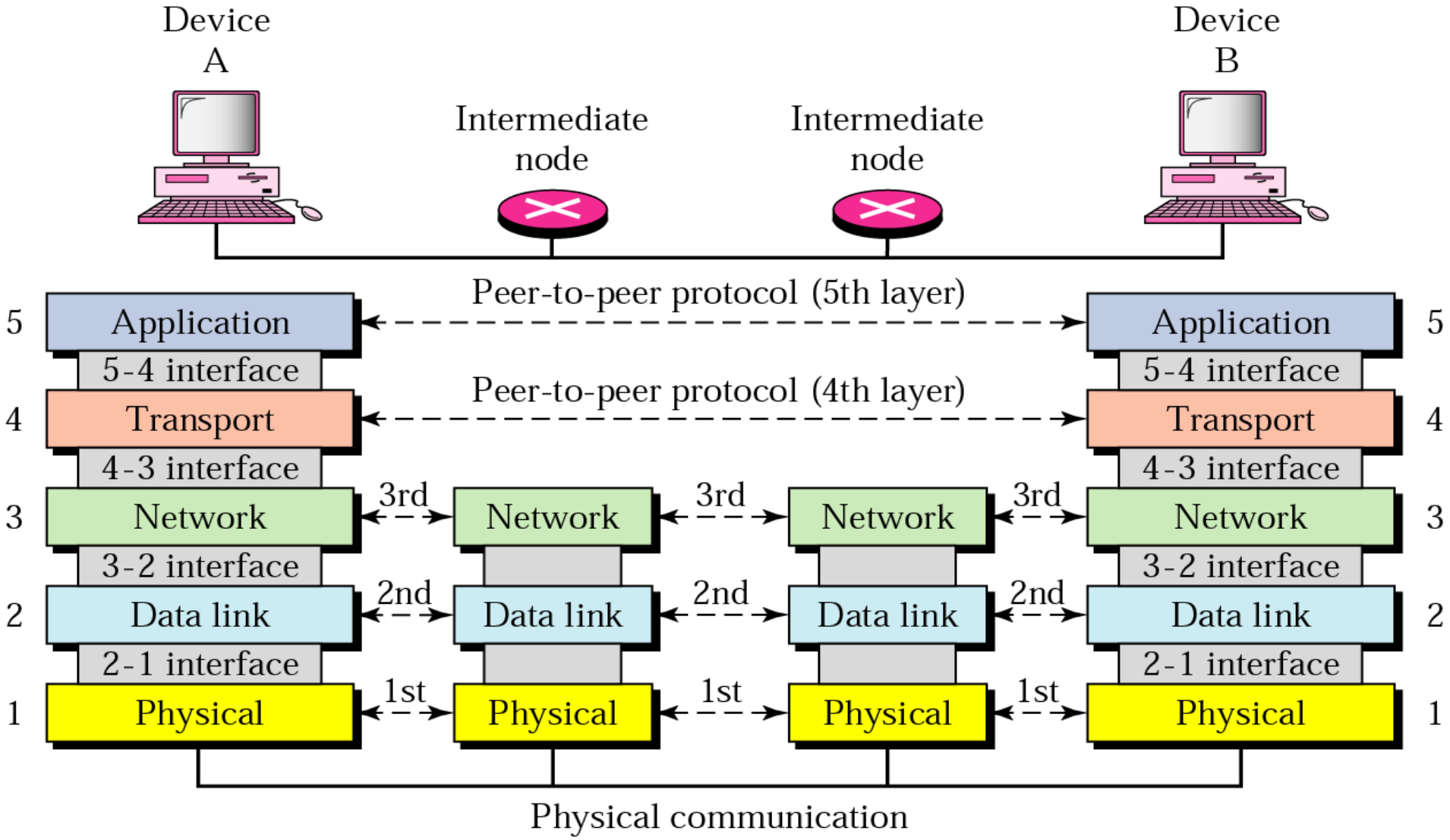
Internet Model

- Dominant model in **data communications and networking**
- 5 ordered layers; often referred to as TCP/IP protocol suite

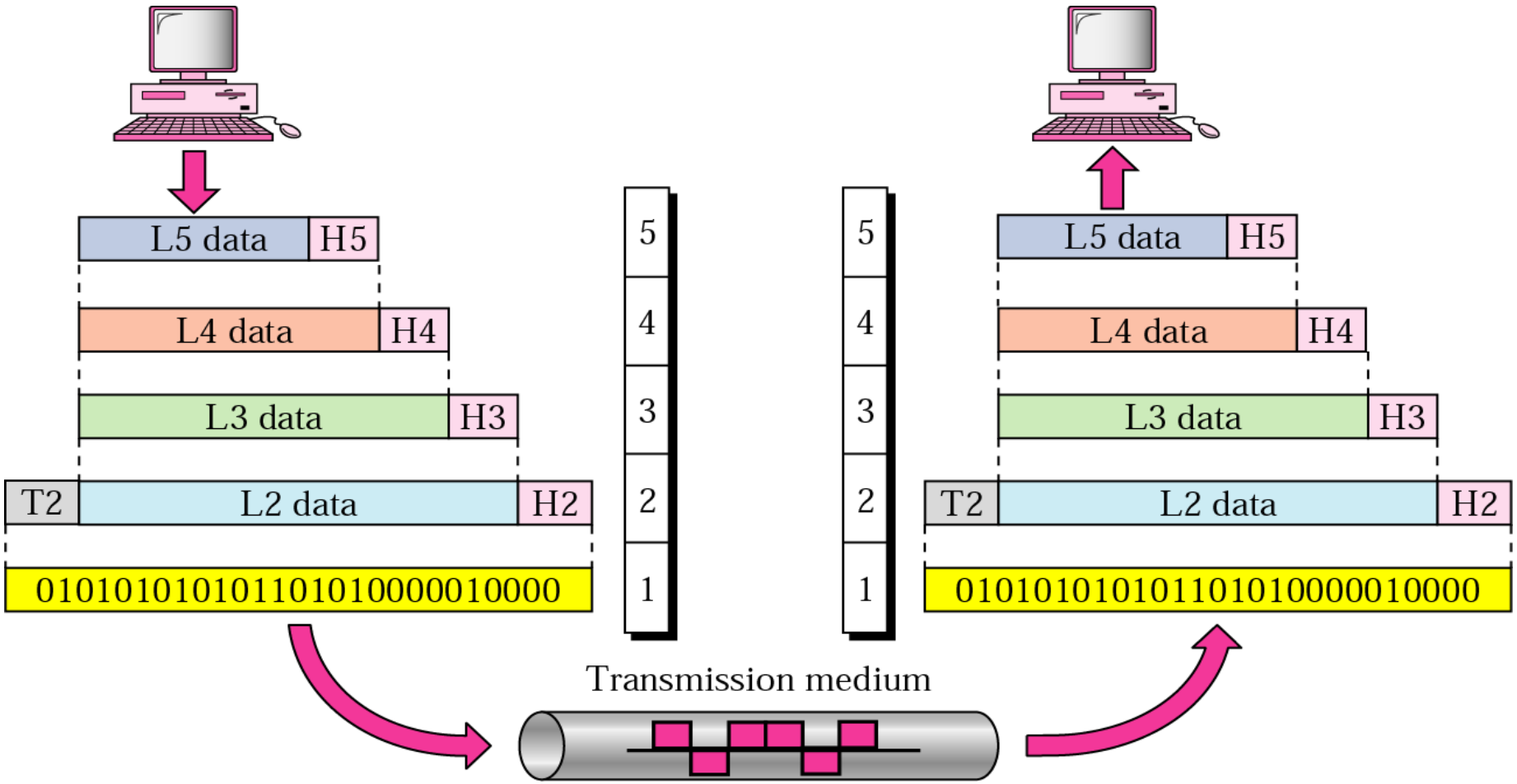



OSI MODEL
vs
TCP/IP MODEL
Data To The Network
UPPER LAYER
APPLICATION
Data
PRESENTATION
Data
SESSION
Data
LOWER LAYER
TRANSPORT
Segments
NETWORK
Packets
DATA LINK
Frames
PHYSICAL
Bits
Data From The Network
Data To The Network
APPLICATION LAYER
APPLICATION
Data
DATA FLOW LAYER
TRANSPORT
Segments
INTERNET
Packets
LINK or NETWORK INTERFACE
Bits & Frames
Data From The Network

Peer-to-peer processes

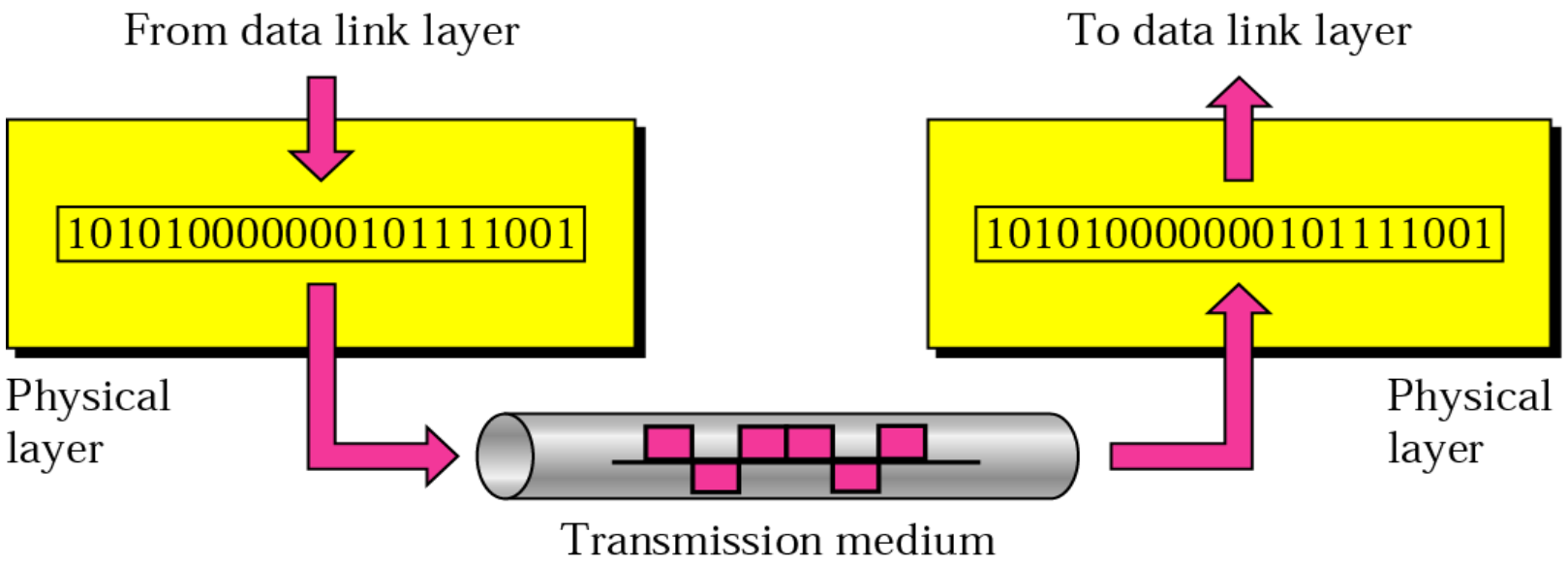


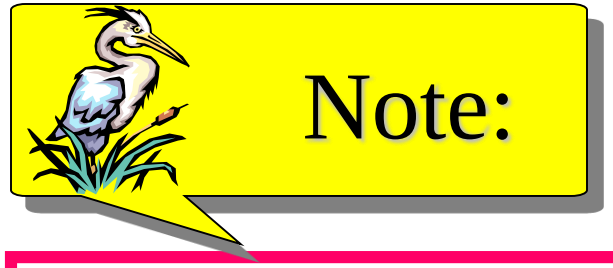
An exchange using the Internet model



Physical characteristics of interfaces and media

- Representation of bits without interpretation
- **Data rate: number of bits per second**
- Synchronization of bits

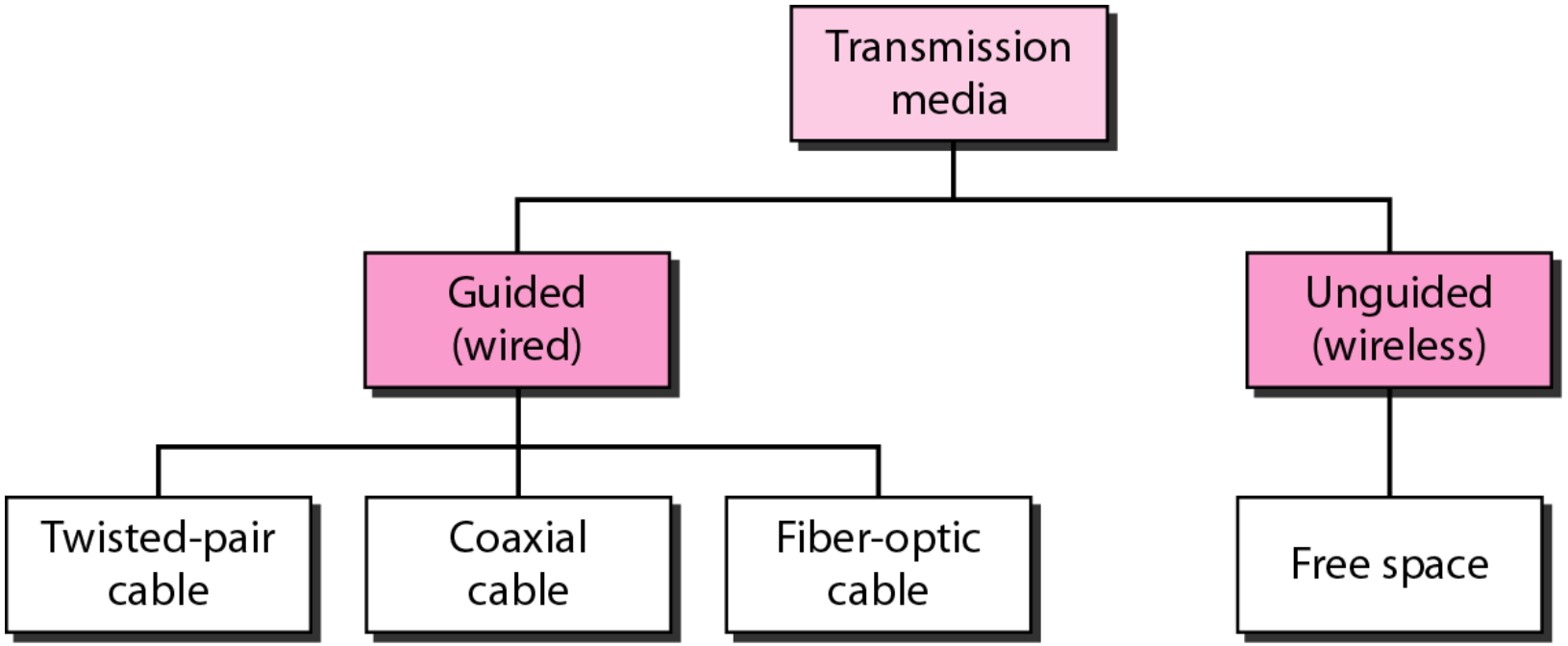




Note:

The physical layer is responsible for transmitting individual bits from one node to the next.

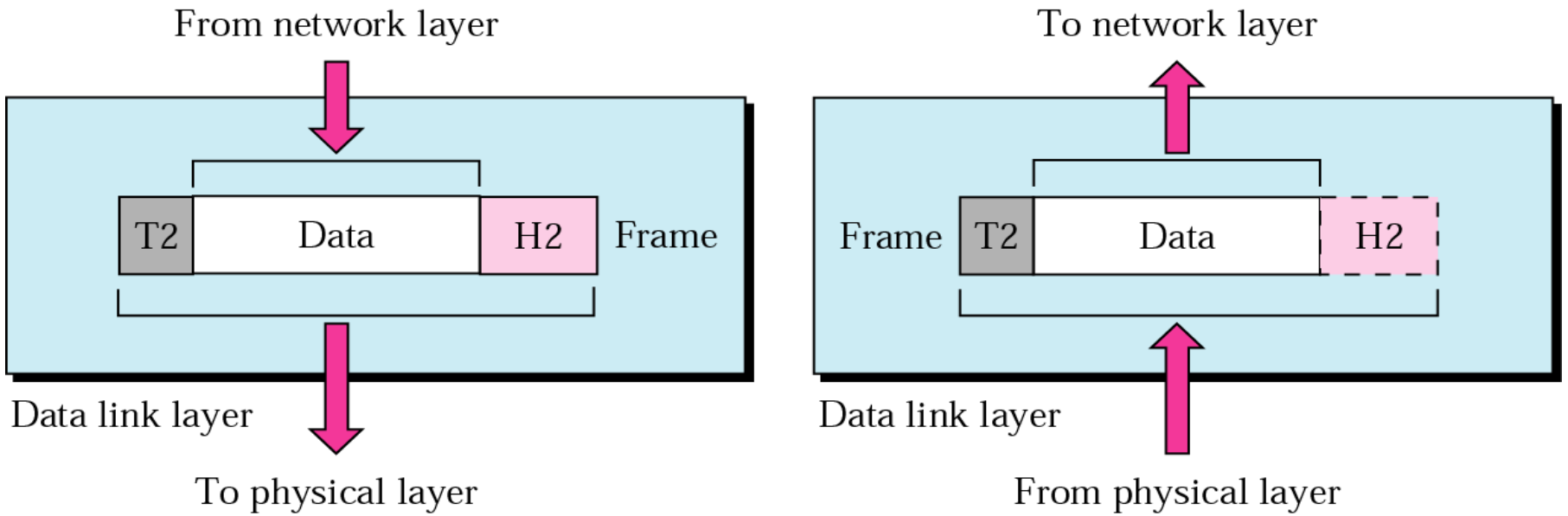
Transmission Media

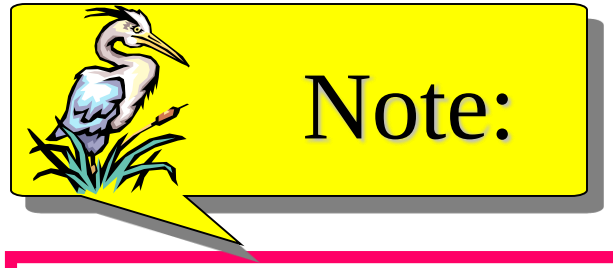


Data Link Layer Responsibilities

Defines frames into manageable **data units**

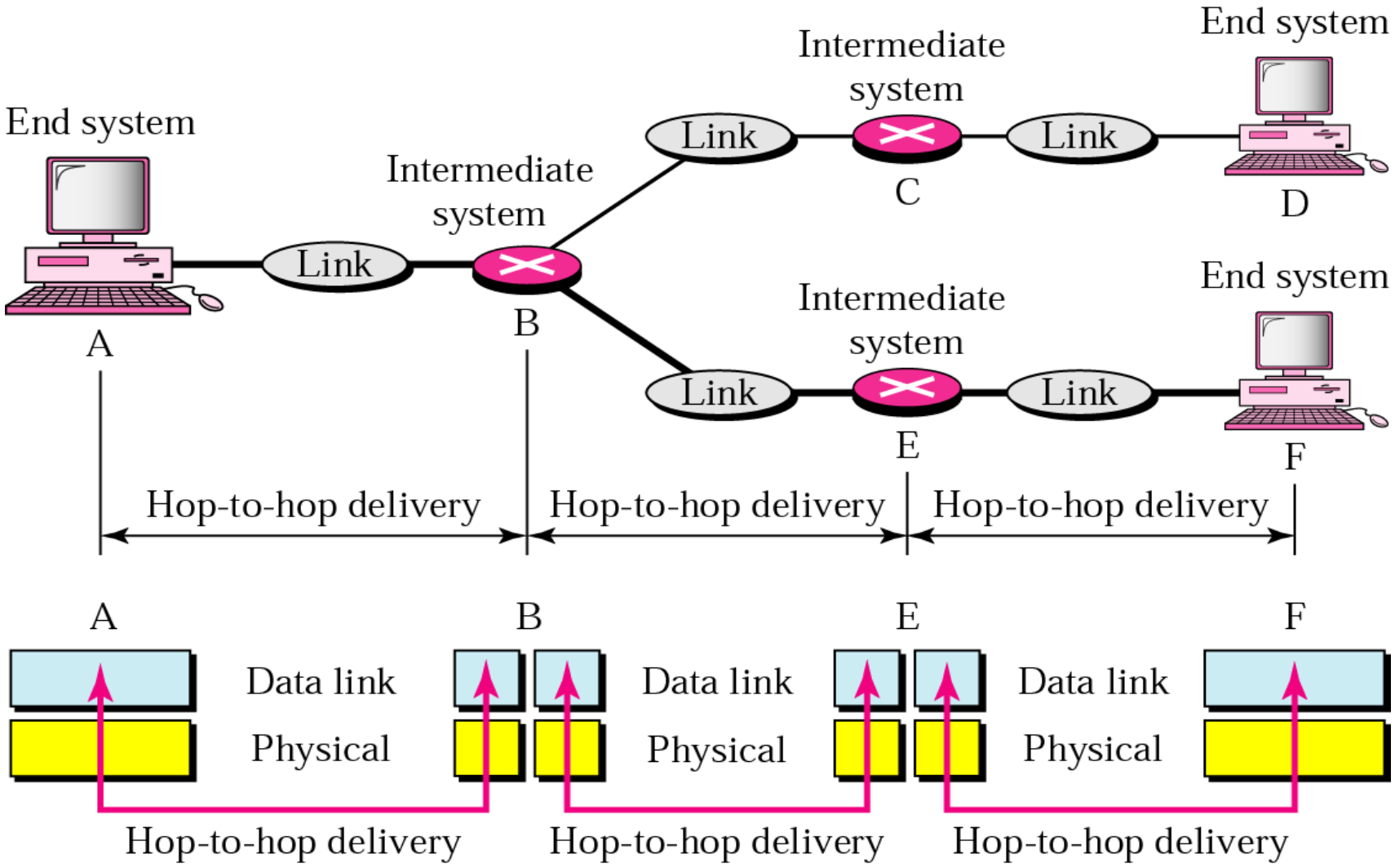
- Physical addressing
- Flow control
- Error control
- Access control



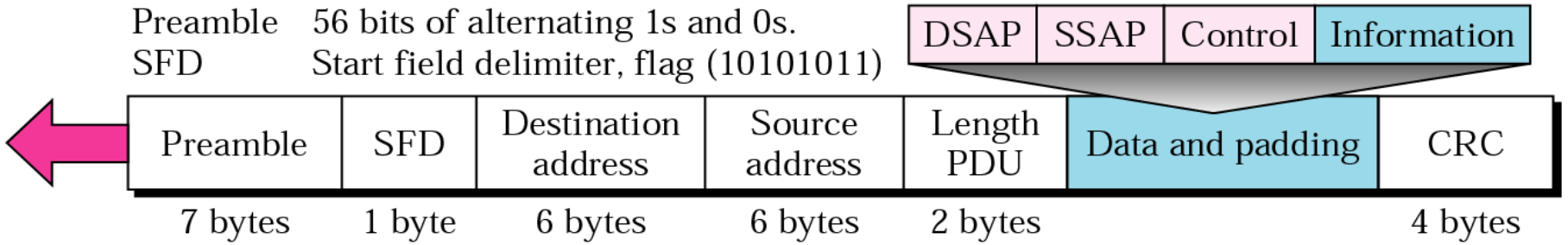


The data link layer is responsible for transmitting frames from one node to the next.

Node-to-node delivery



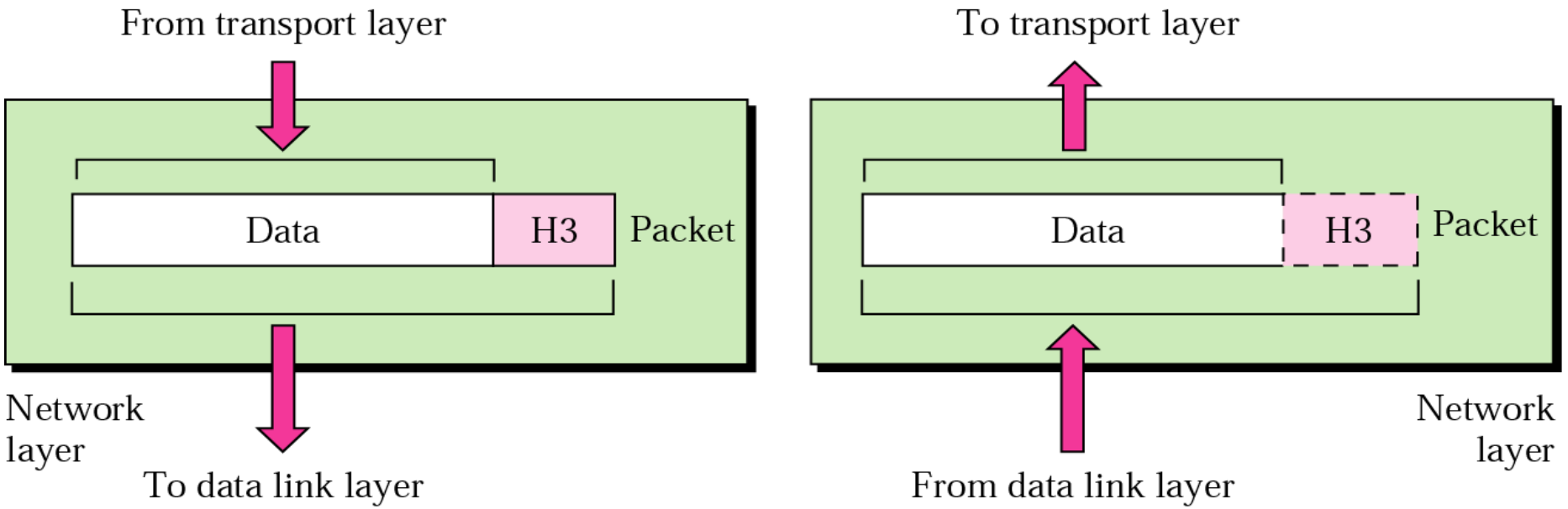
802.3 MAC frame

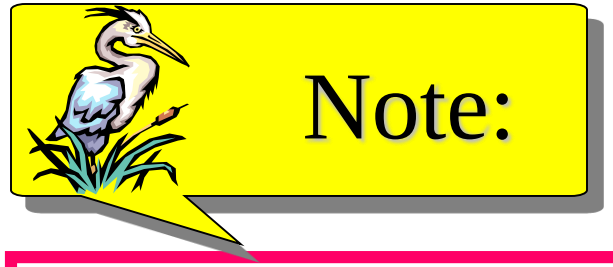


- Preamble – 7 bytes of alternating 0s and 1s to alert the receiver and allow it to synchronize
- Start Frame Delimiter (SFD) – 1 byte – 10101011 signals the beginning of a frame, last chance for synchronization – last 2 bits are 11
- Destination address (DA) – 6 bytes – contains the physical address of the destination station or stations
- Source address (SA) – 6 bytes – contains the physical address of the sender
- Length/type – if less than 1518 then it defines the length of the data field – if more than 1536 then it defines the type of the PDU packet that is encapsulated
- Data – data encapsulated from upper-layer protocols : 46 ~ 1500 bytes
- CRC – CRC-32

Network Layer Responsibilities

- Source-to-destination delivery, possibly across multiple networks
- Logical addressing
- Routing

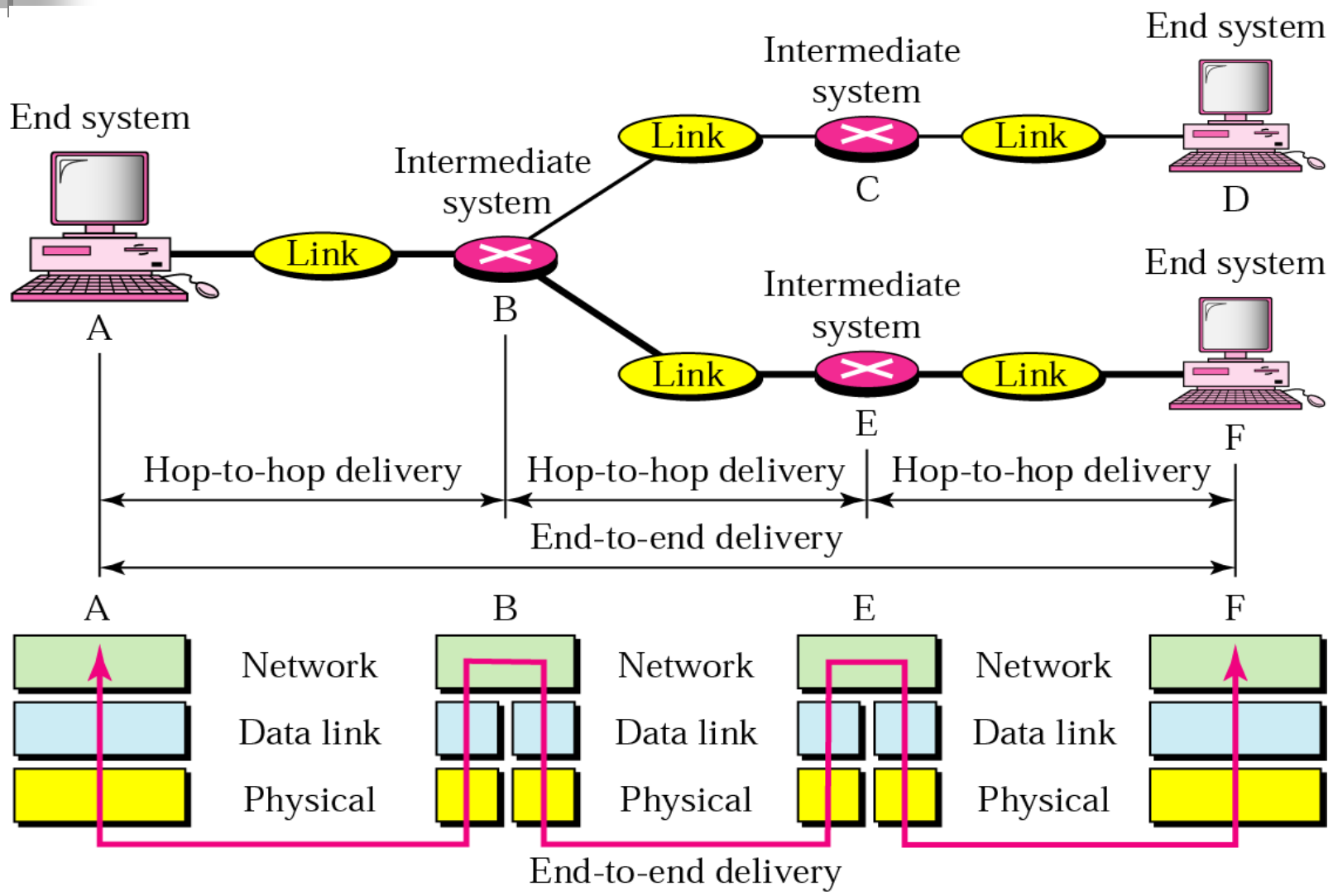




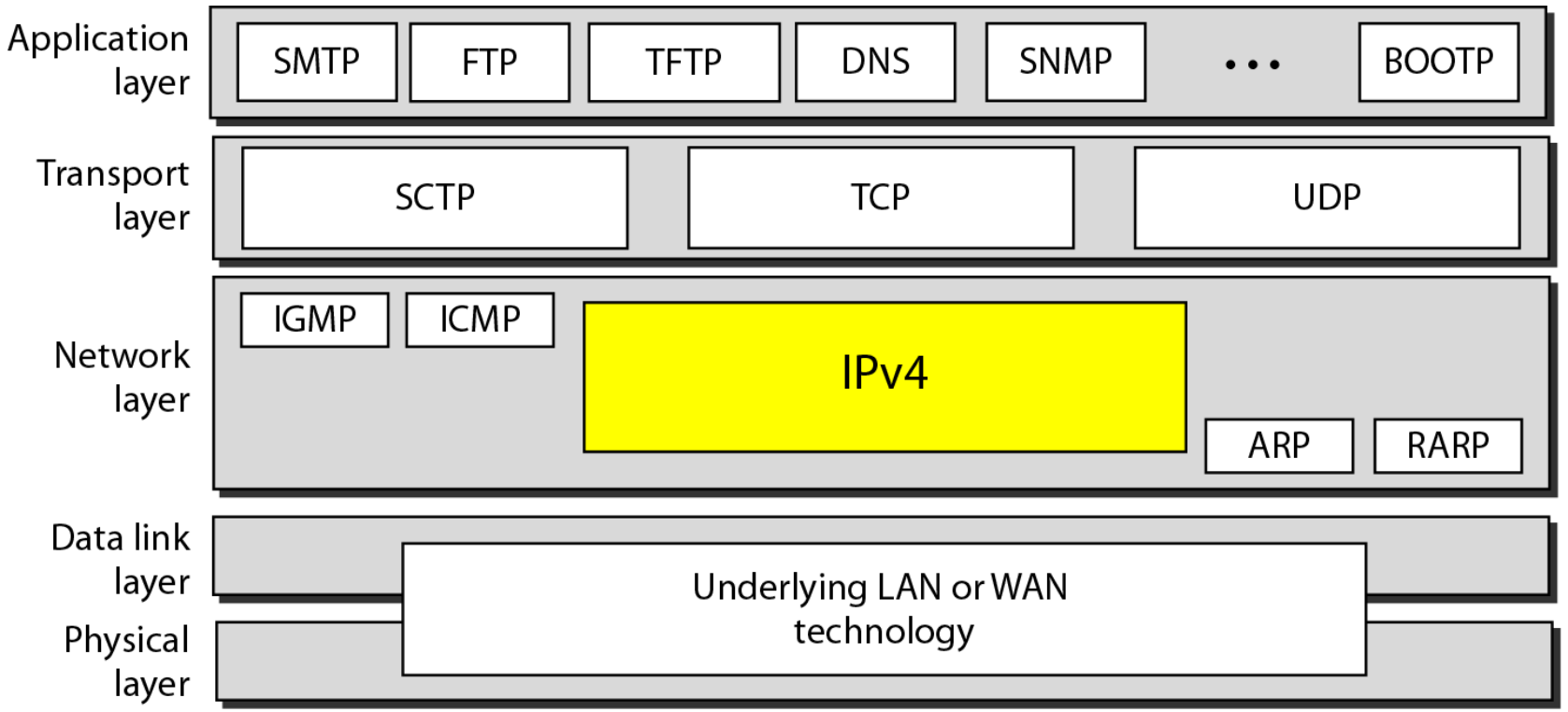
Note:

The network layer is responsible for the delivery of packets from the original source to the final destination.

Source-to-destination delivery



IPv4



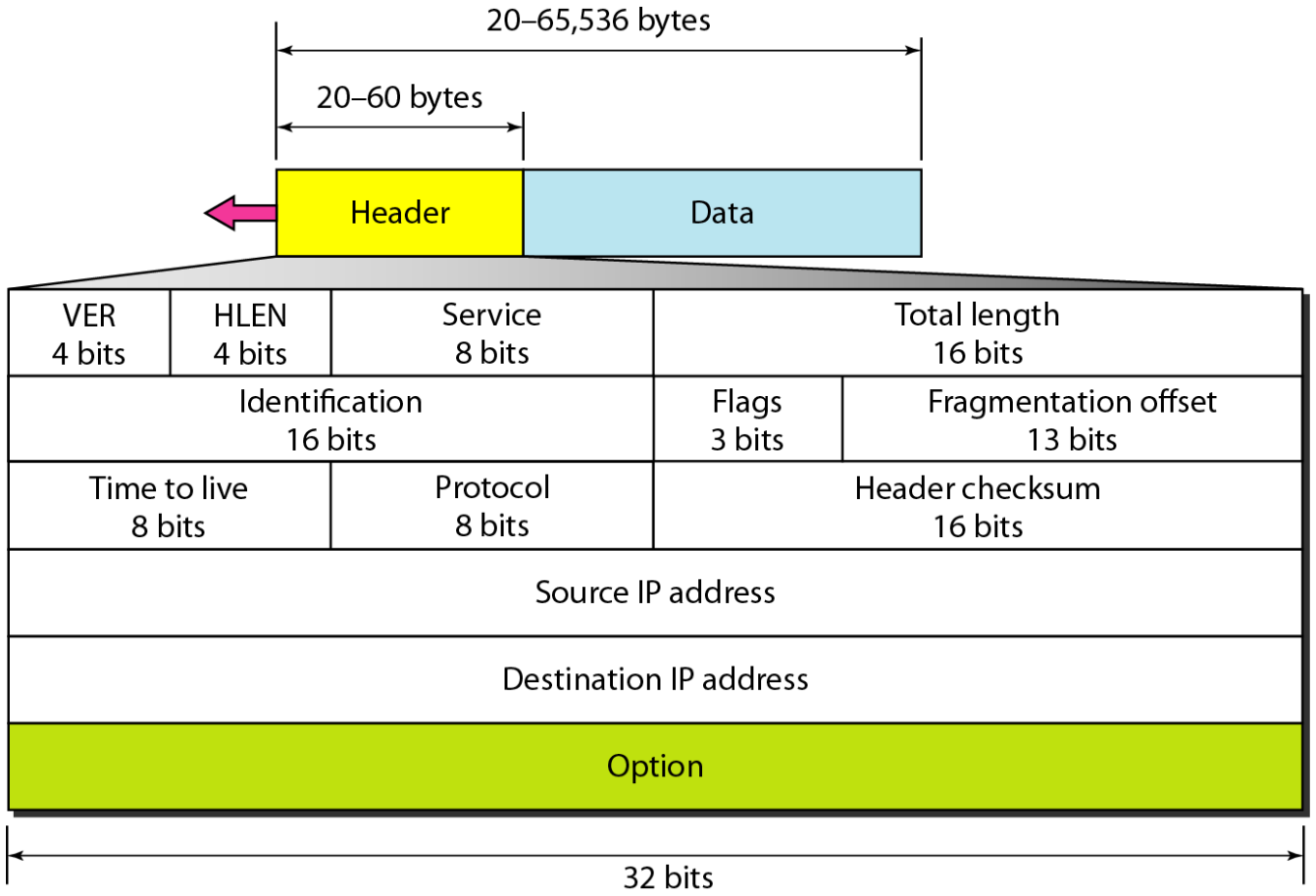
Position of IPv4 in TCP/IP protocol suite

IPv4

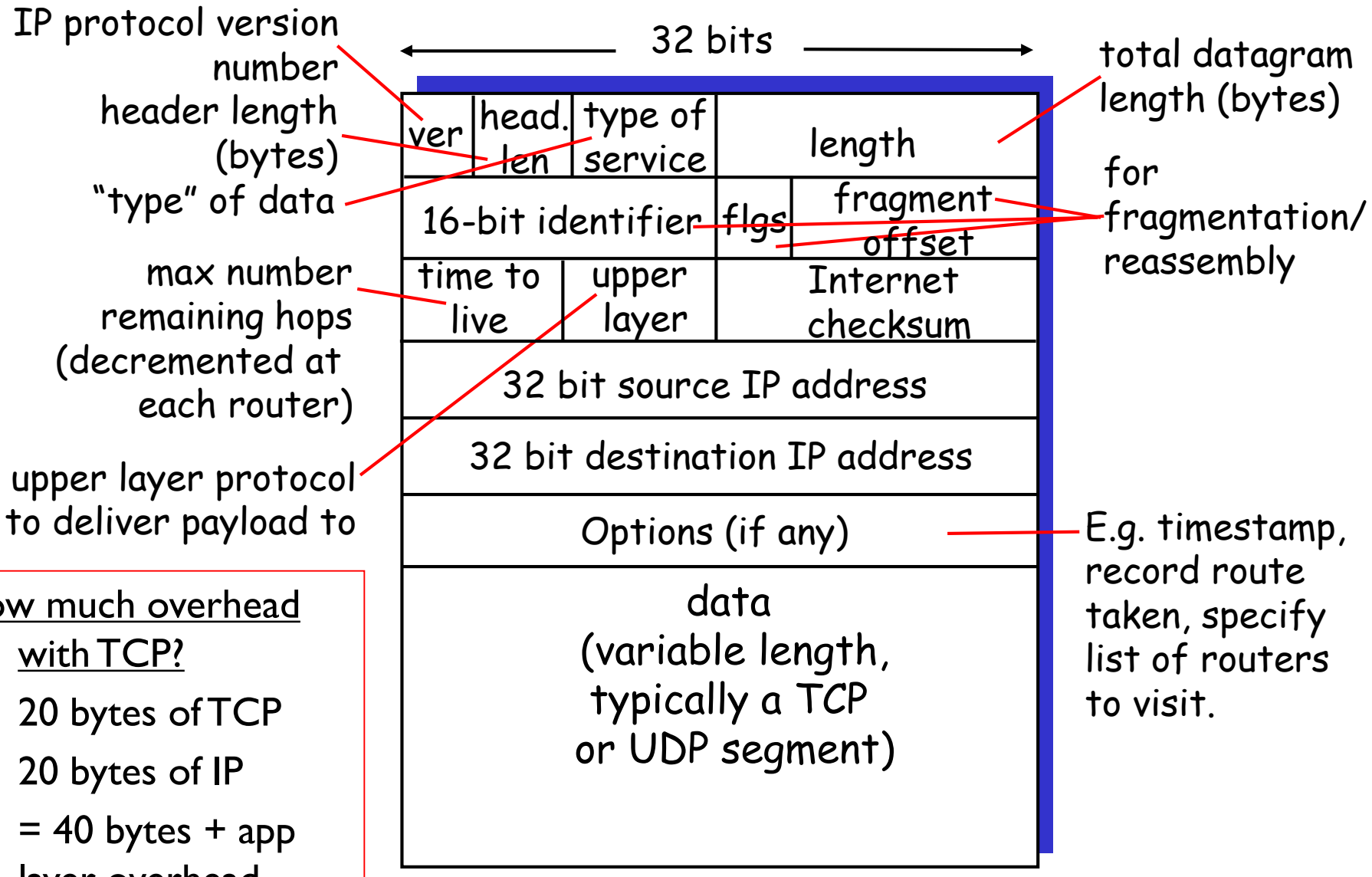
- Best-effort delivery
  - IPv4 is an unreliable and connectionless datagram protocol -
  - The term best-effort means that IPv4 provides no error control or flow control (except for error detection on the header).
- Connectionless protocol
  - Each datagram is handled independently, and diagrams sent by the source to the same destination could arrive out of order.
  - Also, some could be lost or corrupted during transmission.
  - IPv4 relies on a high-level protocol to take of all these problem.

IPv4 Datagram

- Packets in the IPv4 layer are called Datagrams.



IP datagram format



how much overhead with TCP?

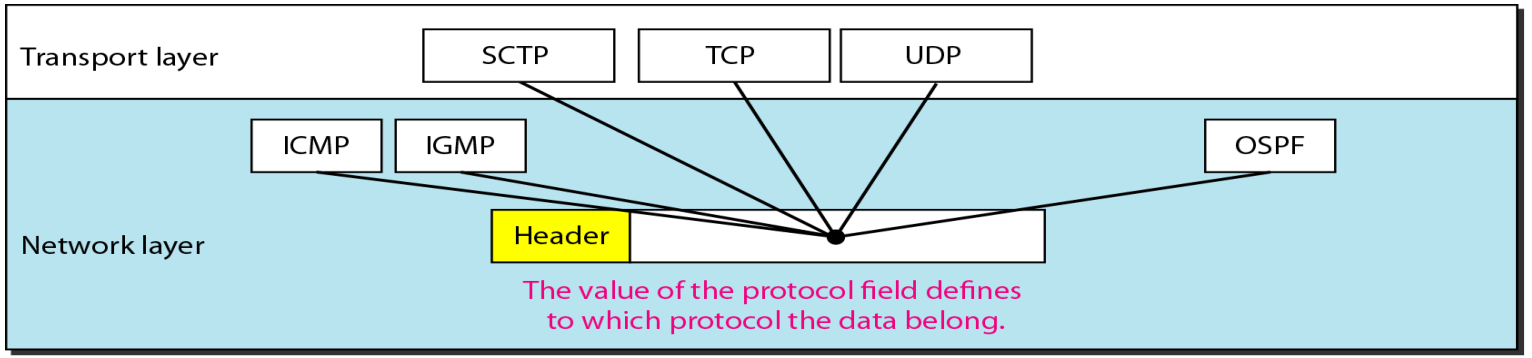
- ❑ 20 bytes of TCP
- ❑ 20 bytes of IP
- ❑ = 40 bytes + app layer overhead

IPv4 Datagram (cont'd)

- A datagram is a variable-length packet consisting of a header and data.
- Header
  - length : 20 - 60 bytes
  - Contains information essential to routing and delivery.
- Version (VER) : It defines the Version of IPv4. it is 4.
- Header Length (HLEN) : Defining the total length of the datagram header in 4byte words.

IPv4 Datagram (cont'd)

- Protocol
  - Defining the higher level protocol that uses the services of the IP layer
    - TCP, UDP, ICMP, and IGMP
    - Multiplexing data from different higher level protocols

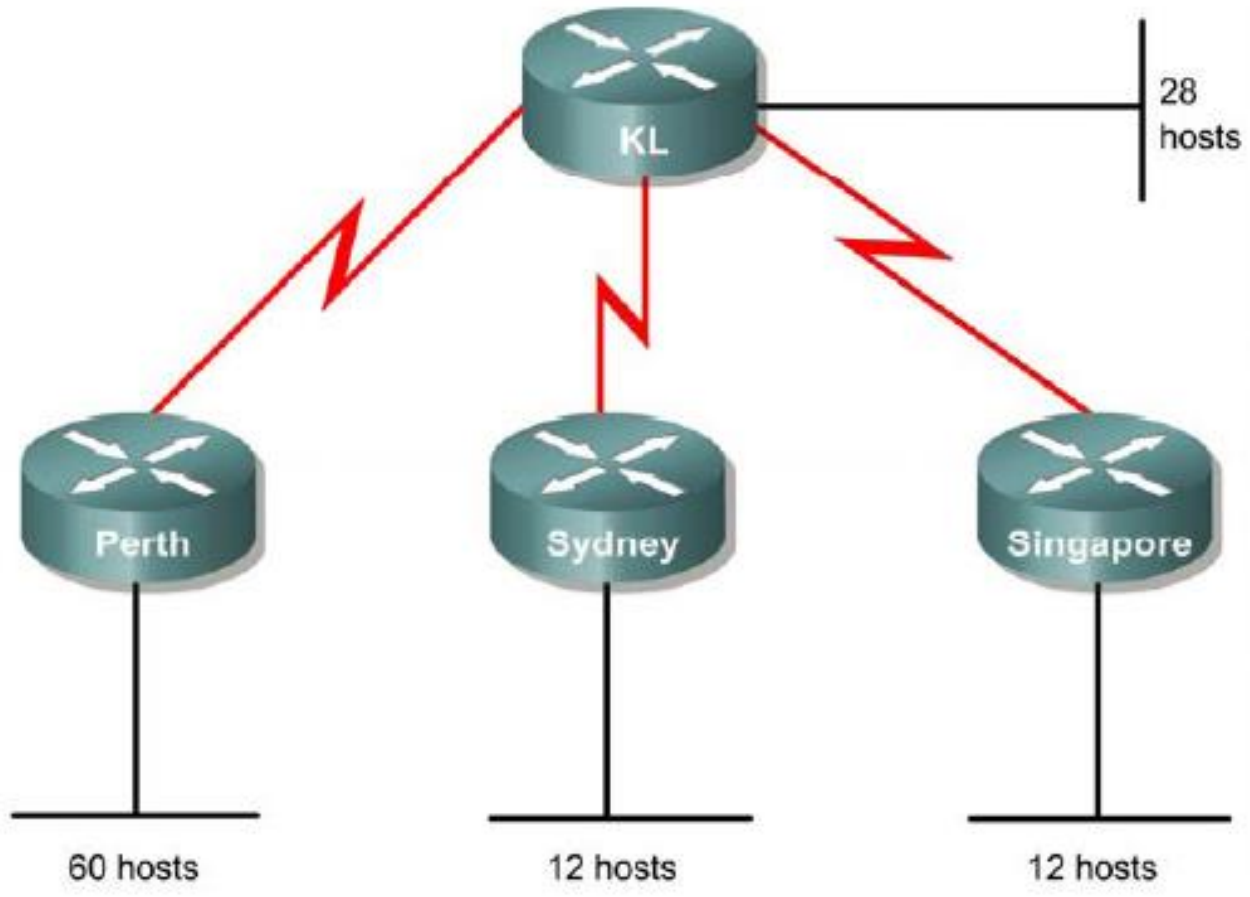


Value	Protocol
1	ICMP
2	IGMP
6	TCP
8	EGP
17	UDP
89	OSPF

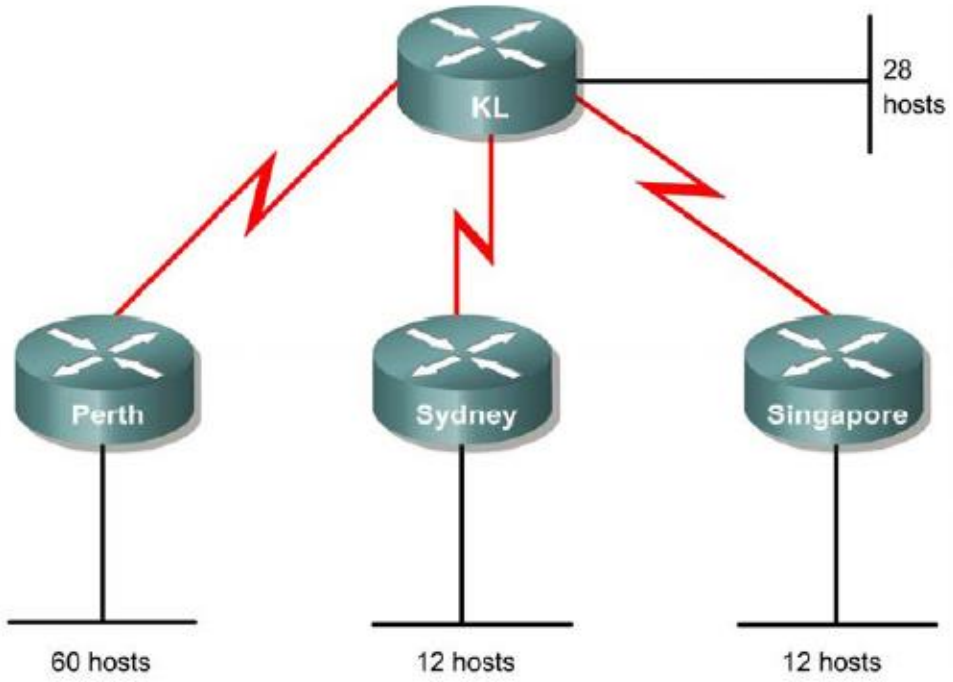
VARIABLE-LENGTH SUBNET MASK (VLSM)

Subnet with VLSM

192.168.10.0/24



Regular Subnet



192.168.10.0/24

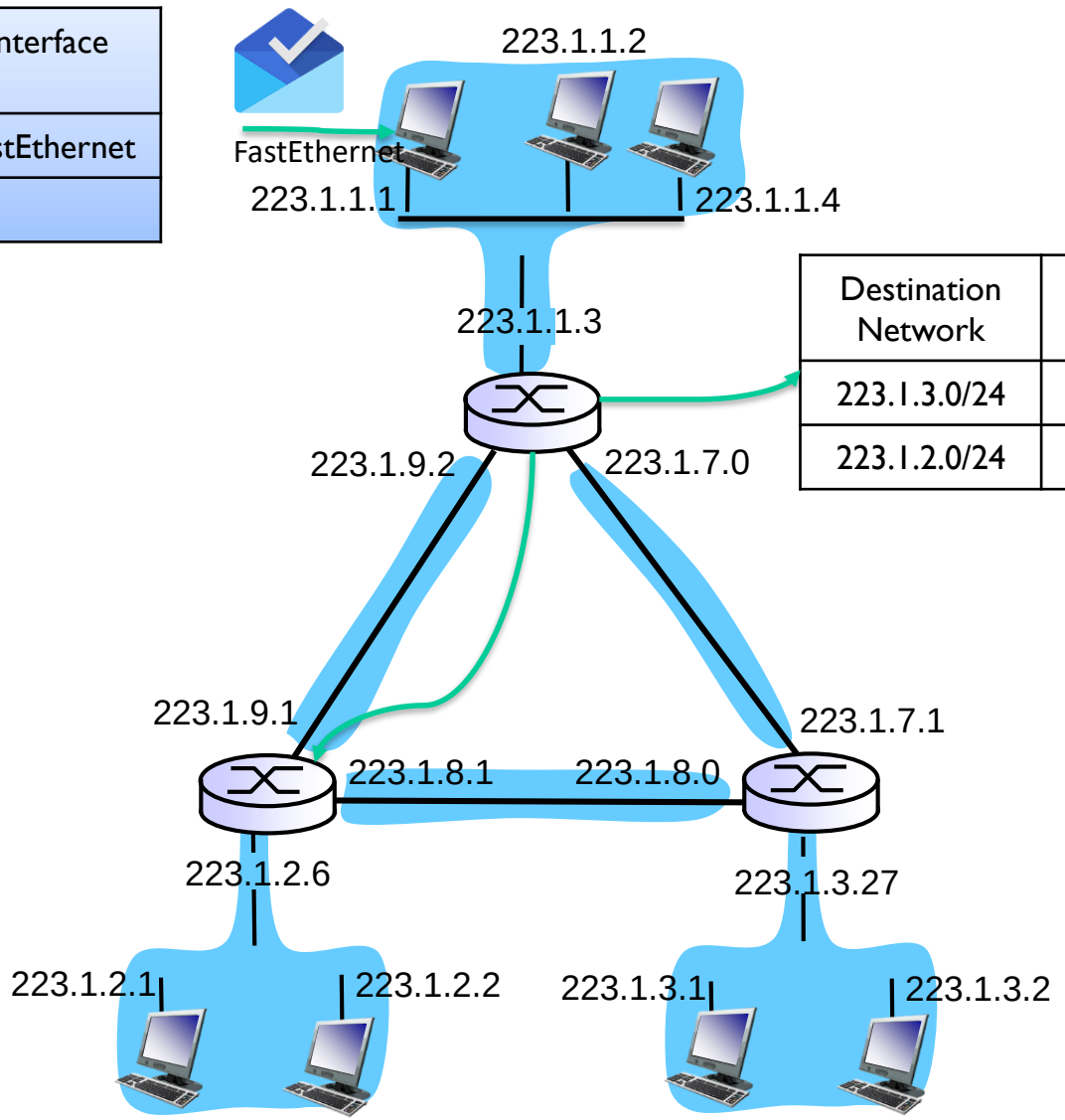
7 subnets; The largest subnet needs **60** hosts

If 3 bits for subnet (8 subnets) ➔ 5 bits for host (32 hosts)

If 6 bits for host (64 hosts) ➔ 2 bits for subnet (4 subnets)

Forwarding - Routing

Destination Network	Next hop	Interface
0.0.0.0	223.1.1.3	FastEthernet

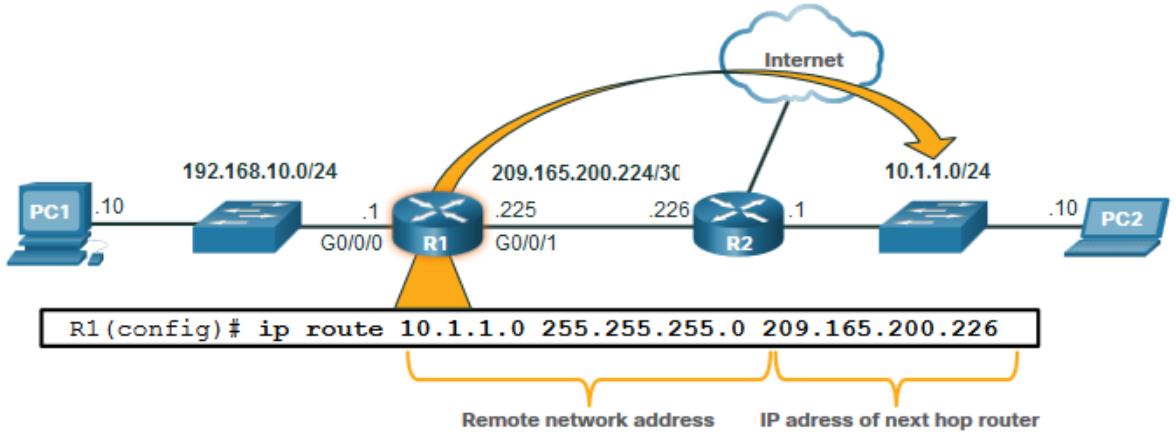


Destination Network	Next hop	Interface
223.1.3.0/24	223.1.7.1	FastEthernet0/0
223.1.2.0/24	223.1.9.1	FastEthernet0/1

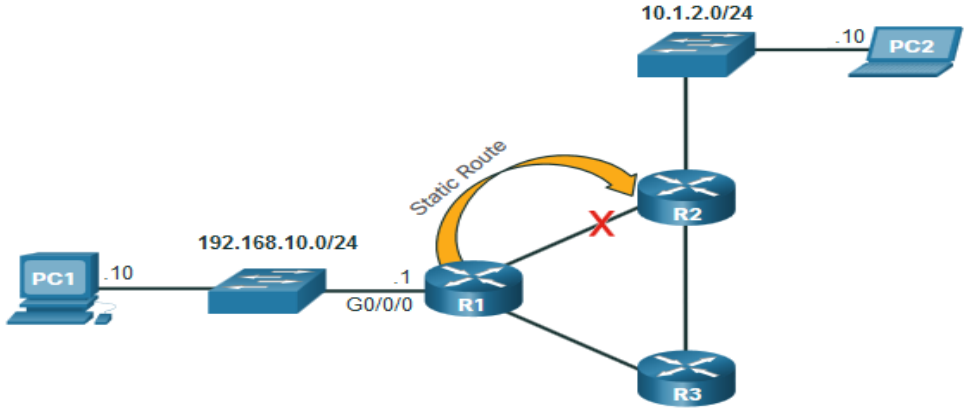
Introduction to Routing

Static Routing

- Static Route Characteristics:
  - Must be configured manually
  - Must be adjusted manually by the administrator when there is a change in the topology
  - Good for small non-redundant networks
  - Often used in conjunction with a dynamic routing protocol for configuring a default route



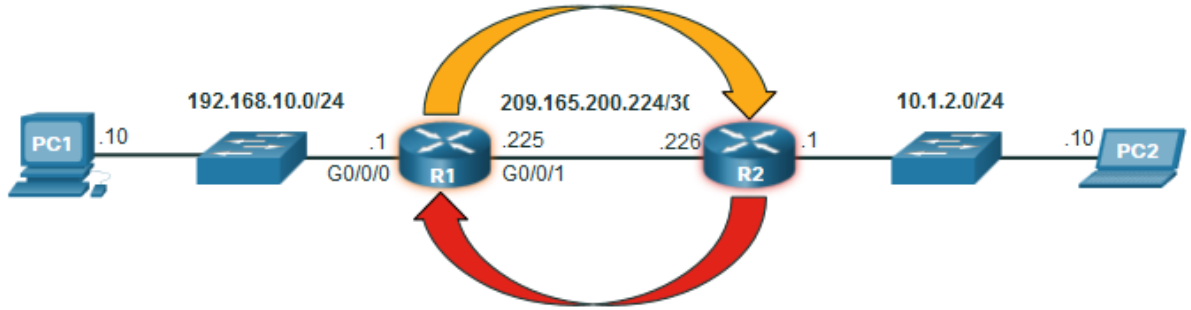
R1 is manually configured with a static route to reach the 10.1.1.0/24 network. If this path changes, R1 will require a new static route.



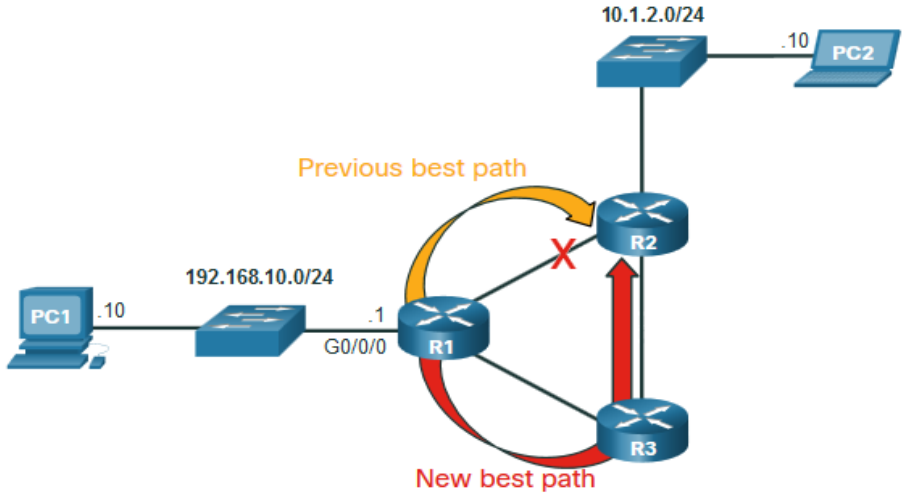
If the route from R1 via R2 is no longer available, a new static route via R3 would need to be configured. Static routes not automatically adjust for topology changes.

Introduction to Routing
Dynamic Routing

- Dynamic Routes Automatically:
  - Discover remote networks
  - Maintain up-to-date information
  - Choose the best path to the destination
  - Find new best paths when there is a topology change
  - Dynamic routing can also share static default routes with the other routers.



- R1 is using the routing protocol OSPF to let R2 know about the 192.168.10.0/24 network.
- R2 is using the routing protocol OSPF to let R1 know about the 10.1.1.0/24 network.



R1, R2, and R3 are using the dynamic routing protocol OSPF. If there is a network topology change, they can automatically adjust to find a new best path.

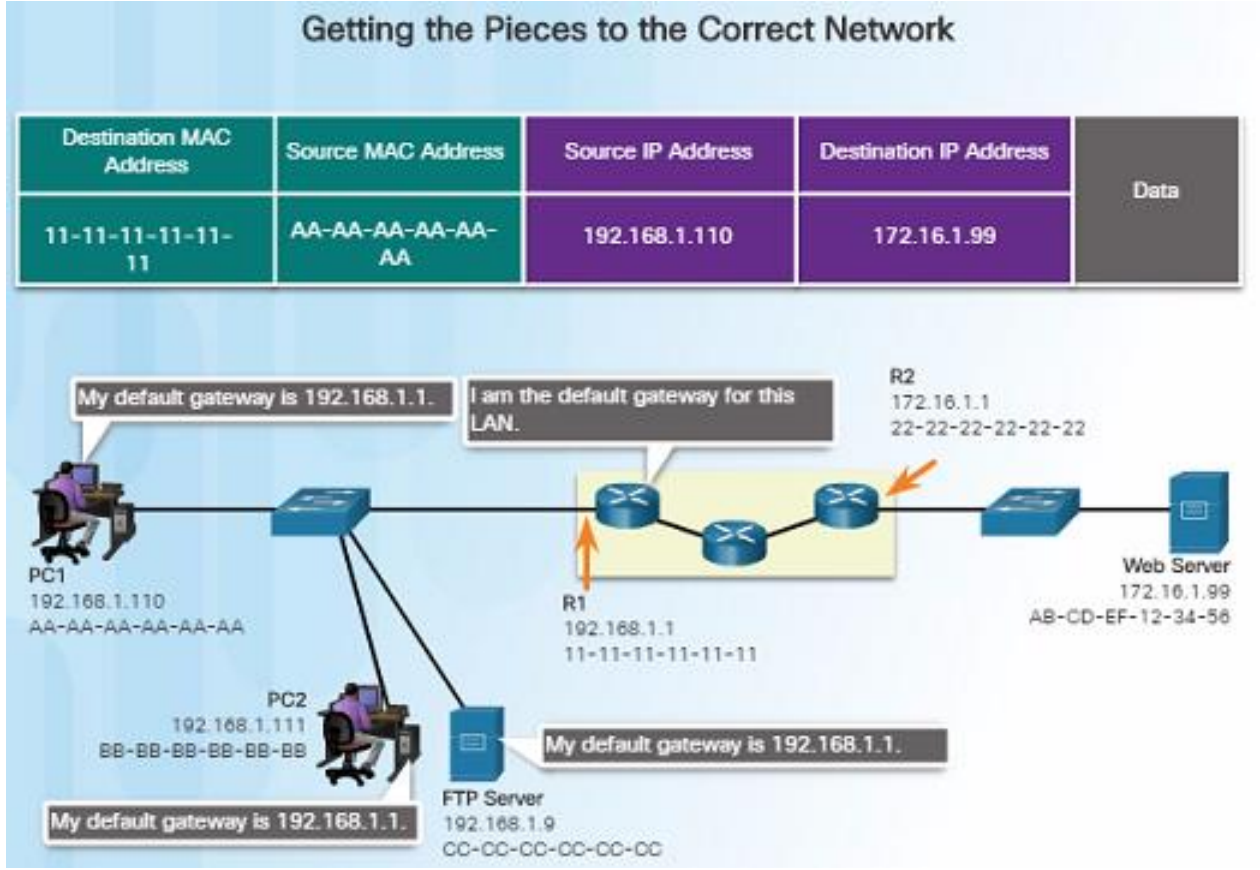
ADDRESS RESOLUTION - ARP

MAC addresses and ARP

- 32-bit IP address:
  - network-layer address for interface
  - used for layer 3 (network layer) forwarding
- MAC (or LAN or physical or Ethernet) address:
  - function: used 'locally" to get frame from one interface to another physically-connected interface (same network, in IP-addressing sense)
  - 48 bit MAC address (for most LANs) burned in NIC ROM, also sometimes software settable
  - e.g.: 1A-2F-BB-76-09-AD

hexadecimal (base 16) notation
(each "numeral" represents 4 bits)

Default Gateways

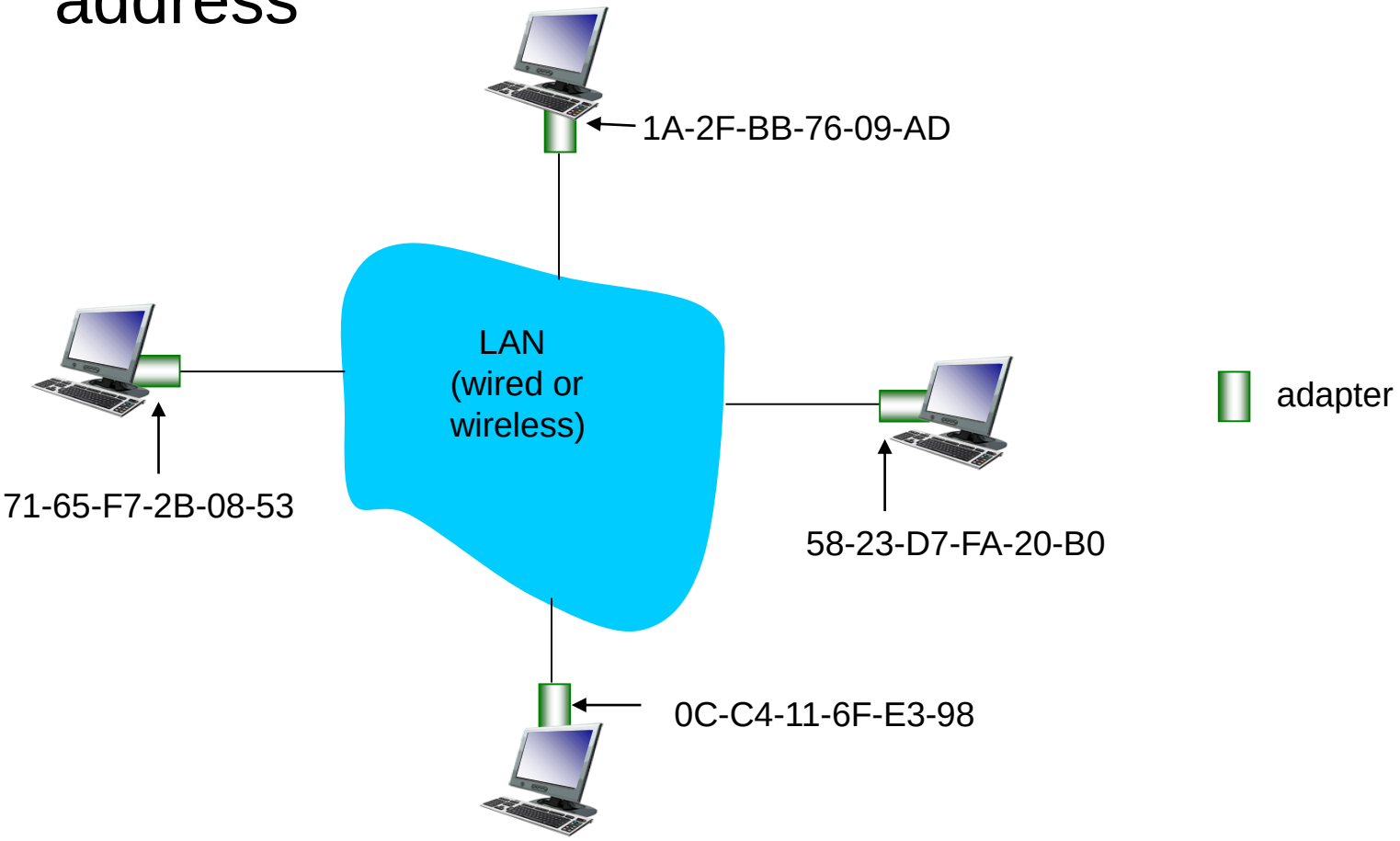


- Routers are also usually configured with their own default gateway.

- Devices need the following information for network access: IP address, subnet mask, and default gateway.
- When a host sends a packet to a device that is on the same IP network, the packet is forwarded out the host interface to the destination device. The router does not need to get involved.
- When a host sends a packet to a device on a different IP network, the packet is forwarded to the default gateway because the host device cannot communicate with devices outside of the local network.
- The default gateway is the device that routes traffic from the local network to devices on remote networks, such as devices on the Internet.

LAN addresses and ARP

each adapter on LAN has unique *LAN* address



LAN addresses (more)

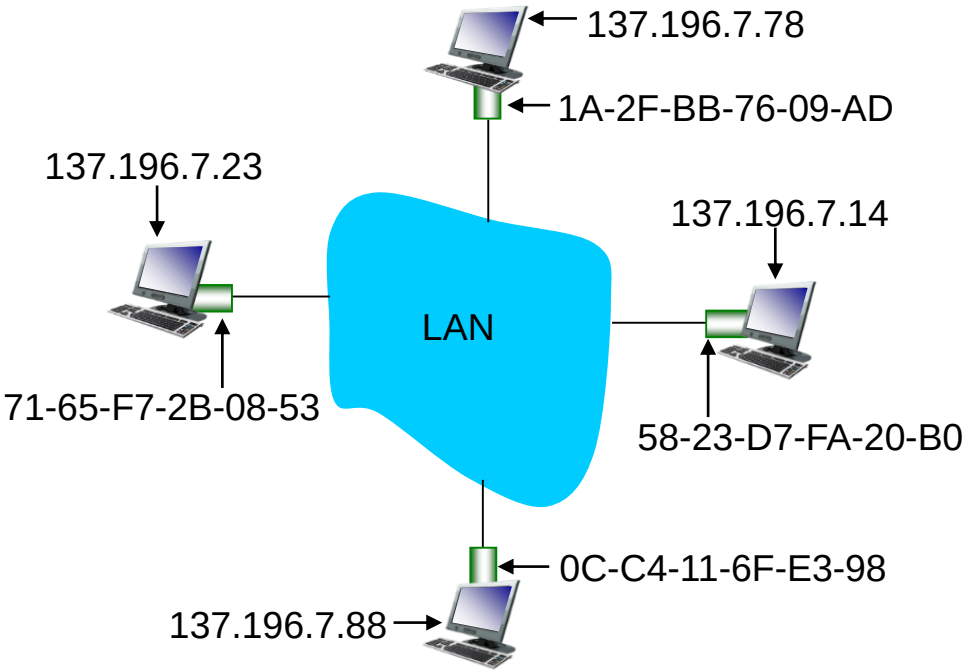
- MAC address allocation administered by IEEE
- manufacturer buys portion of MAC address space (to assure uniqueness)
- analogy:
  - MAC address: like Social Security Number
  - IP address: like postal address
- MAC flat address ➜ portability
  - can move LAN card from one LAN to another
- IP hierarchical address not portable
  - address depends on IP subnet to which node is attached

ARP: address resolution protocol

- ARP table: each IP node (host, router) on LAN has table
  - IP/MAC address mappings for some LAN nodes:
- < IP address; MAC address; TTL>
  - TTL (Time To Live): time after which address mapping will be forgotten (typically 20 min)

IP address	MAC address	TTL
137.196.7.14	58-23-D7-FA-20-B0	20

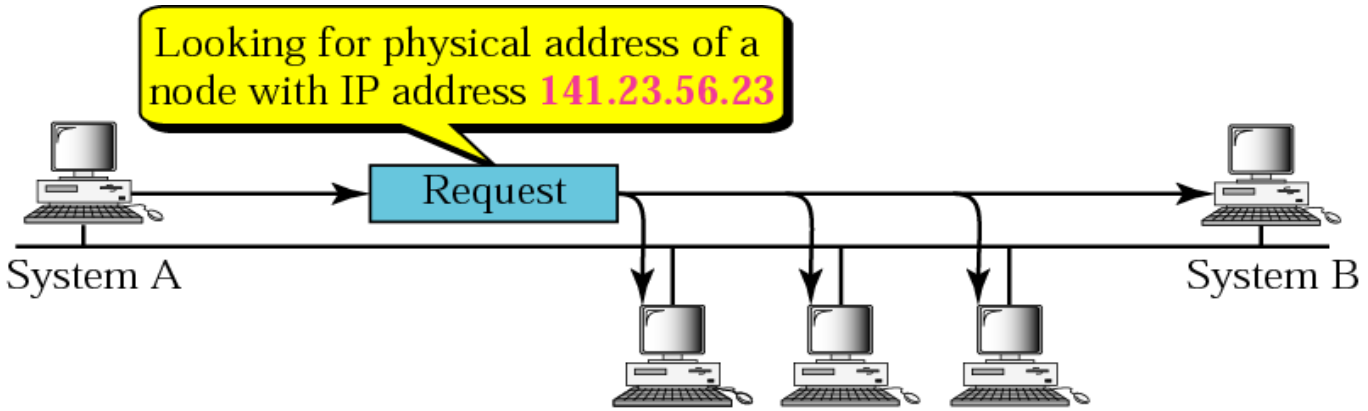
Question: how to determine interface's MAC address, knowing its IP address?



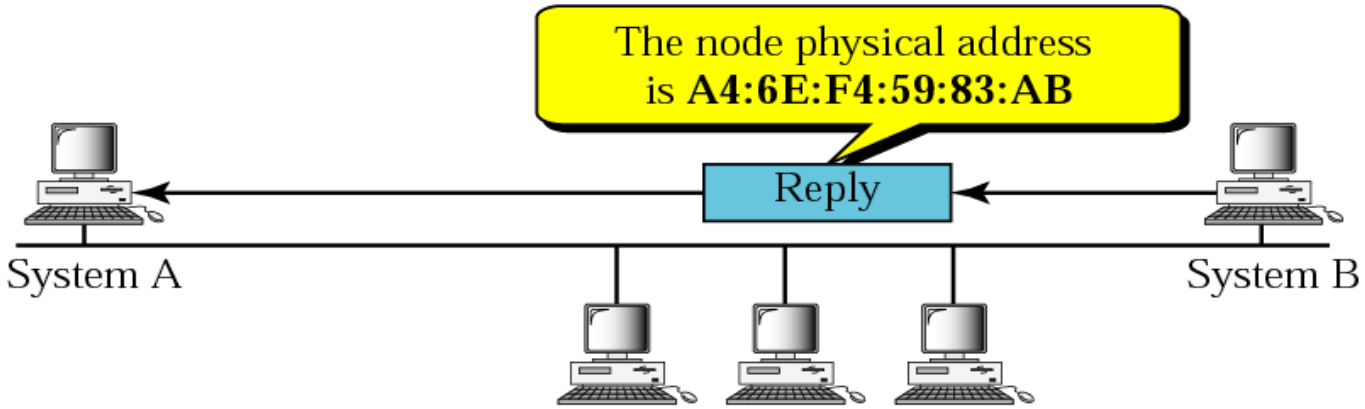
ARP protocol: same LAN

- A wants to send datagram to B
  - B's MAC address not in A's ARP table.
- A broadcasts ARP query packet, containing B's IP address
  - destination MAC address = FF-FF-FF-FF-FF-FF
  - all nodes on LAN receive ARP query
- B receives ARP packet, replies to A with its (B's) MAC address
  - frame sent to A's MAC address (unicast)
- A caches (saves) IP-to-MAC address pair in its ARP table until information becomes old (times out)
  - soft state: information that times out (goes away) unless refreshed
- ARP is "plug-and-play":
  - nodes create their ARP tables without intervention from net administrator

ARP operation

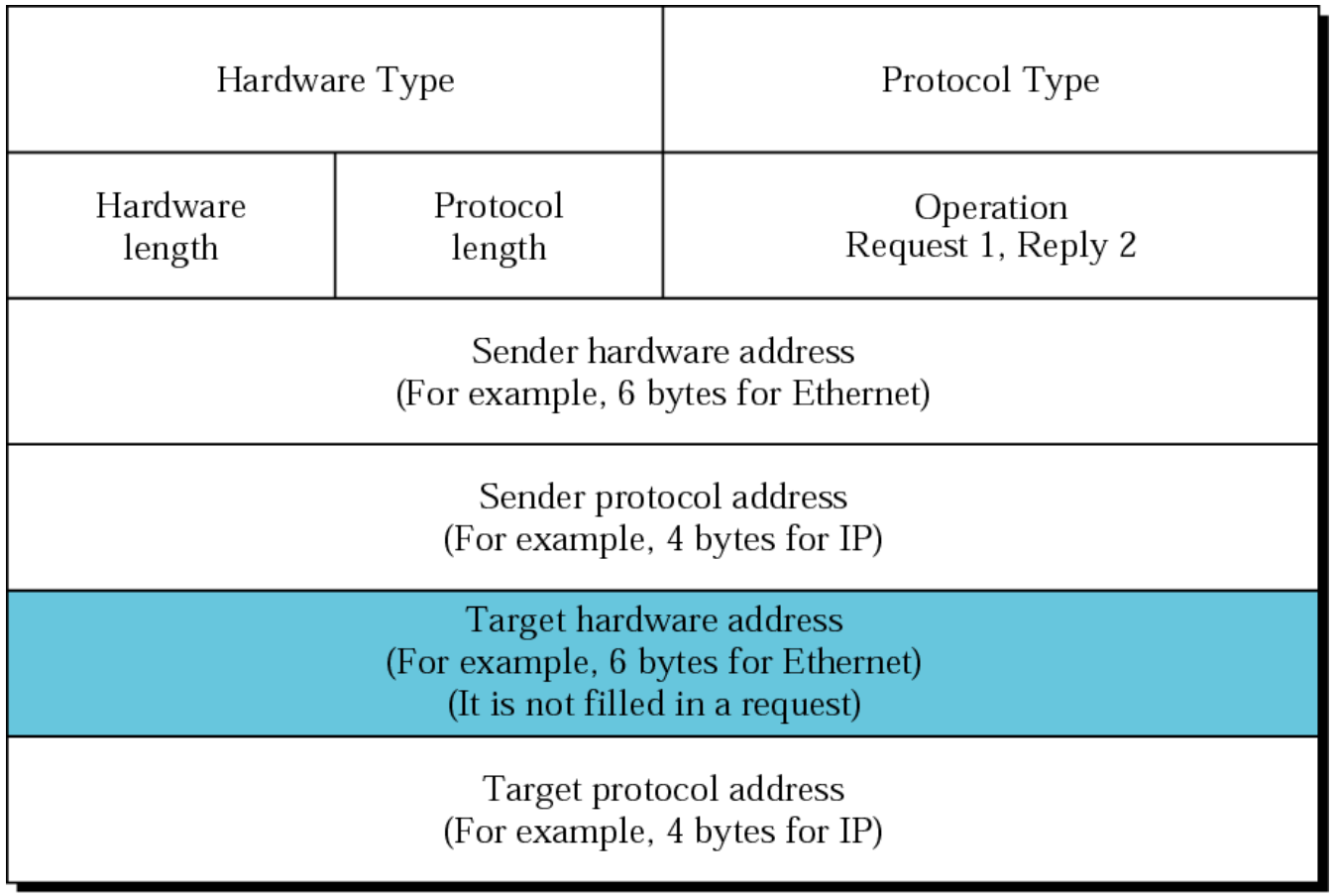


a. ARP request is broadcast



b. ARP reply is unicast

ARP packet



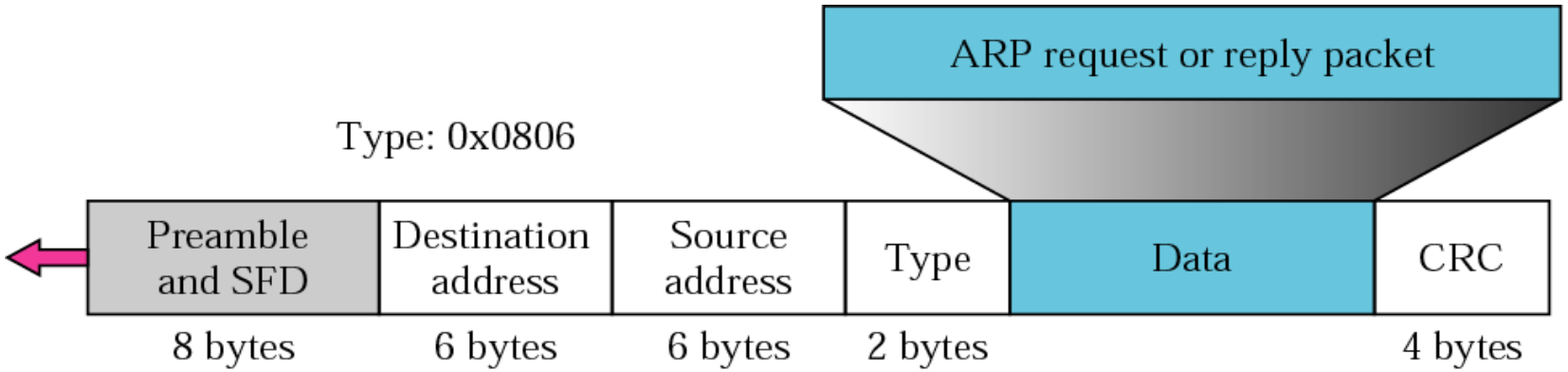
Hardware Type - Ethernet is type 1

Protocol Type- IPv4=x0800

Hardware Length:length of Ethernet Address (6)

Protocol Length:length of IPv4 address (4)

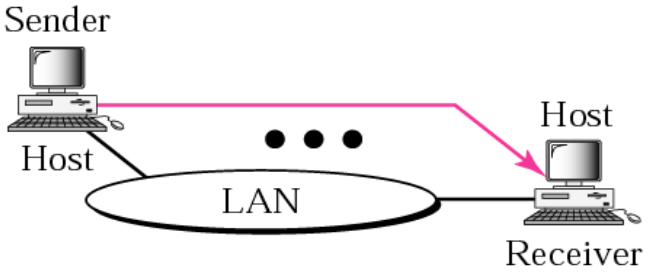
Encapsulation of ARP packet



The ARP packet is encapsulated within an Ethernet packet.
Note: Type field for Ethernet is x0806

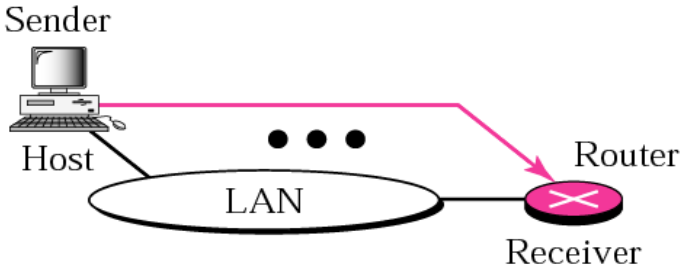
Four cases using ARP

Target IP address:
Destination address in the IP datagram



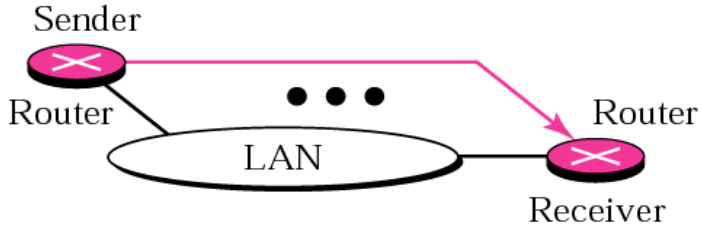
Case 1. A host has a packet to send to another host on the same network.

Target IP address:
IP address of a router



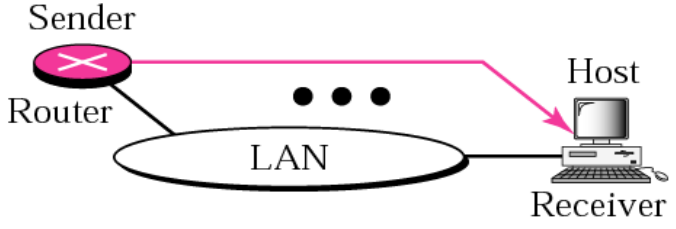
Case 2. A host wants to send a packet to another host on another network.
It must first be delivered to a router.

Target IP address:
IP address of the appropriate router found in the routing table



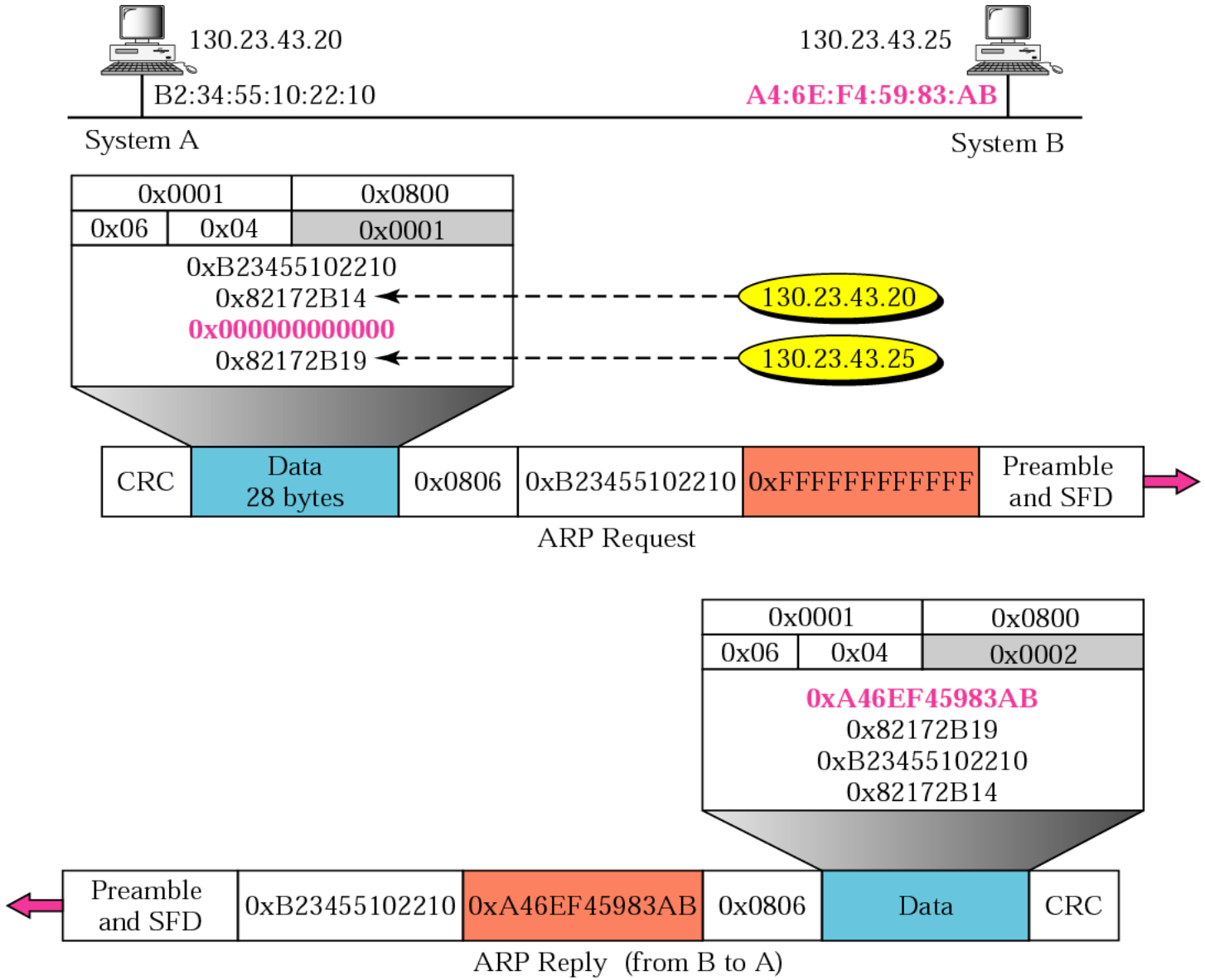
Case 3. A router receives a packet to be sent to a host on another network.
It must first be delivered to the appropriate router.

Target IP address:
Destination address in the IP datagram

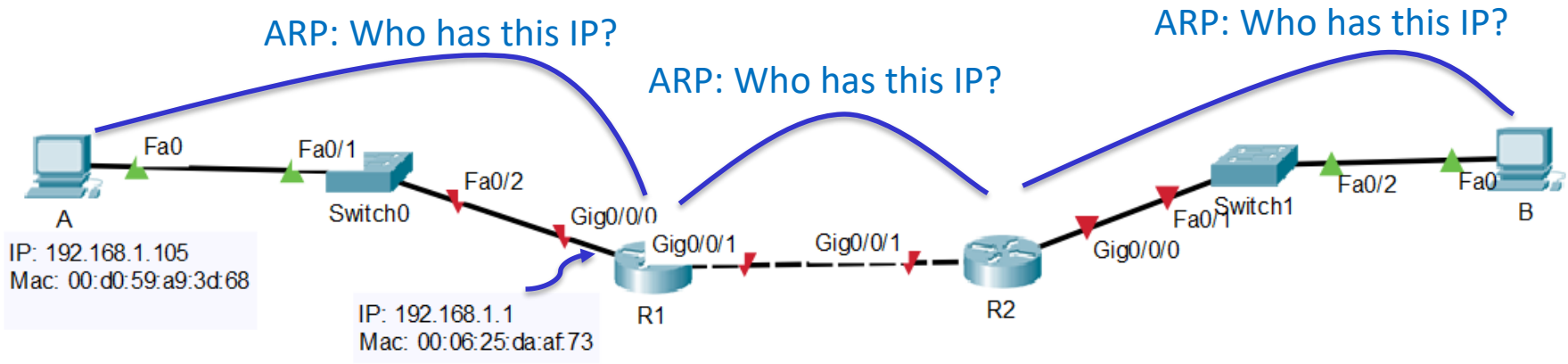


Case 4. A router receives a packet to be sent to a host on the same network.

Example 1


130.23.43.20
130.23.43.25
B2:34:55:10:22:10
A4:6E:F4:59:83:AB
System A
System B
0x0001
0x0800
0x06
0x04
0x0001
0xB23455102210
0x82172B14
0x000000000000
0x82172B19
130.23.43.20
130.23.43.25
CRC
Data
28 bytes
0x0806
0xB23455102210
0xFFFFFFFFFFFF
Preamble
and SFD
ARP Request
0x0001
0x0800
0x06
0x04
0x0002
0xA46EF45983AB
0x82172B19
0xB23455102210
0x82172B14
Preamble
and SFD
0xB23455102210
0xA46EF45983AB
0x0806
Data
CRC
ARP Reply (from B to A)

ARP – A each step



ethernet-ethereal-trace-1

File Edit View Go Capture Analyze Statistics Telephony Wireless Tools Help

arp

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000	AmbitMic_a9:3…	Broadcast	ARP	42	Who has 192.168.1.1? Tell 192.168.1.105
2	0.001018	LinksysG_da:a…	AmbitMic_a9:3d…	ARP	60	192.168.1.1 is at 00:06:25:da:af:73
6	13.542…	CnetTech_73:8…	Broadcast	ARP	60	Who has 192.168.1.117? Tell 192.168.1.104

```
› Frame 1: 42 bytes on wire (336 bits), 42 bytes captured (336 bits)
⌄ Ethernet II, Src: AmbitMic_a9:3d:68 (00:d0:59:a9:3d:68), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
  › Destination: Broadcast (ff:ff:ff:ff:ff:ff)
  › Source: AmbitMic_a9:3d:68 (00:d0:59:a9:3d:68)
    Type: ARP (0x0806)
⌄ Address Resolution Protocol (request)
    Hardware type: Ethernet (1)
    Protocol type: IPv4 (0x0800)
    Hardware size: 6
    Protocol size: 4
    Opcode: request (1)
    Sender MAC address: AmbitMic_a9:3d:68 (00:d0:59:a9:3d:68)
    Sender IP address: 192.168.1.105
    Target MAC address: 00:00:00_00:00:00 (00:00:00:00:00:00)
    Target IP address: 192.168.1.1
```

ethernet-ethereal-trace-1

File Edit View Go Capture Analyze Statistics Telephony Wireless Tools Help

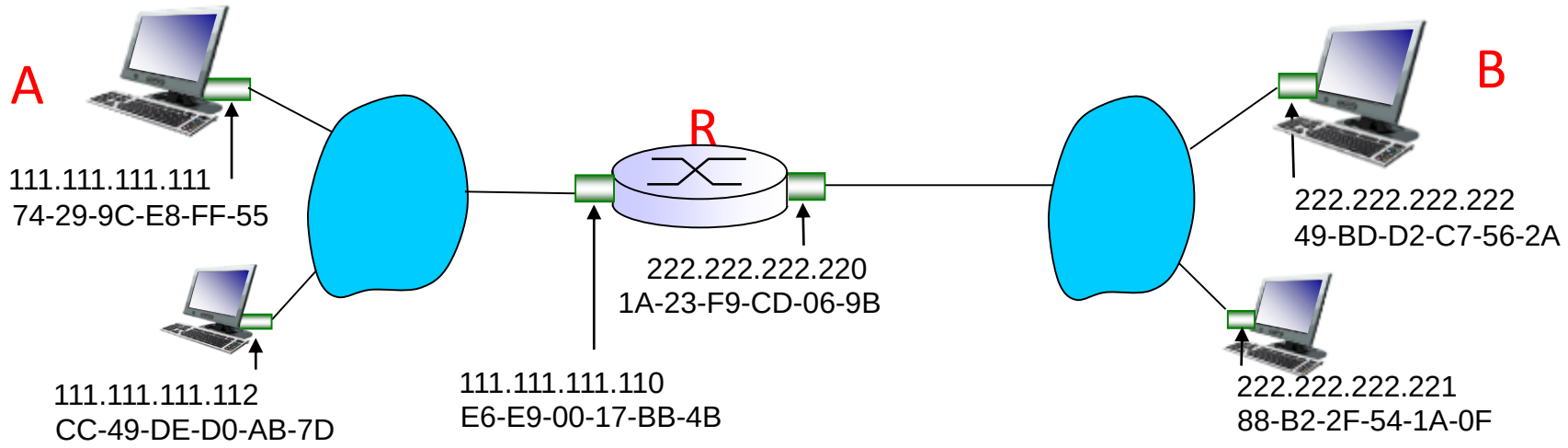
arp

No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000	AmbitMic_a9:3…	Broadcast	ARP	42	Who has 192.168.1.1? Tell 192.168.1.105
2	0.001018	LinksysG_da:a…	AmbitMic_a9:3d…	ARP	60	192.168.1.1 is at 00:06:25:da:af:73
6	13.542…	CnetTech_73:8…	Broadcast	ARP	60	Who has 192.168.1.117? Tell 192.168.1.104

```
› Frame 2: 60 bytes on wire (480 bits), 60 bytes captured (480 bits)
⌄ Ethernet II, Src: LinksysG_da:af:73 (00:06:25:da:af:73), Dst: AmbitMic_a9:3d:68 (00:d0:59:a9:3d:68)
  › Destination: AmbitMic_a9:3d:68 (00:d0:59:a9:3d:68)
  › Source: LinksysG_da:af:73 (00:06:25:da:af:73)
    Type: ARP (0x0806)
    Padding: 000000000000000000000000000000000000
⌄ Address Resolution Protocol (reply)
    Hardware type: Ethernet (1)
    Protocol type: IPv4 (0x0800)
    Hardware size: 6
    Protocol size: 4
    Opcode: reply (2)
    Sender MAC address: LinksysG_da:af:73 (00:06:25:da:af:73)
    Sender IP address: 192.168.1.1
    Target MAC address: AmbitMic_a9:3d:68 (00:d0:59:a9:3d:68)
    Target IP address: 192.168.1.105
```

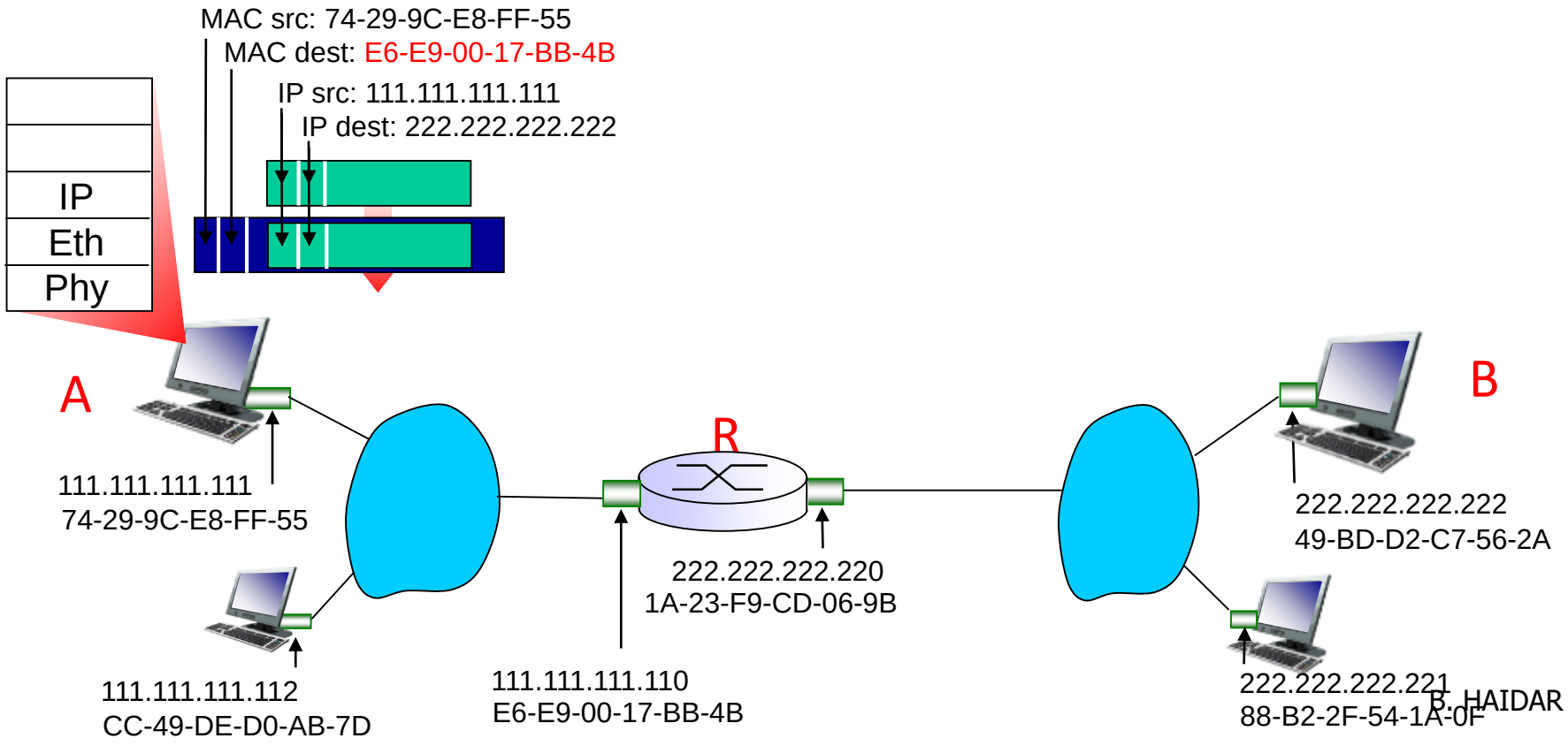
Addressing: routing to another LAN

- walkthrough: send datagram from A to B via R
  - focus on addressing – at IP (datagram) and MAC layer (frame)
  - assume A knows B's IP address
  - assume A knows IP address of first hop router, R (how?)
  - assume A knows R's MAC address (how?)



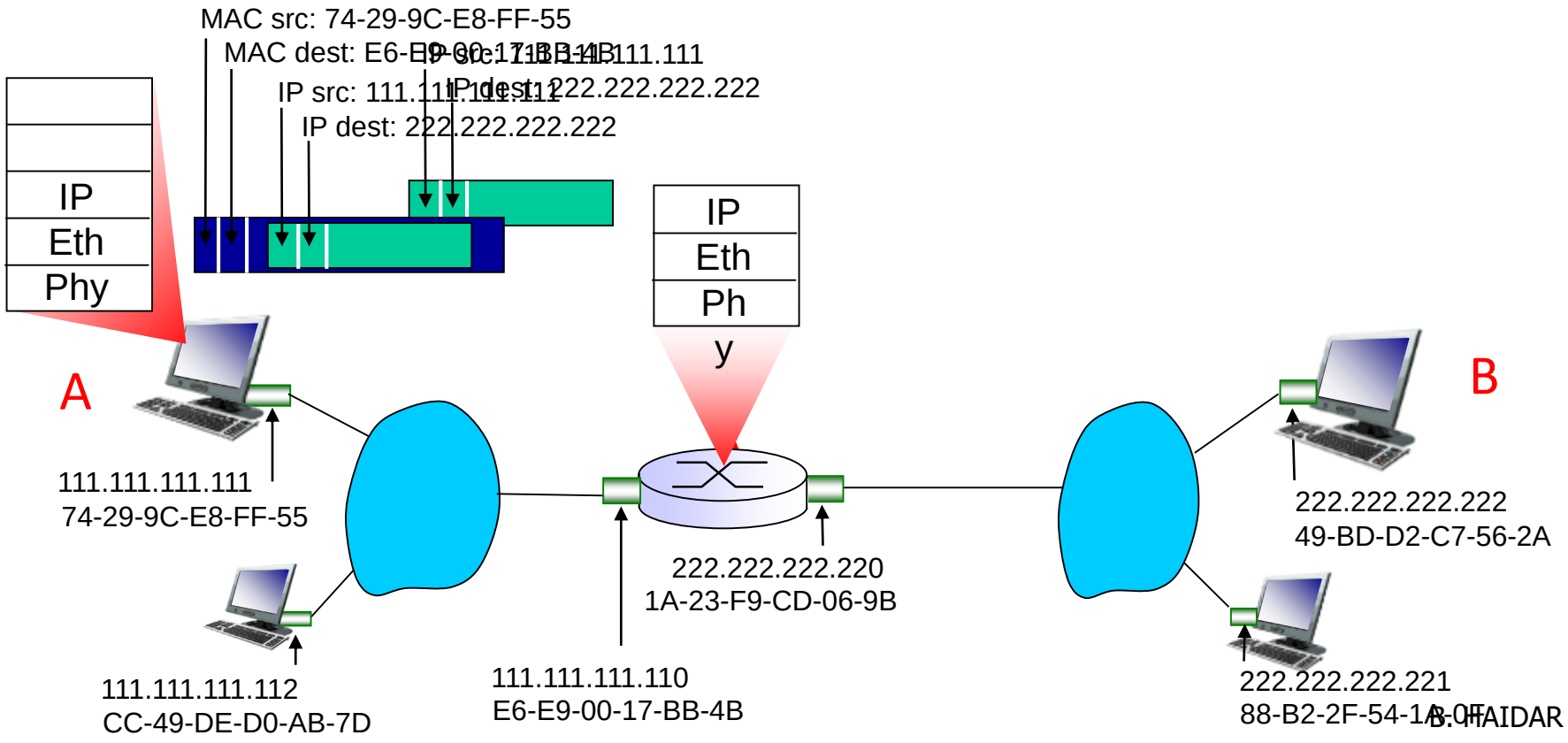
Addressing: routing to another LAN

- A creates IP datagram with IP source A, destination B
- A creates link-layer frame with R's MAC address as destination address, frame contains A-to-B IP datagram



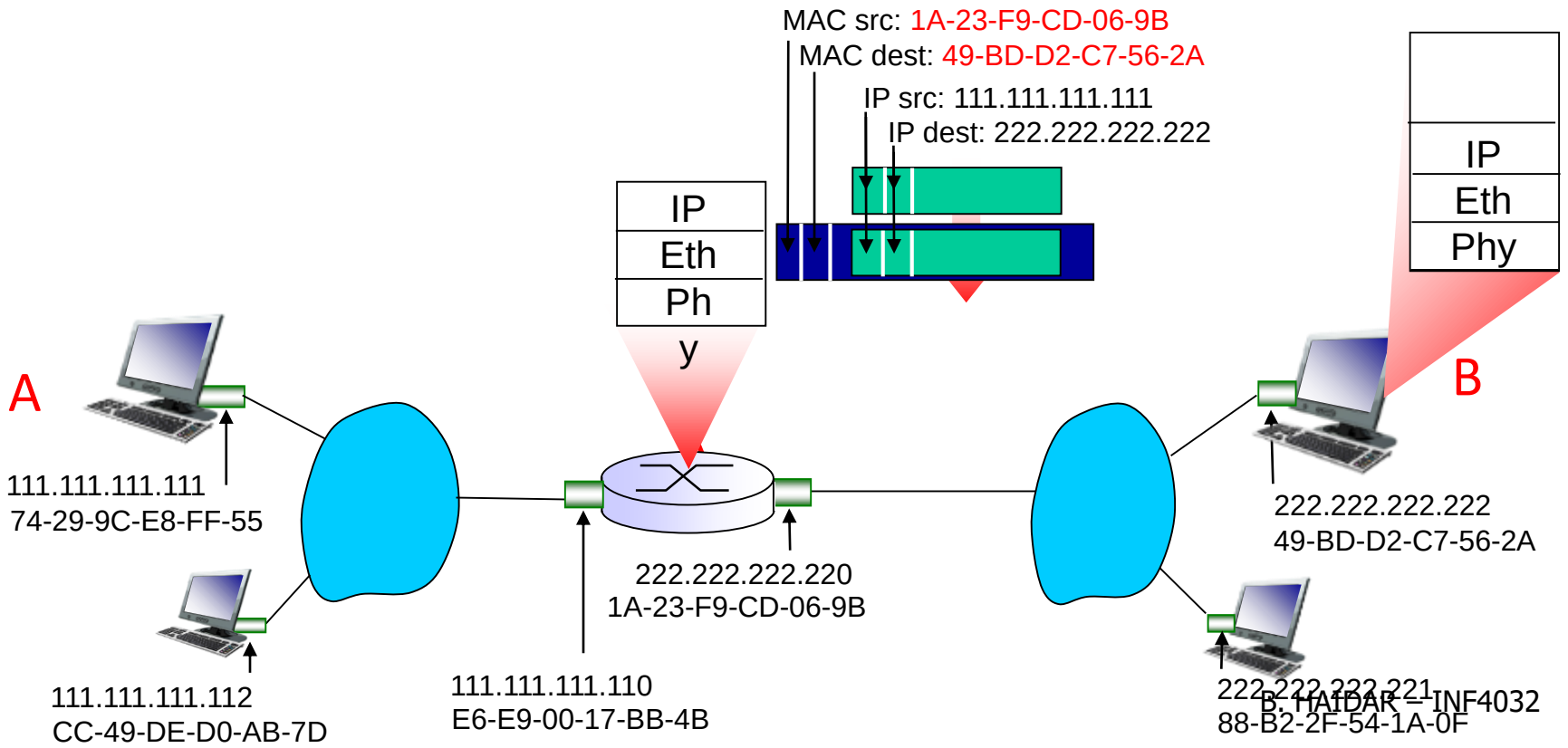
Addressing: routing to another LAN

- frame sent from A to R
- frame received at R, datagram removed, passed up to IP



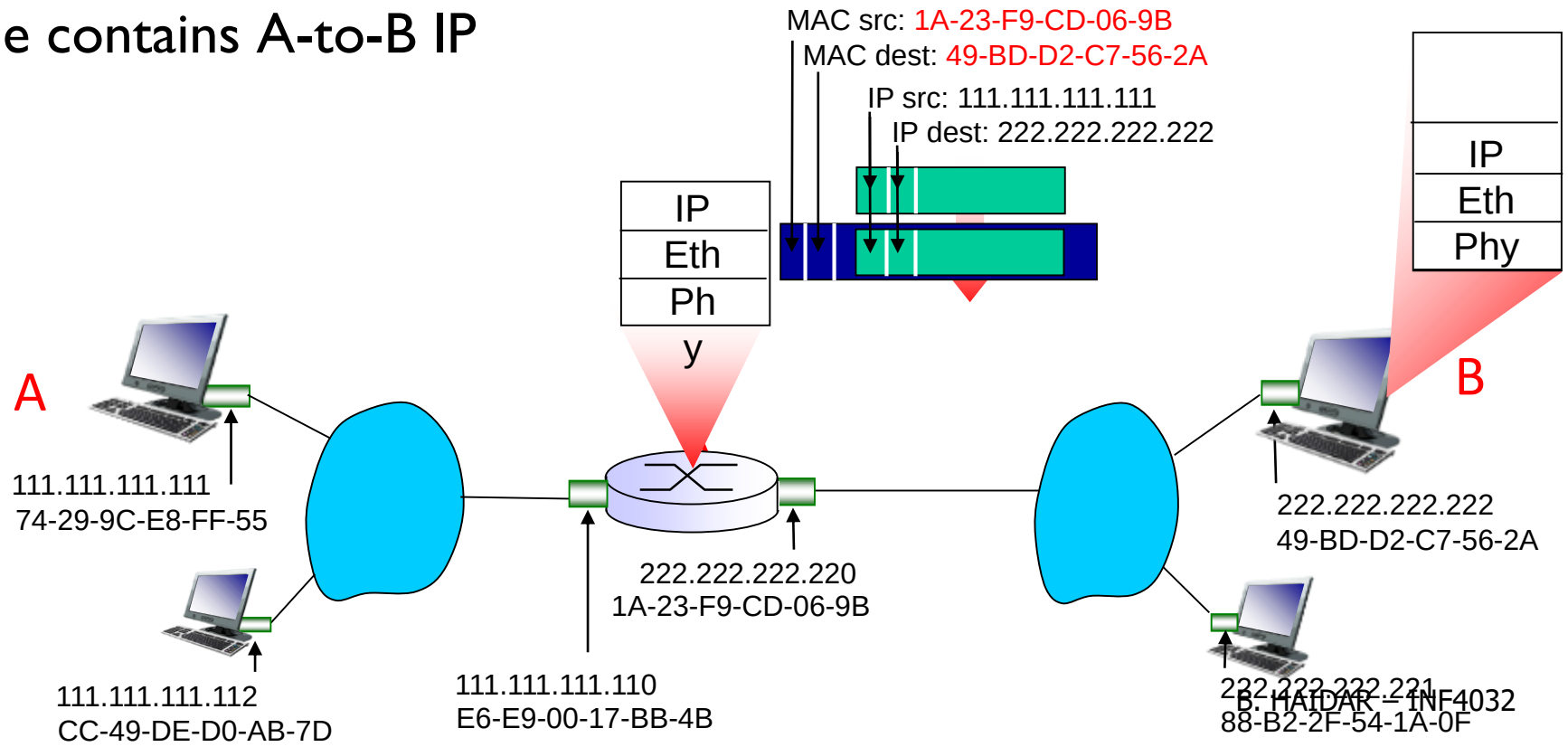
Addressing: routing to another LAN

- ⑩ R forwards datagram with IP source A, destination B
- ⑩ R creates link-layer frame with B's MAC address as destination address, frame contains A-to-B IP datagram



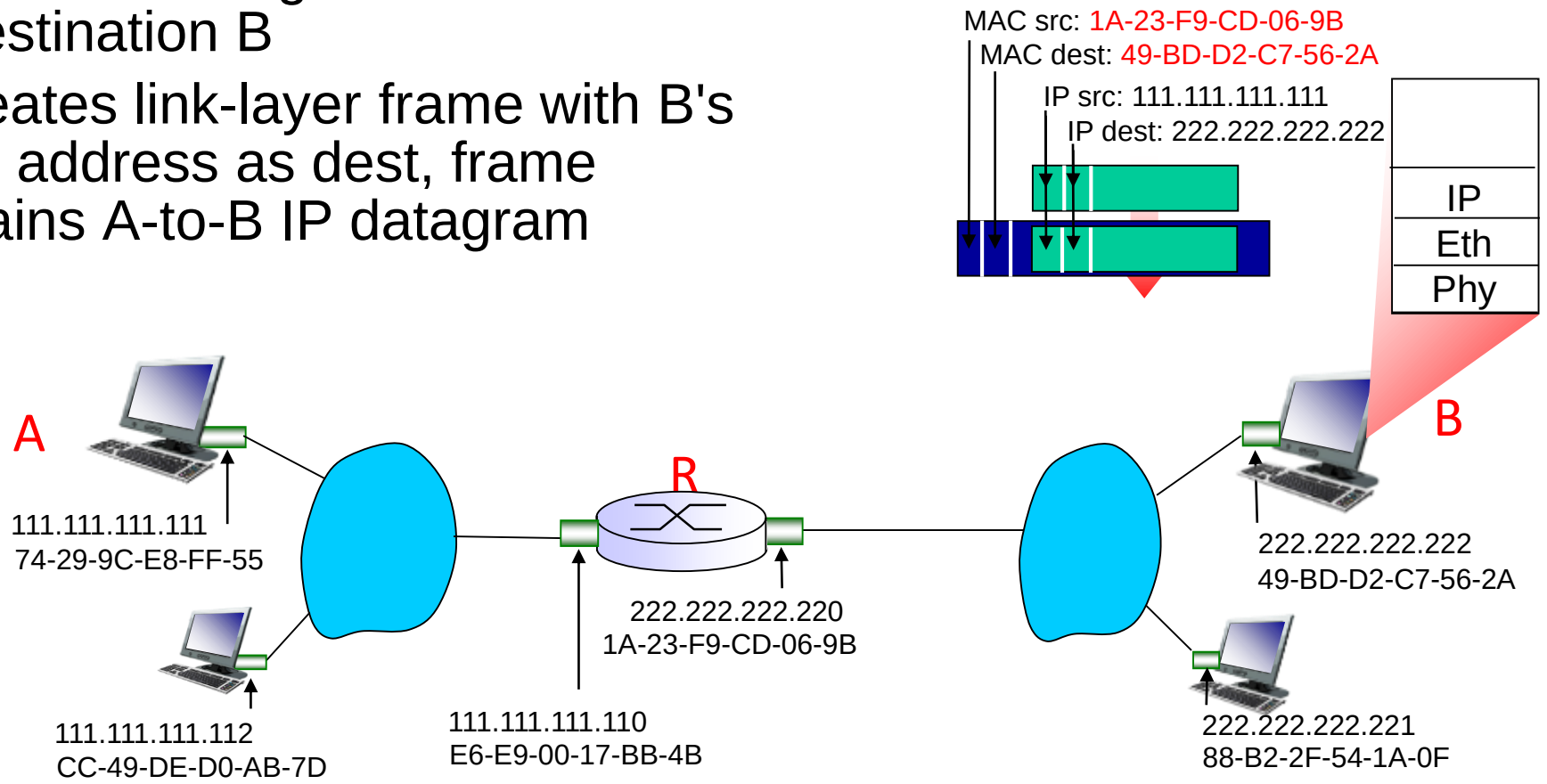
Addressing: routing to another LAN

- ❿ R forwards datagram with IP source A, destination B
- ❿ R creates link-layer frame with B's MAC address as destination address, frame contains A-to-B IP datagram



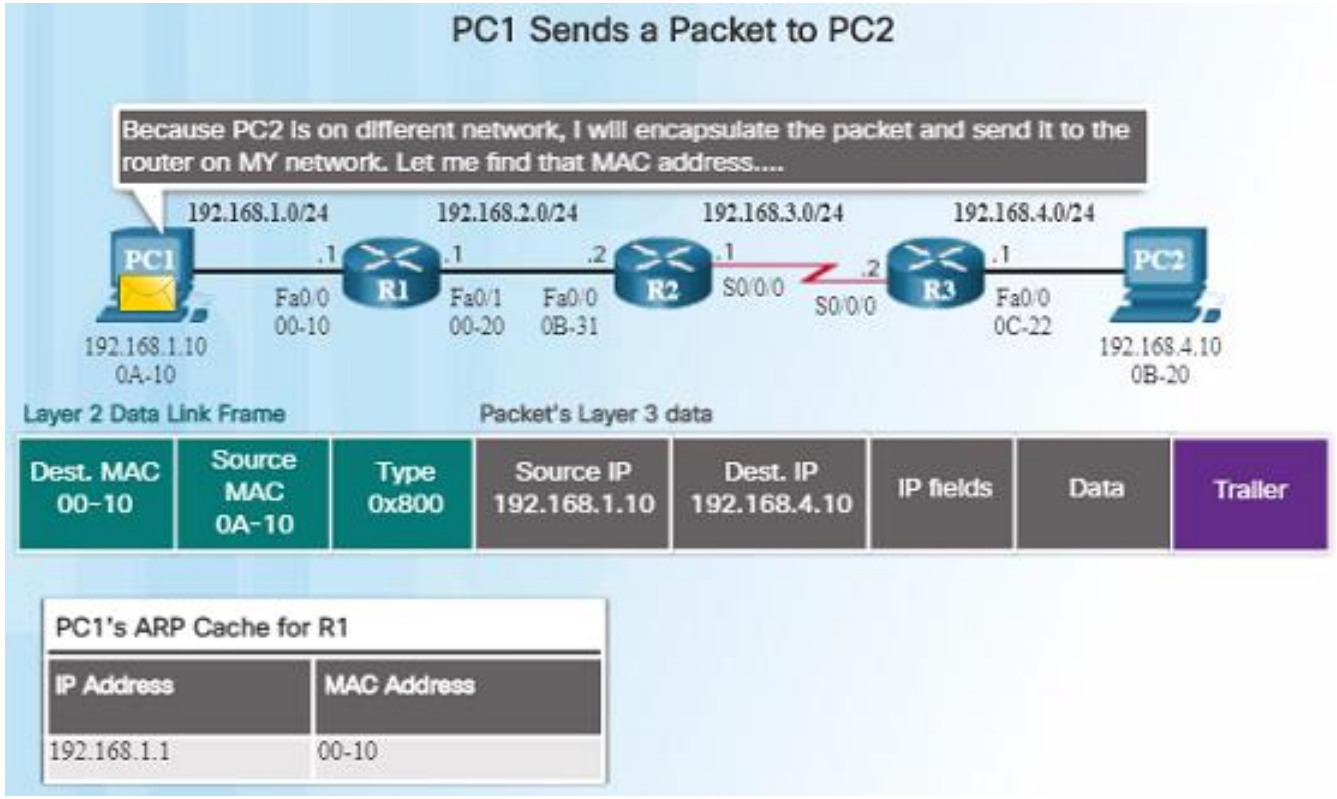
Addressing: routing to another LAN

- R forwards datagram with IP source A, destination B
- R creates link-layer frame with B's MAC address as dest, frame contains A-to-B IP datagram



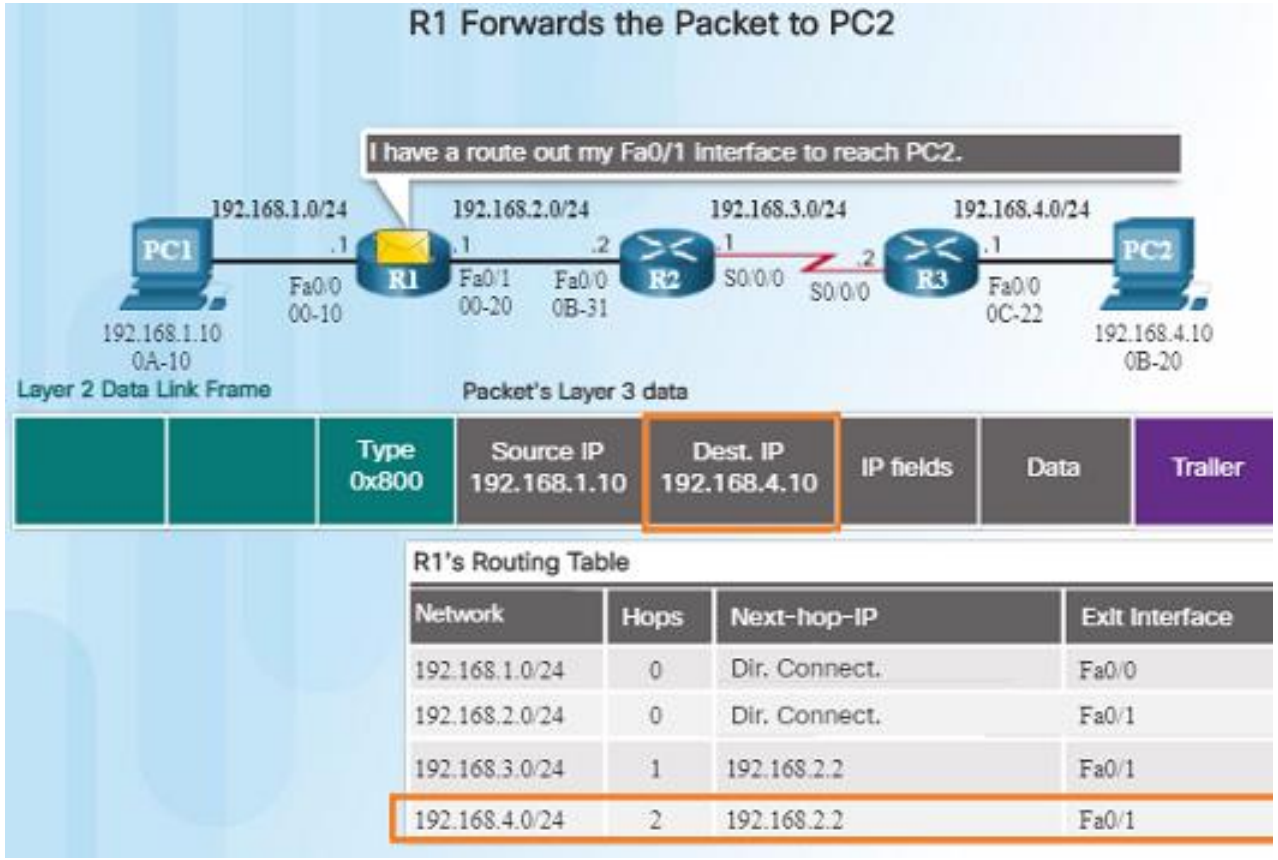
* Check out the online interactive exercises for more examples: http://gaia.cs.umass.edu/kurose_ross/interactive/

Send a Packet



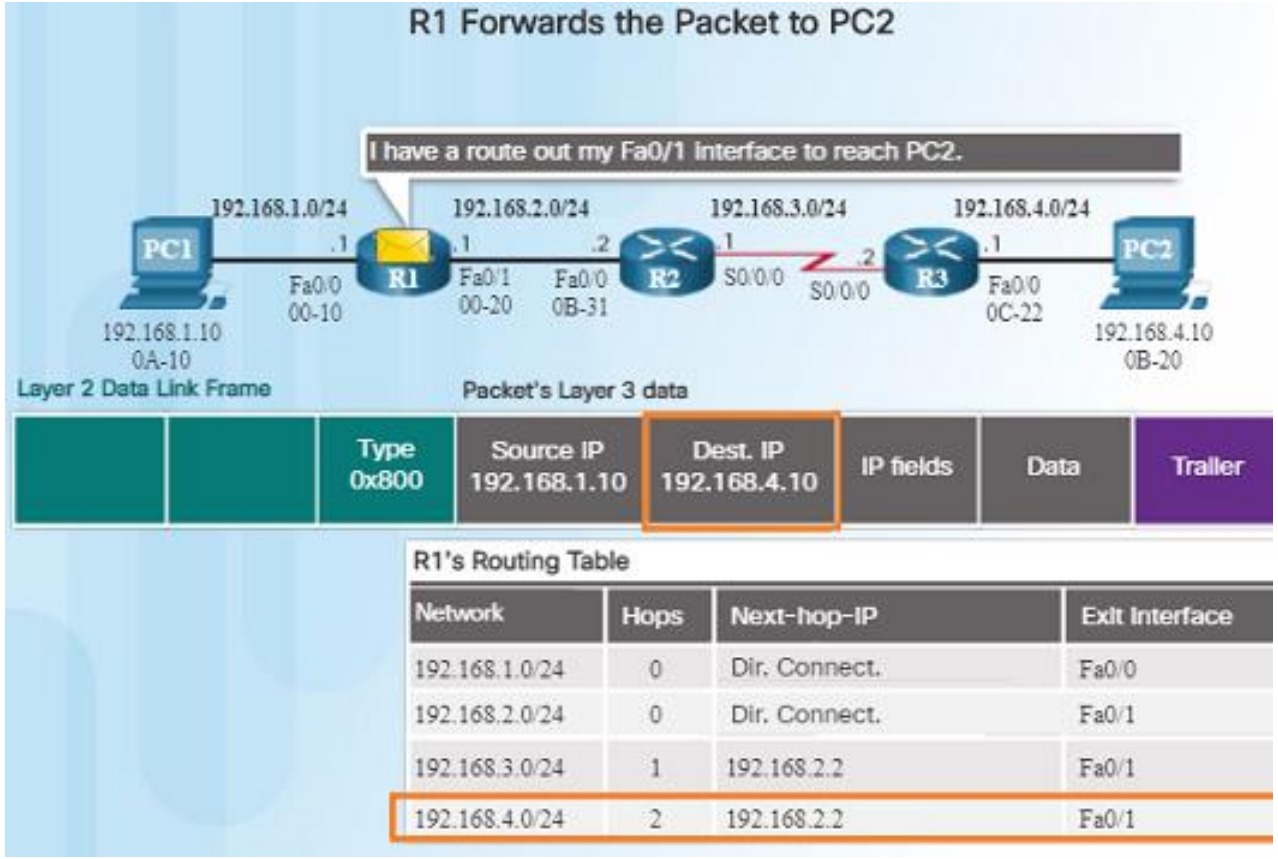
- For PC1 to send a packet to PC2, the following occurs:
  - PC1 must determine if the destination IPv4 address is on the same network. If it is on the same network, PC1 will obtain the destination MAC address from its ARP cache or use an ARP request.
  - Because the destination network is on a different network, PC1 forwards the packet to its default gateway.
  - To determine the MAC address of the default gateway, PC1 checks its ARP table for the IPv4 address of the default gateway and its corresponding MAC address. An ARP request is sent if it is not found.
  - When PC1 has the MAC address of Router R1, it can forward the packet.

Forward to the Next Hop



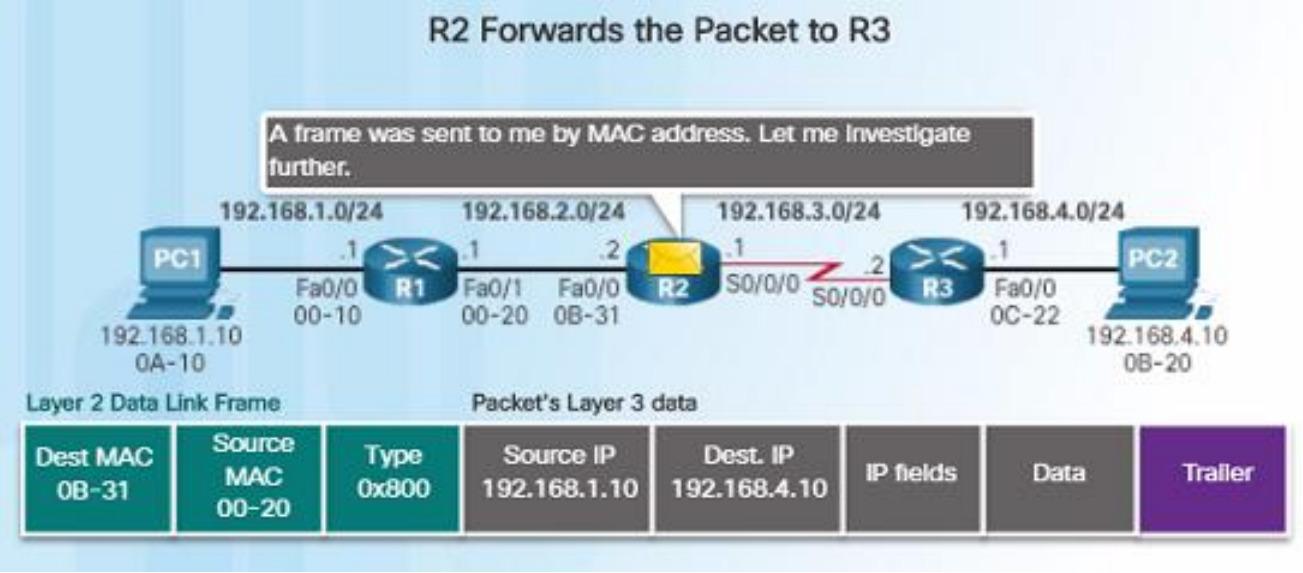
- When R1 receives the Ethernet frame from PC1, the following occurs:
  - R1 examines the destination MAC address which matches the MAC address of the receiving interface and copies the frame into its buffer.
  - R1 identifies the Ethernet Type field as 0x800 which indicates that the Ethernet frame contains an IPv4 packet in the data portion of the frame.
  - R1 de-encapsulates the Ethernet frame.
  - Because the destination IPv4 address of the packet, 192.168.4.10, does not match any of the directly connected networks on R1, R1 searches the routing table for a corresponding route.
    - R1's Routing Table has a route for the 192.168.4.0/24 network.

Forward to the Next Hop (Cont.)



- When R1 receives the Ethernet frame from PC1, the following occurs:
  - The route that R1 finds to the 192.168.4.0/24 network has a next-hop address of 192.168.2.2 and an exit interface of FastEthernet 0/1.
  - This will require that the IPv4 packet be encapsulated in a new Ethernet frame with the destination MAC address of the IPv4 address of the next-hop router, 192.168.2.2
  - Because the exit interface is on an Ethernet network, R1 must resolve the next-hop IPv4 address with a destination MAC address using ARP, assuming it is not in its ARP cache.
  - When R1 has the MAC address for the next-hop, the Ethernet frame is forwarded out of the FastEthernet 0/1 interface of R1.

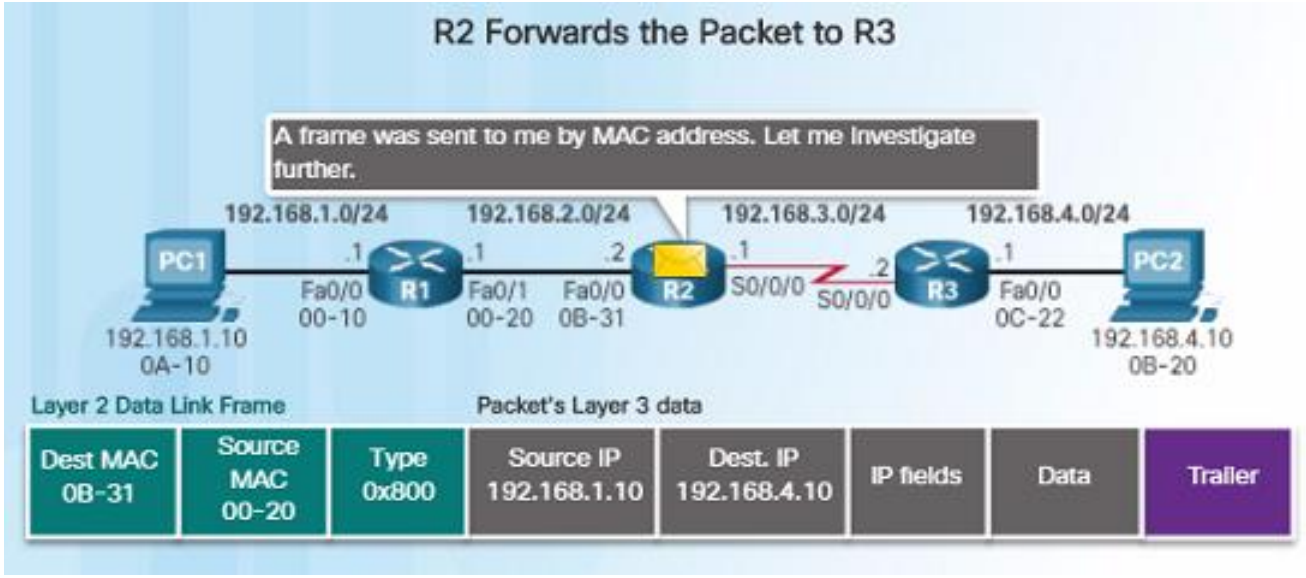
Packet Routing



- The process outlined to the right describes what happens when router R2 receives a frame on its FA0/0 interface that needs to be forwarded to router R3.

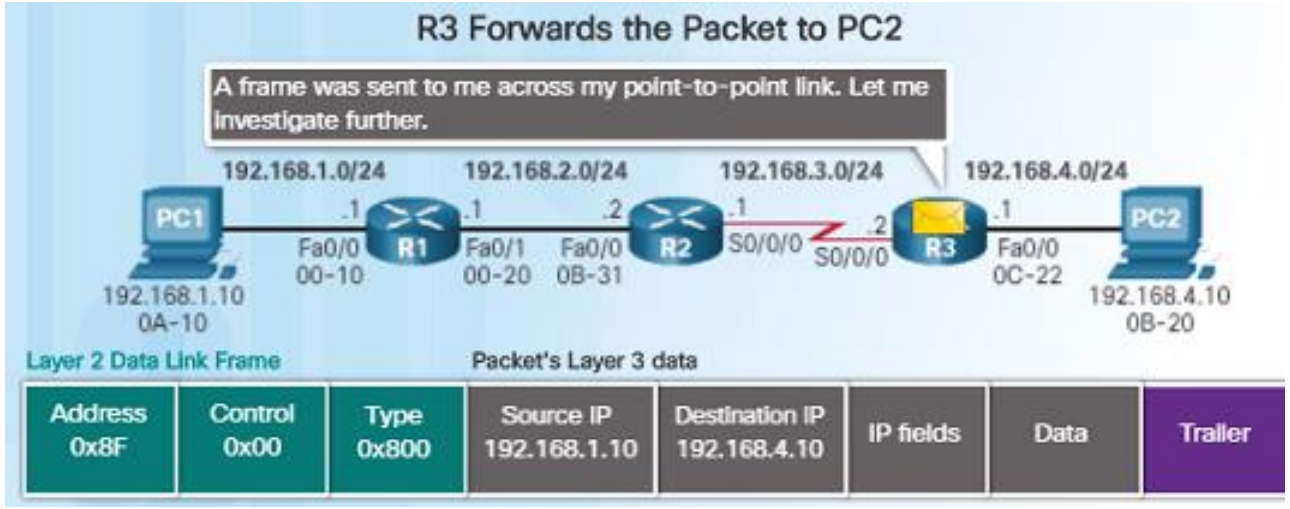
- R2 examines the destination MAC address. Because it matches the MAC address of its receiving interface, R2 copies the frame into its buffer.
- R2 determines that that frame contains an IPv4 packet in the data portion of the frame.
- R2 de-encapsulates the Ethernet frame.
- Because the destination IP address is on a different network, the routing table is searched to find a corresponding route for the destination IPv4 address.

Packet Routing (Cont.)



- The routing table of R2 has a route to the 192.168.4.0/24 network with a next-hop IPv4 address of 192.168.3.2 and an exit interface of Serial 0/0/0.
- Because the exit interface is not Ethernet, R2 does not have to resolve the next-hop IP-v4 address with a destination MAC address.
- The IPv4 packet is encapsulated into a new data link frame used by the exit interface and sent out the Serial 0/0/0 exit interface.
- Because there are no MAC addresses on serial interfaces, R2 sets the data link destination address to an equivalent of a broadcast.

Reach the Destination

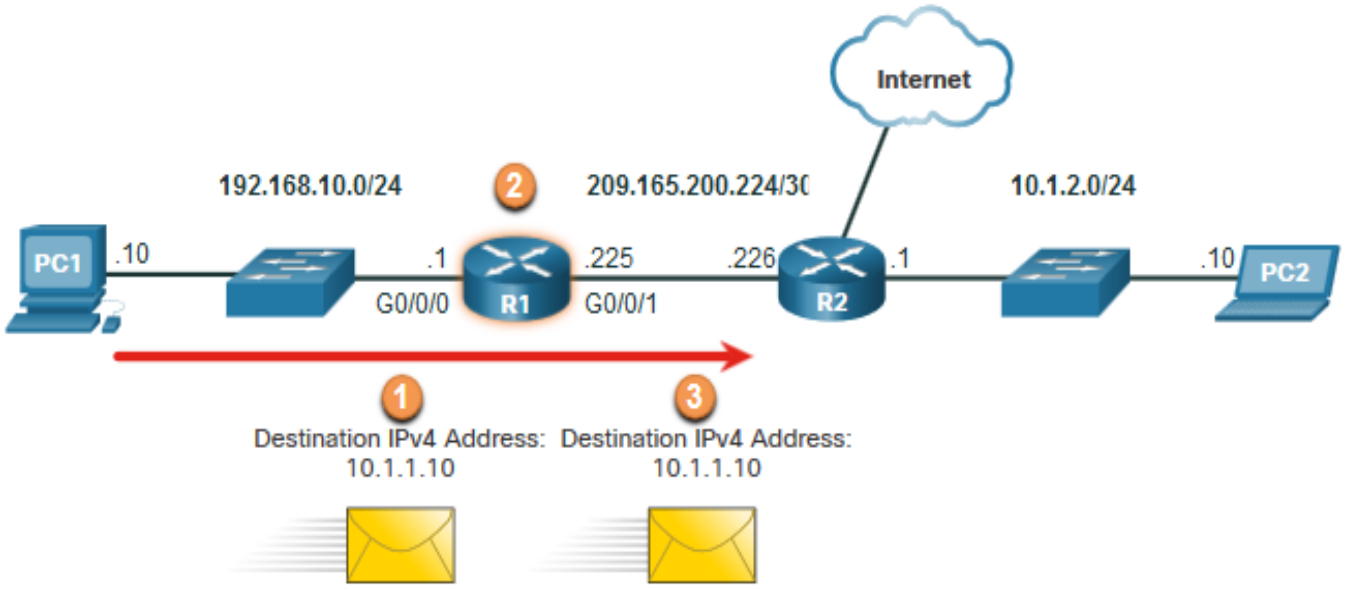


- The process outlined on the right describes what takes place when R3 receives a frame on its serial interface.

- R3 copies the data link PPP frame into its buffer.
  - R3 de-encapsulates the data link PPP frame.
  - R3 searches the routing table for the destination IPv4 address of the packet.
  - Because the destination network is on R3's directly connected network, the packet can be sent directly and does not need to be sent to another router.
  - Because the exit interface is a directly connected Ethernet network, R3 must resolve the destination IPv4 address of the packet with a destination MAC address by either finding it in its ARP cache or send out an ARP request.

What happens when the router receives the frame from the host device?

Router Packet Forwarding Decision



1. Packet arrives on the Gigabit Ethernet 0/0/0 interface of router R1. R1 de-encapsulates the Layer 2 Ethernet header and trailer.
2. Router R1 examines the destination IPv4 address of the packet and searches for the best match in its IPv4 routing table. The route entry indicates that this packet is to be forwarded to router R2.
3. Router R1 encapsulates the packet into a new Ethernet header and trailer, and forwards the packet to the next hop router R2.

R1 Routing Table

Route	Next Hop or Exit Interface
192.168.10.0 /24	G0/0/0
209.165.200.224/30	G0/0/1
10.1.1.0/24	**via R2**
Default Route 0.0.0.0/0	via R2

ICMP Type and Code

Type	Code	Description
0	0	echo reply (ping)
3	0	destination network unreachable
3	1	destination host unreachable
3	2	destination protocol unreachable
3	3	destination port unreachable
3	6	destination network unknown
3	7	destination host unknown
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

Trace Route – ICMP

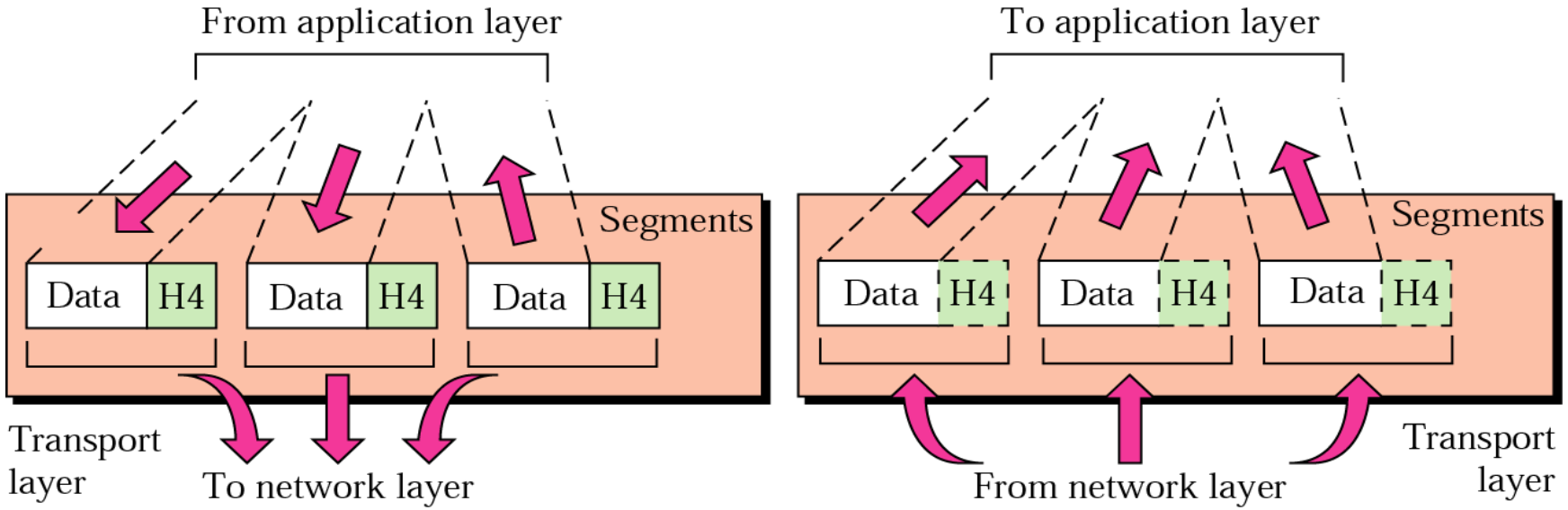
- Source sends series of UDP segments to destination host
  - First has TTL =1
  - Second has TTL=2, etc.
  - Unlikely port number
- When n^{th} datagram arrives to n^{th} router:
  - Router discards datagram
  - Sends to source an ICMP message (type 11, code 0)
  - Datagram includes router IP address. Traceroute does DNS lookup to find name of router (if any)
- When ICMP message arrives, source calculates RTT
- Traceroute repeat the process 3 times (To get average RTT)

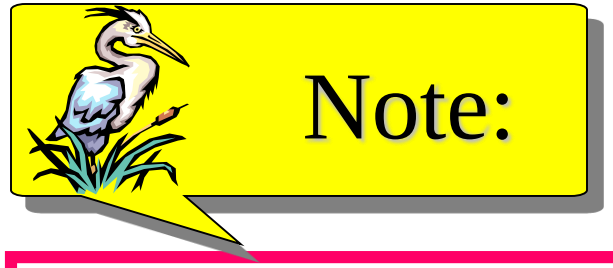
Stopping criterion

- UDP segment eventually arrives at destination host
- Destination returns ICMP "port unreachable" packet (type 3, code 3)

Transport Layer Responsibilities

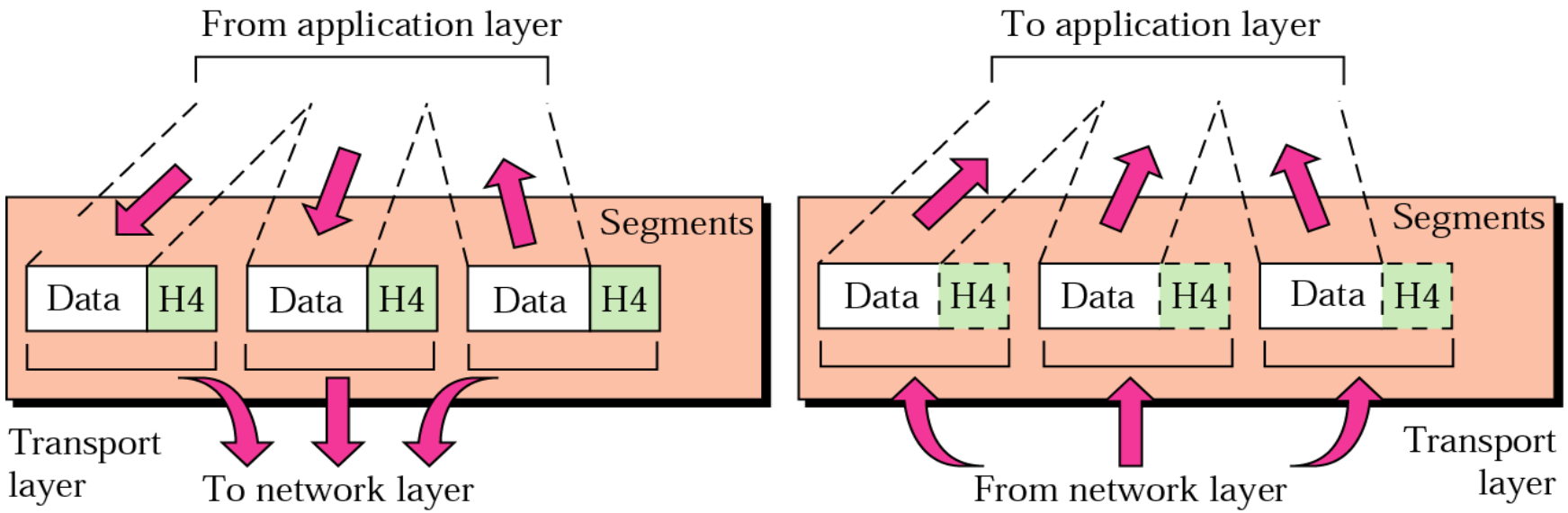
- Process-to-process delivery of entire message
- Port addressing
- Segmentation and reassembly
- Connection control: connectionless or connection-oriented
- End-to-end flow control
- End-to-end error control



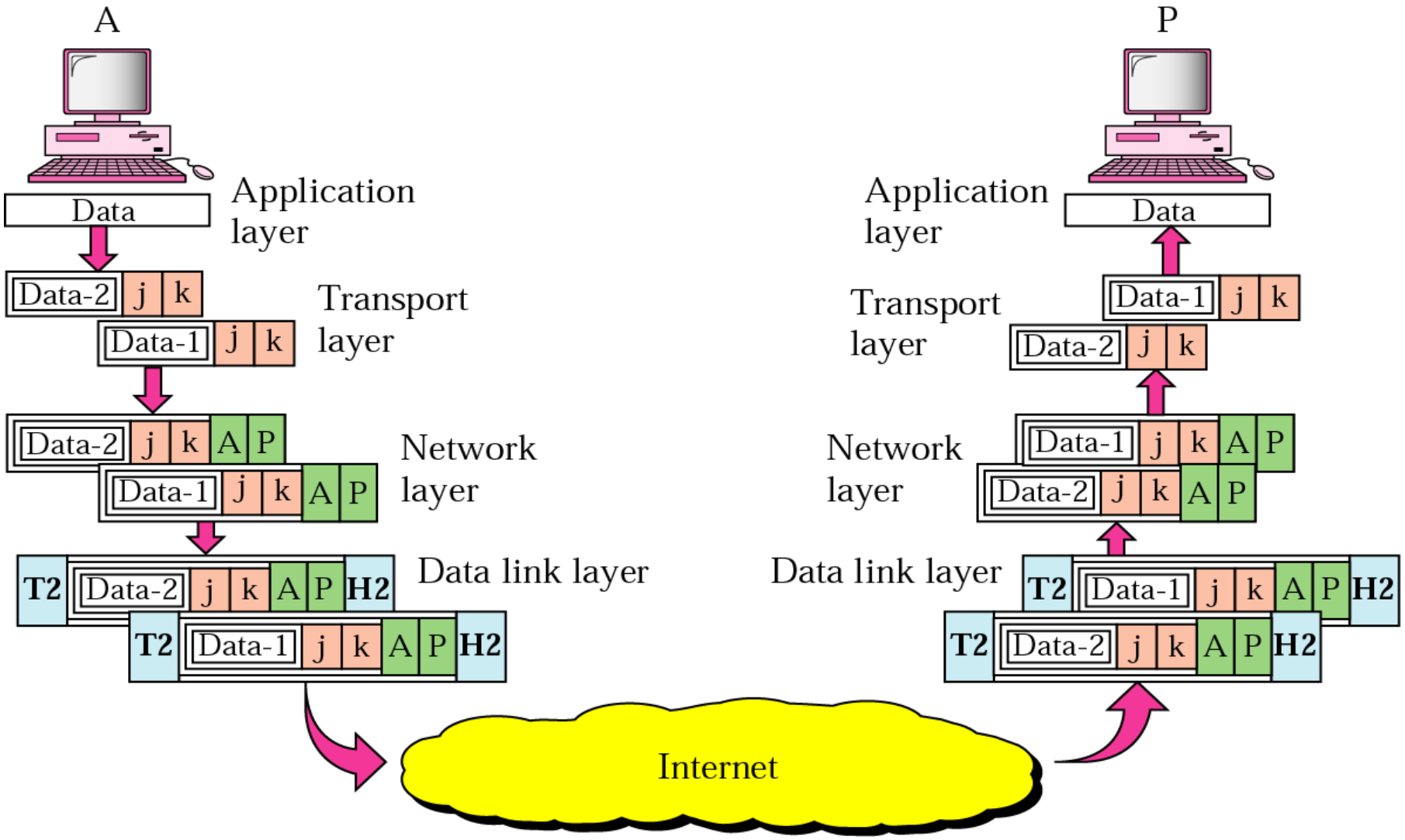


Note:

The transport layer is responsible for delivery of a message from one process to another.

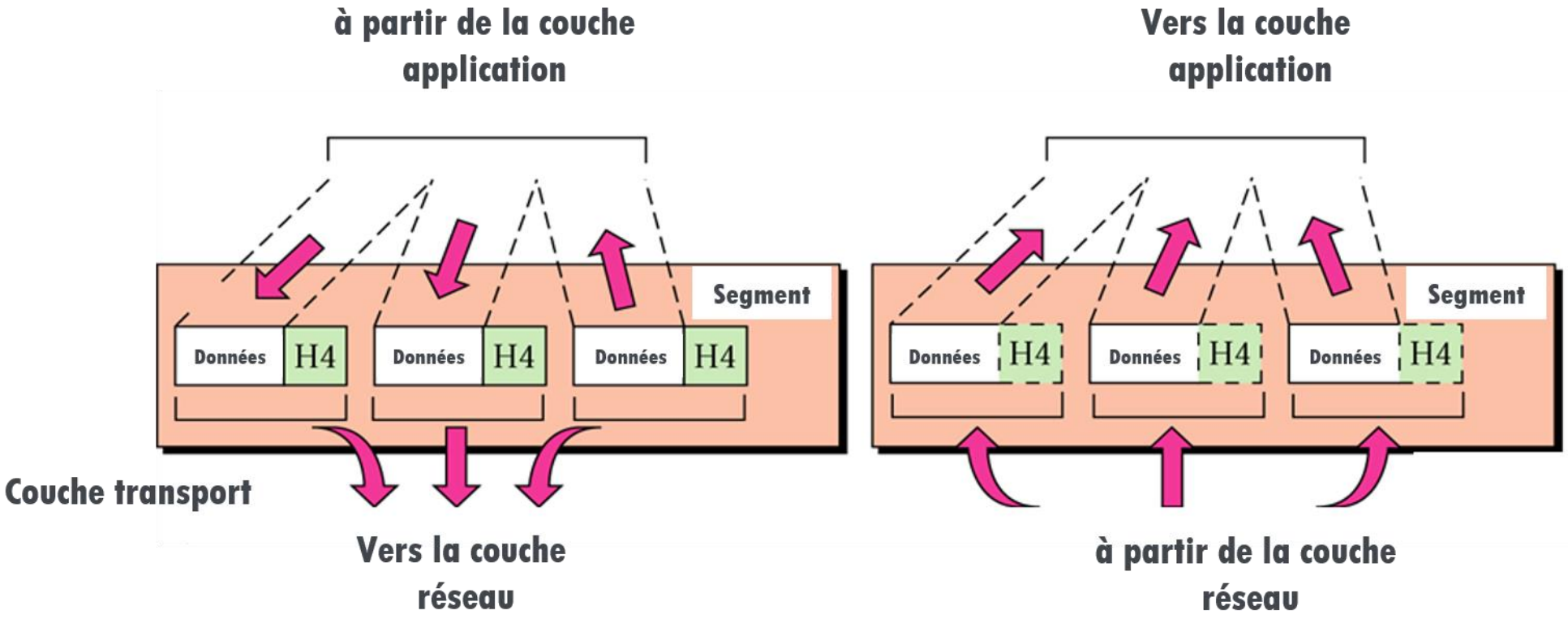

From application layer
Segments
Data
H4
Data
H4
Data
H4
Transport layer
To network layer
To application layer
Segments
Data
H4
Data
H4
Data
H4
From network layer
Transport layer

Example 3

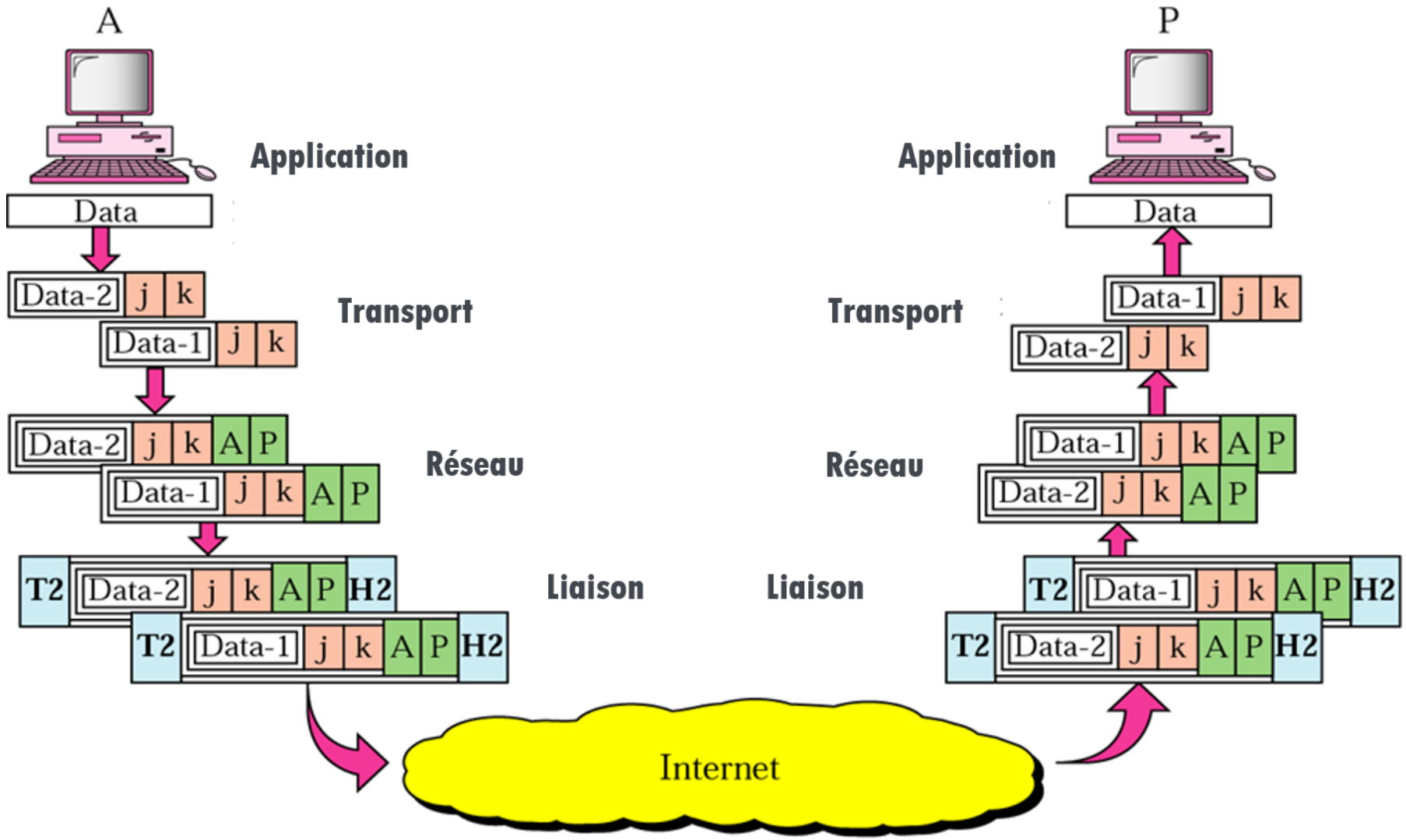

A
P
Application layer
Transport layer
Network layer
Data link layer
Data
Data-2
Data-1
j
k
A
P
T2
H2
Internet

La couche transport

- **La couche transport est responsable de la livraison des segments d'un processus source à un processus destination final.**

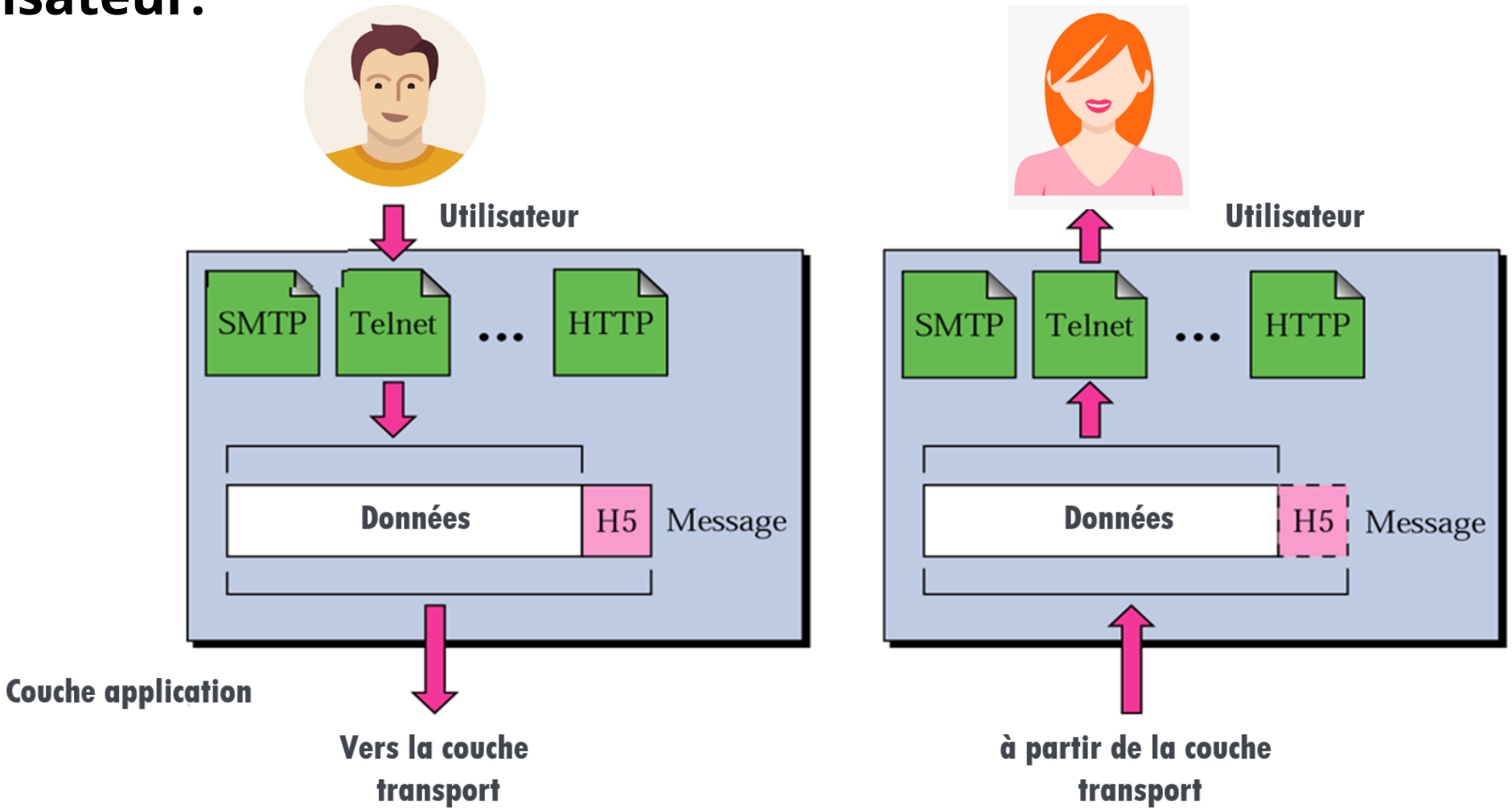


Encapsulation



Couche application

- **La couche application est chargée de fournir des services à l'utilisateur.**



THANKS !