

Problem I

(17+18)

(35 points)

- 1) We would like to simulate the work of a walkie-talkie between two processes, a father and a son. Each of the two processes reads words from the keyboard simulating typing a message when it reads the keyword "over" it stops and sends a signal to the other process to begin reading. The keyword "out" means that the communication has ended. The process that reads the "out" message stops and should stop his partner.

P.S: the father starts the communication

- 2) What is the possible output of the following pieces of programs

Program1	Program2
<pre> void main() { int i, j = 1, p[2]; pipe(p); write(p[1], &j, sizeof(int)); for(i=1; i<5; i++) if(!fork()) { close(p[1]); break; } read(p[0], &j, sizeof(int)); printf("%d\n", i); } </pre>	<pre> void main() { int i, j = 1, p[2]; pipe(p); for(i=1; i<5; i++) if(!fork()) { close(p[1]); break; } wait(0); read(p[0], &j, sizeof(int)); printf("%d\n", i); } </pre>
Program3	
<pre> void main() { int i, j = 1, p[2]; pipe(p); for(i=1; i<5; i++) if(!fork()) { close(p[1]); break; } write(p[1], &j, sizeof(int)); printf("%d\n", i); read(p[0], &j, sizeof(int)); } </pre>	

Problem 2

(19 + 16)

(35) points

- I. Consider a system with paginated main memory. The memory is composed of 4 frames (frames) each has a size of 64 bytes. At a given moment, the memory is empty, then two processes P1 (4 pages) and P2 (2 pages) are launched in the system. The processor sends requests submissions in the following order in the format [hexadecimal logical address, Process]:

[3F, P1]
[4A, P1]
[1D, P1]
[00, P2]
[CA, P1]
[87, P1]
[39, P2]
[2B, P1]
[00, P1]
[11, P2]

19

1. What is in bytes the size of the main memory? 2
2. What is the size of the virtual address space? 3
3. Indicate by figures the evolution of the memory and the number of page faults using the following page replacement algorithms:
 - a) FIFO 4
 - b) LRU 4
 - c) Second chance. 4

- II. We assume the existence of six free space holes in a contiguous allocation memory, named from A to F, as shown in the right table:

Hole	Address	Size
A	360	20 Bytes
B	500	12 Bytes
C	1200	18 Bytes
D	1400	4 Bytes
E	1820	40 Bytes
F	1960	32 Bytes

Three successive R1, R2 and R3 requests of size 12, 10 and 6 bytes must be satisfied. What is the memory address allocated to each of these requests if the allocation strategy is:

- a. First Fit ? 4
- b. Best Fit ? 4
- c. Worst Fit ? 4
- d. Next Fit ? 4

2) Prog 1: (6 pts)

> if the parent reads first, then all child will print also

5 4 3 2 1

0

(because the parent exits and the write side on child's is closed)

> if any child read first, it prints the value of i (ex. 4)
all other processes included parent will block
because they will attempt to read from an empty
pipe with write side opened.

Prog 2: (6 pts)

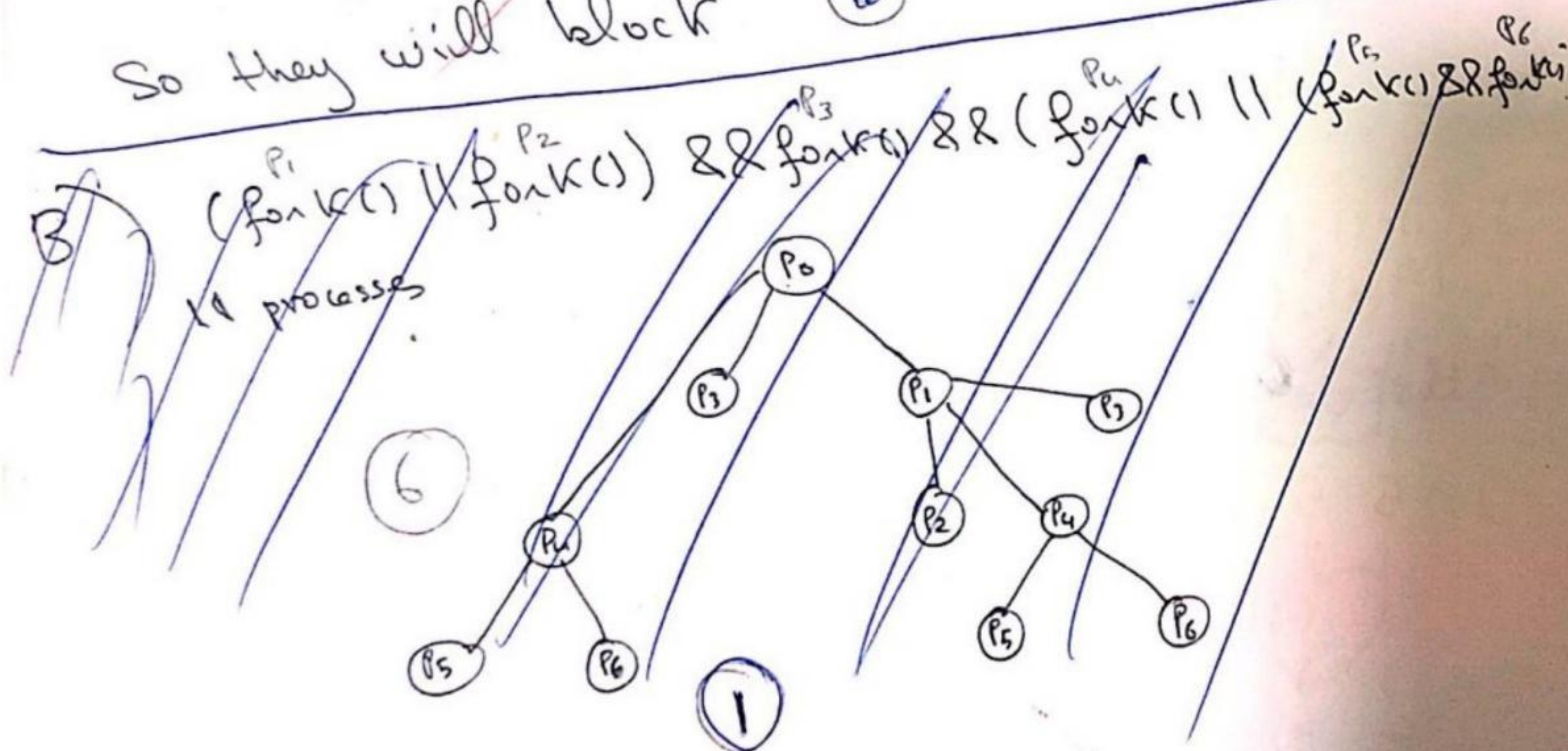
No output

here the parent can't read immediately, he must waiting
at least one child to read from pipe and exit.

Suppose the fourth child read from pipe, it will
print 4 and exit.

Now the parent or other child will attempt to
read from an empty pipe with write side opened

So they will block



Prog 3 (6 pts)

- each child reads from empty pipe with no writer (closed port) \Rightarrow no block
 - then it ~~writes~~ ^{prints} the value of i
- possible outputs :

1	2	3	4	5
2	1	3	4	5
...				
				5

Possible combinations of $\{1, 2, 3, 4\}$

- the parent prints 5 at the end after the ending of childs

~~all the processes~~

if the parent starts first, it can write into the pipe the value 5, prints 5, and read 5, ~~then exit~~

- all other child will print its value of i then they will block on reading the empty block if the parent doesn't exist yet. If not they will not block

if parent first

5	5	...	5
1	2		4
2	1		3
3	3		2
4	4	...	1

if any of the child starts first

2	2	2	2
1	3	4	5
3	1	3	1
5	5	5	1

2: Memory management (15 pts)

$P_1 \rightarrow 4$ pages

$P_2 \rightarrow 2$ pages

3	64B
2	
1	
0	

1) size of memory:

$$64 \times 4 = 256 \text{ Bytes} = 2^8 \text{ Bytes} \quad (2) \quad (3)$$

2) size of the virtual address space

Page size = 64 By = 2^6 Bytes

\Rightarrow 6 bits is needed for the offset (2) (3)

The logical address is in hexadecimal (2 digits)
 - each is encoded on 4 bits \Rightarrow the logical address is encoded on 8 bits \Rightarrow $\frac{\# \text{ offset}}{2} = 6$ (virtual address on 8 bits)

$$\Rightarrow 2^2 \text{ pages} \Rightarrow 4 \times 64 = 256 \text{ Bytes} = 2^8 \quad (3)$$

⇒ 2 page

Page

3)

3F ⇒	0011 1111	63	0	P ₁
4A ⇒	0100 1010	74	1	P ₁
1D ⇒	0001 1101	29	0	P ₁
00 ⇒	0000 0000	0	0	P ₂
CA ⇒	1100 1010	202	3	P ₁
87 ⇒	1000 0111	135	2	P ₁
39 ⇒	0011 1001	57	0	P ₂
2B	0010 1011	43	0	P ₁
00	0000 0000	00	0	P ₁
11	0001 0001	17	0	P ₂

(3)

007

\Rightarrow The $W = \{0, 1, 0, 0, 3, 2, 0, 0, 0, 0\}$

a) FIFO

	0	0	0	0	0	0	0	0	0
0									
1		1	1		1	1	1	1	1
2					3	3	3	3	3
3						2	2	2	2

$W = \{(0; P_1), (1; P_1), (0; P_1), (0; P_2), (3; P_1), (2; P_1), (0; P_2), (0; P_1), (0; P_1), (0; P_2)\}$

4 page faults

(2)

FIFO

Page refs	4 page frames				
	Fault?	Page contents			
(0, P ₁)	✓	(0, P ₁)			
(1, P ₁)	✓	(1, P ₁)	(0, P ₁)		
(0, P ₁)	X	(1, P ₁)	(0, P ₁)		
(0, P ₂)	✓	(0, P ₂)	(1, P ₁)	(0, P ₁)	
(3, P ₁)	✓	(3, P ₁)	(0, P ₂)	(1, P ₁)	(0, P ₁)
(2, P ₁)	✓	(2, P ₁)	(3, P ₁)	(0, P ₂)	(1, P ₁)
(0, P ₂)	X	(2, P ₁)	(3, P ₁)	(0, P ₂)	(1, P ₁)
(0, P ₁)	✓	(0, P ₁)	(2, P ₁)	(3, P ₁)	(0, P ₂)
(0, P ₁)	X	—	—	—	—
(0, P ₂)	X	—	—	—	—

6 page faults

b) LRU

Page refs	Fault?	4 Page Frames			
		Page contents			
(0, P ₁)	✓	(0, P ₁)			
(1, P ₁)	✓	(1, P ₁)	(0, P ₁)		
(0, P ₁)	X	(0, P ₁)	(1, P ₁)		
(0, P ₂)	✓	(0, P ₂)	(0, P ₁)	(1, P ₁)	
(3, P ₁)	✓	(3, P ₁)	(0, P ₂)	(0, P ₁)	(1, P ₁)
(2, P ₁)	✓	(2, P ₁)	(3, P ₁)	(0, P ₂)	(0, P ₁)
(0, P ₂)	X	(0, P ₂)	(2, P ₁)	(3, P ₁)	(0, P ₁)
(0, P ₁)	X	(0, P ₁)	(0, P ₂)	(2, P ₁)	(3, P ₁)
(0, P ₁)	X	—	—	—	—
(0, P ₂)	X	(0, P ₂)	(0, P ₁)	(2, P ₁)	(3, P ₁)

5 page faults

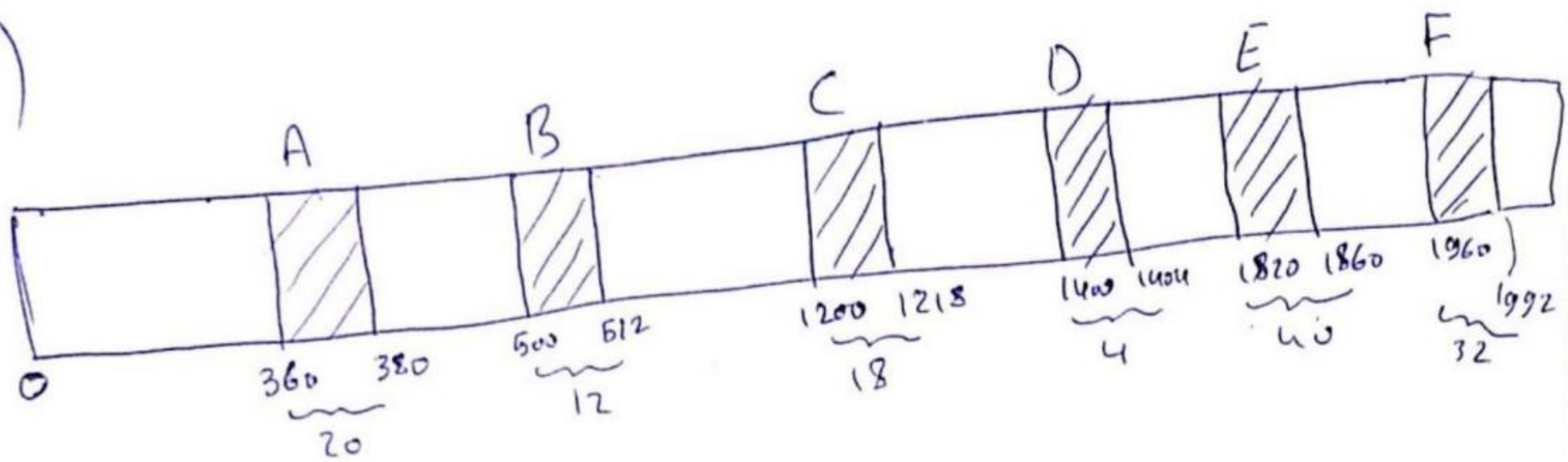
c) Second chance

Page refs	4 Page Frames						out
	Fault?	Page contents					
A (0, P ₁)	✓	(0, P ₁)					
B (1, P ₁)	✓	(1, P ₁)	(0, P ₁)				
A (0, P ₁)	X	(1, P ₁)	(0, P ₁) [*]				
C (0, P ₂)	✓	(0, P ₂)	(1, P ₁)	(0, P ₁) [*]			
D (3, P ₁)	✓	(3, P ₁)	(0, P ₂)	(1, P ₁)	(0, P ₁) [*]		
E (2, P ₁)	✓	(2, P ₁)	(3, P ₁)	(0, P ₂)	(0, P ₁)	(1, P ₁)	
C (0, P ₂)	X	(2, P ₁)	(3, P ₁)	(0, P ₂) [*]	(0, P ₁)		
A (0, P ₁)	X	(2, P ₁)	(3, P ₁)	(0, P ₂) [*]	(0, P ₁) [*]		
A (0, P ₁)	X	—	—	—	—		
C (0, P ₂)	X	—	—	—	—		

~~5~~ page faults

6 page faults

3

~~Scanned output~~~~Scanned output~~

$R_1 \div 12$ bytes
 $R_2 \div 10$ bytes
 $R_3 \div 6$ bytes

a) First Fit

R_1 is allocated on A $\Rightarrow A = 8$ bytes
 R_2 is allocated on B $\Rightarrow B = 2$ bytes
 R_3 is allocated on A $\Rightarrow A = 2$ bytes

b) Best Fit

R_1 is allocated on B $\Rightarrow B = 0$
 R_2 is allocated on C $\Rightarrow C = 8$
 R_3 is allocated on C $\Rightarrow C = 2$

c) Worst Fit

R_1 is allocated on E $\Rightarrow E = 28$
 R_2 is allocated on F $\Rightarrow F = 22$
 R_3 is allocated on E $\Rightarrow E = 22$

d) Next Fit

4

- R_1 on A $\Rightarrow A = 8$
 - R_2 on B $\Rightarrow B = 2$
 - R_3 on C $\Rightarrow C = 12$

Problem 1

2nd session 19-20

File: /home/parallels/Documents/FS1...ms/19-20/Second Sessio/Pb1-1.c Page 1 of 1

(17 pts)

```

#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <signal.h>
#include <sys/types.h>

void handler(){} 2

void main(){
    signal(SIGSTP, handler); 1
    signal(SIGCONT, handler); 1
    char word[10];

    if(!pid=fork()){ 1
        pause(); 1
        while(1){ 1
            scanf("%s", name); 1
            if (!strcmp(name, "over")){
                kill(getppid(), SIGCONT); 1
                pause();
            }
            elseif (!strcmp(name, "out")){
                kill(getppid(), SIGQUIT); 1
                exit(1); 1
            }
        }
    }
    else {
        while(1){ 1
            scanf("%s", name); 1
            if (!strcmp(name, "over")){
                kill(pid, SIGCONT); 1
                pause(); 1
            }
            elseif (!strcmp(name, "out")){
                kill(pid, SIGQUIT); 1
                exit(1); 1
            }
        }
    }
}

```