

Sorting

Guttag Chapter 12

Goals

Understand why and when is sorting useful

Understand different methods of sorting

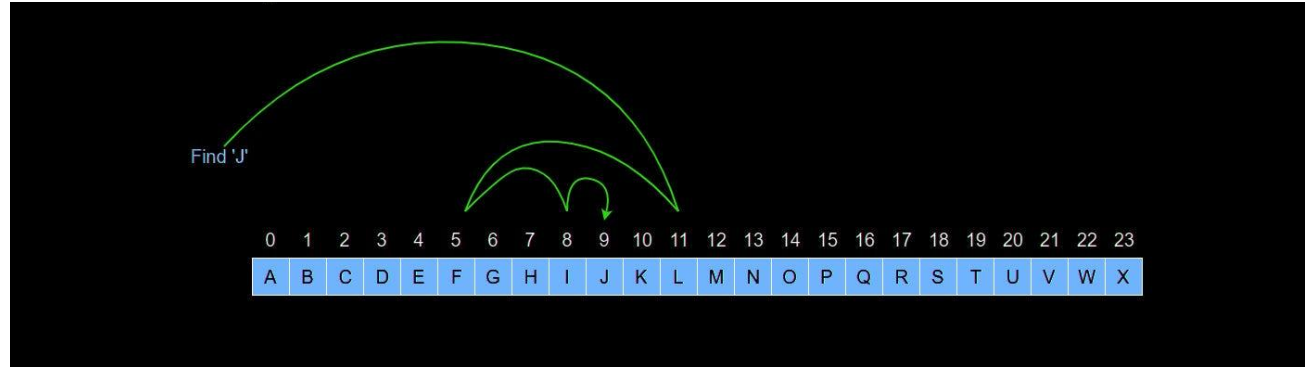
- selection
- bubble
- insert
- merge sort
- python built-in algorithms (quicksort and timsort)

Live coding example

Why and When to Sort

Why is sorting useful?

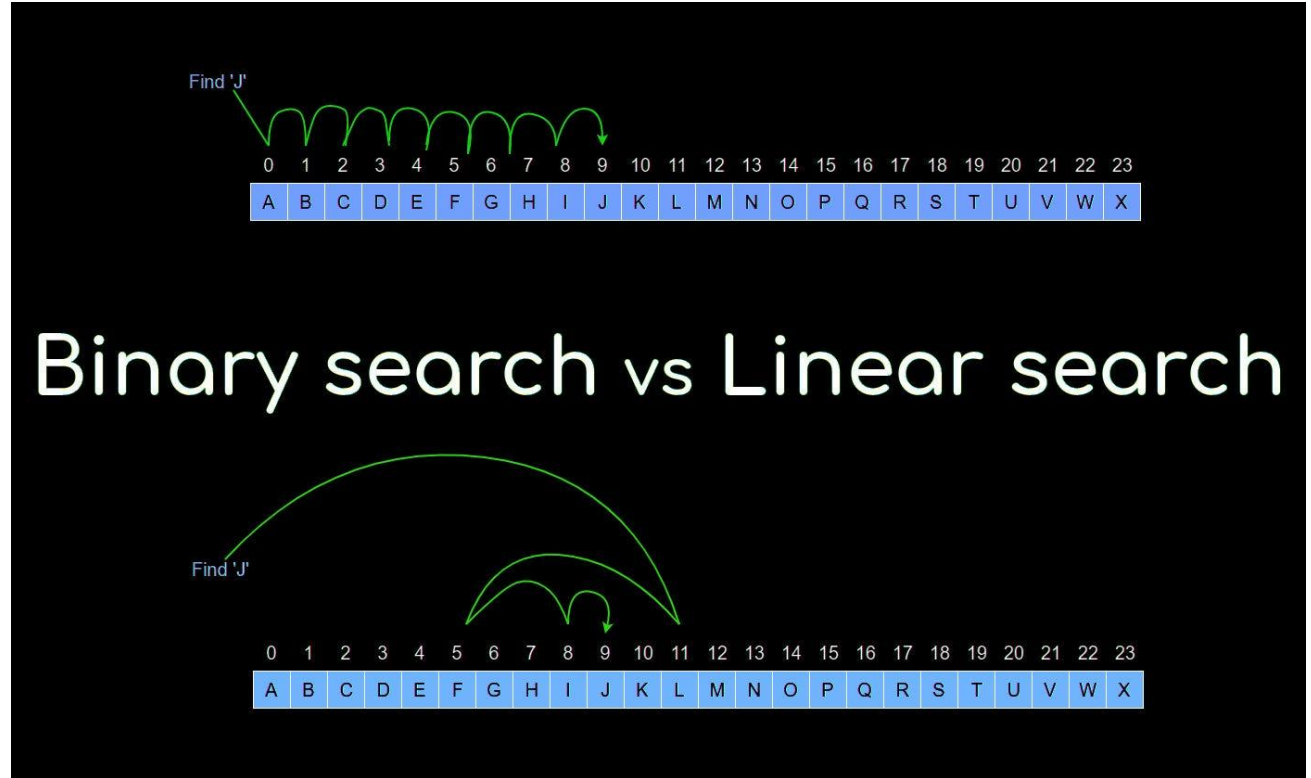
It makes searching easier!



When is sorting useful?

The "cost" can be
amortized

I.e. sorting can
happen just once,
but subsequent
searching could
happen multiple
times, diluting the
cost.



When is sorting useful?

The "cost" can be
amortized

- Consider a pair of shoes costing \$100
- If you wear the shoes 100 times
- that is \$1 per use

You should probably only get the shoes if they will be used a lot

You should only sort if the sorted material will be searched a lot

How to Sort

Algorithm: Selection Sort

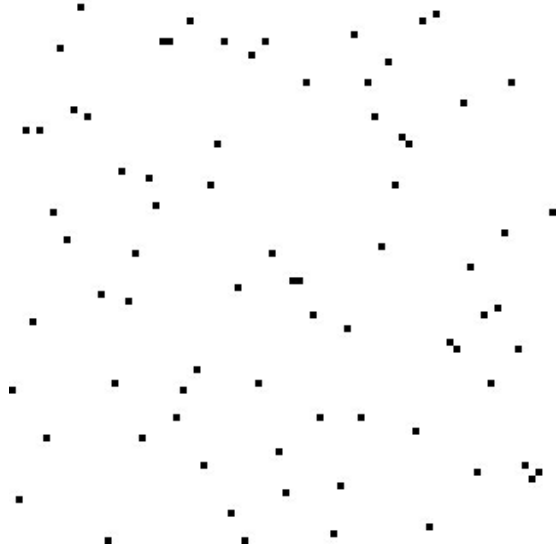
Conceptually, the list can have two parts: **sorted prefix** and **unsorted suffix**

game: select the smallest element in the suffix, and move it to the prefix ending in a loop

- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # unsorted suffix
- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # locate smallest
- [1, 8, 3, 8, 9, 6, 4, 8, 2, 2] # move to prefix ending
- [1, 8, 3, 8, 9, 6, 4, 8, 2, 2] # locate smallest
- [1, 2, 8, 3, 8, 9, 6, 4, 8, 2] # move to prefix ending
- [1, 2, 8, 3, 8, 9, 6, 4, 8, 2] # locate smallest
- [1, 2, 2, 8, 3, 8, 9, 6, 4, 8] # move to prefix ending
- [1, 2, 2, 8, 3, 8, 9, 6, 4, 8] # locate smallest...etc

Algorithm: Selection Sort

Repeatedly, **select** smallest and build up prefix



	8
	5
	2
	6
	9
	3
	1
	4
	0
	7

Characteristics

- selecting the smallest takes $O(n)$
- selection repeats n times!
- if sorting occurs in place, all the shuffling takes even more time!
- Complexity: $O(n^2)$

Algorithm: Bubble Sort

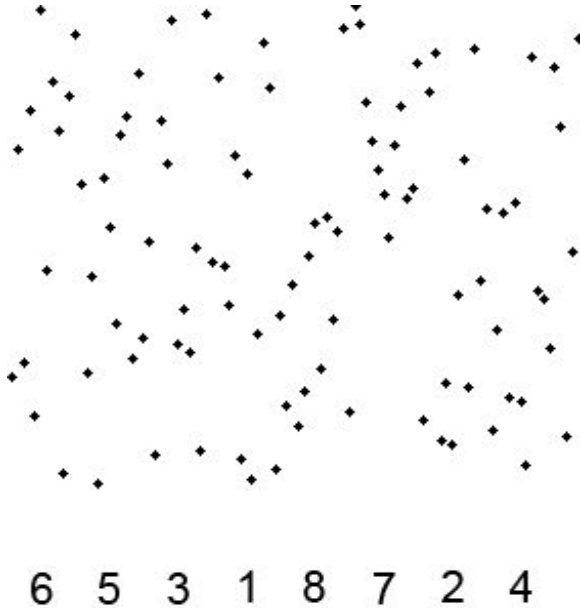
Conceptually, the list can have two parts: **unsorted prefix** and **sorted suffix**

game: find the largest element in prefix via "bubbling" through to the end, in a loop

- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # unsorted prefix
- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # locate largest
- [8, 3, 8, 6, 4, 8, 2, 2, 1, 9] # bubble to the suffix beginning
- [8, 3, 8, 6, 4, 8, 2, 2, 1, 9] # locate largest
- [3, 8, 6, 4, 8, 2, 2, 1, 8, 9] # bubble to the suffix beginning
- [3, 8, 6, 4, 8, 2, 2, 1, 8, 9] # locate largest
- [3, 6, 4, 8, 2, 2, 1, 8, 8, 9] # bubble to the suffix beginning
- [3, 6, 4, 8, 2, 2, 1, 8, 8, 9] # locate largest...etc

Algorithm: Bubble Sort

Repeatedly, **bubble** up the largest to suffix



Characteristics

- bubbling requires traversing the list to make decisions about what is largest $\rightarrow O(n)$
- bubbling repeats n times!
- Complexity: $O(n^2)$

Algorithm: Insert Sort

Conceptually, the list can have two parts: **sorted prefix** and **unsorted suffix**

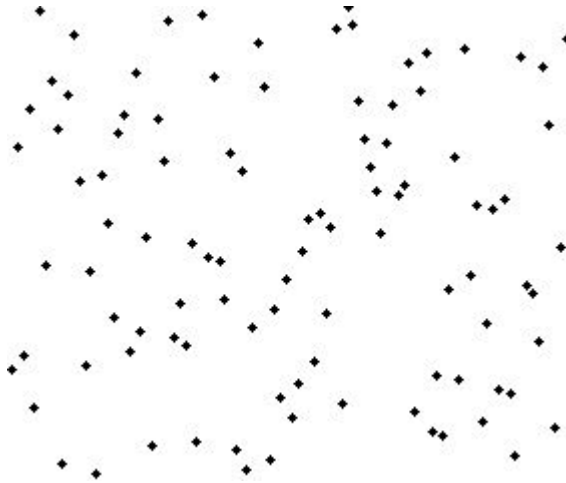
game: get the next element and insert it into the prefix at correct location in a loop

- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # unsorted suffix
- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # get next element
- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # insert in prefix after spot found
- [8, 3, 8, 9, 6, 4, 8, 2, 2, 1] # get next element
- [3, 8, 8, 9, 6, 4, 8, 2, 2, 1] # insert in prefix after spot found
- [3, 8, 8, 9, 6, 4, 8, 2, 2, 1] # get next element
- [3, 8, 8, 9, 6, 4, 8, 2, 2, 1] # insert in prefix after spot found
- [3, 8, 8, 9, 6, 4, 8, 2, 2, 1] # get next element...etc

Algorithm: Insert Sort

Repeatedly, get next element and **insert** it in prefix

Characteristics



- inserting requires traversing the prefix to make decision about correct location $\rightarrow O(n)$
- inserting repeats n times!
- Complexity: $O(n^2)$

Algorithm: Merge Sort

RECURSIVE!

but first....

Think about whether two sorted lists could be merged easily computationally.

[1,4,6,8,10,11] and [2,3,4,7,9,10]

Algorithm: Merge Sort

RECURSIVE!

Conceptually

#1. if a list is of length 0 or 1, it is already sorted.

#2. merging two sorted lists is fast

game: split large lists in half repeatedly until sorted, then sort and merge neighboring lists, repeatedly

Algorithm: Merge Sort

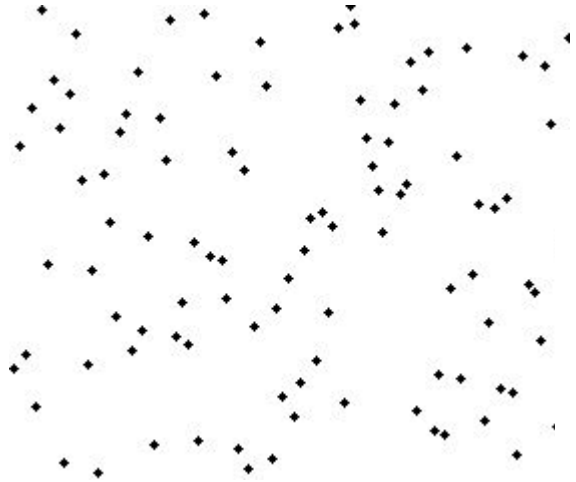
Merging two sorted lists:

Remaining in L_1	Remaining in L_2	Result
[1, 5, 12, 18, 19, 20]	[2, 3, 4, 17]	[]
[5, 12, 18, 19, 20]	[2, 3, 4, 17]	[1]
[5, 12, 18, 19, 20]	[3, 4, 17]	[1, 2]
[5, 12, 18, 19, 20]	[4, 17]	[1, 2, 3]
[5, 12, 18, 19, 20]	[17]	[1, 2, 3, 4]
[12, 18, 19, 20]	[17]	[1, 2, 3, 4, 5]

Algorithm: Merge Sort

RECURSIVE!

split up until **sorted**, then **sort** and **merge** neighboring pairs



6 5 3 1 8 7 2 4

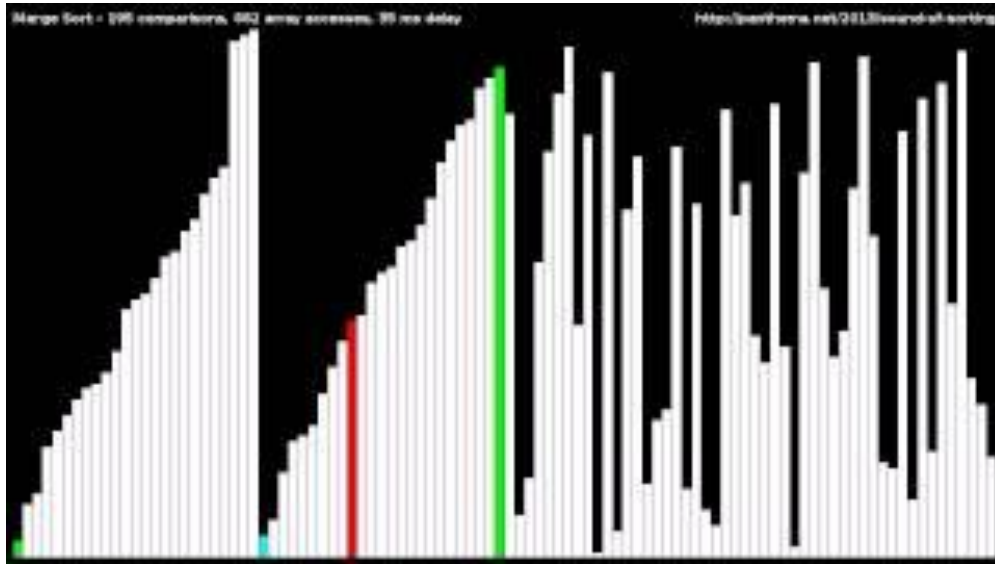
Characteristics

- merging lists takes $O(n)$ time
- merging reduces the number of lists by factor of 2
- total number of merges is $\log(n)$
- Complexity is $O(n \log(n))$

Algorithm: Merge Sort

RECURSIVE!

split up until **sorted**, then **sort** and **merge** neighboring pairs



Characteristics

- merging lists takes $O(n)$ time
- merging reduces the number of lists by factor of 2
- total number of merges is $\log(n)$
- Complexity is $O(n \log(n))$

Algorithm: built-in sorting algorithm

```
import random

# make a list filled with random ints, hopefully not in order :)
li = [random.randint(0,10) for _ in range(10)]

# sort in place
li.sort()
```

```
import random

# make a list filled with random ints, hopefully not in order :)
li = [random.randint(0,10) for _ in range(10)]

# sort into a new list
li_new = sorted(li)
```

Critical Thinking

- Why is sorting useful?
- What are commonalities in the sorting algorithms?
- What are differences between the algorithms?