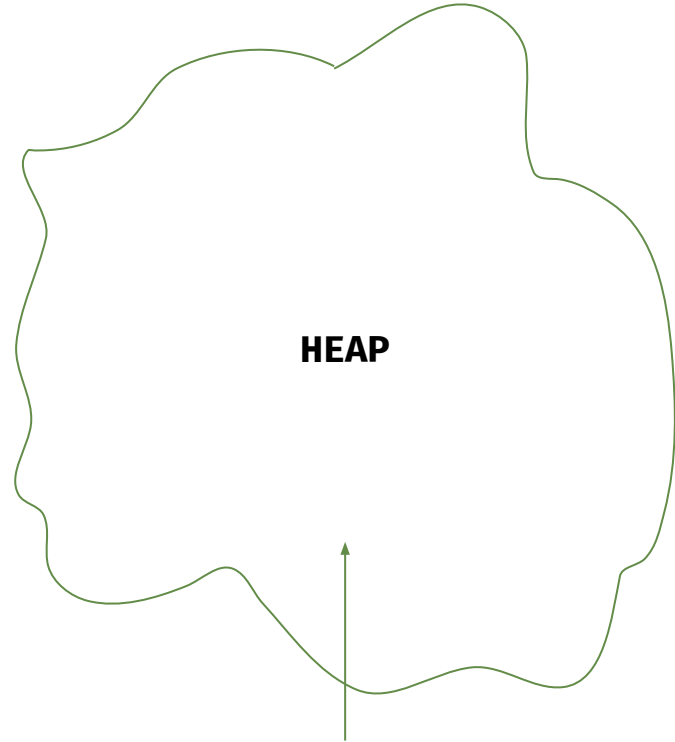




# Stack and Heap

STACK
0xdef45ea5
0x134a48fd
0x6785efae

↑  
Lawful good

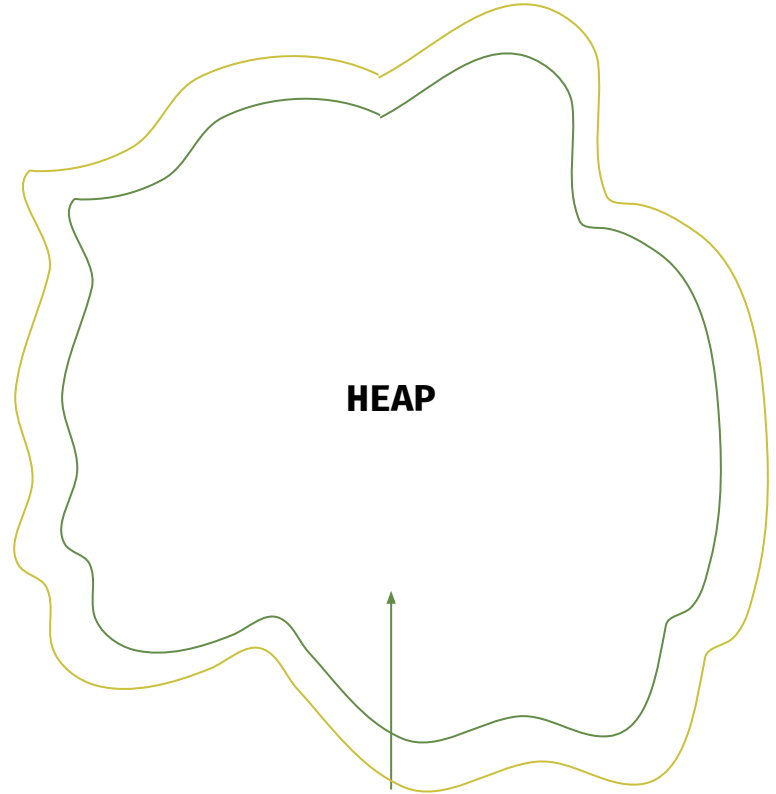


Chaotic neutral

# Stack and Heap

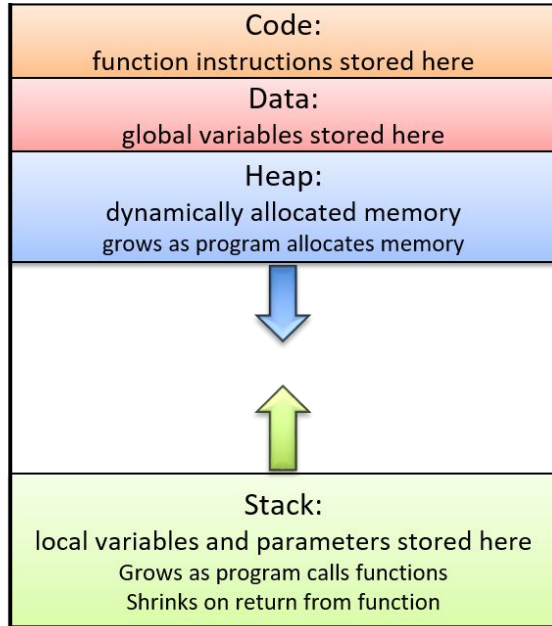
STACK
0xdef45ea5
0x134a48fd
0x6785efae

↑  
Lawful good



Chaotic neutral

# Stack and heap



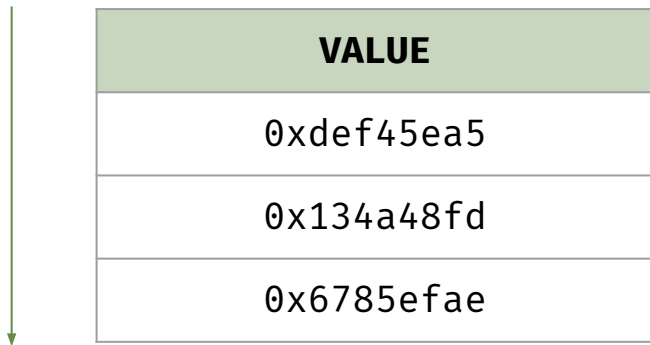
Dynamically allocated; can grow or shrink depending on usage.

(Can collide!)

# Brief introduction to the “stack”

There exists another quasi-temporary spot for memory which can get us out of a tricky jam: the *stack*.

LIFO  
Last In, First Out

A diagram illustrating the stack structure. On the left, a vertical green arrow points downwards, indicating the direction of growth. To the right of the arrow is a table with a header row labeled 'VALUE' and three data rows containing hexadecimal values. The text '“grows down”' is positioned below the arrow.

VALUE
0xdef45ea5
0x134a48fd
0x6785efae

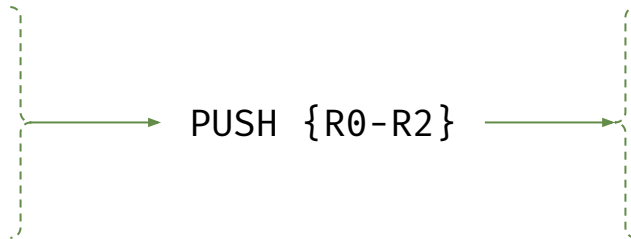
“grows down”

# Brief introduction to the “stack”

## PUSH

Places a value in the  
stack *in order*

REGISTERS	
R0	0xdef45ea5
R1	0x134a48fd
R2	0x6785efae



STACK
0xdef45ea5
0x134a48fd
0x6785efae

# Brief introduction to the “stack”

## POP

Extracts values from  
the stack *in order*

REGISTERS	
R0	0x6785efae
R1	
R2	

}

POP {R0}

}

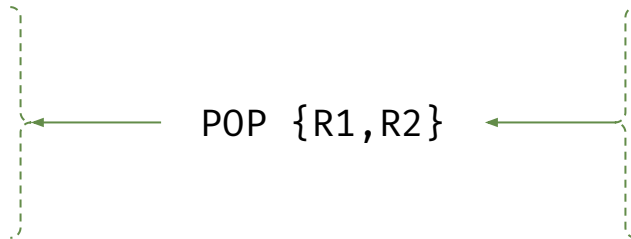
STACK
0xdef45ea5
0x134a48fd
0x6785efae

# Brief introduction to the “stack”

## POP

Extracts values from  
the stack *in order*

REGISTERS	
R0	0x6785efae
R1	0x134a48fd
R2	0xdef45ea5



STACK
0xdef45ea5
0x134a48fd
0x6785efae



# Brief introduction to the “stack”

PUSH       $\{R_M, R_D, \dots\}$

POP       $\{R_M, R_D, \dots\}$

PUSH       $\{R_M - R_D\}$

POP       $\{R_M - R_D\}$

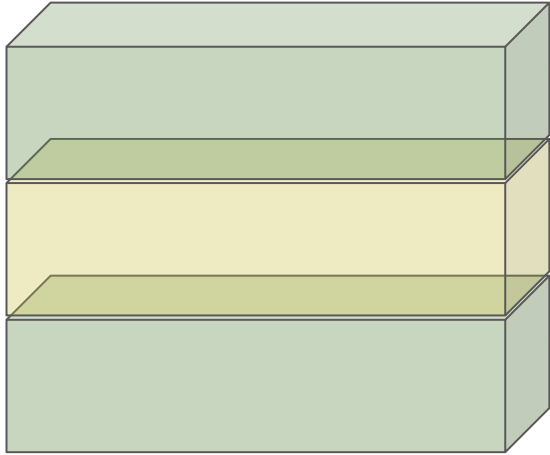
# Getting framed



Individual stack “frame”

Contains all PUSH'd,  
POP'd values for a  
given subroutine

# Getting framed



Implies: each subroutine  
can have its own stack  
“frame”

# Getting framed

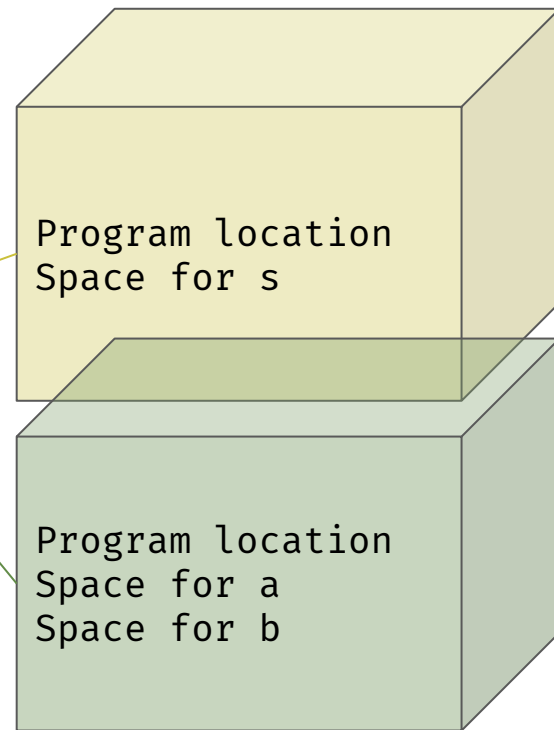
```
def add(a: int, b: int) -> int:
```

```
    c = a + b
```

```
    return c
```

```
s = add(2, 3)
```

```
print(s)
```



# Getting framed

```
def add(a: int, b: int) -> int:
```

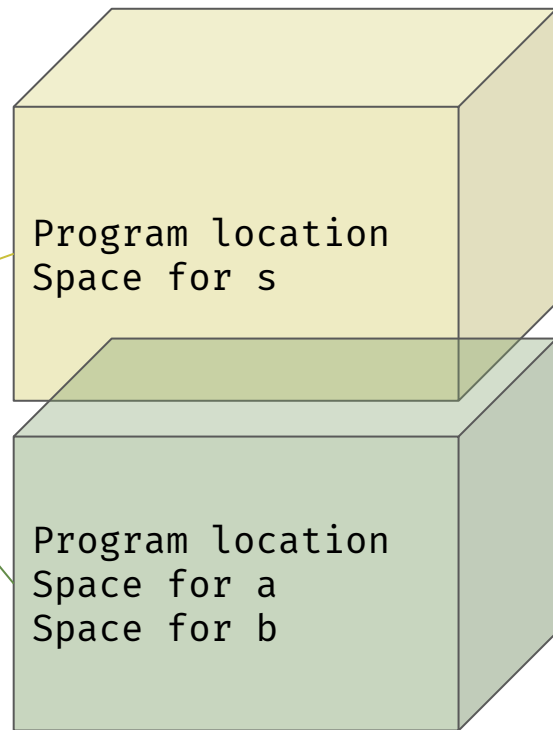
```
    c = a + b
```

```
    return c
```

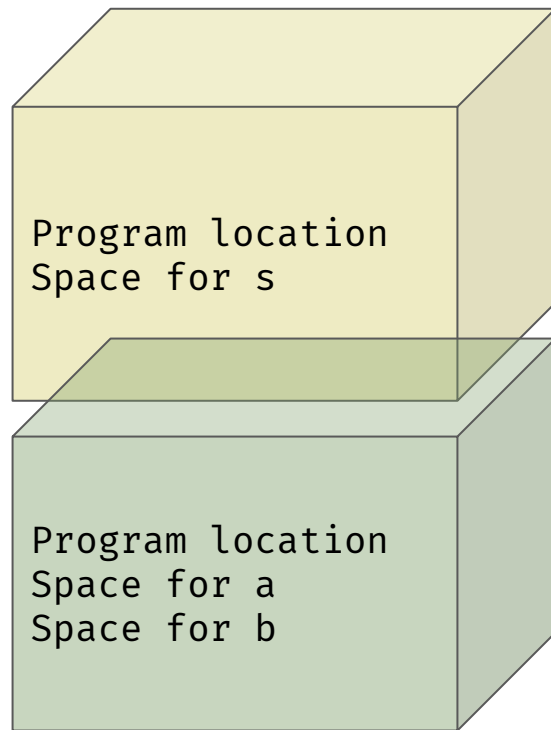
But we can't do  
this, can we?

```
s = add(2, 3)
```

```
print(s)
```



# Getting framed

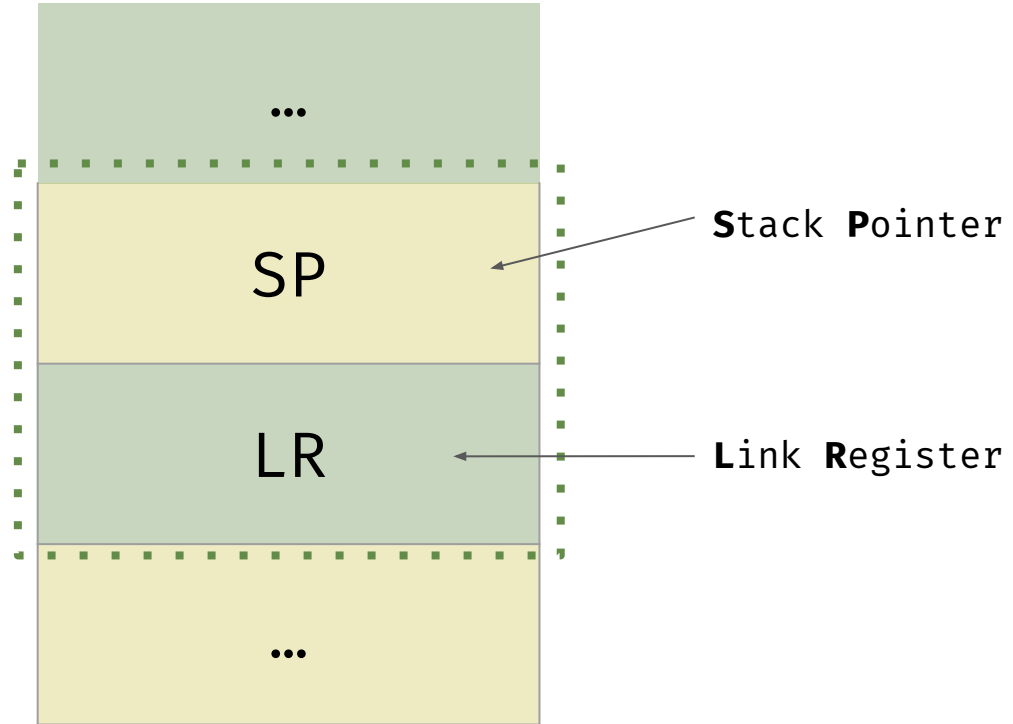
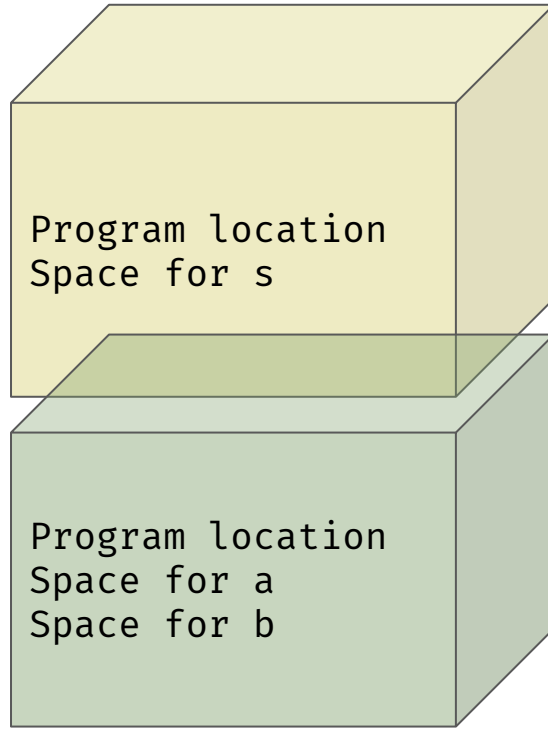


Back to here?

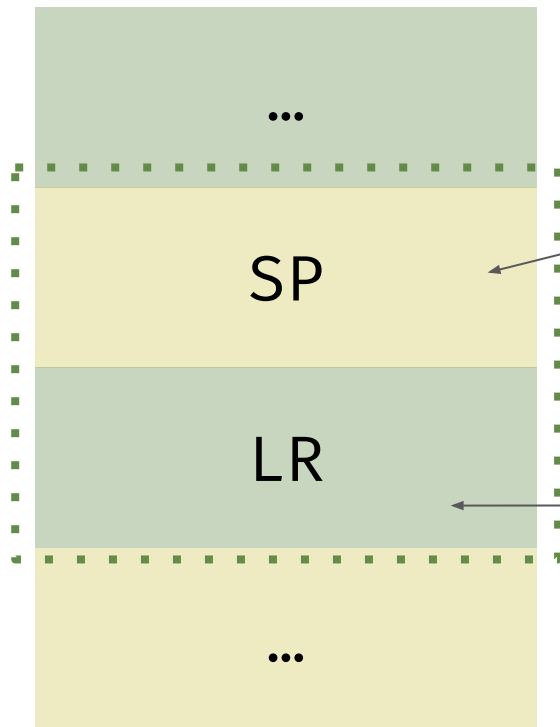


How do we get  
from here...

# Getting framed



# Getting framed



Seems like the CARDIAC had something like this...



**Stack Pointer**

**Link Register:** stores memory location of the previous point in memory that we branched from.



## Going out on a limb...

BL	<b>B</b> branch and <b>L</b> ink	Jump to a label; store location in <b>LR</b>
BX	<b>B</b> branch and e <b>X</b> ecute	Jump back to a memory location and continue to execute

# Going out on a limb...

BL

LABEL

BX

$R_D$

Has to be a memory  
location stored in  
a register...

But can't POP {LR}

# Getting shifty

Logical Shift Right



LSR             $R_D, R_X, \#_{BITS}$

LSR        R5, R5, #24

~~0xff000000~~ → ~~0x000000ff~~

LSL             $R_D, R_X, \#_{BITS}$



Logical Shift Left

LSL        R5, R5, #24

0x000000ff → 0xff000000